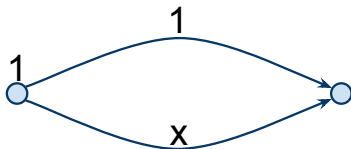


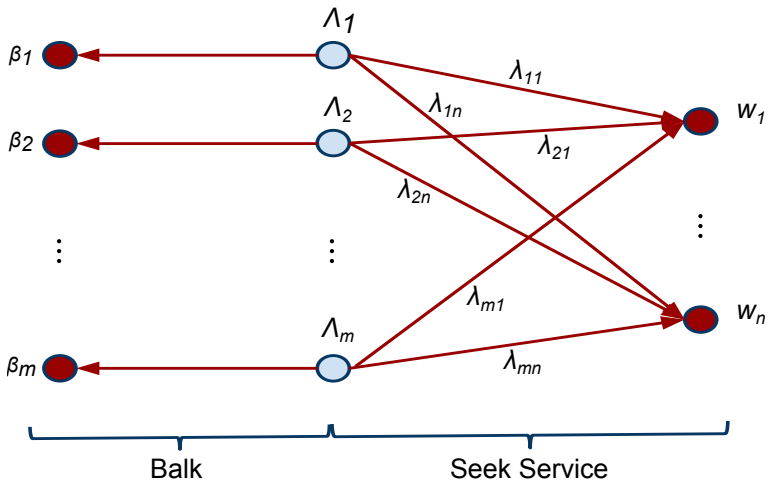
$$\begin{pmatrix} (2, 2) & (5, 0) \\ (0, 5) & (4, 4) \end{pmatrix}$$

- ▶ $k = 1$
- ▶ $\mathcal{P}_1 = \{1, 2\}$
- ▶ $c_1 = 1$ and $c_2 = x$
- ▶ $r = 1$

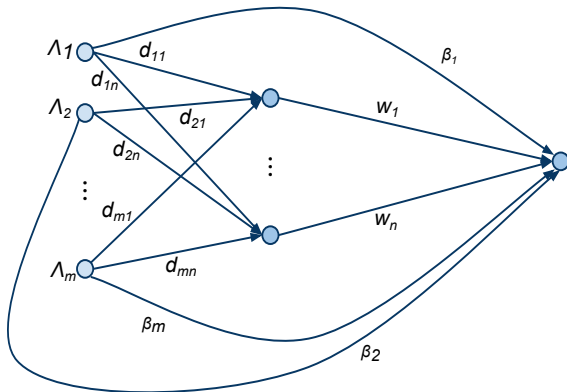


The Nash flow minimises:

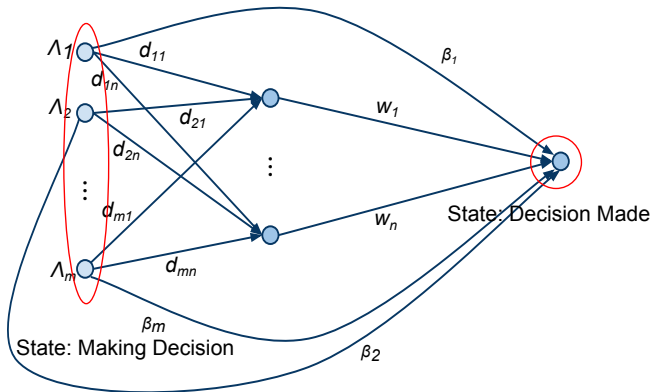
$$\begin{aligned}
 \Phi(y, 1-y) &= \sum_{e=1}^2 \int_0^{f_e} c_e(x) dx = \int_0^y 1 dx + \int_0^{1-y} x dx \\
 &= y + \frac{(1-y)^2}{2} = \frac{1}{2} + \frac{y^2}{2} \\
 &\Rightarrow \tilde{f} = (0, 1)
 \end{aligned}$$



- $k = m$
- $|\mathcal{P}_i| = n + 1$
- $r_i = \Lambda_i$



- $k = m$
- $|\mathcal{P}_i| = n + 1$
- $r_i = \Lambda_i$



Theorem Assuming $\sum_{i=1}^m \Lambda_i < \sum_{j=1}^n c_j \mu_j$ we have:

$$\lim_{\beta_i \rightarrow \infty} PoA(\beta) < \infty \text{ for all } i \in [m]$$

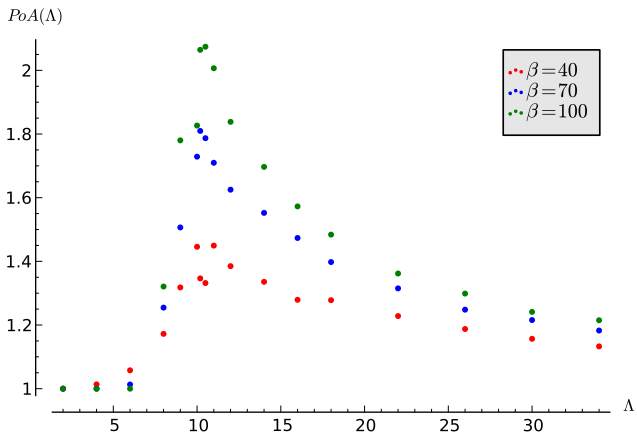
The price of anarchy increases with worth of service, up to a point.

