# Proof of Convergence for Kernelized M3L

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We now describe the convergence of the algorithm for kernelized M3L [2]. It closely follows the proof of convergence of SMO [3].

#### 1 Notation

We denote vectors in bold small letters, for example  $\mathbf{v}$ . If  $\mathbf{v}$  is a vector of dimension d, then  $v_k, k \in \{1, \ldots, d\}$  is the kth component of  $\mathbf{v}$ , and  $\mathbf{v}_I, I \subseteq \{1, \ldots, d\}$  denotes the vector with components  $v_k, k \in I$  (with the  $v_k$ 's arranged in the same order as in  $\mathbf{v}$ ). Similarly, matrices will be written in bold capital letters, for example  $\mathbf{A}$ . If  $\mathbf{A}$  is an  $m \times n$  matrix, then  $A_{ij}$  represents the ijth entry of  $\mathbf{A}$ , and  $\mathbf{A}_{IJ}$  represents the matrix with entries  $A_{ij}, i \in I, j \in J$ .

A sequence is denoted as  $\{a^n\}$ , and  $a^n$  is the nth element of this sequence. If  $\hat{a}$  is a limit point of the sequence, we write  $a^n \to \hat{a}$ .

## 2 The Optimization Problem

The dual that we are trying to solve is:

$$\max_{\alpha} 2 \sum_{l=1}^{L} \alpha_{l}^{t} \mathbf{1} - 2 \sum_{l=1}^{L} \sum_{k=1}^{L} R_{lk} \alpha_{l}^{t} \mathbf{Y}_{l} \mathbf{K} \mathbf{Y}_{k} \alpha_{k}(1)$$

s.t

$$0 < \alpha < C1$$

where  $\alpha_l = [\alpha_{1l}, \dots \alpha_{Nl}]$ ,  $\mathbf{Y}_l = diag([y_{1l}, \dots y_{Nl}])$  and  $\mathbf{K} = \phi(\mathbf{X})^t \phi(\mathbf{X})$ . Making the substitution  $\theta_{il} = 2y_{il}\alpha_{il}$ , we get the following optimization problem:

$$\max_{\theta} \sum_{i=1}^{N} \sum_{l=1}^{L} \theta_{il} y_{il} - \frac{1}{2} \sum_{i=1}^{N} \sum_{j=1}^{N} \sum_{k=1}^{L} \sum_{l=1}^{L} \theta_{ik} \theta_{jl} R_{kl} K_{ij}(2)$$

s.t

$$0 \le \theta_{ik} \le 2C$$
  $\forall i, k \text{ s.t } y_{ik} > 0$ 

$$-2C \le \theta_{ik} \le 0 \quad \forall i, k \text{ s.t } y_{ik} < 0$$

This can be written as the following optimization problem:

Problem:

$$\max_{\boldsymbol{\theta}} f(\boldsymbol{\theta}) = -\frac{1}{2} \boldsymbol{\theta}^t \mathbf{Q} \boldsymbol{\theta} + \mathbf{y}^t \boldsymbol{\theta}$$
 (3)

s.1

$$1 \le \theta \le u$$

Here the vector  $\boldsymbol{\theta} = [\theta_{11} \dots \theta_{1L}, \theta_{21}, \dots \theta_{NL}]^t$ ,  $\mathbf{y} = [\mathbf{y}_1^t \dots \mathbf{y}_N^t]^t$  and  $\mathbf{Q} = \mathbf{K} \otimes \mathbf{R}$  where  $\otimes$  is the Kronecker product. I and  $\mathbf{u}$  are NL- dimensional vectors with entries:

$$l_{ik} = \begin{cases} 0 & \text{if } y_{ik} > 0\\ -2C & \text{if } y_{ik} < 0 \end{cases}$$
 (4)

$$u_{ik} = \begin{cases} 2C & \text{if } y_{ik} > 0\\ 0 & \text{if } y_{ik} < 0 \end{cases}$$
 (5)

#### 3 The algorithm

We will prove the convergence result for the algorithm given in Algorithm 1. We start with  $\boldsymbol{\theta}^1 = \mathbf{0}$ , and generate a sequence of vectors  $\boldsymbol{\theta}^n$ , where  $\boldsymbol{\theta}^n$  is the vector after the (n-1)th iteration. The gradient  $\nabla f(\boldsymbol{\theta}^n)$ , denoted by  $\mathbf{g}^n$ , is stored. The projected gradient  $\nabla^P f(\boldsymbol{\theta}^n)$ , denoted by  $\tilde{\mathbf{g}}^n$ , is computed from the gradient on the fly.  $\tau > 0$  is a user-defined parameter that determines when we stop. The algorithm terminates when all projected gradients are less than  $\tau$ .

