osX Assembly (asm)

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If you find a bug in the osX Assembler please contact Dr. Welborn (Curtis.Welborn@uvu.edu) to have it fixed. If you have questions, please come and talk to me— I would love to help you. If you don't like the error messages, formatting of instructs, or just want additional features write your own assembler. You're a UVU Graduate student in the MCS if you are using this program and as such you should be able to duplicate this entire program in a matter of days as the original version was written in one weekend. The documentation I've provided took much longer to produce than the code itself.

Note on Language usage: The osX Assember was written using C. No complex data structures or advanced libraries were needed. I find writing an assembler and vm in C to be the simplest solution (students who have tried this agree with me). **This doesn't mean you should write your programs in C**, you should use the language you feel most comfortable using and the language you feel is best for the problem. In an actual commercial product C/C++ would be a great choice but you're not being asked to implement a commercial product you are in a classroom setting where topics on software design, development, and testing will be every bit as important as the topics on operating systems. Don't get bogged down in language issues. Focus on solving the problem, your instructor will introduce optimization concept as they see fit. Make sure your program works and that sound design and development practices have been used before wasting time worrying about efficiency and optimization.

Instruction Notation

address(label) – returns the address of a label byte(register) – returns the first byte in a register memory[address] – a location in memory

Comment

; This is a comment

Instruction Formats

All osX instructions are of fixed length (6 bytes long). The addressing mode of all operands can always be determined by checking the op code of the instruction. The <op code> of an instruction is always encoded as a single unsigned byte (e.g., a char in C). An operand which is a register is always encoded as a single unsigned byte (e.g., a char in C). Immediate operands and addresses are encoded as 4 byte integers. There are 7 different instruction formats:

Encoding

Two data types (char, int) were used to encode the instructions. When only a single byte was required the C data type of char or more specifically unsigned char. When 4 bytes were needed a simple int (4 byte signed integer) was used. The means you only need to worry about reading two data types from the .osx file, chars and integers that is it. No nibbles, no bit arrays, nothing overly complex.

- 1. 1 byte op code | 1 byte register | 4 bytes unused
- 2. 1 byte op code | 1 byte register | 4 byte immediate
- 3. 1 byte op code | 1 byte register | 4 byte address
- 4. 1 byte op code | 4 byte address | 1 byte unused
- 5. 1 byte op code | 4 byte immediate | 1 byte unused
- 6. 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused
- 7. 1 byte op code | 1 byte register | 1 byte register | 2 bytes unused

Arithmetic

[< label>] ADD < reg1> < reg2> < reg3> ; Add

Encode: 1 byte op code | 1 byte register | 1 byte register | 1 byte register | 2 bytes unused ADD = 16

Suggested Usage: <reg1> ← < reg2> + < reg3>

$[\langle label \rangle]$ SUB $\langle reg1 \rangle \langle reg2 \rangle \langle reg3 \rangle$; Subtract

Encode: 1 byte op code | 1 byte register | 1 byte register | 1 byte register | 2 bytes unused SUB = 17

[< label>] MUL < reg1> < reg2> < reg3> ; Multiple

Encode: 1 byte op code | 1 byte register | 1 byte register | 1 byte register | 2 bytes unused MUL = 18

Suggested Usage: <reg1> ← < reg2> * < reg3>

[<label>] DIV < reg1> < reg2> < reg3> ; Divide

Encode: 1 byte op code | 1 byte register | 1 byte register | 1 byte register | 2 bytes unused DIV = 19

Suggested Usage: <reg1> ← < reg2> / < reg3>

Move Data

[<label>] MOV <reg1> < reg2> ; Load data from register

Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused MOV = 1

Suggested Usage: <reg1> ← < reg2>

[<label>] MVI <reg1> <imm>; Load register with immediate value

Encode: 1 byte op code | 1 byte register | 4 byte immediate

MVI = 22

Suggested Usage: <reg1> ← <imm>

[<label>] ADR <reg1> <label> ; Get address of label

Encode: 1 byte op code | 1 byte register | 4 byte address ADR = 0

Suggested Usage: <reg1> ← address(<label>)

$\begin{tabular}{l} $\tt STR < \tt reg1 > \tt reg2 > \tt ; Store word (int) using register indirect addressing \\ \end{tabular}$

Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused STR=2

Suggested Usage: memory[<reg2>] ← <reg1>

[<label>] STRB <reg1> < reg2> ; Store byte using register indirect addressing

Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused STRB=3

Suggested Usage: memory[< reg2>] ← byte(memory[<reg1>])

[<label>] LDR <reg1> < reg2> ; Load word (int) using register indirect addressing

Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused LDR=4

Suggested Usage: <reg1> ← memory[<reg2>]

[<label>] LDRB <reg1> < reg2> ; Load byte using register indirect addressing

Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused LDRB=5

Suggested Usage: < reg1> ← byte(memory[<reg2>]

Branch

[<label>] B <label>; Jump to label

Encode: 1 byte op code | 4 byte address | 1 byte unused B = 7

Suggested Usage: PC ← address(<label>)

[<label>] BL <label>; Jump to label and link (place the PC of the next instruction into a register)

