暑期課程 基本影像處理 day6-2

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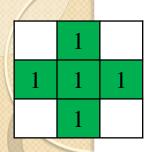
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2015.07.08

Course Outline

- 型態影像學
 - 。Dilation (膨脹)
 - · Erosion (侵蝕)
 - Opening and Closing
- Connected Component Labeling (連通元件標記法)
 - Flood fill
 - Two-pass



Connected Component Labeling

- 對4-connected或8-connected的foreground pixel標示相同的Label,形成連通元件(connected component)
- 。未連接則以不同label標示

1	0	0	0	0	0	0	0	1
1	1	0	0	0	0	0	0	1
1	1	0	1	1	0	0	1	1
0	0	0	1	1	1	0	0	1
0	0	0	1	1	1	0	0	0
1	1	0	0	1	1	0	1	1
0	1	1	0	0	1	1	0	1
0	1	1	0	0	0	0	0	1

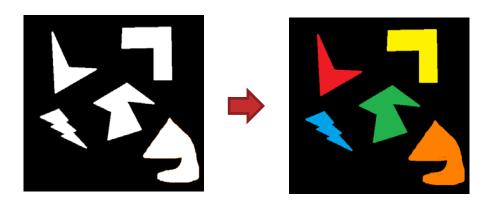


1	0	0	0	0	0	0	0	2
1	1	0	0	0	0	0	0	2
1	1	0	3	3	0	0	2	2
0	0	0	3	3	3	0	0	2
0	0	0	3	3	3	0	0	0
4	4	0	0	3	3	0	5	5
0	4	4	0	0	3	3	0	5
0	4	4	0	0	0	0	0	5

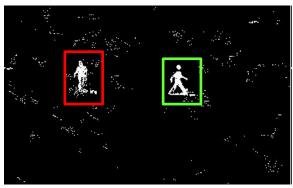
binary map

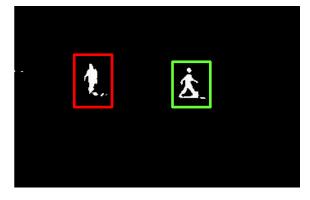
4-connected component map

• Why Connected Component Labeling?











- Flood fill
- Two-pass

Flood fill

void flood_fill(int x, int y, int color) {

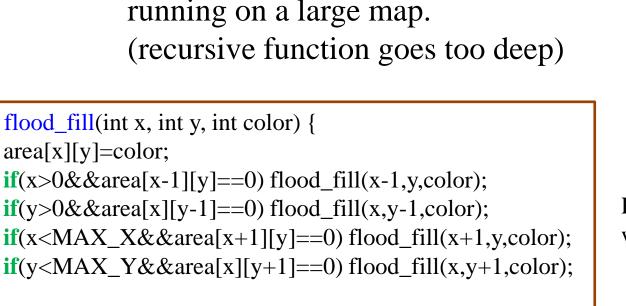
area[x][y]=color;

Recursive implementation

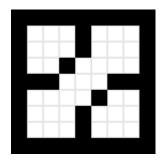
 $if(x>0&&area[x-1][y]==0) flood_fill(x-1,y,color);$

 $if(y>0\&\&area[x][y-1]==0) flood_fill(x,y-1,color);$

- drawback:
 - May cause stack overflow while running on a large map. (recursive function goes too deep)

