

Module ModularArithmetic.

Definition F (m : positive) := { z : Z | z = z mod (Z.pos m) }.

Program Definition of\_Z m (a:Z) : F m := a mod (Z.pos m).

Definition to\_Z {m} (a:F m) : Z := proj1\_sig a.

Context {m : positive}.

Definition zero : F m := of\_Z m 0.

Definition one : F m := of\_Z m 1.

Definition add (a b:F m) : F m := of\_Z m (to\_Z a + to\_Z b).

Definition mul (a b:F m) : F m := of\_Z m (to\_Z a \* to\_Z b).

Definition opp (a :F m) : F m := of\_Z m (0 - to\_Z a).

Definition sub (a b:F m) : F m := add a (opp b).

Definition inv\_with\_spec : { inv : F m → F m  
| inv zero = zero ∧ ( prime (Z.pos m) →  
    ∀ a, a ≠ zero → mul (inv a) a = one )  
} := ModularArithmeticPre.inv\_impl.

Definition inv : F m → F m := Eval hnf in proj1\_sig inv\_with\_spec.

Definition div (a b:F m) : F m := mul a (inv b).

End ModularArithmetic.

Module EdwardsCurve. (\* <<https://eprint.iacr.org/2008/013.pdf>> \*)

Context `{field:@Algebra.Hierarchy.field F Feq Fzero Fone Fopp Fadd Fsub Fmul Finv Fdiv}  
    {char\_ge\_3 : @Ring.char\_ge F Feq Fzero Fone Fopp Fadd Fsub Fmul 3}  
    {a : F} {nonzero\_a : a ≠ 0} {square\_a : ∃ sqrt\_a, sqrt\_a^2 = a}  
    {d : F} {nonsquare\_d : ∀ x, x^2 ≠ d} {Feq\_dec:Decidable.DecidableRel Feq}.

Definition point := { xy | let '(x, y) := xy in a\*x^2 + y^2 = 1 + d\*x^2\*y^2 }.

Definition coordinates (P:point) : (F\*F) := let (xy, xy\_onCurve\_proof) := P in xy.

Program Definition zero : point := (0, 1).

Program Definition add (P1 P2:point) : point :=  
    match coordinates P1, coordinates P2 return (F\*F) with  
    (x1, y1), (x2, y2) =>  
    (((x1\*y2 + y1\*x2)/(1 + d\*x1\*x2\*y1\*y2)) , ((y1\*y2 - a\*x1\*x2)/(1 - d\*x1\*x2\*y1\*y2)))  
    end.

Fixpoint mul (n:N) (P : point) : point :=  
    match n with 0 => zero | S n' => add P (mul n' P) end.

End EdwardsCurve.

Module EdDSA. (\* <<https://eprint.iacr.org/2015/677.pdf>> \*)

Class EdDSAParametersOK

{E Eeq Eadd Ezero Eopp} {EscalarMult} {c:N} {B:E} {l : positive} {n b:N}

{H : ∀ {t}, word t → word (b + b)} {Eenc:E→word b} {Senc:F l→word b}

:= { EdDSA\_group :> @Algebra.Hierarchy.group E Eeq Eadd Ezero Eopp;  
    EdDSA\_scalarmult:@Algebra.ScalarMult.is\_scalarmult E Eeq Eadd Ezero EscalarMult;  
    EdDSA\_c\_valid : c = 2 ∨ c = 3; EdDSA\_n\_ge\_c : n ≥ c; EdDSA\_n\_le\_b : n ≤ b;  
    EdDSA\_B\_not\_identity : not (Eeq B Ezero); EdDSA\_l\_prime : Znumtheory.prime l;  
    EdDSA\_l\_odd : Z.lt 2 l; EdDSA\_l\_order\_B :> Eeq (EscalarMult (Z.to\_nat l) B) Ezero }.

Context `{prm:EdDSAParametersOK}. Notation signature := (word (b + b)).

Notation secretkey := (word b). Notation publickey := (word b).

Infix "^" := Nat.pow. Local Infix "mod" := Nat.modulo. Infix "++" := Word.combine.

Infix "+" := Eadd. Local Infix "\*" := EscalarMult.

Coercion wordToNat : word → N. Coercion Z.to\_nat : Z → N.

