

CS32 Week 5: Sorting

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Sorting

- Sorting is useful.
 - Better retrieval efficiency.
 - Better data organization.
 - Sorting midterm papers by last name, first name.

Sorting

- Selection Sort

- Find the smallest item in the unsorted portion, then place it in the front.

- Time complexity:

- Worst case?

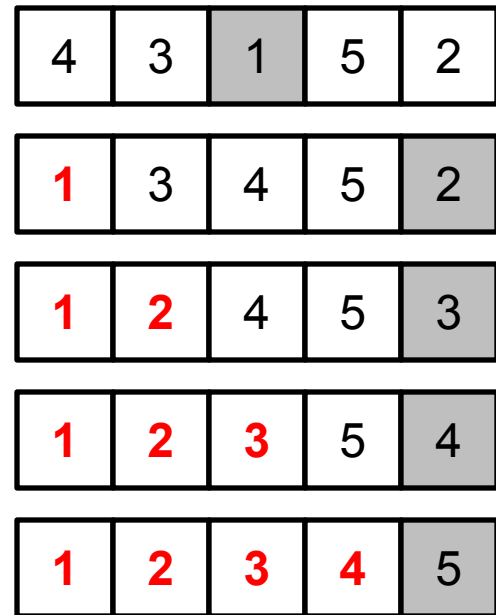
- $O(n^2)$

- Best case?

- $O(n^2)$

- Average case?

- $O(n^2)$



Sorting

- Insertion Sort

- Pick one from unsorted portion, and place it in the right position in the sorted part.

- Time complexity:

- Worst case?

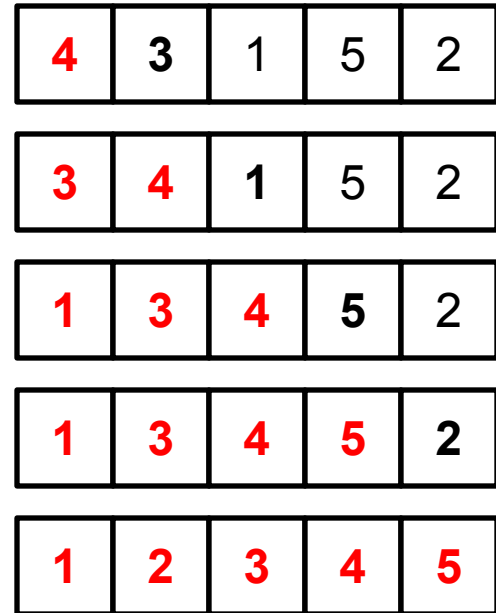
- $O(n^2)$

- Best case?

- $O(n)$

- Average case?

- $O(n^2)$



Sorting

- Bubble Sort

- Bubble up the largest item to the end of unsorted part, until all is sorted. Loop through the array, if not in the correct order, swap.

4	3	1	5	2
3	4	1	5	2
3	1	4	5	2
3	1	4	5	2
3	1	4	2	5

Sorting

- Bubble Sort

- Bubble up the largest item to the end of unsorted part, until all is sorted. Loop through the array, if not in the correct order, swap.

3	1	4	2	5
---	---	---	---	---

1	3	4	2	5
---	---	---	---	---

1	3	4	2	5
---	---	---	---	---

1	3	2	4	5
---	---	---	---	---

Sorting

- Bubble Sort

- Bubble up the largest item to the end of unsorted part, until all is sorted. Loop through the array, if not in the correct order, swap.

- Time complexity:

- Worst case?

- $O(n^2)$

- Best case?

- $O(n)$

- Average case?

- $O(n^2)$

1	3	2	4	5
---	---	---	---	---

1	3	2	4	5
---	---	---	---	---

1	2	3	4	5
---	---	---	---	---

1	2	3	4	5
---	---	---	---	---

Sorting

- Merge Sort

- Break the array into half, sort each half, then merge them together.

