Scrap your Boilerplate with Object Algebras!

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Abstract. This is the abstract ...

1 Introduction

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2 Overview

We start by considering the company structure introduced in Figure 1.
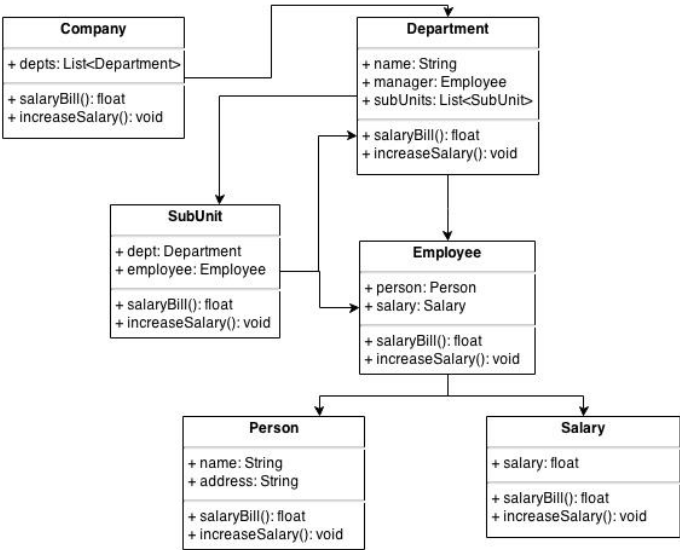

Fig. 1. Company Structure

```
public class Company {
  private List<Department> depts;
  public Company(List<Department> depts){this.depts = depts;}
  public Float salaryBill(){
    Float r = 0f;
    for (Department dept: depts) r+= dept.salaryBill();
    return r;
  }
  public void increaseSalary(){
    for (Department dept: depts) dept.increaseSalary();
  }
}
```

Fig. 2. Company Class of OOP style

```
public class Salary {
  private Float salary;
  public Salary(Float salary){this.salary = salary;}
  public Float salaryBill(){return this.salary;}
  public void increaseSalary(){this.salary *= 1.1f;}
}
```

Fig. 3. Salary Class of OOP style

2.1 Object Oriented Solution

A very natural Object-Oriented way to model the company structure is as illustrated in Figure 2 and Figure 3. Similar code can be applied to Department, Employee, SubUnit and Person. A Company contains a list of Departments. each Department is managed by an Employee as the manager and contains a list of SubUnits. The SubUnit can be either a department or an Employee. An Employee is a Person with the Salary Information.

Now consider adding two operations to our company structure: query the salary bill for the whole company and increase the salary of each employee by 10%. The methods Float salaryBill() and void increaseSalary() in Figure 2 and Figure 3 is an easy solution.

This way of Object Oriented style representation of tree structures can become cumbersome and inflexible due to the bound relationship between classes. For instance, adding a new operation such as pretty printing of the company structure requires a lot of changes on the existing code and violates the no modification rule.

2.2 Modeling Company Structure with Object Algebras

Object Algebras is a good solution to solve the extensibility problem. Figure 4 shows the approach to model the Company structure with Object Algebras.

Hence different operations can be realized by inheriting object algebras from the object algebra interface. To implement query bill operation for the whole

```
@Algebra
public interface SybAlg<Company, Dept, SubUnit, Employee, Person, Salary>{
  public Company C(List<Dept> depts);
  public Dept D(String name, Employee manager, List<SubUnit> subUnits);
  public SubUnit PU(Employee employee);
  public SubUnit DU(Dept dept);
  public Employee E(Person p, Salary s);
  public Person P(String name, String address);
  public Salary S(float salary);
}
```

Fig. 4. Company Structure represented by Object Algebra Interface

Company structure, we need to implement the Company interface with specific operation for each component.

```
public class SalaryQuerySybAlg implements SybAlg<Float,Float,Float,Float,
    Float,Float> {
  public Float C(List<Float> depts){
    Float r = 0f;
    for (Float f: depts) r += f;
    return r;
  }
  //... The code for Department, Employee, SubUnit and Person are omitted
      here.
  public Float S(float salary){
    return salary;
  }
}
```

