# Principles of Programming Languages Lecture 09 Subtyping Bounded Universal and Existential Types

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1 Subtyping Quantified Types

## Recall Subtyping for Function Types

$$\frac{\vdash T_1 <: S_1 \qquad \vdash S_2 <: T_2}{\vdash S_1 \to S_2 <: T_1 \to T_2}$$

- contravariant in argument position: flip types
- **covariant** in result position: keep order of types

## Bounded Universal Types

#### Syntax

$$T ::= \cdots \mid \forall X <: T.T$$

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New: May have subtyping assumptions for type variables X

$$\frac{X<: T\in \Gamma}{\Gamma\vdash X<: T}$$

$$T ::= \cdots \mid \forall X <: T.T$$

### New: May have subtyping assumptions for type variables X

$$\frac{X <: T \in \Gamma}{\Gamma \vdash X <: T}$$

#### Consider $\forall$ as a function

$$\frac{\Gamma \vdash T_1 <: S_1 \qquad \Gamma, X <: T_1 \vdash S_2 <: T_2}{\Gamma \vdash \forall X <: S_1.S_2 <: \forall X <: T_1.T_2}$$

## Syntax

$$T ::= \cdots \mid \exists X <: T.T$$

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$$T ::= \cdots \mid \exists X <: T.T$$

#### Covariant in bound of exists

$$\frac{\Gamma \vdash S_1 <: T_1 \qquad \Gamma, X <: S_1 \vdash S_2 <: T_2}{\Gamma \vdash \exists X <: S_1.S_2 <: \exists X <: T_1.T_2}$$

### Source



■ Benjamin Pierce. Types in Programming Languages. MIT Press 2002.