Program Definition curveKey (sk:secretkey) : N :=

    let x := wfirstn n (H sk) in (\* hash the key, use first "half" for secret scalar \*)  
    let x := x - (x mod (2^c)) in (\* it is implicitly 0 mod (2^c) \*)  
    PeanoNat.Nat.setbit x n. (\* and the high bit is always set \*)

Definition public (sk:secretkey) : publickey := Eenc (curveKey sk\*B).

Definition prngKey (sk:secretkey) : word b := Word.split2 b b (H sk).

Program Definition sign (A\_:publickey) sk {n} (M : word n) :=

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let r :  $\mathbb{N}$  := H (prngKey sk ++ M) in (* secret nonce *)
let R : E := r * B in (* commitment to nonce *)
let s :  $\mathbb{N}$  := curveKey sk in (* secret scalar *)
let S : F l := F.nat_mod l (r + H (Eenc R ++ A_ ++ M) * s) in
  Eenc R ++ Senc S.
Definition valid {n} (message : word n) pubkey signature :=
   $\exists$  A S R, Eenc A = pubkey  $\wedge$  Eenc R ++ Senc S = signature  $\wedge$ 
  Eq (F.to_nat S * B) (R + (H (Eenc R ++ Eenc A ++ message) mod l) * A).
End EdDSA.

Section Ed25519.
Definition q : positive := 2255 - 19.
Definition a : F q := F.opp 1. Definition d : F q := F.opp(F.of_Z _ 121665)/(F.of_Z _ 121666).
Definition E := E.point(F:=F q)(Feq:=eq)(Fone:=F.one)(Fadd:=F.add)(Fmul:=F.mul)(a:=a)(d:=d).
Definition b :  $\mathbb{N}$  := 256. Definition n :  $\mathbb{N}$  := b - 2. Definition c :  $\mathbb{N}$  := 3.
Definition l : positive := 2252 + 27742317777372353535851937790883648493.
Program Definition B : E :=
  (F.of_Z q 15112221349535400772501151409588531511454012693041857206046113283949847762202,
   F.of_Z q 4 / F.of_Z q 5)%F.
Definition Fencode {len} {m} : F m → word len :=
   $\lambda$  x : F m ⇒ (NToWord _ (Z.to_N (F.to_Z x))).
Definition signum (x : F q) := Z.testbit (F.to_Z x) 0.
Definition Eenc : E → word b :=  $\lambda$  P ⇒
  let '(x,y) := E.coordinates P in Word.combine (Fencode (len:=b-1) y) (bit (signum x)).
Definition Senc : F l → word b := Fencode (len:=b).
End Ed25519.

Module WeierstrassCurves. (* <https://hyperelliptic.org/EFD/glp/auto-shortw.html> *)
Context {field:@Algebra.Hierarchy.field F Feq Fzero Fone Fopp Fadd Fsub Fmul Finv Fdiv}
  {Feq_dec:Decidable.DecidableRel Feq} {a b:F}
  {char_ge_3:@Ring.char_ge F Feq Fzero Fone Fopp Fadd Fsub Fmul 3}.
Notation "' $\infty$ '" := unit : type_scope. Notation "' $\infty$ '" := (inr tt) : core_scope.
Notation "( x , y )" := (inl (pair x y)). Open Scope core_scope.

Definition point := { P | match P with (x, y) ⇒ y2 = x3 + a*x + b |  $\infty$  ⇒ T end }.
Definition coordinates (P:point) : (F*F +  $\infty$ ) := proj1_sig P.
Program Definition zero : point :=  $\infty$ .
Program Definition add (P1 P2:point) : point :=
  match coordinates P1, coordinates P2 return F*F +  $\infty$  with
  | (x1, y1), (x2, y2) ⇒
    if x1 =? x2
    then
      if y2 =? -y1
      then  $\infty$ 
      else let k := (3*x12+a)/(2*y1) in
        let x3 := k2-x1-x1 in
        let y3 := k*(x1-x3)-y1 in
        (x3, y3)
    else let k := (y2-y1)/(x2-x1) in
      let x3 := k2-x1-x2 in
      let y3 := k*(x1-x3)-y1 in
      (x3, y3)
  |  $\infty$ ,  $\infty$  ⇒  $\infty$ 
  |  $\infty$ , _ ⇒ coordinates P2
  | _,  $\infty$  ⇒ coordinates P1
end.
Fixpoint mul (n: $\mathbb{N}$ ) (P : point) : point :=
  match n with 0 ⇒ zero | S n' ⇒ add P (mul n' P) end.
End WeierstrassCurves.

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