We assume that  $\mathbf{R}$  and  $\mathbf{K}$  are both positive definite matrices. The eigenvalues of  $\mathbf{Q}$  are then  $\lambda_i \mu_j$ , where  $\lambda_i$  are the eigenvalues of  $\mathbf{K}$  and  $\mu_j$  are the eigenvalues of  $\mathbf{R}$ . Because all eigenvalues of both  $\mathbf{R}$  and  $\mathbf{K}$  are positive, so are the eigenvalues of  $\mathbf{Q}$  and thus  $\mathbf{Q}$  is positive definite. Therefore, using Lemma 8 (refer to Appendix A), we have that  $\boldsymbol{\theta}$  is optimal iff  $\boldsymbol{\theta}$  is feasible and  $\nabla^P f(\boldsymbol{\theta}) = \mathbf{0}$ .

Our algorithm will not reach an exact solution, and therefore the KKT conditions will not be satisfied exactly. Call a (feasible) solution  $\boldsymbol{\theta}$   $\tau$ -optimal if the KKT conditions are satisfied to a precision of  $\tau$ , i.e,  $|\nabla_{ik}^P f(\boldsymbol{\theta})| < \tau$  for each i, k. Then our algorithm terminates when it reaches a  $\tau$ -optimal solution. Lemma 10 in Appendix A shows that by making  $\tau$  arbitrarily close to 0, we can make our solution arbitrarily close to the optimum. In the following sections we will show that our algorithm will reach the  $\tau$ -optimal solution in a finite number of steps.

## 4 Convergence

In this section we prove that the sequence of vectors  $\boldsymbol{\theta}^n$  converges.

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Algorithm 1 Our algorithm
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\theta_{ik} \leftarrow 0 \quad \forall i, k

repeat

Pick i, k such that \tilde{g}_{ik} \mid \geq \tau

if \exists j such that K_{ij} \neq \sqrt{K_{ii}K_{jj}} and \tilde{g}_{jk} \neq 0

then

Pick one such j.

Optimize w.r.t \theta_{ik} and \theta_{jk}

else

Optimize w.r.t \theta_{ik}

end if.

Update g_{ik} \quad \forall i, k

until \mid \tilde{g}_{ik} \mid < \tau \quad \forall i, k
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Note the following:

- In each iteration of the algorithm, we optimize over a set of variables, which may either be a single variable  $\theta_{ik}$  or a pair of variables  $\{\theta_{ik}, \theta_{jl}\}.$
- The projected gradient of all the chosen variables is non zero at the start of the iteration.
- At least one of the chosen variables has projected gradient with magnitude greater than
  τ.

Consider the *n*th iteration. Denote by *B* the set of indices of the variables chosen:  $B = \{(i,k)\}$  or  $B = \{(i,k),(j,k)\}$ . Without loss of generality, reorder variables so that the variables in *B* occupy the first few indices. In the *n*th iteration, we optimize f over the variables in *B* keeping the rest of the variables constant. Thus we have to maximize  $h(\Delta_B) = f(\theta^n + [\Delta_B^t, 0]^t) - f(\theta^n)$ . This amounts to solving the optimization

problem:

$$\max_{\mathbf{\Delta}_{B}} h(\mathbf{\Delta}_{B}) = -\frac{1}{2} \mathbf{\Delta}_{B}^{t} \mathbf{Q}_{BB} \mathbf{\Delta}_{B} - \mathbf{\Delta}_{B}^{t} (\mathbf{Q} \boldsymbol{\theta}^{n})_{B} \qquad \qquad \tilde{\mathbf{g}}_{B}^{n+1} = \nabla^{P} h(\mathbf{\Delta}_{B}^{*}) \qquad (1)$$

$$+ \mathbf{y}_{B}^{t} \mathbf{\Delta}_{B} \qquad (6) \quad \text{Using (7)}$$

s.t

$$l_B - \theta_B \le \Delta_B \le u_B - \theta_B$$

Assuming that  $R_{kk} > 0 \ \forall k \text{ and } K_{ii} > 0 \ \forall i, \mathbf{Q}_{BB}$ is positive definite. This can be seen as follows. B has either one or two elements. In the first case,  $B = \{(i,k)\}$  and  $\mathbf{Q}_{BB} = R_{kk}K_{ii} > 0$  and hence  $\mathbf{Q}_{BB}$  is positive definite. In the second case, suppose  $B = \{(i, k), (j, k)\}$ . The matrix  $\mathbf{Q}_{BB}$  is given by  $\begin{bmatrix} K_{ii}R_{kk} & K_{ij}R_{kk} \\ K_{ij}R_{kk} & K_{jj}R_{kk} \end{bmatrix}$ . Because  $K_{ij}^2 \neq K_{ii}K_{ij}$  (refer Algorithm 1) we have that  $\mathbf{Q}_{BB}$  must be positive definite.