Proof.

- ▶ $\lim_{\beta_i \rightarrow \infty} \lambda^* = k^*$ and $\lim_{\beta_i \rightarrow \infty} \tilde{\lambda} = \tilde{k}$
- ▶ As $\beta_i \rightarrow \infty$:

$$\sum_{i=1}^m \Lambda_i = \sum_{i=1}^m \sum_{j=1}^n \lambda_{ij}^* = \sum_{i=1}^m \sum_{j=1}^n \tilde{\lambda}_{ij}$$

- ▶ $PoA(\beta) < \infty$

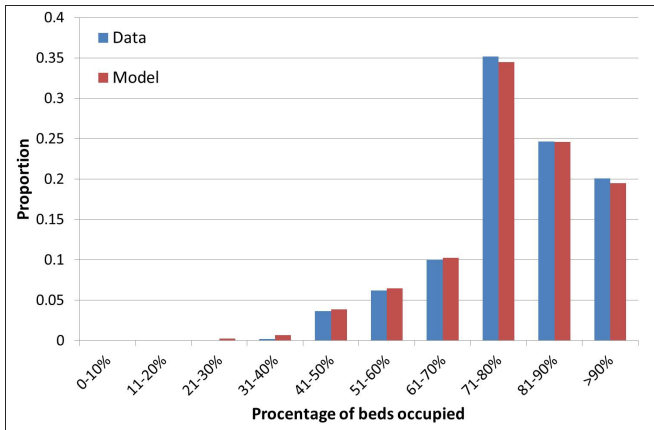


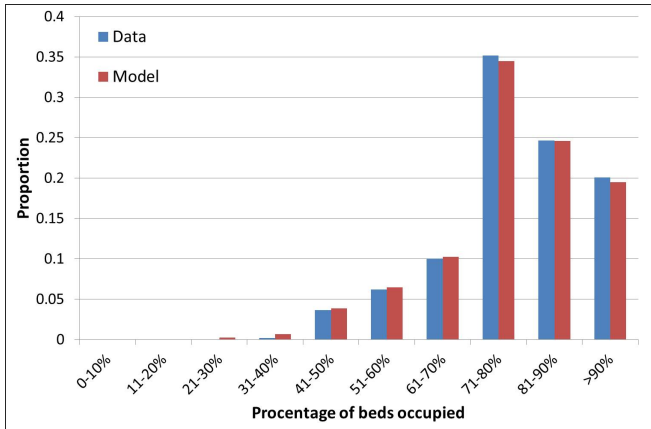
Price of Anarchy in Public Services *EJORS*, 2013.

What about the controllers?

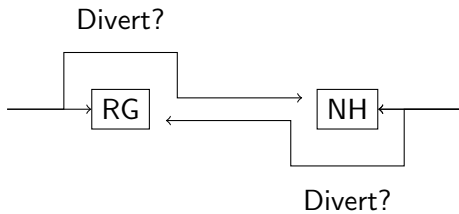
What about the controllers?

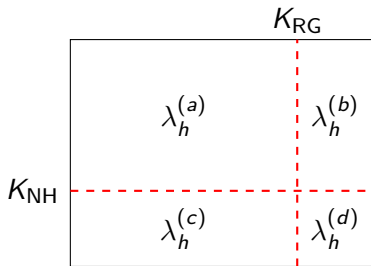
S. Deo and I. Gurvich. **Centralized vs. Decentralized Ambulance Diversion: A Network Perspective.** *Management Science*, May 2011.

