

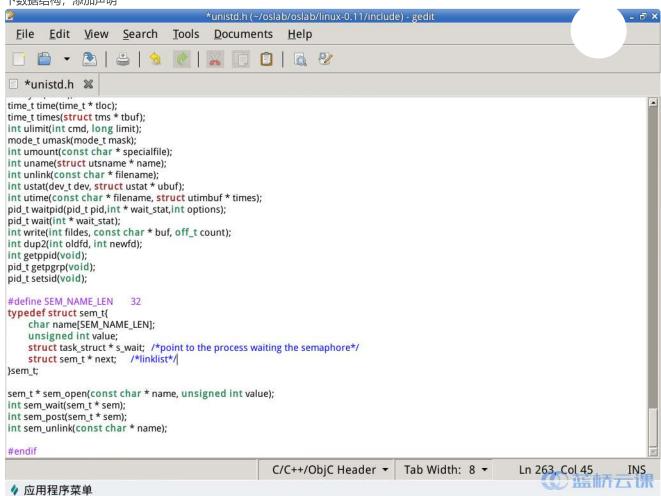
"操作系统原理与实践"实验报告

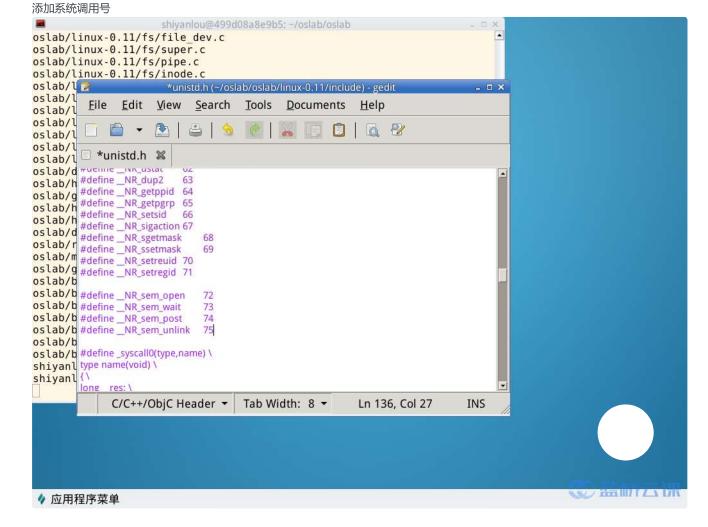
信号量的实现和应用

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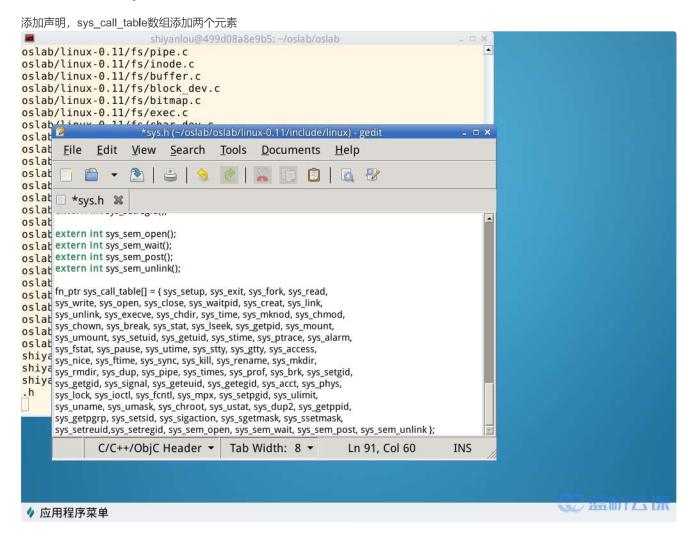
在linux 0.11上实现信号量