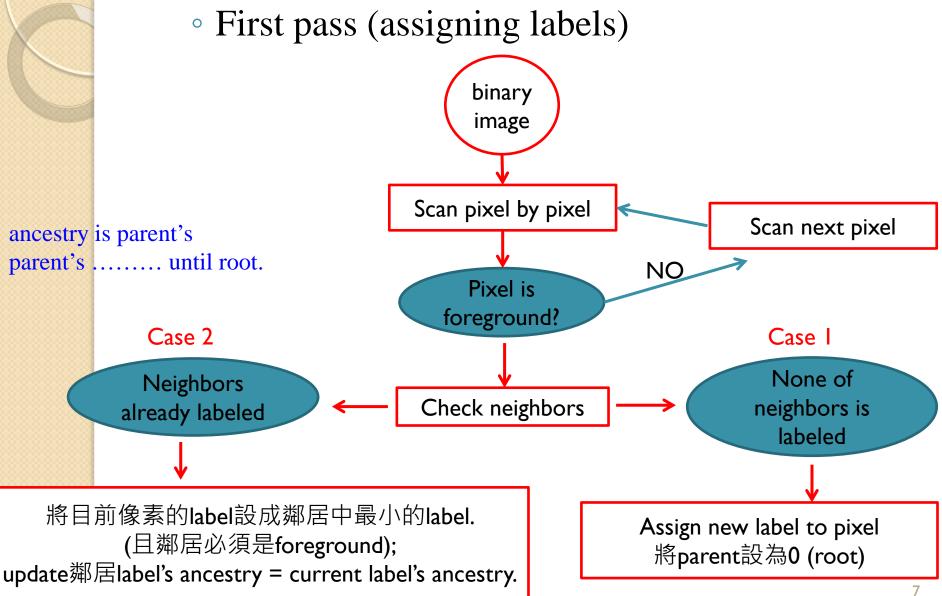


Recursive flood fill with 4 directions



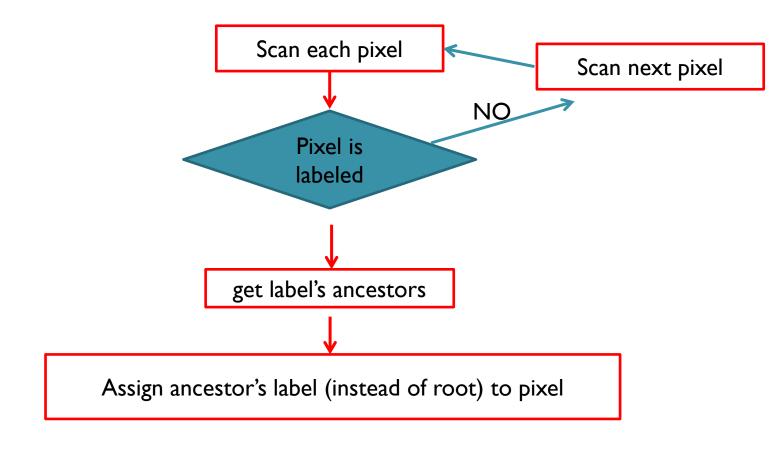
Recursive flood fill with 8 directions

Two-pass Algorithm



Two-pass Algorithm

Second pass (aggregation)



ancestry or ancestor is parent's parent's until root.



ancestor(x, parent);

//parent array

ancestor label

Find the parent label of a set.

X is a label of the set.

PARENT is the array containing the union-find data structure.

Algorithm 2: Find

Two-pass Algorithm

update_ancestor(x, y, parent)

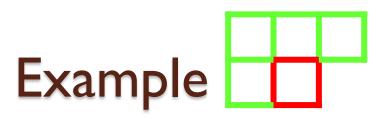
```
Construct the union of two sets.
X is the label of the first set.
Y is the label of the second set.
PARENT is the array containing the union-find data structure.
     procedure union(X, Y, PARENT);
     j := X;
     k := Y;
     while PARENT[j] <> 0
                                //find parent of set 1 until root
       j := PARENT[j];
                               //find parent of set 2 until root
     while PARENT[k] <> 0
       k := PARENT[k];
     if j <> k then PARENT[k] := j; //if(j != k) parent[k] = j;
                                       //將某一subtree接到另一個subtree之下
```

• First pass (assigning labels)

input: binary map

row-wise scanning

0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	1	1	0	Û	ĺ	1	0	0	1	1	0	0	1	1	0
0	1	1	1	1	1	1	1	1	0	0	1	1	1	1	0	0
0	0	0	1	1	1	1	0	0	0	1	1	1	1	0	0	0
0	0	1	1	1	1	0	0	0	1	1	1	0	0	1	1	0
0	1	1	1	0	0	1	1	0	0	0	1	1	1	0	0	0
0	0	1	1	0	0	0	0	0	1	1	0	0	0	1	1	0
0	0	0	0	0	0	1	1	1	1	0	0	1	1	1	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

