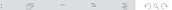
Software Foundations Introduction

Albert-Ludwigs-Universität Freiburg

Luminous Fennell

2012-10-24





- Which semester?
- Experience:
 - Logic courses, Th. comp. science
 - Verification, Hoare Calculus
 - Functional Programming
 - Formal Systems
- Coq:
 - Proof Assistant
 - Programming language
 - (show live)

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Software Foundations (Benjamin Pierce et al.)

- Self study course
- Chapters: Commented source code with exercises
- http://www.cis.upenn.edu/~bcpierce/sf/ Version 2012-7-25
- Work the chapters at home
- Meeting once a week for questions/discussion
- Exercises may be submitted
- Course Homepage: http://proglang.informatik.unifreiburg.de/teaching/softwarefoundations/2012ws/

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- Chapter Exercises
 - Edited versions on course website
 - (* EXPECTED *) Exercise is strongly recommended
 - (* NO SOLUTION *) Solution on demand
 - Sample solution 1-2 weeks later
- Graded Exercises
 - 4 graded exercises, distributed throughout the semester
 (2 before, 2 after christmas)
 - Each 25% of final grade
 - 2 weeks time to submit

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Departement of Programming Languages Building 079, Rooms 00-013 and 00-014

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- (Manuel Geffken, Robert Jakob, Matthias Keil)

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http://coq.inria.fr/

```
Tactic Notation "aevalR cases" tactic(first) ident(c) :=
 [ Case aux c "E ANum" | Case aux c "E APlus"
 I Case aux c "E AMinus" | Case aux c "E AMult" |.
Theorem aeval iff aevalR : forall a n.
Proof
  aevalR cases (induction H) SCase: simpl.
  SCase "E ANum".
  SCase "E APlus".
   rewrite IHaevalR1. rewrite IHaevalR2. reflexivity.
   rewrite IHaevalR1, rewrite IHaevalR2,
  SCase "E AMult".
    Imp.v
               35% (685,29) <N> Git-master (Coq Script(4) Holes)--1:55PM------
 IHaevalR1 : aeval e1 = n1
 IHaevalR2 : aeval e2 = n2
subgoal 2 is:
```

Stating and Proving formal theorems

Informal

"Clearly, zero is the smallest natural number!"

Formal (Coq)

```
Inductive nat : Set :=
| 0 : nat
| S : nat -> nat.
Inductive le : nat -> nat -> Prop :=
| le_n : forall n : nat, le n n
| le_S : forall ni n2 : nat,
le ni n2 -> le ni (S n2).
```

```
Theorem le_nat_total: forall n : nat, le 0 n.

Proof. intros n. induction n as [| n'].

(* Case n = 0 *)

apply le_n.

(* Case n = S n' *)

apply le_S. apply IHn'.

Qed.

(* Or with automation *)

Theorem le_nat_total: forall n : nat, le 0 n.

Proof. intros n; induction n as [| n']; auto.

Qed.
```

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Formalization of Programming Languages

While Programs

$$e ::= k \mid \mathtt{True} \mid \mathtt{False} \mid x \mid e + e \mid e - e$$

$$\mid x := e \mid e; e \mid \mathtt{IF} \ e \ \mathtt{THEN} \ e \ \mathtt{ELSE} \ e \mid \mathtt{WHILE} \ e \ \mathtt{DO} \ e$$

Lambda Calculus

$$e := k \mid \texttt{True} \mid \texttt{False} \mid x \mid \texttt{IF} \ e \ \texttt{THEN} \ e \ \texttt{ELSE} \ e$$
 $\mid \lambda x. \ e \mid e \ e$

- Precise definition of semantics
- Type systems
- Proving properties about programs (e.g. Correctness)

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Implementation of Certified Functional Programs if you're interested

```
Define sortedlist := { 1 : natlist | sorted 1 }.
insertSorted : nat -> sortedlist -> sortedlist := ...
myFancySort : forall 1 : natlist,
              { 1' : natlist | 1' = simpleSort 1 } := ...
```

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