Increasing Salary is trickier. We need the visitable classes to help reconstruct the tree structures of the Company with increased salary.

```
public interface G_Company {
  <Company,Dept,SubUnit,Employee,Person,Salary> Company accept(SybAlg<
      Company,Dept,SubUnit,Employee,Person,Salary> alg);
}
public interface G_Salary {
  <Company,Dept,SubUnit,Employee,Person,Salary> Salary accept(SybAlg<
      Company,Dept,SubUnit,Employee,Person,Salary> alg);
}
public class IncreaseSalarySybAlg implements SybAlg<G_Company, G_Dept,
    G_SubUnit, G_Employee, G_Person, G_Salary>{
  public G_Company C(List<G_Dept> p0) {
    return new G_Company() {
      public <Company,Dept,SubUnit,Employee,Person,Salary> Company accept(
          SybAlg<Company,Dept,SubUnit,Employee,Person,Salary> alg) {
        List<Dept> gp0 = new ArrayList<Dept>();
        for (G_Dept s: p0) {
          gp0.add(s.accept(alg));
        }
        return alg.C(gp0);
      }
    };
```

```
public class FloatQuerySybAlgebra extends SybAlgQuery<Float> {
  public FloatQuerySybAlgebra(Monoid<Float> m) {
    super(m);
  }
  public Float S(float p0){return p0;}
}
```

Fig. 5. Query Salary Class with Object Algebra Framework

```
  }
 //... The code for Department, Employee, SubUnit and Person are omitted
      here.
 public G_Salary S(float p0) {
   return new G_Salary() {
     public <Company,Dept,SubUnit,Employee,Person,Salary> Salary accept(
         SybAlg<Company,Dept,SubUnit,Employee,Person,Salary> alg) {
       return alg.S(1.1f*p0);
     }
   };
 }
}
```

However, although we solved the problem of extensibility with object algebras, the traversal code become so long and most of the time we are writing boilerplate routine code, which is to call the methods of its child leaves. The only code we are really interested in is the Salary S(Float salary) method to return or to increase the salary. It will be great if the boilerplate code can be generated automatically every time we want to traverse the tree structure.

2.3 Object Algebra Framework

Motivated by this problem of generating generic code for tree structure traversals, more specifically, queries and transformations in object algebras, we introduce monoids in queries and generic visitable interfaces in transformations, and write generic query and transformation which can be easily inherited by real cases of queries and transformations. Furthermore, we designed an object algebra framework with great features. With our framework, even the generic query and transformation code can be generated automatically by adding an "@Algebra" annotation.

Now with our Object Algebra Framework, the code we need to write for Salary Bill and Increase Salary will be much shorter. A Generic query code will be as short as Figure 5. Transformation code will be like Figure 6 The classes SybAlgQuery<R>, SybAlgTransform and G_SybAlg_Salary are all generated by the framework automatically.

3 Queries

As a specific type of object algebras, queries allow users to define new operations handling a user-defined data structure. A *query algebra* is a class implementing

```
public interface SybIncSalary extends SybAlgTransform {
  default G_SybAlg_Salary S(float p0) {
    return new G_SybAlg_Salary() {
      public <Company, Dept, SubUnit, Employee, Person, Salary> Salary
          accept(SybAlg<Company, Dept, SubUnit, Employee, Person, Salary>
          alg) {
        return alg.S(1.1f * p0);
      }
    };
  }
}
```

Fig. 6. Increase Salary Class with Object Algebra Framework

an object algebra interface by a top-down traversal throughout the hierarchy. It is something supporting the program to gather information from the substructures of a data type recursively, and make a response at the root node to the query.

3.1 FreeVars: a simple query algebra

An example is shown here to discuss about query algebras in a clearer way. The object algebra interface is related to an expression, where it can be treated as a literal, a string or composed of two smaller ones. Specifically, the structure is defined as follows:

```
public interface ExpAlg<Exp> {
   Exp Var(String s);
   Exp Lit(int i);
   Exp Add(Exp e1, Exp e2);
}
```

Based on the interface above, a query might be raised on collecting all the names of free variables defined in an expression. More precisely, an array of strings would be used to store the names of those variables. In that case, a $Var(s)$ would simply return an one-element array of s, and a $Lit(i)$ corresponds to an empty set, whereas two arrays would be merged into one if we are combining two expressions with the $Add()$ method.

