

Computational Higher Inductive Types

Computing with Custom Equalities

Jason Gross
`jgross@mit.edu`

MIT CSAIL Student Workshop

August 28, 2021

Properties of Equality

Warm Up: Linked Lists

Example: Unordered Sets
Canonical Inhabitants
Higher Inductive Types

Computing with Higher Inductive Types

Thank you

Properties of Equality

Properties of Equality

- ▶ Reflexivity: $x = x$

Properties of Equality

- ▶ Reflexivity: $x = x$
- ▶ Symmetry: if $x = y$ then $y = x$

Properties of Equality

- ▶ Reflexivity: $x = x$
- ▶ Symmetry: if $x = y$ then $y = x$
- ▶ Transitivity: if $x = y$ and $y = z$, then $x = z$

Properties of Equality

- ▶ Reflexivity: $x = x$
- ▶ Symmetry: if $x = y$ then $y = x$
- ▶ Transitivity: if $x = y$ and $y = z$, then $x = z$
- ▶ Leibniz rule: if $x = y$, then $f(x) = f(y)$

Warm Up: Linked Lists

- ▶ Two constructors:

Warm Up: Linked Lists

- ▶ Two constructors: `nil`, or `[]`, and `cons`

Warm Up: Linked Lists

- ▶ Two constructors: `nil`, or `[]`, and `cons`
- ▶ Two accessors on non-nil lists:

Warm Up: Linked Lists

- ▶ Two constructors: `nil`, or `[]`, and `cons`
- ▶ Two accessors on non-`nil` lists: `head` and `tail`

Warm Up: Linked Lists

- ▶ Two constructors: `nil`, or `[]`, and `cons`
- ▶ Two accessors on non-nil lists: `head` and `tail`
- ▶ Equality is defined on an element-by-element basis
 - ▶ $[] = []$
 - ▶ $[] \neq [a, \dots]$
 - ▶ $[a, \dots] \neq []$
 - ▶ $[x_0, x_1, \dots, x_n] = [y_0, y_1, \dots, y_m]$ iff $[x_1, \dots, x_n] = [y_1, \dots, y_m]$ and $x_0 = y_0$

Warm Up: Linked Lists

- ▶ Two constructors: `nil`, or `[]`, and `cons`
- ▶ Two accessors on non-`nil` lists: `head` and `tail`
- ▶ Equality is defined on an element-by-element basis
 - ▶ $[] = []$
 - ▶ $[] \neq [a, \dots]$
 - ▶ $[a, \dots] \neq []$
 - ▶ $[x_0, x_1, \dots, x_n] = [y_0, y_1, \dots, y_m]$ iff $[x_1, \dots, x_n] = [y_1, \dots, y_m]$ and $x_0 = y_0$
- ▶ Fairly easy to prove the properties of equality
 - ▶ In Coq, Agda, and Idris, you get all of these properties for free

Example: Unordered Sets

Example: Unordered Sets

- ▶ `nil`, or \emptyset

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`
- ▶ `contains`

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`
- ▶ `contains`
- ▶ Often implemented internally as a list or a tree

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`
- ▶ `contains`
- ▶ Often implemented internally as a list or a tree
- ▶ Equality is then implemented as “is one a permutation of the other?”

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`
- ▶ `contains`
- ▶ Often implemented internally as a list or a tree
- ▶ Equality is then implemented as “is one a permutation of the other?”
- ▶ Fairly easy to prove that it’s an equivalence relation

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`
- ▶ `contains`
- ▶ Often implemented internally as a list or a tree
- ▶ Equality is then implemented as “is one a permutation of the other?”
- ▶ Fairly easy to prove that it’s an equivalence relation
- ▶ Leibniz rule (if $x = y$, then $f(x) = f(y)$) is harder
- ▶ In Haskell, Agda, Coq, and Idris, the Leibniz rule is false!

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`
- ▶ `contains`
- ▶ Often implemented internally as a list or a tree
- ▶ Equality is then implemented as “is one a permutation of the other?”
- ▶ Fairly easy to prove that it’s an equivalence relation
- ▶ Leibniz rule (if $x = y$, then $f(x) = f(y)$) is harder
- ▶ In Haskell, Agda, Coq, and Idris, the Leibniz rule is false! (or at least not internally provable)

Example: Unordered Sets

- ▶ `nil`, or \emptyset
- ▶ `add`
- ▶ `remove`
- ▶ `contains`
- ▶ Often implemented internally as a list or a tree
- ▶ Equality is then implemented as “is one a permutation of the other?”
- ▶ Fairly easy to prove that it’s an equivalence relation
- ▶ Leibniz rule (if $x = y$, then $f(x) = f(y)$) is harder
- ▶ In Haskell, Agda, Coq, and Idris, the Leibniz rule is false! (or at least not internally provable)
 - ▶ The problem is that either you don’t have private fields, or you can’t make use of the fact that everything is defined in terms of your public methods.