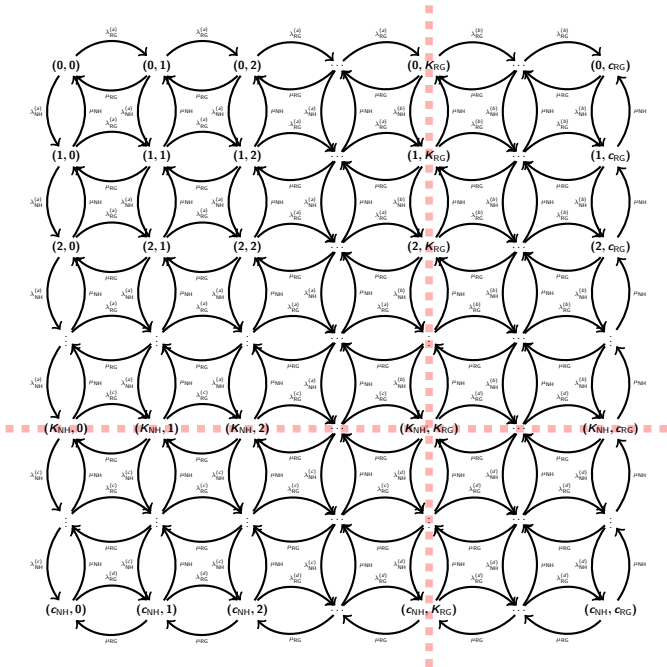


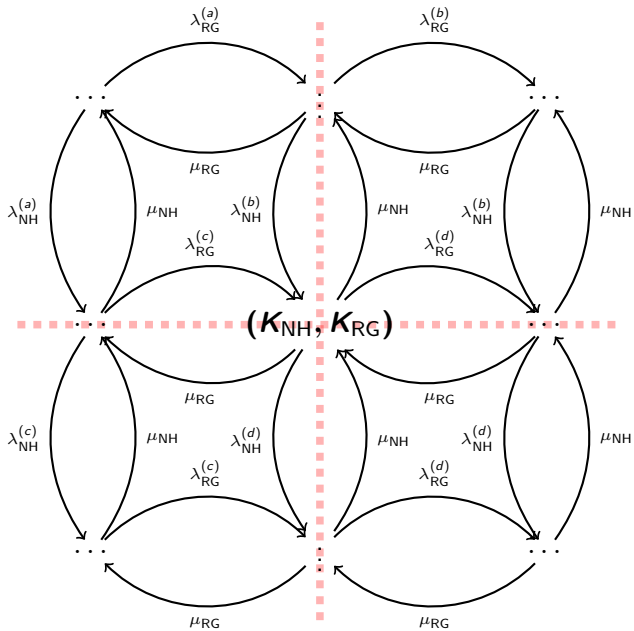


Mathematical modelling of patient flows to predict critical care capacity required following the merger of two District General Hospitals into one., *Submitted to Anaesthesia*







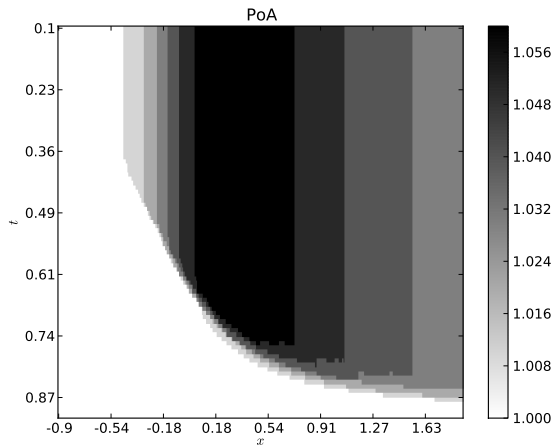


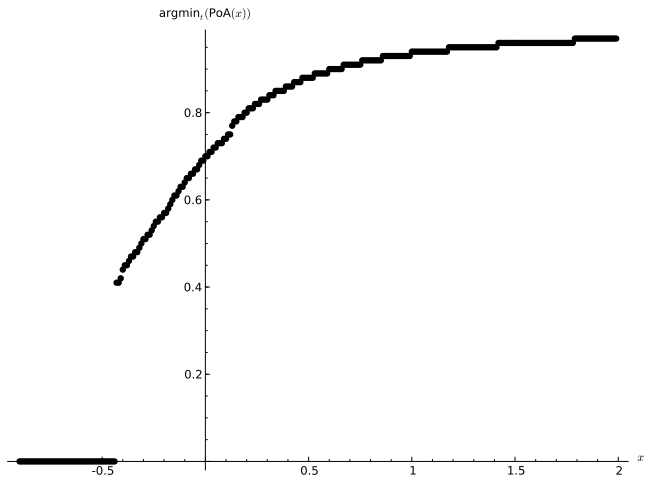
$$A = \begin{pmatrix} (U_{\text{NH}}(1, 1) - t)^2 & \dots & (U_{\text{NH}}(1, c_{\text{RG}}) - t)^2 \\ (U_{\text{NH}}(2, 1) - t)^2 & \dots & (U_{\text{NH}}(2, c_{\text{RG}}) - t)^2 \\ \vdots & \ddots & \vdots \\ (U_{\text{NH}}(c_{\text{NH}}, 1) - t)^2 & \dots & (U_{\text{NH}}(c_{\text{NH}}, c_{\text{RG}}) - t)^2 \end{pmatrix}$$