Hence by Lemma 8 in Appendix A,  $\Delta_B^*$  optimizes (6) iff it is feasible and

$$\nabla^P h(\mathbf{\Delta}_B^*) = \mathbf{0} \tag{7}$$

Then we have that  $\boldsymbol{\theta}^{n+1} = \boldsymbol{\theta}^n + \boldsymbol{\Delta}^*$ , where  $\Delta^* = [\Delta_B^{*t}, 0^t]^t$ . Now note that since  $\mathbf{g}_B^n =$  $-(\mathbf{Q}\boldsymbol{\theta}^n)_B + \mathbf{y}_B$ 

$$h(\mathbf{\Delta}_B) = -\frac{1}{2}\mathbf{\Delta}_B^t \mathbf{Q}_{BB}\mathbf{\Delta}_B + \mathbf{\Delta}_B^t \mathbf{g}_B^n \qquad (8)$$

$$\Rightarrow \nabla h(\mathbf{\Delta}_B) = -\mathbf{Q}_{BB}\mathbf{\Delta}_B + \mathbf{g}_B^n \qquad (9)$$

Also.

$$\mathbf{g}_{B}^{n+1} = -(\mathbf{Q}\boldsymbol{\theta}^{n+1})_{B} + \mathbf{y}_{B}$$

$$= -(\mathbf{Q}(\boldsymbol{\theta}^{n} + [\boldsymbol{\Delta}_{B}^{*t}, \mathbf{0}^{t}]^{t}))_{B} + \mathbf{y}_{B}$$
(10)

$$\Rightarrow \mathbf{g}_B^{n+1} = \mathbf{g}_B^n - \mathbf{Q}_{BB} \mathbf{\Delta}_B^* = \nabla h(\mathbf{\Delta}_B^*) \quad (11)$$

Then (11) means that:

$$\tilde{\mathbf{g}}_B^{n+1} = \nabla^P h(\mathbf{\Delta}_B^*) \tag{12}$$

$$\tilde{\mathbf{g}}_{B}^{n+1} = \nabla^{P} h(\mathbf{\Delta}_{B}^{*}) = \mathbf{0}$$
 (13)

This leads us to the following lemma:

**Lemma 1.** Let  $\theta^n$  be the solution at the start of the nth iteration. Let B be the set of indices of the variables over which we optimize. Let the updated solution be  $\theta^{n+1}$ . Then

1. 
$$\tilde{\mathbf{g}}_{B}^{n+1} = \mathbf{0}$$

2. 
$$\boldsymbol{\theta}^{n+1} \neq \boldsymbol{\theta}^n$$

3. If 
$$l_{jk} < \theta_{jk}^{n+1} < u_{jk}$$
 then  $g_{jk}^{n+1} = 0$   $\forall (j,k) \in B$ 

*Proof.* 1. This follows directly from (13).

- 2. If  $\boldsymbol{\theta}^{n+1} = \boldsymbol{\theta}^n$ , then  $\boldsymbol{\Delta}_B^* = \mathbf{0}$  and so, from (11),  $\mathbf{g}_B^{n+1} = \nabla h(\mathbf{0}) = \mathbf{g}_B^n$ . This means that from (13)  $\tilde{\mathbf{g}}_B^n = \tilde{\mathbf{g}}_B^{n+1} = \mathbf{0}$ . But this is a contradiction since we required that all variables in the chosen set have non zero projected gradient before the start of the iteration.
- 3. Since the final projected gradients are 0 for all variables in the chosen set (from (13)), if  $l_{jk} < \theta_{jk}^{n+1} < u_{jk} \text{ then } g_{jk}^{n+1} = 0 \ \forall (j,k) \in B$

**Lemma 2.** In the same setup as the previous lemma,  $f(\boldsymbol{\theta}^{n+1}) - f(\boldsymbol{\theta}^n) \ge \sigma \|\boldsymbol{\theta}^{n+1} - \boldsymbol{\theta}^n\|^2$ , for some fixed  $\sigma > 0$ .