Encode: 1 byte op code | 4 byte address | 1 byte unused BL = 21

Suggested Usage: PC \leftarrow address(<|abel>) R5 \leftarrow PC+6

[<label>] BX < reg>; Jump to address in register

Encode: 1 byte op code | 1 byte register | 4 bytes unused BX = 6

Suggested Usage: PC ← <reg>

[<label>] BNE <label>; Jump to label if Z register is not zero

Encode: 1 byte op code | 4 byte address | 1 byte unused

Suggested Usage: PC ← address(<label>) if Z != 0

[<label>] BGT <label>; Jump to label if Z register is greater than zero

Encode: 1 byte op code | 4 byte address | 1 byte unused BGT = 9

Suggested Usage: PC \leftarrow address(<|abel>) if Z > 0

[<label>] BLT <label>; Jump to label if Z register is less than zero

Encode: 1 byte op code | 4 byte address | 1 byte unused BLT = 10

BLT = 10

Suggested Usage: $PC \leftarrow address(\langle label \rangle)$ if Z < 0

[<label>] BEQ <label>; Jump to label if Z register is zero

Encode: 1 byte op code | 4 byte address | 1 byte unused BEQ = 11

Suggested Usage: $PC \leftarrow address(\langle label \rangle)$ if Z == 0

Logical

[<label>] CMP < reg1> < reg2> ; Compare < reg1> and < reg2> place result in Z register

Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused CMP = 12

Suggested Usage: Z ← < reg1> - < reg2>

[<label>] AND < reg1> < reg2> ; Perform an AND operation on <reg1> and <reg2> result in Z register

Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused AND = 13

Note: The AND operation can be logical or bitwise either will work

Suggested Usage: Z ← < reg1> & < reg2> bitwise AND operation

[<label>] ORR < reg1> < reg2> ; Perform an OR operation on <reg1> and <reg2> result in Z register Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused ORR = 14

Note: The ORR operation can be logical or bitwise either will work

Suggested Usage: Z ← < reg1> | < reg2> bitwise OR operation

Suggested Usage: Z ← <reg1> | | <reg2> logical OR operation

[<label>] EOR < reg1> < reg2> ; Exclusive OR operation on <reg1> and <reg2> result in Z register Encode: 1 byte op code | 1 byte register | 1 byte register | 3 bytes unused EOR = 15

Note: The EOR operation can be logical or bitwise, there is little actual need for this instruction

Suggested Usage: Z ←< reg1> ^ < reg2> bitwise Exclusive OR operation

Suggested Usage: Z ← (<reg1> | | <reg2>) && !(<reg1> && <reg2>)

Interrupts

SWI <imm>; **Software Interrupt**

Encode: 1 byte op code | 4 byte immediate | 1 byte unused SWI = 20

Suggested Usage: Execute interrupt <imm>

Directives

[<label>] .WORD <integer>; Allocate space for a positive or negative 4 byte integer

[<label>] .BYTE '<char>'; Allocate space for printable ASCII character

[<label>] .BYTE \<value>; Allocate space for a byte value 0...255

[<label>] .SPACE <value> ; Allocate 1...1024 bytes of space

Registers

R0; General Purpose Register, Encode as 0

Suggested Usage: General purpose Register

R1; General Purpose Register, Encode as 1

Suggested Usage: general purpose Register

R2; General Purpose Register, Encode as 2

Suggested Usage: general purpose Register

R3; General Purpose Register, Encode as 3

Suggested Usage: general purpose Register

R4; General Purpose Register, Encode as 4

Suggested Usage: general purpose Register

R5; General Purpose Register, Encode as 5

Suggested Usage: general purpose Register

SP; Space Pointer Register, Encode as 6

Suggested Usage: Point at top of run-time stack, used with SL register to test for Stack-Overflow and Out-Of-Memory.

FP; Frame Pointer Register, Encoded as 7

Suggested Usage: Point at current frame on run-time stack

SL; Stack Limit Register, Encode as 8

Suggested Usage: Limit the size of the run-time stack, set this register in your VM to the byteSize value stored as the first value in the osX byte file. Can be used to test for Stack-Overflow and Out-Of-Memory

Z; Flag Register, Encode as 9

Suggested Usage: Set when CMP instruction is run, used by conditional branch instructions

SB; Stack Base Register, Encode as 10

Suggested Usage: Limit the size of memory in your VM, set this register in your VM to the maximum size of memory. Can be used to test for Stack-Underflow.

PC; Program Counter, Encode as 11

Suggested Usage: Point to next instruction to execute. Never directly set this register using assembly code.

osX byte code file

The osX Assembler will read an osX .asm file and convert this into byte code. The osX byte code file contains a three integer header before the first line of assembly code in encoded.

- Integer One Is the size of the byte code file, minus the three integer header.
- Integer Two Is the location of the first instruction found in the assembly file. You should set the PC of your VM to this value. As before this value doesn't account for the three integer header.
- Integer Three Is a loader address you can specify at the command-line of the osX Assembler.

 Your professor will explain this value and how to use as it; not really related to the assembler or vm directly, it just provide you and your instructor some flexibility when loading byte code.

NOTE: It is suggested you write your osX Assembly code as directives followed by instructions, the osX Assembler will convert your even if you write it instructions followed by directives. You should never write assembly code with instructions and directives intermixed.