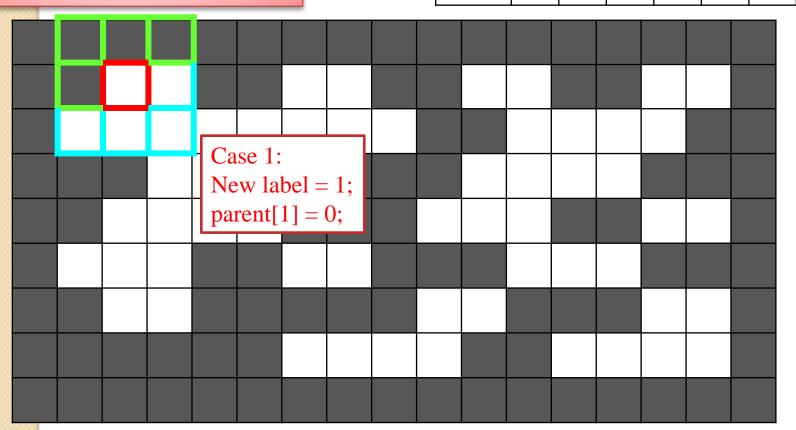


step1

8-connected

start with currentLabelCount = 1

Label				
parent				

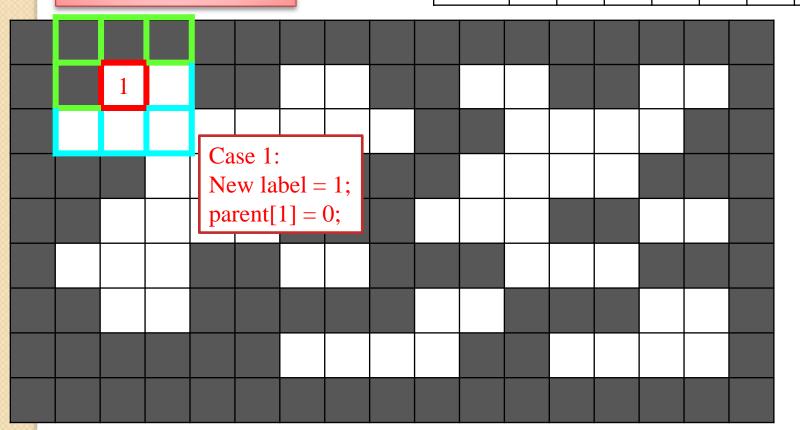


root 1

Example

step1

Label	1			
parent	0			



1

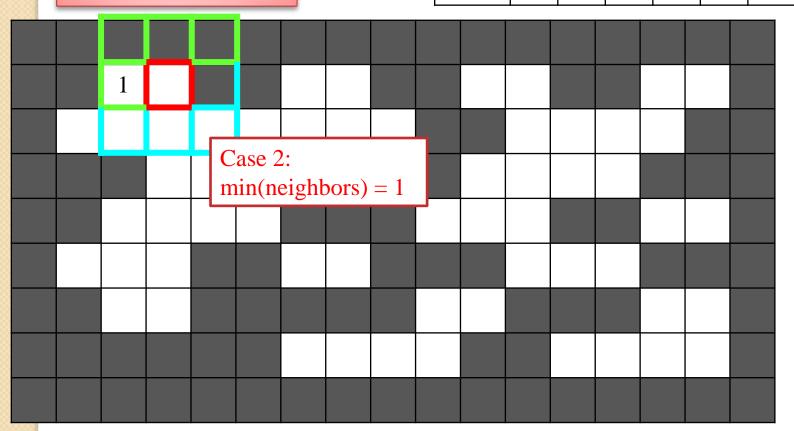
Example

step2

currentLabelCount = 1

Label	1			
parent	0			

root



root 1

Example

step2

Label	1			
parent	0			

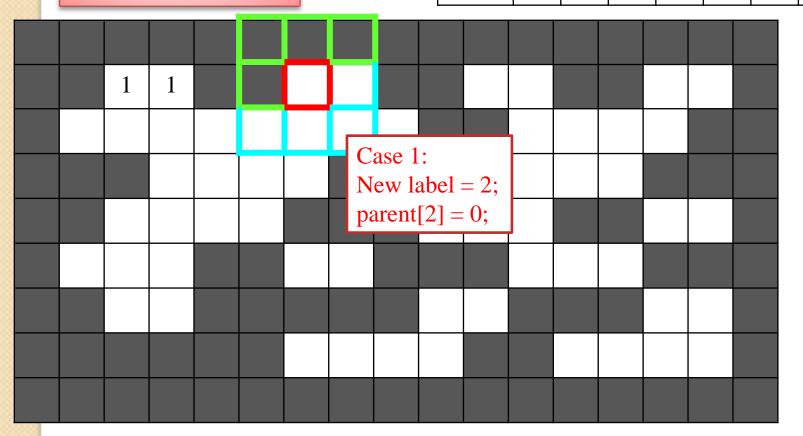


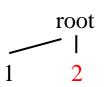
root 1

Example

step3

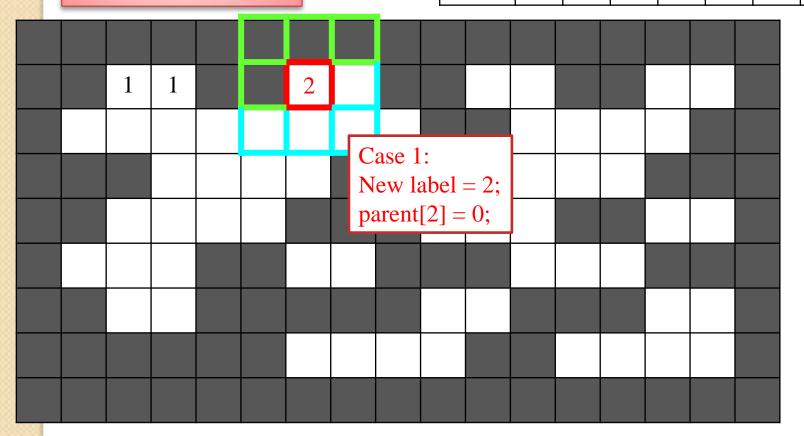
Label	1			
parent	0			

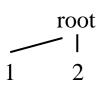




step3

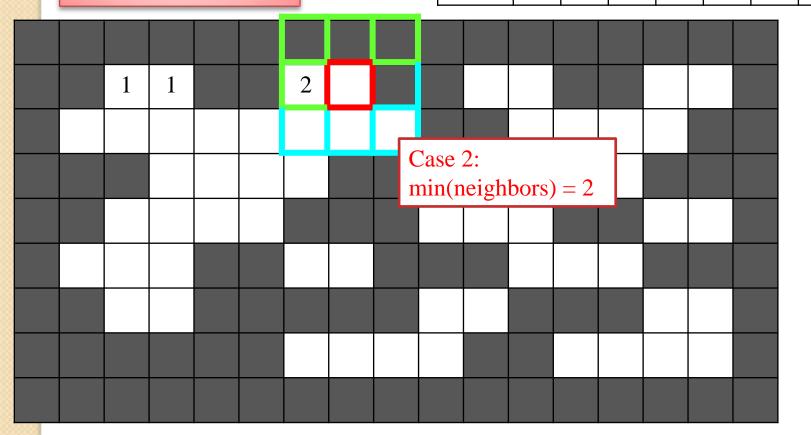
Label	1	2			
parent	0	0			

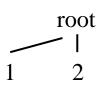




step4

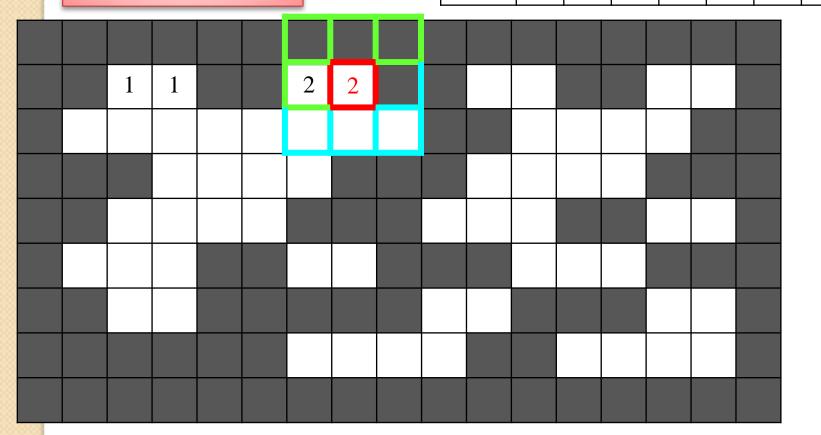
Label	1	2			
parent	0	0			

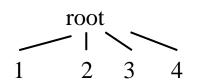




step4