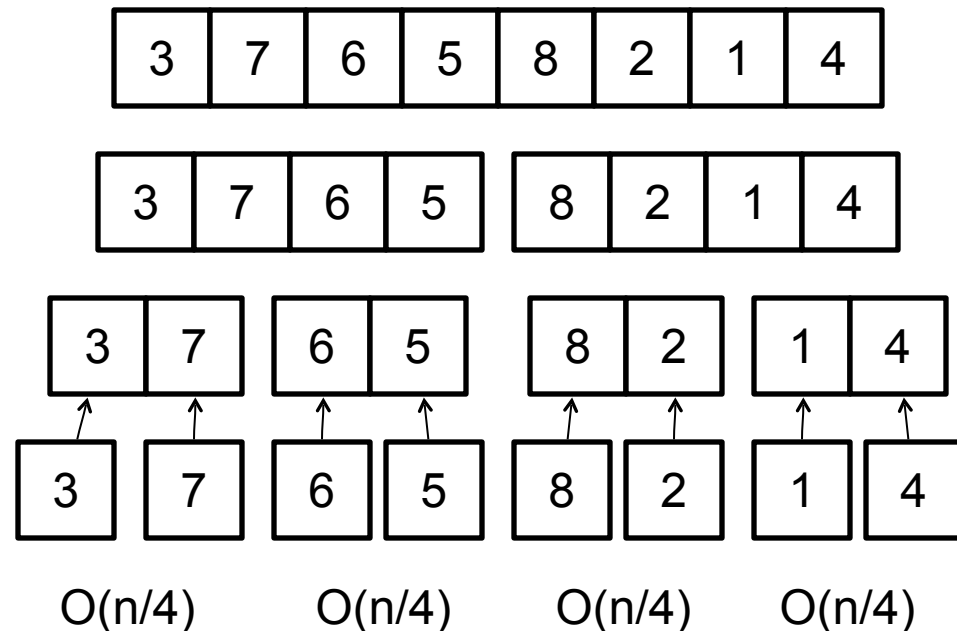
- Recursion

- Base case?
- How to merge results?

- Merge complexity?

- n is array size
- Each level is $O(n)$

Level sum: $O(n)$



Sorting

- Merge Sort

- Break the array into half, sort each half, then merge them together.

- Merge complexity?

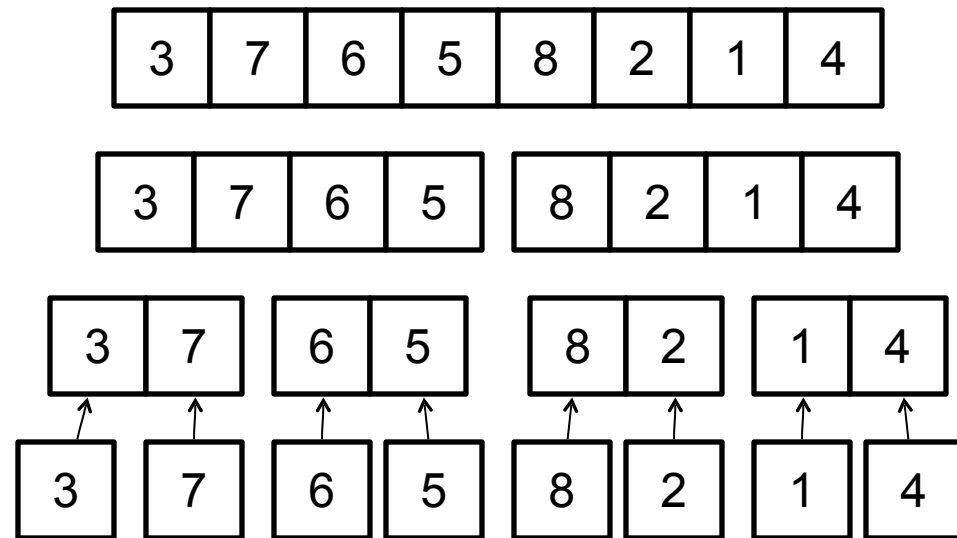
- n is array size
- Each level is $O(n)$