Proof.

$$f(\boldsymbol{\theta}^{n+1}) - f(\boldsymbol{\theta}^n) = h(\boldsymbol{\Delta}_B^*)$$
  
=  $-\frac{1}{2}\boldsymbol{\Delta}_B^{*t}\mathbf{Q}_{BB}\boldsymbol{\Delta}_B^* + \boldsymbol{\Delta}_B^{*t}\mathbf{g}_B^n(14)$ 

where  $\Delta_B^*$  is the optimum solution of Problem (6). Now, note that since  $\Delta_B^*$  is feasible and **0** is feasible, we have, using Lemma 9 in Appendix A,

$$\Delta_B^{*t} \nabla h(\Delta_B^*) \ge 0 \tag{15}$$

$$\Rightarrow -\mathbf{\Delta}_{B}^{*t}\mathbf{Q}_{BB}\mathbf{\Delta}_{B}^{*} + \mathbf{\Delta}_{B}^{*t}\mathbf{g}_{B}^{n} \ge 0 \qquad (16)$$

$$\Rightarrow \mathbf{\Delta}_{B}^{*t} \mathbf{Q}_{BB} \mathbf{\Delta}_{B}^{*} \le \mathbf{\Delta}_{B}^{*t} \mathbf{g}_{B}^{n} \tag{17}$$

This gives us that

$$-\frac{1}{2}\boldsymbol{\Delta}_{B}^{*t}\mathbf{Q}_{BB}\boldsymbol{\Delta}_{B}^{*}+\mathbf{g}_{B}^{nt}\boldsymbol{\Delta}_{B}^{*}\geq\frac{1}{2}\boldsymbol{\Delta}_{B}^{*t}\mathbf{Q}_{BB}\boldsymbol{\Delta}_{B}^{*}$$
(18)

$$\Rightarrow f(\boldsymbol{\theta}^{n+1}) - f(\boldsymbol{\theta}^n) \ge \frac{1}{2} \boldsymbol{\Delta}_B^{*t} \mathbf{Q}_{BB} \boldsymbol{\Delta}_B^* \quad (19)$$

$$\Rightarrow f(\boldsymbol{\theta}^{n+1}) - f(\boldsymbol{\theta}^n) \ge \nu_B \frac{1}{2} \boldsymbol{\Delta}_B^{*t} \boldsymbol{\Delta}_B^* \quad (20)$$

where  $\nu_B$  is the minimum eigenvalue of the matrix  $\mathbf{Q}_{BB}$ . Since  $\mathbf{Q}_{BB}$  is positive definite always, this value is always greater than zero, and bounded below by the minimum eigenvalue among all  $2 \times 2$  positive definite submatrices of  $\mathbf{Q}$ . Thus

$$f(\boldsymbol{\theta}^{n+1}) - f(\boldsymbol{\theta}^n) \ge \sigma \boldsymbol{\Delta}_B^{*t} \boldsymbol{\Delta}_B^*$$
  
=  $\sigma \|\boldsymbol{\theta}^{n+1} - \boldsymbol{\theta}^n\|^2$  (21)

for some fixed  $\sigma \geq 0$ .

**Theorem 1.** The sequence  $\{\theta^n\}$  generated by Algorithm 1 converges.

Proof. From Lemma 2, we have that  $f(\boldsymbol{\theta}^{n+1}) - f(\boldsymbol{\theta}^n) \geq 0$ . Thus the sequence  $\{f(\boldsymbol{\theta}^n)\}$  is monotonically increasing. Since it is bounded from above (by the optimum value) it must converge. Since convergent sequences are Cauchy, this sequence is also Cauchy. Thus for every  $\epsilon$ ,  $\exists n_0$  s.t  $f(\boldsymbol{\theta}^{n+1}) - f(\boldsymbol{\theta}^n) \leq \sigma \epsilon^2 \ \forall n \geq n_0$ . Again using Lemma 2, we get that

$$\|\boldsymbol{\theta}^{n+1} - \boldsymbol{\theta}^n\|^2 \le \epsilon^2 \tag{22}$$

for every  $n \geq n_0$ . Hence the sequence  $\{\boldsymbol{\theta}^n\}$  is Cauchy. The feasible set of  $\boldsymbol{\theta}$  is closed and compact, so Cauchy sequences are also convergent. Hence  $\{\boldsymbol{\theta}^n\}$  converges.

#### 5 Finite termination

We have shown that  $\{\boldsymbol{\theta}^n\}$  converges. Let  $\hat{\boldsymbol{\theta}}$  be a limit point of  $\{\boldsymbol{\theta}^n\}$ . We will start from the assumption that the algorithm runs for an infinite number of iterations and then prove a contradiction.

Call the variable  $\theta_{ik}$  as  $\tau$ -violating if the magnitude of the projected gradient  $\tilde{g}_{ik}$  is greater than  $\tau$ . Note that at every iteration, the chosen set of variables contains at least one that is  $\tau$ -violating. Now suppose the algorithm runs for an infinite number of iterations. Then it means that the sequence of iterates  $\boldsymbol{\theta}^k$  contains an infinite number of  $\tau$ -violating variables. Since there are only a finite number of distinct variables, we have that at least one variable figures as a  $\tau$ -violating variable in the chosen set B an infinite number of times. Suppose that  $\theta_{il}$  is one such variable, and let  $\{k_{il}\}$  be the subsequence in which this variable is chosen as a  $\tau$ -violating variable.