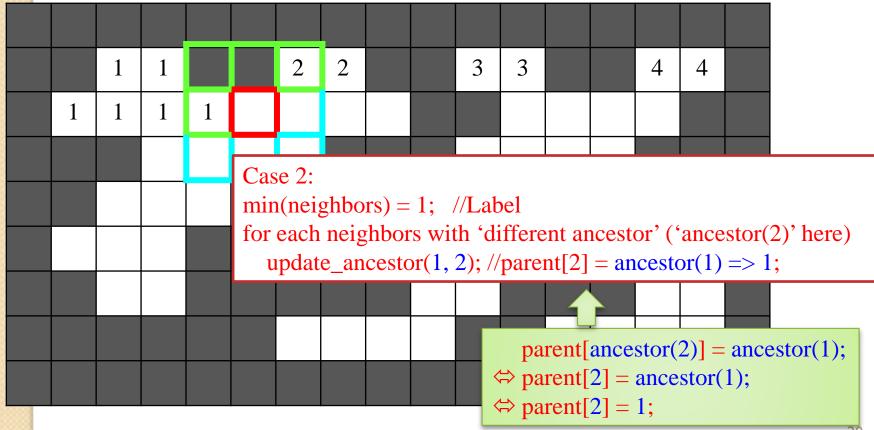
Label	1	2			
parent	0	0			





step13

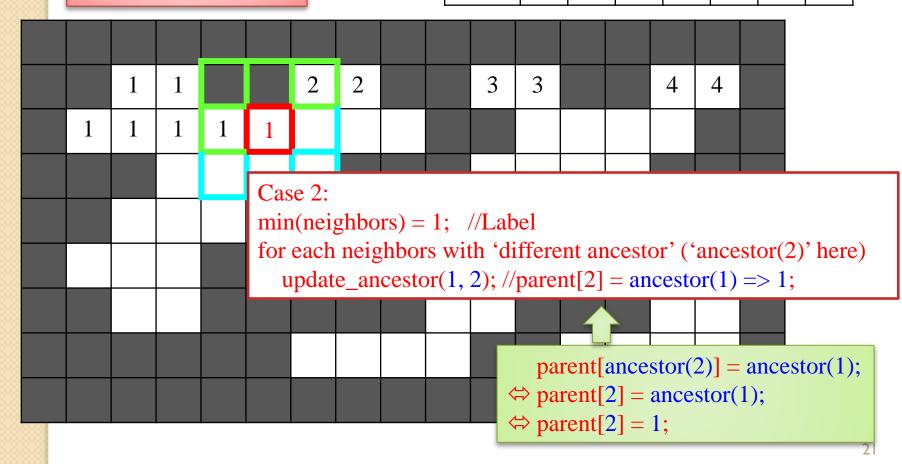
Label	1	2	3	4		
parent	0	0	0	0		



root 1 3 4 1 2

step13

Label	1	2	3	4		
parent	0	1	0	0		



step14

 Label
 1
 2
 3
 4

 parent
 0
 1
 0
 0

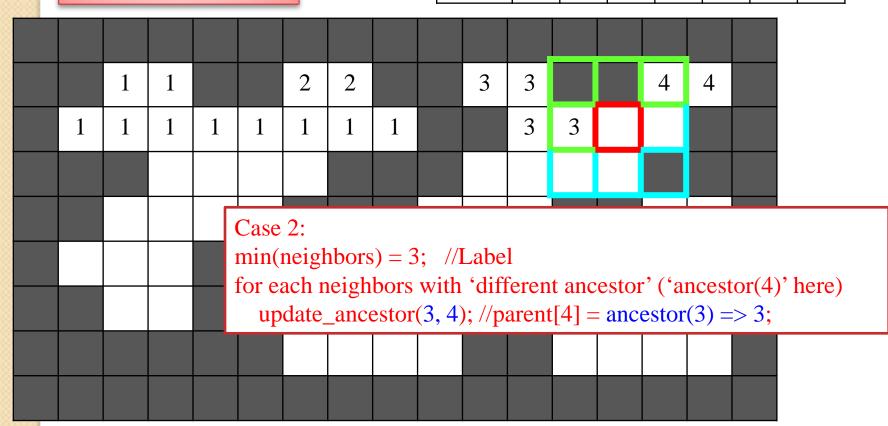
root

	1	1			2	2			3	3			4	4		
1	1	1	1	1	1											
				mir for	Case 2: min(neighbors) = 1; //Label for each neighbors with 'different ancestor' //since ancestor(2) == ancestor(1), no change;											

root
1 3 4
1
2

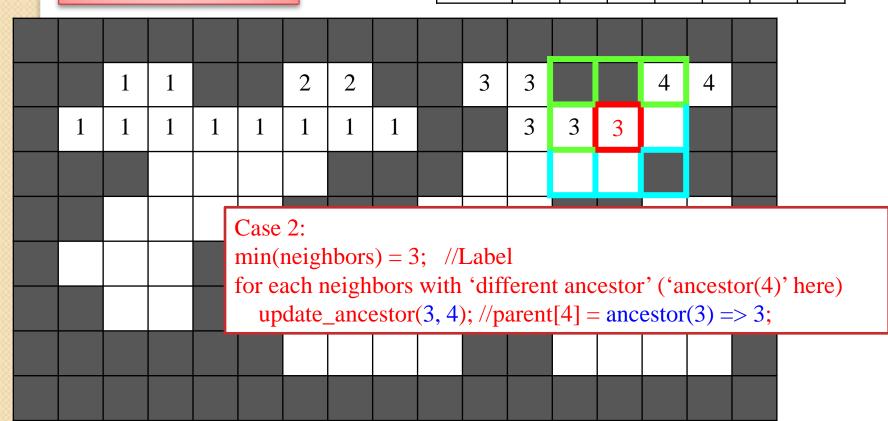
step19

Label	1	2	3	4		
parent	0	1	0	0		



step19

Label	1	2	3	4		
parent	0	1	0	3		

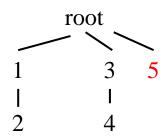


step47

Label	1	2	3	4		
parent	0	1	0	3		

	1	1			2	2			3	3			4	4	
1	1	1	1	1	1	1	1			3	3	3	3		
		1	1	1	1				3	3	3	3			
	1	1	1	1				3	3	3			3	3	