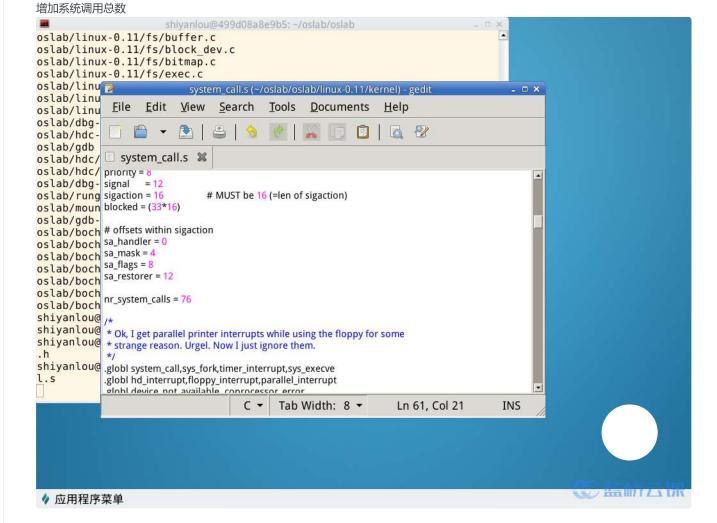
在实现的过程中将运用到系统调用的实验知识,通过此次实验也对前面做了一个比较系统的复习###在include/unistd.h文件中 定义如下数据结构,添加声明





在include/linux/sys.h文件中





在kernel/中实现sem.c

在此文件内实现4个山寨信号量函数。 首先包含了必要的头文件, 然后定义信号量队列的头指针。

```
#include <linux/kernel.h>
#include <asm/system.h>
#include <linux/sched.h>
#include <asm/segment.h>
#include <unistd.h>

sem_t *sem_head = &((sem_t *){"", 0, NULL, NULL}); /* head of the linklist */
```

自己新增的一个函数,负责把信号量的名字从用户栈拷贝到内核栈

下面进入正题,实现四个山寨的信号量函数

```
sem_t *sys_sem_open(const char *name, unsigned int value)
   sem_t *sem_cur, *sem_pre;
   char pname[SEM_NAME_LEN];
   str_u2k(name, pname, SEM_NAME_LEN);
   /*whether the sem existes in the linklist */
   for(sem_pre=sem_head, sem_cur=sem_head->next; sem_cur && strcmp(pname, sem_cur->name);
           sem_pre=sem_cur, sem_cur=sem_cur->next);
   /*if not exists, new a sem */
   if(!sem_cur)
       printk("semaphore %s no found. created a new one. \n", pname);
       sem_cur = (sem_t *)malloc(sizeof(sem_t));
       strcpy(sem_cur->name, pname);
       sem_cur->value = value;
       sem_cur->next = NULL;
       sem_pre->next = sem_cur;
   printk("pid %d opens semaphore %s(value %u) OK. \n", current->pid, pname, sem_cur->value);
   return sem_cur;
```

```
int sys_sem_wait(sem_t *sem)
{
    cli();
    while(sem->value<=0)
        sleep_on(&(sem->s_wait));
    sem->value--;
    sti();
    return 0;
}
```

```
int sys_sem_post(sem_t *sem)
{
    cli();
    sem->value++;
    /* if still have process waiting for sem */
    if(sem->s_wait)
    {
        wake_up(&(sem->s_wait));
        sti();
        return 0;
    }
    sti();
    return -1;
}
```

```
int sys_sem_unlink(const char *name)
{
    sem_t *sem_cur, *sem_pre;
    char pname[SEM_NAME_LEN];
    int i;

    str_u2k(name, pname, SEM_NAME_LEN);

    for(sem_pre=sem_head, sem_cur=sem_head->next; sem_cur && strcmp(pname, sem_cur->name);
        sem_pre=sem_cur, sem_cur=sem_cur->next);

/*if not found, return -1 */
    if(!sem_cur)
        return -1;

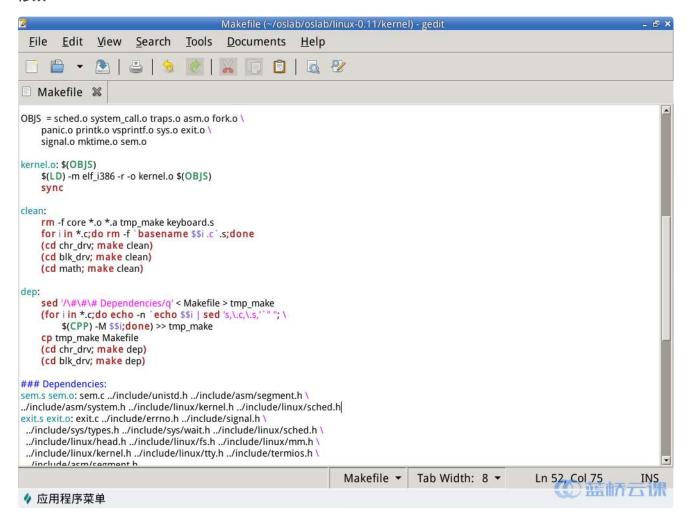
/*if found, free it */
    sem_pre->next = sem_cur->next;
    free(sem_cur);
    printk("unlink semaphore %s OK. \n", pname);
    return 0;
}
```