**Lemma 3.** For every  $\epsilon \quad \exists k_{il}^0 \quad s.t$   $\mid \theta_{il}^{k_{il}+1} - \theta_{il}^{k_{il}} \mid \leq \epsilon \quad \forall k_{il} > k_{il}^0.$ 

*Proof.* We have that since  $\theta^k \to \hat{\theta}$ ,  $\theta^{k_{il}} \to \hat{\theta}$ , Lemma 6. There can be only a finite number of and  $\boldsymbol{\theta}^{k_{il}+1} \to \hat{\boldsymbol{\theta}}$ . Thus, for any given  $\epsilon \exists k_{il}^0$  such that

$$\mid \theta_{il}^{k_{il}} - \hat{\theta}_{il} \mid \le \epsilon/2 \qquad \forall k_{il} > k_{il}^0 \qquad (23)$$

$$|\theta_{il}^{k_{il}+1} - \hat{\theta}_{il}| \le \epsilon/2 \qquad \forall k_{il} + 1 > k_{il}^0$$
 (24)

This gives, by triangle inequality,

**Lemma 4.**  $|\hat{g_{il}}| \geq \tau$ , where  $\hat{g_{il}}$  is the derivative of f w.r.t  $\theta_{il}$  at  $\boldsymbol{\theta}$ .

*Proof.* This is simply because of the fact that  $|g_{il}^{k_{il}}| \ge \tau$  for every  $k_{il}$ , and the absolute value of the derivative w.r.t  $\theta_{il}$  is a continuous function of  $\boldsymbol{\theta}$ , and  $\boldsymbol{\theta}^{k_{il}} \rightarrow \hat{\boldsymbol{\theta}}$ .

We use some notation. If  $\theta_{il}^{k_{il}} \in (l_{il}, u_{il})$  and if  $\theta_{il}^{k_{il}+1} = l_{il}$  or  $\theta_{il}^{k_{il}+1} = u_{il}$ , then we say that " $k_{il}$  is int  $\rightarrow$  bd", where "int" stands for interior and "bd" stands for boundary. Similar interpretations are assumed for "bd  $\rightarrow$  bd" and "int  $\rightarrow$ int". Thus each iteration  $k_{il}$  can be of one of only four possible kinds: int  $\rightarrow$  int,int  $\rightarrow$  bd, bd  $\rightarrow$  int and bd  $\rightarrow$  bd. We will prove that each of these kinds of iterations can only occur a finite number of times.

Lemma 5. There can be only a finite number of  $int \rightarrow int \ and \ bd \rightarrow int \ transitions.$ 

*Proof.* Suppose not. Then we can construct an infinite subsequence  $\{s_{il}\}$  of the sequence  $\{k_{il}\}$  that consists of these transitions. Then we have that  $g_{il}^{s_{il}+1}=0$ , using Lemma 1. Hence  $g_{il}^{s_{il}+1} \to 0$ . Since the gradient is a continuous function of  $\boldsymbol{\theta}$ , and since  $\boldsymbol{\theta}^{s_{il}+1} \to \hat{\boldsymbol{\theta}}$ , we have that  $g_{il}^{s_{il}+1} \rightarrow \hat{g}_{il}$ . But this means  $\hat{g}_{il} = 0$ , which contradicts Lemma 4.

 $int \rightarrow bd \ transitions.$ 

*Proof.* Suppose that we have completed sufficient number of iterations so that all int  $\rightarrow$  int and  $bd \rightarrow int transitions have completed.$  The next int  $\rightarrow$  bd transition will place  $\theta_{il}$  on the boundary. Since there are no bd  $\rightarrow$  int transitions anymore,  $\theta_{il}$  will stay on the boundary henceforth. Hence there can be no more int  $\rightarrow$ bd transitions.

**Lemma 7.** There can only be a finite number of  $bd \rightarrow bd \ transitions.$ 

*Proof.* Suppose not, i.e there are an infinite number of bd  $\rightarrow$  bd transitions. Let  $t_{il}$  be the subsequence of  $k_{il}$  consisting of bd  $\rightarrow$  bd transitions. Now, the sequence  $\theta_{il}^{t_{il}} \rightarrow \hat{\theta}_{il}$  and is therefore Cauchy. Hence  $\exists n_1$  s.t

$$\mid \theta_{il}^{t_{il}} - \theta_{il}^{t_{il}+1} \mid \leq \epsilon \ll u_{il} - l_{il} \qquad \forall t_{il} \geq n$$
 (26)

Similarly, because the gradient is a continuous function of  $\theta$ , the sequence  $\{g_{il}^{t_{il}}\}$  is convergent and therefore Cauchy. Hence  $\exists n_2$  s.t

$$|g_{il}^{t_{il}} - g_{il}^{t_{il}+1}| \le \frac{\tau}{2} \qquad \forall k_{il} \ge n_2$$
 (27)

Also, from the previous lemmas,  $\exists n_3 \text{ s.t } t_{il} \text{ is not}$ int  $\rightarrow$  int, bd  $\rightarrow$  int or int  $\rightarrow$  bd  $\forall t_{il} \geq n_3$ .