1	1	1			1	1				3	3	3			
	1	1													

step48



Label	1	2	3	4	5	
parent	0	1	0	3	0	

	1	1			2	2			3	3			4	4	
1	1	1	1	1	1	1	1			3	3	3	3		
		1	1	1	1				3	3	3	3			
	1	1	1	1				3	3	3			3	3	
1	1	1			1	1				3	3	3			
	1	1						5			1				
										Case 1: New label = 5;			5;		
										parent[5] = 0;					

step49

root

1 3 5

1 1
2 4

Label	1	2	3	4	5	
parent	0	1	0	3	0	

		1	1		Case	2:										
	1	1	1		min(_					t on o	astat	·' ('o	naast	ton(5)), homo)
			1	1		date_	_)' here)
		1	1	1	1			3	3	3			3	3		
	1	1	1			1	1			3	3	3				
		1	1					5								

step49

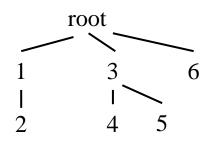
root

1 3
1 1
2 4 5

Label	1	2	3	4	5	
parent	0	1	0	3	3	

Γ																
		1	1		Case	2:										
	1	1	1		min(_					4	4	2 (6		(F)	\2.1 \
			1	1			_			teren arent)' here)
		1	1	1	1			3	3	3			3	3		
	1	1	1			1	1			3	3	3				
		1	1					5	3							