- Number of Level

- $\log_2 n$

- Total complexity

- $O(n \log n)$ Level sum: $O(n)$ $O(n/4)$ $O(n/4)$ $O(n/4)$ $O(n/4)$



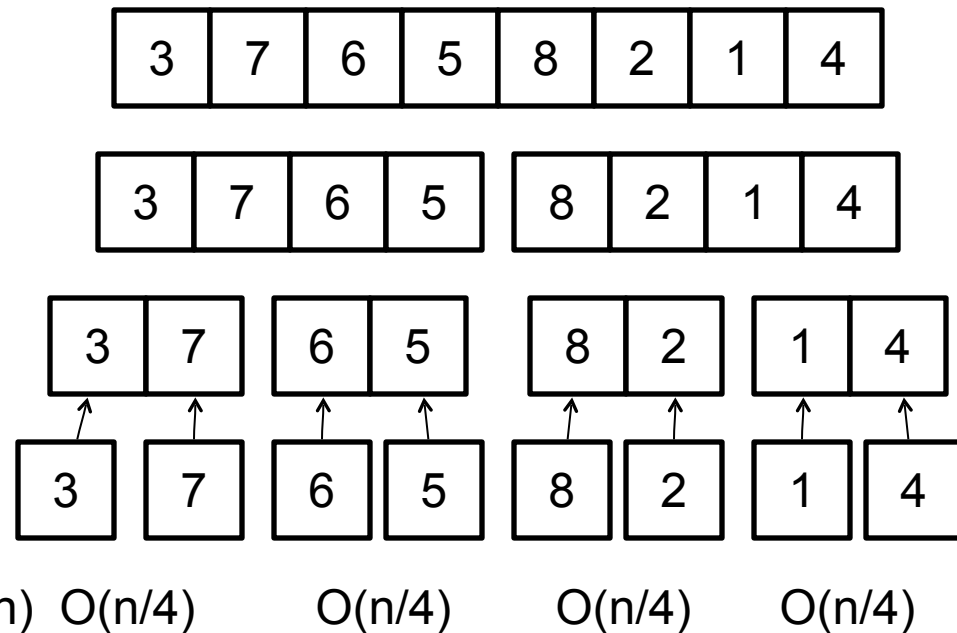
Sorting

- Merge Sort

- Break the array into half, sort each half, then merge them together.

- Disadvantage?

- Need additional memory for merging.



Sorting

- Quick Sort

- Pick a pivot, move items that are smaller than pivot to the front, move items that are larger than pivot to the end.

- #Level

- Expected to be $\log_2 n$
- Depending on pivot

- Complexity:

- Worst case?
 - $O(n^2)$ for already sorted array
- Average case?
 - $O(n \log n)$

4	3	1	5	2
---	---	---	---	---

3	1	2	4	5
---	---	---	---	---

3	1	2
---	---	---

1	2	3
---	---	---

Sorting

```
int Partition(int a[], int low, int high)
{
    int pi = low;
    int pivot = a[low];
    do
    {
        while ( low <= high && a[low] <= pivot )
            low++;
        while ( a[high] > pivot )
            high--;
        if ( low < high )
            swap(a[low], a[high]);
    }
    while ( low < high );
    swap(a[pi], a[high]);
    pi = high;
    return(pi);
}
```

4	3	1	5	2
---	---	---	---	---

4	3	1	2	5
---	---	---	---	---

2	3	1	4	5
---	---	---	---	---

Sorting

- Quick Sort
 - ❑ Pick a pivot, move items that are smaller than pivot to the front, move items that are larger than pivot to the end.
 - ❑ Partition function
 - ❑ How to pick pivot?
 - First? Last? Or random?
 - It's important for the efficiency
 - ❑ In-place sorting

4	3	1	5	2
---	---	---	---	---

3	1	2	4	2
---	---	---	---	---

3	1	2
---	---	---

1	2	3
---	---	---

Sorting

- What's the best algorithm to sort 1,000,000 random numbers that are all between 1 and 5?
- Bucket sort
 - Partition the whole into multiple buckets, sort each buckets individually (could recursively call bucket sort or other sort), and gather together.
 - Bucket: 1, 2, 3, 4, 5

Sorting

- Bucket sort
 - Partition the whole into multiple buckets, sort each buckets individually (could recursively call bucket sort or other sort), and gather together.
 - Average case: $O(n+k)$