Take  $n_0 = \max(n_1, n_2, n_3)$ . Now, consider  $t_{il} \geq n_0$ . Without loss of generality, assume that  $\theta_{il}^{t_{il}} = l_{il}$ . Then, since  $|\tilde{g}_{il}^{t_{il}}| \geq \tau$ , we must have that  $g_{il}^{t_{il}} \geq \tau$ . From (26), and using the fact that this is a  $bd \rightarrow bd$  transition, we must have that

$$\theta_{il}^{t_{il}+1} = l_{il} \tag{28}$$

From (27), we have that

$$g_{il}^{t_{il}+1} \ge \frac{\tau}{2} \tag{29}$$

From (28) and (29), we have that  $\tilde{g}_{il}^{t_{il}+1}$ which contradicts Lemma 1.

But if all int  $\rightarrow$  int, int  $\rightarrow$  bd, bd  $\rightarrow$  int and bd  $\rightarrow$  bd transitions are finite, then  $\theta_{il}$  cannot be  $\tau$ -violating an infinite number of times and hence we have a contradiction. This gives us the following theorem:

**Theorem 2.** Algorithm 1 terminates in finite number of steps

### A Optimality conditions

This section proves three results on quadratic programs that have been used in the proof above. Consider the optimization problem:

$$\max_{\mathbf{x}} -\frac{1}{2}\mathbf{x}^t \mathbf{H} \mathbf{x} + \mathbf{p}^t \mathbf{x} \tag{30}$$

s.t

$$1 \le x \le u$$

The feasible set of the above problem is defined by box constraints and is therefore convex. If the Hessian **H** is positive definite, then this optimization problem will be convex. The gradient of f is denoted as  $\nabla f$ . The projected gradient of f, denoted as  $\nabla^P f$  is given by:

$$\nabla_i^P f(\mathbf{x}) = \begin{cases} \nabla_i f(\mathbf{x}) & \text{if } l_i < x_i < u_i \\ \max(0, \nabla_i f(\mathbf{x})) & \text{if } x_i = l_i \\ \min(0, \nabla_i f(\mathbf{x})) & \text{if } x_i = u_i \end{cases}$$
(31)

**Lemma 8.** Consider the optimization problem (30). Suppose  $\mathbf{H}$  is positive definite. Then,  $\mathbf{x}$  is optimum for (30) iff

1. **x** is feasible

2. 
$$\nabla_i^P f(\mathbf{x}) = 0 \quad \forall i$$

*Proof.* First let's write it as a minimization problem:

$$\min_{\mathbf{x}} h(\mathbf{x}) = -f(\mathbf{x}) \tag{32}$$

 $_{
m s.t}$ 

$$1 \le x \le u$$

The positive definiteness of  $\mathbf{H}$  and the fact that the constraints are box constraints imply that the problem is convex. Because the problem is convex, and because strong duality holds, a point  $\mathbf{x}$  is optimal iff the KKT conditions hold at  $\mathbf{x}$ . The Lagrangian of the above problem is:

$$\mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b}) = h(\mathbf{x}) + \mathbf{a}^t(\mathbf{l} - \mathbf{x}) + \mathbf{b}^t(\mathbf{x} - \mathbf{u})(33)$$

The KKT conditions are then:

$$\nabla_{\mathbf{x}} \mathcal{L} = \nabla h(\mathbf{x}) - \mathbf{a} + \mathbf{b} = \mathbf{0} \tag{34}$$

$$a_i(l_i - x_i) = 0 (35)$$

$$b_i(u_i - x_i) = 0 (36)$$

$$\mathbf{a} \ge 0 \tag{37}$$

$$\mathbf{b} \ge 0 \tag{38}$$

$$1 \le x \le u \tag{39}$$

From (34), (35) and (36), we get that:

$$\frac{\partial h}{\partial x_i}(\mathbf{x}) = a_i - b_i \qquad \forall i \tag{40}$$

$$\Rightarrow \frac{\partial h}{\partial x_i}(\mathbf{x}) = \begin{cases} 0 & \text{if } l_i < x_i < u_i \\ a_i \ge 0 & \text{if } l_i = x_i \\ -b_i \le 0 & \text{if } x_i = u_i \end{cases} \quad \forall i (41)$$