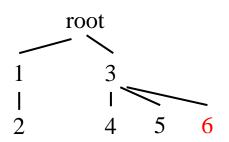
step54



Label	1	2	3	4	5	6	
parent	0	1	0	3	3	0	

	1	1			2	2			3	3			4	4			
1	1	1	1	Ca	se 2:												
		1	1					= 5;				40 0 0 0	ton? (6040	22404	(6)? h ana	`
	1	1	1	11			•				ent a ent[6]					(6)' here 3;)
1	1	1			1	1				3	3	3					
	1	1						5	7	nai	rent[s	ances	L ₂	$\begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$	ance	stor(5);	
					6	6				⇒ paı	rent[6	5] = a	ances	- T			
									¢	⇒ paı	rent[([5] = 3	3;				

step54



Label	1	2	3	4	5	6	
parent	0	1	0	3	3	3	

	1	1			2	2			3	3			4	4		
1	1	1	1	Ca	se 2:											
		1	1					= 5;					ton? (6040	22404	(6)? h ana)
	1	1	1	I I			_				ent a					(6)' here) 3;
1	1	1			1	1				3	3	3				
	1	1						5	7	nai	rent[s	ances	L ₂	$\begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$	ance	stor(5);
					6	6	5			⇒ paı	rent[6	5] = a	ances	- T		
									¢	⇒ paı	rent[([5] = 3	3;			

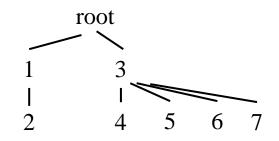
step55

root

1 3
1 1
2 4 5 6

Label	1	2	3	4	5	6	
parent	0	1	0	3	3	3	

	1	1			2	2			3	3			4	4	
1	1	1	1	Ca	se 2:										
		1	1				ors) =				ant a	n 000	ton?		
	1	1	1				ghbo cesto							ige;	
1	1	1			1	1				3	3	3			
	1	1						5	3				3	3	
					6	6	5	3							

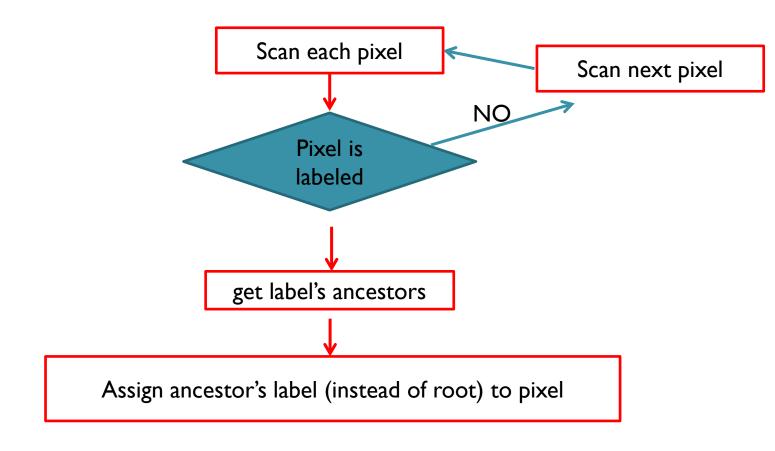


Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3

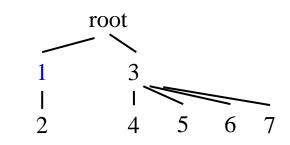
	1	1			2	2			3	3			4	4	
1	1	1	1	1	1	1	1			3	3	3	3		
		1	1	1	1				3	3	3	3			
	1	1	1	1				3	3	3			3	3	
1	1	1			1	1				3	3	3			
	1	1						5	3				3	3	
					6	6	5	3			7	3	3	3	

Two-pass Algorithm

Second pass (aggregation)



ancestry or ancestor is parent's parent's until root.

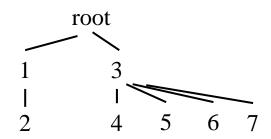


ancestry is parent's parent's until root.

step1 (Second pass)

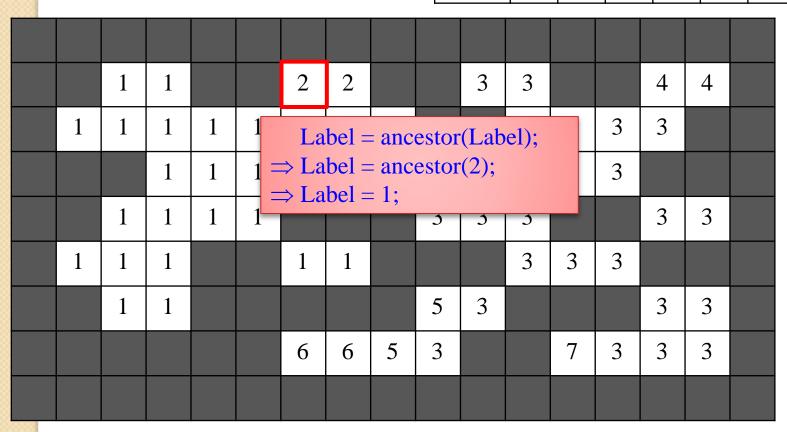
Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3