18.2.4: 之前大括号打错了,以至于fork了新的消费者而没参与到进程同步,现进行修正。此程序有时会出现死锁情况,原因不明。不过输出也是正确的,遇到死锁情况采取放置策略,重新启动系统运行程序即可

```
#define __LIBRARY__
#include <unistd.h>
#include <fcntl.h>
#include <stdio.h>
_syscall2(sem_t *,sem_open,const char *,name,unsigned int,value)
_syscall1(int,sem_wait,sem_t *,sem)
_syscall1(int,sem_post,sem_t *,sem)
_syscall1(int,sem_unlink,const char*,name)
const char *FILENAME = "/usr/root/buffer_file";
const int NR_CONSUMERS = 5;
                                                   /* number of consumers */
const int NR_ITEMS = 50;
                                                /\star the maximum of item \star/
const int BUFFER_SIZE = 10;
sem_t *mutex, *full, *empty;
unsigned int item_pro, item_used;
                                                 /* the NO. of item which is produced or consumed just now */
                                                   /* the I/O pointer*/
int fi, fo;
int main(int argc, char *argv[])
   char *filename;
    int pid;
    int i, flag = 1;
    filename = argc > 1 ? argv[1] : FILENAME;
    /* O_TRUNC means open the file in read-only or write-only, if exists, wipe the file*/
    fi = open(filename, O_CREAT| O_TRUNC| O_WRONLY, 0222);     /* producer write-only */
    fo = open(filename, O_TRUNC| O_RDONLY, 0444);
                                                             /* consumer read-only */
   mutex = sem_open("MUTEX", 1);
   full = sem_open("FULL", 0);
    empty = sem_open("EMPTY", BUFFER_SIZE);
    item_pro = 0;
    if ((pid = fork()))  /* father process is producer */
        printf("pid %d:\tproducer created....\n", pid);
        fflush(stdout);
        while (item_pro < NR_ITEMS) /* while not finish produce the item */</pre>
            sem_wait(empty);
           sem_wait(mutex);
            /\star if the buffer is full, reset the pointer to the head of file \star/
            if(!(item_pro % BUFFER_SIZE))
                lseek(fi, 0, 0);
            write(fi, (char *) &item_pro, sizeof(item_pro));
                                                                   /* produce */
            printf("pid %d:\tproduces item %d\n", pid, item_pro);
            fflush(stdout);
            item_pro++;
            sem_post(mutex);
                                 /* wake up the consumer */
            sem_post(full);
    }
    else
            /* child process is consumer */
        i = NR_CONSUMERS;
        while(i--)
            if(!(pid=fork())) /* new the consumer */
                pid = getpid();
                printf("pid %d:\tconsumer %d created....\n", pid, NR_CONSUMERS-i);
                fflush(stdout);
                while(1)
                    sem_wait(full);
                    sem_wait(mutex);
                    /*when read() finish, return 0, reset the pointer to the head of file */
```

```
if(!read(fo, (char *)&item_used, sizeof(item_used)))
                        lseek(fo, 0, 0);
                        read(fo, (char *)&item_used, sizeof(item_used));
                    printf("pid %d:\tconsumer %d consumes item %d\n", pid, NR_CONSUMERS-i+1, item_used);
                    fflush(stdout);
                    sem_post(mutex);
                    sem_post(empty);
                                      /* wake up the producer */
                    if(item_used == NR_ITEMS){    /* if all items have been consumed */
                        flag = 0;
                        goto OK;
           }
   }
OK:
   while(flag);
   sem_unlink("MUTEX");
   sem_unlink("FULL");
   sem_unlink("EMPTY");
   close(fi):
   close(fo);
    return 0;
```