$$\Rightarrow \frac{\partial f}{\partial x_i}(\mathbf{x}) = \begin{cases} 0 & \text{if } l_i < x_i < u_i \\ -a_i \le 0 & \text{if } l_i = x_i \\ b_i \ge 0 & \text{if } x_i = u_i \end{cases} \quad \forall i(42)$$

$$\Rightarrow \nabla_i^P f(\mathbf{x}) = 0 \qquad \forall i \tag{43}$$

Conversely, suppose  $\nabla_i^P f(\mathbf{x}) = 0$  and  $\mathbf{x}$  is feasible. Taking  $a_i = \max(\frac{\partial h}{\partial x_i}(\mathbf{x}), 0)$  and  $b_i =$  $-\min(\frac{\partial h}{\partial x_i}(\mathbf{x}),0)$ , we have that

$$l_i < x_i < u_i \tag{44}$$

$$\Rightarrow \frac{\partial h}{\partial x_i}(\mathbf{x}) = 0 \tag{45}$$

$$\Rightarrow a_i = 0 \text{ and } b_i = 0$$
 (46)

$$x_i = l_i \tag{47}$$

$$\Rightarrow \frac{\partial h}{\partial x_i}(\mathbf{x}) = -\frac{\partial f}{\partial x_i}(\mathbf{x}) \ge 0 \tag{48}$$

$$\Rightarrow a_i \ge 0 \text{ and } b_i = 0$$
 (49)

$$x_i = u_i \tag{50}$$

$$\Rightarrow \frac{\partial h}{\partial x_i}(\mathbf{x}) = -\frac{\partial f}{\partial x_i}(\mathbf{x}) \le 0 \tag{51}$$

$$\Rightarrow a_i = 0 \text{ and } b_i > 0$$
 (52)

Thus (34), (35), (36), (37) and (38) are satisfied by this choice of a and b. Thus the KKT conditions are equivalent to the following conditions, which are therefore necessary and sufficient for  $\mathbf{x}$  to be optimal:

$$\nabla_i^P f(\mathbf{x}) = 0 \qquad \forall i \tag{53}$$

$$1 \le \mathbf{x} \le \mathbf{u} \tag{54}$$

 $\Rightarrow \frac{\partial f}{\partial x_i}(\mathbf{x}) = \begin{cases} 0 & \text{if } l_i < x_i < u_i \\ -a_i \le 0 & \text{if } l_i = x_i \\ b_i \ge 0 & \text{if } x_i = u_i \end{cases} \quad \forall i (42) \quad \text{(30)}. \quad \begin{array}{l} \text{Lemma 9. Consider the optimization problem} \\ \text{(30)}. \quad \text{Suppose $\mathbf{H}$ is positive definite. If $\mathbf{x}^*$ is optimal, and $\mathbf{x} \ne \mathbf{x}^*$ is feasible, then $(\mathbf{x}^* - \mathbf{x}^*)$ is the sum of th$  $(\mathbf{x})^t \nabla f(\mathbf{x}^*) \ge 0$ 

> *Proof.* Because the feasible set is convex and because both  $\mathbf{x}$  and  $\mathbf{x}^*$  are feasible, we have that  $\mathbf{x}(\lambda) = \mathbf{x} + \lambda(\mathbf{x}^* - \mathbf{x})$  is feasible for all  $\lambda \in [0, 1]$ . Consider the optimization problem:

$$\max_{\lambda} f(\mathbf{x}(\lambda)) - f(\mathbf{x}) \tag{55}$$

$$0 < \lambda < 1$$

From a Taylor series expansion, it can be seen

$$f(\mathbf{x}(\lambda)) - f(\mathbf{x}) = -\frac{\lambda^2}{2} (\mathbf{x}^* - \mathbf{x})^t H(\mathbf{x}^* - \mathbf{x})$$
$$+\lambda (\mathbf{x}^* - \mathbf{x})^t \nabla f(\mathbf{x})$$
(56)

Since **H** is positive definite,  $(\mathbf{x}^* - \mathbf{x})^t \mathbf{H} (\mathbf{x}^* - \mathbf{x}) >$ 0. Hence, the problem (55) satisfies the conditions of Lemma 8. Also, because  $\mathbf{x}^*$  is optimal for (30),  $\lambda = 1$  is optimal for (55). Hence, by Lemma 8, the projected gradient of (55) at  $\lambda = 1$ is 0. This means that

$$\left. \frac{\partial f(\mathbf{x}(\lambda))}{\partial \lambda} \right|_{\lambda=1} \ge 0 \tag{57}$$