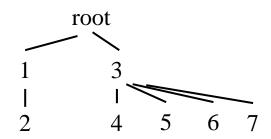
	1	1			2	2			3	3			4	4	
1		Labe	el = a	ances	stor(I	Label	l);			3	3	3	3		
		Labo		nces	stor(1);			3	3	3	3			
				1				3	3	3			3	3	
1	1	1			1	1				3	3	3			
	1	1						5	3				3	3	
					6	6	5	3			7	3	3	3	



step3

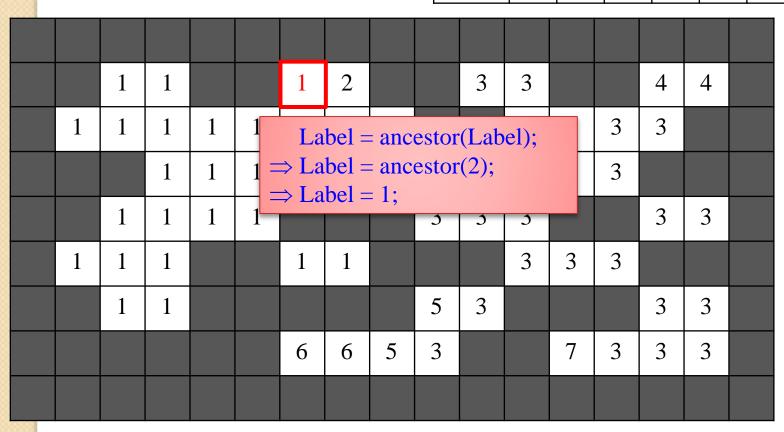
Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3

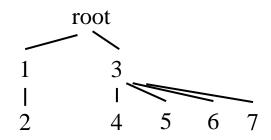




step3

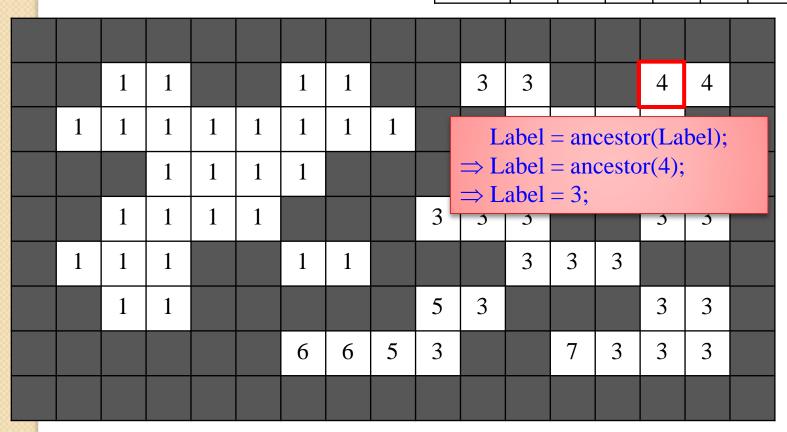
Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3

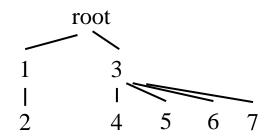




step7

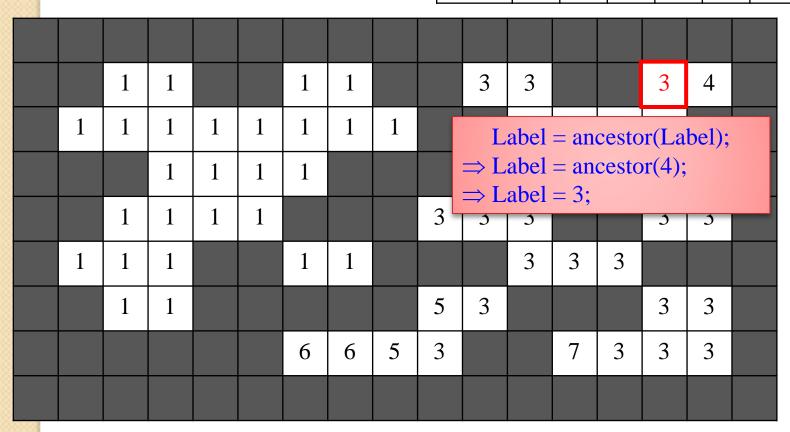
Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3