修改kernel/Makefile

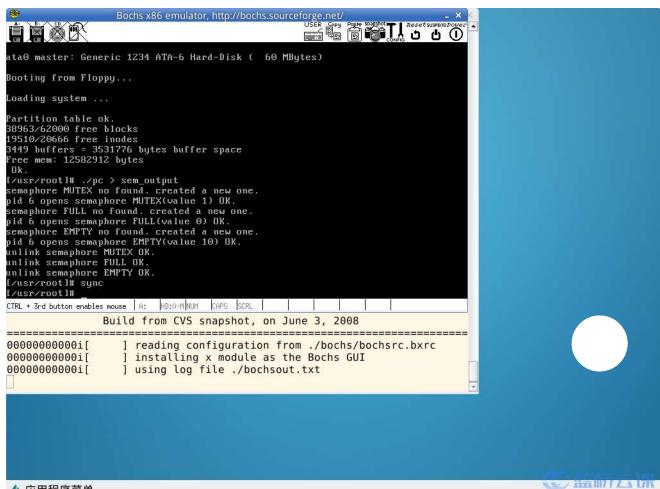


运行程序

手动替换库文件

在oslab下输入sudo ./mount-hdc,将做过修改的库文件和pc.c拷贝进虚拟镜像之中,执行sudo umount hdc取消挂载

```
gcc -o pc pc.c -Wall
./pc > sem_output
sync
```



♦ 应用程序菜单

退出前务必执行sync

在Ubuntu下执行sudo less hdc/usr/root/sem_output, 得到输出 shiyanlou@d0e9574fc318: ~/oslab/oslab pid 7: producer created.... pid 7: produces item 0 pid 7: produces item 1 pid 7: produces item 2 pid 7: produces item 3 pid 7: produces item 4 pid 7: produces item 5 pid 7: produces item 6 pid 7: produces item pid 7: produces item 8 pid 7: produces item 9 pid 12: consumer 5 created.... pid 12: consumer 6 consumes item 0 pid 12: consumer 6 consumes item 1 pid 12: consumer 6 consumes item 2 pid 12: consumer 6 consumes item 3 pid 12: consumer 6 consumes item pid 12: consumer 6 consumes item 5 pid 12: consumer 6 consumes item 6 pid 12: consumer 6 consumes item pid 12: consumer 6 consumes item 8 pid 12: consumer 6 consumes item 9 pid 11: consumer 4 created.... pid 10: consumer 3 created.... pid 9: consumer 2 created.... pid 8: consumer 1 created.... pid 7: produces item 10 pid 7: produces item 11 pid 7: produces item 12 17 LA 1711 ♦ 应用程序菜单 shiyanlou@d0e9574fc318: -/oslab/oslab pid 7: produces item 12 pid 7: produces item 13 pid 7: produces item 14 produces item 15 pid 7: pid 7: produces item 16 pid 7: produces item 17 pid 7: produces item 18 pid 7: produces item 19 consumer 2 consumes item 10 pid 8: pid 8: consumer 2 consumes item 11 pid 8: consumer 2 consumes item 12 pid 8: consumer 2 consumes item 13 pid 8: consumer 2 consumes item 14 pid 8: consumer 2 consumes item 15 pid 8: consumer 2 consumes item 16 pid 8: consumer 2 consumes item 17 pid 8: consumer 2 consumes item 18 pid 8: consumer 2 consumes item 19 pid 7: produces item 20 pid 7: produces item 21 pid 7: produces item 22 pid 7: produces item 23 pid 7: produces item 24 pid 7: produces item 25 pid 7: produces item 26 pid 7: produces item 27 pid 7: produces item 28 pid 7: produces item 29 pid 12: consumer 6 consumes item 20

