

MATHS 315 Assignment 4

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(a)

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|----|---|--------------------|
| 1. | $(\forall x \neg Ax \rightarrow \neg Ax)$ | K4 |
| 2. | $((\forall x \neg Ax \rightarrow \neg Ax) \rightarrow (Ax \rightarrow \neg \forall x \neg Ax))$ | Tautology Instance |
| 3. | Ax | Hypothesis |
| 4. | $(Ax \rightarrow \neg \forall x \neg Ax)$ | 1,2 Modus Ponens |
| 5. | $\neg \forall x \neg Ax$ | 3,4 Modus Ponens |
| 6. | $\exists x Ax$ | 5, Abbreviation |

(b)

We have shown that $\{Ax\} \vdash_{K_L} \exists x Ax$. Therefore we may add a hypothesis Bx such that $\{Ax, Bx\} \vdash_{K_L} \exists x Ax$. As x was not generalised over as a free variable (no generalisation was used at all in the derivation), we have by the deduction theorem that $\{Ax\} \vdash_{K_L} (Bx \rightarrow \exists x Ax)$

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As $\Sigma \cup \{(A \vee (B \wedge C))\}$ is a consistent set of closed wffs, by theorem 4.20, there is some model \mathcal{I} such that $\mathcal{I} \models D$ for every $D \in \Sigma \cup \{(A \vee (B \wedge C))\}$, by definition 4.10.

Therefore the model $\mathcal{I} \models (A \vee (B \wedge C))$, this being an abbreviation for $\mathcal{I} \models_v (\neg A \rightarrow (B \wedge C))$ for every \mathcal{I} -assignment v . By Definition 3.10, 5, we have that $\mathcal{I} \not\models_v \neg A$ or $\mathcal{I} \models_v (B \wedge C)$

iff $\mathcal{I} \models_v A$, due to A being derived from a closed, consistent set of wffs, or $\mathcal{I} \models_v B$ and $\mathcal{I} \models_v C$.

So either \mathcal{I} is a model for $\Sigma \cup \{A\}$ or $\Sigma \cup \{B\}$.

As $\Sigma \cup \{A\}$ has a model, and is a closed set of wffs, it is consistent, by proposition 4.18. Conversely, the same is true of $\Sigma \cup \{B\}$. So either $\Sigma \cup \{A\}$ or $\Sigma \cup \{B\}$ is consistent.

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(a)

- G1. $\forall x(Px \vee Lx)$
G2. $\neg \exists x(Px \wedge Lx)$

(b)

- G3. $\exists x Px$
G4. $\exists y Ly$

(c)

- G5. $(((((Ll_1 \wedge Ll_2) \wedge \neg El_1 l_2) \wedge Pp_1) \wedge Pp_2) \wedge Il_1 p_1) \wedge Il_1 p_2) \wedge Il_2 p_1) \wedge Il_2 p_2) \rightarrow (Ep_1 p_2))$

(d)

- G6. $((Pp_1 \wedge Ll_2) \wedge \neg Ep_1 p_2) \rightarrow \exists l((Ll \wedge Ilp_1) \wedge Ilp_2))$

(e)

If \mathcal{I} is a normal model of S with domain D and p_1 and p_2 are points of \mathcal{I} (in other words, $p_1, p_2 \in D$ and $P^{\mathcal{I}}p_1, P^{\mathcal{I}}p_2$ both hold and $p_1 \neq p_2$) then there is a unique $l \in D$ such that $L^{\mathcal{I}}, I^{\mathcal{I}}lp_1$ and $I^{\mathcal{I}}lp_2$ all hold.

As \mathcal{I} is a normal model, it must satisfy axioms G1 - G6, as well as K1 - K5 and E1 - E3, with E interpreted as equality.

Given two unique points p_1 and p_2 , by axiom G6 we must have under the normal model at least one line l incident with both points, such that $L^{\mathcal{I}}, I^{\mathcal{I}}lp_1$ and $I^{\mathcal{I}}lp_2$ all hold.