Example: Unordered Sets

Solution 1: Canonical Inhabitants

- ▶ Give up private fields, but use element-wise equality

Example: Unordered Sets

Solution 1: Canonical Inhabitants

- ▶ Give up private fields, but use element-wise equality
- ▶ Define a type of “sorted lists without duplication”, and call them sets

Example: Unordered Sets

Solution 1: Canonical Inhabitants

- ▶ Give up private fields, but use element-wise equality
- ▶ Define a type of “sorted lists without duplication”, and call them sets
- ▶ Now we can use element-wise equality, and get Leibniz (and other properties) for free

Example: Unordered Sets

Solution 1: Canonical Inhabitants

- ▶ Give up private fields, but use element-wise equality
- ▶ Define a type of “sorted lists without duplication”, and call them sets
- ▶ Now we can use element-wise equality, and get Leibniz (and other properties) for free
- ▶ What if we don't have an ordering on the elements, only equality?

Example: Unordered Sets

Solution 1: Canonical Inhabitants

- ▶ Give up private fields, but use element-wise equality
- ▶ Define a type of “sorted lists without duplication”, and call them sets
- ▶ Now we can use element-wise equality, and get Leibniz (and other properties) for free
- ▶ What if we don't have an ordering on the elements, only equality?
- ▶ Is this really what we wanted? We asked for unordered sets, and instead made sorted lists.

Example: Unordered Sets

Solution 2: Higher Inductive Types

- ▶ Higher Inductive Types

Example: Unordered Sets

Solution 2: Higher Inductive Types

- ▶ Higher Inductive Types
- ▶ Keep the built-in equality (so we get the properties for free), but turn it into equality up to permutation

Example: Unordered Sets

Solution 2: Higher Inductive Types

- ▶ Higher Inductive Types
- ▶ Keep the built-in equality (so we get the properties for free), but turn it into equality up to permutation
- ▶ How do we get that it's an equivalence relation for free?

Example: Unordered Sets

Solution 2: Higher Inductive Types

- ▶ Higher Inductive Types
- ▶ Keep the built-in equality (so we get the properties for free), but turn it into equality up to permutation
- ▶ How do we get that it's an equivalence relation for free?
 - ▶ Take the reflexive symmetric transitive closure of the given relation

Example: Unordered Sets

Solution 2: Higher Inductive Types

- ▶ Higher Inductive Types
- ▶ Keep the built-in equality (so we get the properties for free), but turn it into equality up to permutation
- ▶ How do we get that it's an equivalence relation for free?
 - ▶ Take the reflexive symmetric transitive closure of the given relation
- ▶ How do we get Leibniz for free?

Example: Unordered Sets

Solution 2: Higher Inductive Types

- ▶ Higher Inductive Types
- ▶ Keep the built-in equality (so we get the properties for free), but turn it into equality up to permutation
- ▶ How do we get that it's an equivalence relation for free?
 - ▶ Take the reflexive symmetric transitive closure of the given relation
- ▶ How do we get Leibniz for free?
 - ▶ Require proving it each time you define a particular function
 - ▶ To define a function that deals with unordered sets, you have to simultaneously prove that your function is invariant under permutations

Computing with Higher Inductive Types

- ▶ It seems simple enough, so what's the problem?

Computing with Higher Inductive Types

- ▶ It seems simple enough, so what's the problem?
- ▶ Having higher inductive types gives you functional extensionality (if $f(x) = g(x)$ for all x , then $f = g$), which doesn't yet have a good computational interpretation in Coq nor Agda nor Idris

Computing with Higher Inductive Types

- ▶ It seems simple enough, so what's the problem?
- ▶ Having higher inductive types gives you functional extensionality (if $f(x) = g(x)$ for all x , then $f = g$), which doesn't yet have a good computational interpretation in Coq nor Agda nor Idris
- ▶ Equality in Coq and Agda (`--without-K`) actually has a rich structure

Computing with Higher Inductive Types

- ▶ It seems simple enough, so what's the problem?
- ▶ Having higher inductive types gives you functional extensionality (if $f(x) = g(x)$ for all x , then $f = g$), which doesn't yet have a good computational interpretation in Coq nor Agda nor Idris
- ▶ Equality in Coq and Agda (`--without-K`) actually has a rich structure
- ▶ If you look at proofs of equality, and equality of these proofs, and you iterate this process, you get enough math to do topology!

Computing with Higher Inductive Types

- ▶ It seems simple enough, so what's the problem?
- ▶ Having higher inductive types gives you functional extensionality (if $f(x) = g(x)$ for all x , then $f = g$), which doesn't yet have a good computational interpretation in Coq nor Agda nor Idris
- ▶ Equality in Coq and Agda (`--without-K`) actually has a rich structure
- ▶ If you look at proofs of equality, and equality of these proofs, and you iterate this process, you get enough math to do topology!
- ▶ This is Homotopy Type Theory

Thank you

Thanks!

Thank you

Thanks!

Questions?

Example: Unordered Sets

Solution 3: Parametricity

- ▶ Make use of the fact that private fields are private

Example: Unordered Sets

Solution 3: Parametricity

- ▶ Make use of the fact that private fields are private
- ▶ Very hard to do!

Example: Unordered Sets

Solution 3: Parametricity

- ▶ Make use of the fact that private fields are private
- ▶ Very hard to do!
- ▶ Can probably be done by way of parametricity (aka “theorems for free”), or a generalization of it

Example: Unordered Sets

Solution 3: Parametricity

- ▶ Make use of the fact that private fields are private
- ▶ Very hard to do!
- ▶ Can probably be done by way of parametricity (aka “theorems for free”), or a generalization of it
- ▶ Parametricity can be given a computational interpretation, but it’s very non-trivial to do so