This gives, using the chain rule

$$(\mathbf{x} - \mathbf{x}^*)^t \nabla f(\mathbf{x}(\lambda)) \Big|_{\lambda=1} \ge 0$$
 (58)

$$\Rightarrow (\mathbf{x}^* - \mathbf{x})^t \nabla f(\mathbf{x}^*) \ge 0 \tag{59}$$

where the last step uses the fact that  $\mathbf{x}(1) =$  $\square$   $\mathbf{x}^*$ .  **Lemma 10.** Consider the optimization problem (30). Suppose **H** is positive definite. If **x** is feasible and  $|\nabla_{ik}^P f(\mathbf{x})| \leq \tau$ , and if  $f^*$  is the optimal value, then

$$f^* - f(\mathbf{x}) \le \frac{d^2 \tau^2}{2m} \tag{60}$$

where d is the dimensionality of  $\mathbf{x}$  and m is the least eigenvalue of  $\mathbf{H}$ .

*Proof.* This proof is the extension of the proof in [1, p. 459]. Again, take h = -f and define the Lagrangian. Then, for the given  $\mathbf{x}$ , we can define

$$\hat{a}_i = \max(\frac{\partial h}{\partial x_i}(\mathbf{x}), 0) \tag{61}$$

$$a_i = \begin{cases} \hat{a}_i & \text{if } \hat{a}_i \ge \tau \\ 0 & ow \end{cases} \tag{62}$$

$$\hat{b}_i = -\min(\frac{\partial h}{\partial x_i}(\mathbf{x}), 0) \tag{63}$$

$$b_i = \begin{cases} \hat{b}_i & \text{if } \hat{b}_i \ge \tau \\ 0 & ow \end{cases} \tag{64}$$

Then, equations (35), (36), (37) and (38) are satisfied by this choice of  $\bf a$  and  $\bf b$ . Equations (35) and (36) mean that this choice of  $\bf a$  and  $\bf b$  satisfies:

$$h(\mathbf{x}) = \mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b}) \tag{65}$$

(34) is only satisfied to a precision of  $\tau$ , which means that

$$\|\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b})\| \le d\tau \tag{66}$$

Also note that:

$$h^* = -f^* = \max_{\mathbf{a}, \mathbf{b}} \min_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b})$$

$$\geq \min_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b})$$

$$= \mathcal{L}(\mathbf{x}^0, \mathbf{a}, \mathbf{b})$$
(67)

for some  $\mathbf{x}^0$ .

Hence,

$$h(\mathbf{x}) - h^* \le h(\mathbf{x}) - \mathcal{L}(\mathbf{x}^0, \mathbf{a}, \mathbf{b})$$
 (68)

$$\Rightarrow h(\mathbf{x}) - h^* \le \mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b}) - \mathcal{L}(\mathbf{x}^0, \mathbf{a}, \mathbf{b})$$
 (69)

Using a Taylor series expansion of  $g(\mathbf{x}) = \mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b})$  we see that:

$$g(\mathbf{x}^0) - g(\mathbf{x}) = (\mathbf{x}^0 - \mathbf{x})^t \nabla g(\mathbf{x}) + \frac{1}{2} (\mathbf{x}^0 - \mathbf{x})^t \nabla^2 g(\mathbf{x}) (\mathbf{x}^0 - \mathbf{x}) (70)$$

The Hessian  $\nabla^2 g(\mathbf{x})$  is nothing but  $\mathbf{H}$ . The quantity  $(\mathbf{x}^0 - \mathbf{x})^t \nabla^2 g(\mathbf{x}) (\mathbf{x}^0 - \mathbf{x})$  is lowerbounded by  $m \|\mathbf{x}^0 - \mathbf{x}\|^2$ , m > 0 being the least eigenvalue of  $\mathbf{H}$ . Hence,

(64) 
$$g(\mathbf{x}^0) - g(\mathbf{x}) \ge (\mathbf{x}^0 - \mathbf{x})^t \nabla g(\mathbf{x}) + \frac{m}{2} ||\mathbf{x}^0 - \mathbf{x}||^2 (71)$$

The right hand side is a convex function of  $\mathbf{x}^0$  for fixed  $\mathbf{x}$ . Setting the gradient = 0, we find that the right hand side attains its minimum value of  $-\frac{1}{2m}\|\nabla g(\mathbf{x})\|^2$  at  $\mathbf{x}^0 = \mathbf{x} - \frac{1}{m}\nabla g(\mathbf{x})$ . Thus,

$$g(\mathbf{x}^0) - g(\mathbf{x}) \ge -\frac{1}{2m} \|\nabla g(\mathbf{x})\|^2 \qquad (72)$$

Combining (66), (69) and (72), and noting that f = -h and  $g(\mathbf{x}) = \mathcal{L}(\mathbf{x}, \mathbf{a}, \mathbf{b})$  we get the desired inequality.

### References

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