Generally speaking, it is natural to deal with the traversal in an algebra-based approach like this:

```
public interface FreeVarsExpAlg extends ExpAlg<String[]> {
   default String[] Var(String s) {
      return new String[]{s};
   }
   default String[] Lit(int i) {
      return new String[]{};
   }
```

```
    default String[] Add(String[] e1, String[] e2) {
        int e1len = e1.length;
        int e2len = e2.length;
        String[] res = new String[e1len+e2len];
        System.arraycopy(e1, 0, res, 0, e1len);
        System.arraycopy(e2, 0, res, e1len, e2len);
        return res;
    }
}
```

Information on our query is collected by traversal and passed on to a higher-level structure. Nonetheless, a programmer has to write a lot of boring code handling the traversals, and it could be even worse for a more complicated data structure. Moreover, it is a query-based approach: you still have to write a bunch of similar stuff with a different query raised, for instance, a pretty printer.

3.2 Generic query algebra with a monoid

Queries are so similar actually: a user has to indicate the rules in which the program may address cases on primitive types and "append" the information. With these two issues, everything becomes simple in the traversal. Hence we introduce the concept of monoid and generic traversal here in our query algebras.

```
public interface Monoid<R> {
    R join(R x, R y);
    R empty();
    default R fold(List<R> lr){
      R res = empty();
      for (R r: lr){
        res = join(res, r);
      }
      return res;
    }
}
```

The interface of a monoid is defined above. Intuitively, the join() method implies how we gather the information from substructures during merging, and the empty() is just an indicator of "no information". Hence now we are able to write a "generic traversal" manually based on monoids as follows:

And now we find everything goes in an easier way: we don't care about what kind of query it is any more during the traversal. Despite whether it asks for all the names of free variables or a printer showing the hierarchy of an expression, all we need to do is to traverse a monoid to the generic traversal class, and the monoid is exactly constructed by implementing the previous interface. This is the progress, once we prepare such a template dealing with the traversal, all query algebras can be addressed in a more concise way, which is called the *generic query algebra.*

```java
public class QueryExpAlg<Exp> implements ExpAlg<Exp> {
 private Monoid<Exp> m;
 public Monoid<Exp> m() { return m; }
 public QueryExpAlg(Monoid<Exp> m) {
   this.m = m;
 }
 public Exp Add(Exp p0,Exp p1) {
  Exp res = m.empty();
  res = m.join(res, p0);
  res = m.join(res, p1);
  return res;
 }
 public Exp Lit(int p0) {
  Exp res = m.empty();
  return res;
 }
 public Exp Var(String p0) {
  Exp res = m.empty();
  return res;
 }
}
```

Fig. 7. Generic Query by hand with Monoid

3.3 Solving freeVars with generic query algebra

As an alternative way to handle the freeVars query, the query algebra is going
to be a subclass of QueryExpAlg, the generic algebra, with generic type to be
String[]. To use the generic traversal code, a monoid is defined as follows:

```java
public class FreeVarsMonoid implements Monoid<String[]> {
   public String[] empty() {
      return new String[]{};
   }
   public String[] join(String[] e1, String[] e2) {
      int e1len = e1.length;
      int e2len = e2.length;
      String[] res = new String[e1len+e2len];
      System.arraycopy(e1, 0, res, 0, e1len);
      System.arraycopy(e2, 0, res, e1len, e2len);
      return res;
   }
}
```

But the result for an expression can only be a null array based on the monoid.
Thus in the freeVars query, furthermore, we expect the variables to store their
names into an array, and by using the monoid, the query algebra will be like:

```java
public class FreeVarsQueryExpAlg extends QueryExpAlg<String[]> {
   public FreeVarsQueryExpAlg(Monoid<String[]> m) {super(m);}
   public String[] Var(String s) {return new String[]{s};}
}
```

```
public interface G_Exp {
  <Exp> Exp accept(ExpAlg<Exp> alg);
}
```

Fig. 8. Generic Visitable Interface

When the class FreeVarsQueryExpAlg is used, an object of the FreeVarsMonoid should be created and traversed to the constructor. As we can see, it is needless for a user to address the traversals in a data structure. Nothing but a monoid is required together with a few methods being overwritten. And furthermore, a monoid can usually be shared among query algebras with the same data type.