验 胎剂// 1/1

```
shiyanlou@d0e9574fc318: ~/oslab/oslab
pid 12: consumer 6 consumes item 20
pid 12: consumer 6 consumes item 21
pid 12: consumer 6 consumes item 22
pid 12: consumer 6 consumes item 23
pid 12: consumer 6 consumes item 24
pid 12: consumer 6 consumes item 25
pid 12: consumer 6 consumes item 26
pid 12: consumer 6 consumes item 27
pid 12: consumer 6 consumes item 28
pid 12: consumer 6 consumes item 29
pid 7:
       produces item 30
pid 7:
       produces item 31
pid 7:
       produces item 32
pid
   7:
       produces item 33
pid 7:
       produces item 34
pid 7:
       produces item 35
pid 7:
        produces item 36
pid 7:
       produces item 37
pid 7:
        produces item 38
pid 7:
        produces item 39
pid 8:
        consumer 2 consumes item 30
        consumer 2 consumes item 31
pid 8:
pid 8:
       consumer 2 consumes item 32
pid 8:
       consumer 2 consumes item 33
pid 8:
        consumer 2 consumes item 34
pid 8:
       consumer 2 consumes item 35
pid 8:
        consumer 2 consumes item 36
       consumer 2 consumes item 37
pid 8:
pid 8: consumer 2 consumes item 38
                                                                                     in || 1 / 1 / 1/10

か 应用程序菜单

                    shiyanlou@d0e9574fc318: ~/oslab/oslab
pid 8: consumer 2 consumes item 33
pid 8:
        consumer 2 consumes item 34
       consumer 2 consumes item 35
pid 8:
pid 8: consumer 2 consumes item 36
pid 8:
        consumer 2 consumes item
pid 8: consumer 2 consumes item 38
pid 8: consumer 2 consumes item 39
pid 7:
        produces item 40
pid 7:
       produces item 41
       produces item 42
pid 7:
pid 7:
       produces item 43
pid 7:
       produces item 44
pid
   7:
       produces item 45
pid 7:
       produces item 46
       produces item 47
pid 7:
pid 7:
        produces item 48
pid 7:
       produces item 49
pid 12: consumer 6 consumes item 40
pid 12: consumer 6 consumes item 41
pid 12: consumer 6 consumes item 42
pid 12: consumer 6 consumes item 43
pid 12: consumer 6 consumes item 44
pid 12: consumer 6 consumes item 45
pid 12: consumer 6 consumes item 46
pid 12: consumer 6 consumes item 47
pid 12: consumer 6 consumes item 48
pid 12: consumer 6 consumes item 49
pid 7: produces item 50
pid 8:
        consumer 2 consumes item 50
(END)
                                                                                     W B面III71口 [7]
♦ 应用程序菜单
```

与实验所要求的输出基本一致

实验报告问题

执行效果变化很大,打印出的消费序列完全是乱的。因为没有信号量对临界区进行保护,加上进程调度是不确定的,导致生产者和消费者们的动作都无法做到互斥。也就是说,在临界区中完成一个动作之前,它们都有可能被另一个进程抢占,导致在缓冲区中的数据是混乱的。比如,消费者C正准备在缓冲区的起始处取出产品0时,就被生产者P抢占,而P已经完成了一轮的生产,也处于缓冲区的起始处,它往这里写入一个1。如果有轮到C执行。那么C取出的产品就是1,而不是0。

2.实验的设计者在第一次编写生产者——消费者程序的时候,是这么做的:

这样可行吗?如果可行,那么它和标准解法在执行效果上会有什么不同?如果不可行,那么它有什么问题使它不可行

- 1. 假设Producer刚生产完一件商品,释放了Mutex,Mutex为1,此时缓存区满了,Empty为0;
- 2. 然后OS执行调度,若被Producer拿到CPU,它拿到Mutex,使Mutex为0,而Empty为0,Producer让出CPU,等待Consumer执行V(Empty);
- 3. 而Consumer拿到CPU后,却要等待Producer执行V(Mutex);
- 4. 两者相互持有对方需要的资源,造成死锁。



Έ.



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LOU2979148015 L19 回复 LOU2979148015 L19

直接在这回复也ok的

2019-05-21 22:51:48

⊕ 回复



LOU2979148015 L19

你好!我有一处不明,在pc.c中你好像只对item_pro变量进行了控制,而未对item_used进行控制。没想通是为什么。还望告知一下!谢谢你!qq2280343053

2019-05-21 22:51:26

⊕ 回复