

# Advanced Union-Find

Algorithms: Design and Analysis, Part II

The Ackermann Function

## Tarjan's Bound

Theorem: [Tarjan 75] With Union by Rank and path compression, m UNION+FIND operations take  $O(m\alpha(n))$  time, where  $\alpha(n)$  is the inverse Ackerman function (will define in this video)

Proof in next video.

#### The Ackermann Function

Aside: Many different definitions, all more or less equivalent.

Will define  $A_k(r)$  for all integers k and  $r \ge 1$ . (recursively)

Base case: 
$$A_0(r) = r + 1$$
 for all  $r \ge 1$ .

In general: For  $k, r \ge 1$ :

$$A_k(r) = \text{Apply } A_{k-1} \ r \text{ times to } r$$

$$= (A_{k-1} \circ A_{k-1} \circ \dots \circ A_{k-1})(r)$$

*r*-fold composition

#### Quiz: $A_1$

Quiz:  $A_1(r)$  corresponds to what function of r?

- A) Successor  $(r \mapsto r + 1)$
- B) Doubling  $(r \mapsto 2r)$
- C) Exponentation  $(r \mapsto 2^r)$ D) Tower function  $(r \mapsto 2^{2\cdots r \text{ times } \dots^2})$

$$A_1(r) = (A_0 \circ A_0 \circ \ldots \circ A_0)(r) = 2r$$
  
(r-fold composition, add 1 each time)

## Quiz: A<sub>2</sub>

Quiz: What function does  $A_2(r)$  correspond to?

A) 
$$r \mapsto 4r$$

B) 
$$r \mapsto 2^r$$

B) 
$$r \mapsto r2^r$$

D) 
$$r \mapsto 2^{2\cdots r \text{ times } \dots^2}$$

$$A_2(r) = (A_1 \circ A_1 \circ \ldots \circ A_1)(r) = r2^r$$

(r-fold composition, doubles each time)

## Quiz: $A_3$

Quiz: What is  $A_3(2)$ ? Recall  $A_2(r) = r2^r$ 

- A) 8
- B) 1024
- B) 2048 «
- D) Bigger than 2048

$$A_3(2) = A_2(A_2(2)) = A_2(8) = 82^8 = 2^{11} = 2048$$

In general:  $A_3(r) = (A_2 \circ A_2 \circ \dots (r \text{ times}) \dots \circ A_2)(r) \ge a$  tower of r 2's =  $2^{2 \dots r \text{ times} \dots^2}$ 

#### $A_4$

$$A_4(2) = A_3(A_3(2)) = A_3(2048) \geq 2^{2 \cdot \cdot \cdot \cdot \text{ height 2048 } \cdot \cdot \cdot \cdot^2}$$

In general:  $A_4(r) = (A_3 \circ \dots r \text{ times } \dots \circ A_3)(r) \approx \text{iterated tower function (aka "wowzer" function)}$ 

#### The Inverse Ackermann Function

Definition: For every  $n \ge 4$ ,  $\alpha(n) = \min \max$  value of k such that  $A_k(2) \ge n$ .

$$\alpha(n) = 1, \ n = 4 \ (A_1(2) = 4)$$
  $\log^* n = 1, \ n = 2$   $\alpha(n) = 2, \ n = 5, \dots, 8 \ (A_2(2) = 8)$   $\log^* n = 2, \ n = 3, 4$   $\alpha(n) = 3, \ n = 9, 10, \dots, 2048$   $\log^* n = 3, \ n = 5, \dots, \underline{16}$   $\alpha(n) = 4, \ n \text{ up to roughly a tower}$   $\log^* n = 4, \ n = 17, \underline{65536}$  of 2's of height 2048  $\log^* n = 5, \ n = 65537, \underline{2^{65536}}$   $\alpha(n) = 5 \text{ for } n \text{ up to } ???$   $\log^* n = 2048 \text{ for such } n$