$$B = \begin{pmatrix} (U_{\text{RG}}(1, 1) - t)^2 & \dots & (U_{\text{RG}}(1, c_{\text{RG}}) - t)^2 \\ (U_{\text{RG}}(2, 1) - t)^2 & \dots & (U_{\text{RG}}(2, c_{\text{RG}}) - t)^2 \\ \vdots & \ddots & \vdots \\ (U_{\text{RG}}(c_{\text{RG}}, 1) - t)^2 & \dots & (U_{\text{RG}}(c_{\text{RG}}, c_{\text{RG}}) - t)^2 \end{pmatrix}$$

Theorem.

Let $f_h(k) : [1, c_{\bar{h}}] \rightarrow [1, c_h]$ be the best response of player $h \in \{\text{NH}, \text{RG}\}$ to the diversion threshold of $\bar{h} \neq h$ ($\bar{h} \in \{\text{NH}, \text{RG}\}$). If $f_h(k)$ is a non-decreasing function in k then the game has at least one Nash Equilibrium in Pure Strategies.

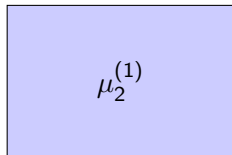




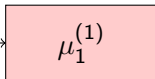
Measuring the Price of Anarchy in Critical Care Unit Interactions, *Submitted to OMEGA*



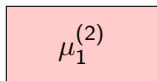
$$C^{(1)} - c^{(1)}$$



$$1 - p^{(1)}$$

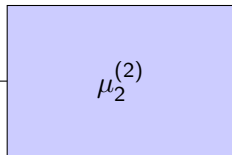


$$c^{(1)}$$

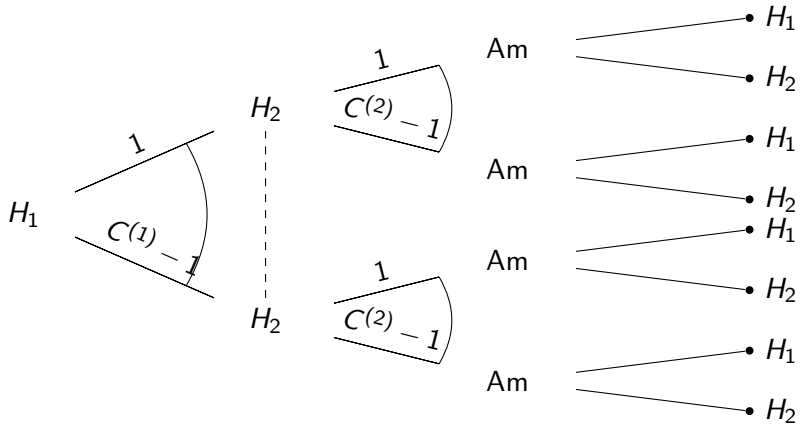


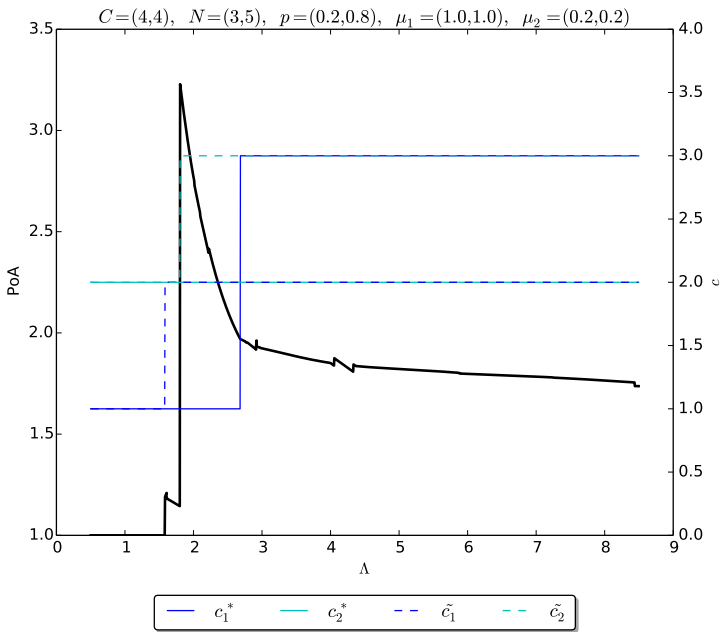
$$c^{(2)}$$

$$1 - p^{(2)}$$



$$C^{(1)} - c^{(1)}$$





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vincent-knight.com/Talks