4 Transformations

Before writing generic transformation interfaces, we first introduce *generic visitable interface* as Figure 8. Implementing the generic visitable interface allows the user to construct objects from the passed in algebra.

The generic transform interface is constructed by inheriting from the *Object Algebra Interface* with *Generic Visitable Interfaces* as Figure 9. Note that the returned *Generic Visitable Interface* will contain all the information for the tree structure.

Now to create a specific transformation, e.g., substitute one variable with another name, it can be easily implemented by inheriting from the generic transform interface with the implementation of interesting cases. Here only the method G_Exp Var(String ss) needs to be overridden as Figure 10.

5 Object Algebras Framework

Writing generic queries and transformations with monoids and generic visitable interfaces can help users write tree structure traversal code with more extensibility and flexibility. However, writing the generic query and transformation interfaces is still painful experience by itself. It will be even better if these boilerplate code can be generated automatically. To address this problem, we provide an *Object Algebra Framework*, which utilizes *Java Annotation* to generate generic query and transformation interfaces based on the *Object Algebra Interface*, as illustrated below:

```
@Algebra
public interface ExpAlg<Exp> {
  Exp Var(String s);
  Exp Lit(int i);
  Exp Add(Exp e1, Exp e2);
}
```

With the annotation "@Algebra", the framework will generate the boilerplate codes for us automatically. As for our ExpAlg example, the following directory structure will be generated by the library.

```java
public interface ExpAlgTransform extends ExpAlg<G_Exp> {
  @Override
  default G_Exp Add(G_Exp e1, G_Exp e2) {
    return new G_Exp() {
      @Override
      public <Exp> Exp accept(ExpAlg<Exp> alg) {
        return alg.Add(e1.accept(alg), e2.accept(alg));
      }
    };
  }
  @Override
  default G_Exp Lit(int i) {
    return new G_Exp() {
      @Override
      public <Exp> Exp accept(ExpAlg<Exp> alg) {
        return alg.Lit(i);
      }
    };
  }
  @Override
  default G_Exp Var(String s) {
    return new G_Exp() {
      @Override
      public <Exp> Exp accept(ExpAlg<Exp> alg) {
        return alg.Var(s);
      }
    };
  }
}
```

Fig. 9. Generic Transformation by hand with acceptor interface

```
src/
├── generic/
│   └── G_ExpAlg_Exp
├── query/
│   └── ExpAlgQuery
├── transform/
    └── ExpAlgTransform
```

Furthermore, the monoid interface is also included in the Object Algebras Framework

Hence when programming with query, one only needs to focus on the interesting case, for instance, the return free variable names, and the specific monoid needed, which in Free Variables case will be a String List monoid. While programming with transformation will be implementing interesting cases by inheriting from the AlgNameTransform interface with *generic visitable interfaces*.

```
public class SubstVarsTransform implements ExpAlgTransform{
  private String s;
  private G_Exp gExp;
  public SubstVarsTransform(String s, G_Exp gExp){
    this.s = s;
    this.gExp = gExp;
  }
  @Override
  public G_Exp Var(String ss) {
    return new G_Exp() {
      @Override
      public <Exp> Exp accept(ExpAlg<Exp> alg) {
        if (ss.equals(s)) return gExp.accept(alg);
        else return alg.Var(ss);
      }
    };
  }

}
```

Fig. 10. Substitute Variables Transformation

6 Other Features

One More section

7 Case Study

Case study section

8 Related Work

Finally related work.

9 Conclusion

And conclusion.

Acknowledgements We should thank someone!

References