step7

Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3



root

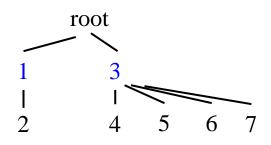
1 3
1 1
2 4 5 6 7

End process

Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3

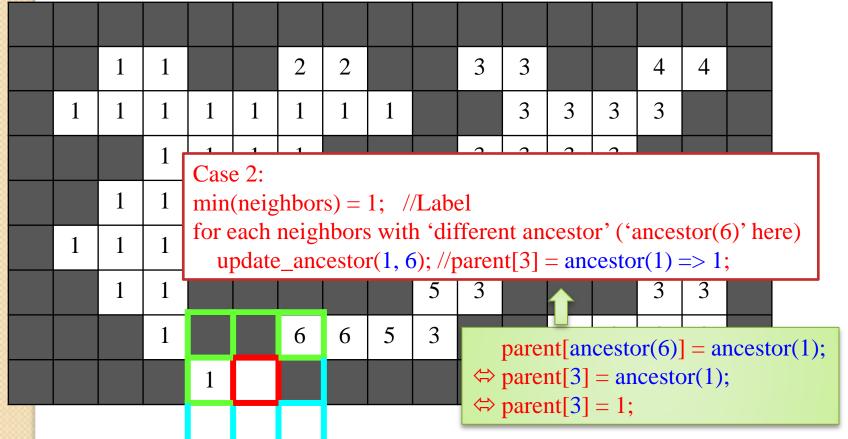
	1	1			1	1			3	3			3	3	
1	1	1	1	1	1	1	1			3	3	3	3		
		1	1	1	1				3	3	3	3			
	1	1	1	1				3	3	3			3	3	
1	1	1			1	1				3	3	3			
	1	1						3	3				3	3	
					3	3	3	3			3	3	3	3	

EXAMPLE 2



(First pass)

Label	1	2	3	4	5	6	7
parent	0	1	0	3	3	3	3

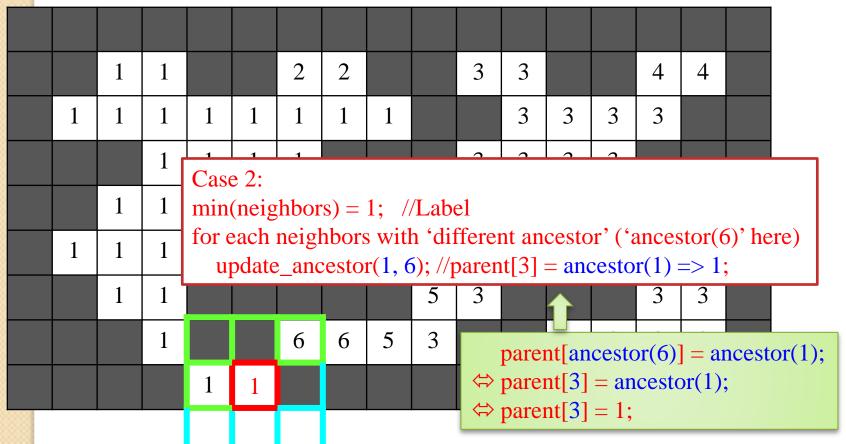


root

1
2
3
4
5
6
7

(First pass)

Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3



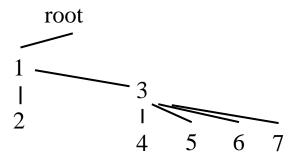
root

1
2
4
5
6
7

(First pass)

Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3

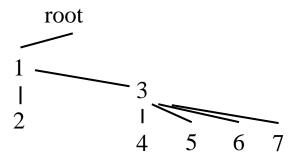
	1	1			2	2			3	3			4	4	
1	1	1	1	1	1	1	1			3	3	3	3		
		1	1	1	1				3	3	3	3			
	1	1	1	1				3	3	3			3	3	
1	1	1			1	1				3	3	3			
	1	1						5	3				3	3	
		1			6	6	5	3			7	3	3	3	
			1	1											



step1 (Second pass)

Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3

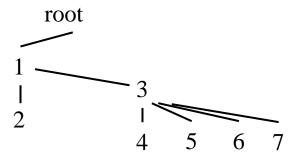
	1	1			2	2			3	3			4	4	
1	1	1	1	1	1	1	1			3	3	3	3		
		1	1	1	1				3	3	3	3			
	1	1	1	1				3	3	3			3	3	
1	1	1			1	1				3	3	3			
	1	1						5	3				3	3	
		1			6	6	5	3			7	3	3	3	
			1	1											



step5 (Second pass)

Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3

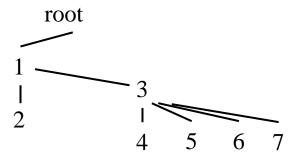
	1	1			1	1			3	3			4	4	
1	1	1	1	1	1	1	1		L	abel	= an	cesto	or(La	bel);	
		1	1	1	1				$\Rightarrow L$ $\Rightarrow L$			cesto	or(3);		
	1	1	1	1				3		3	– 1,		3	5	
1	1	1			1	1				3	3	3			
	1	1						5	3				3	3	
		1			6	6	5	3			7	3	3	3	
			1	1											



step5 (Second pass)

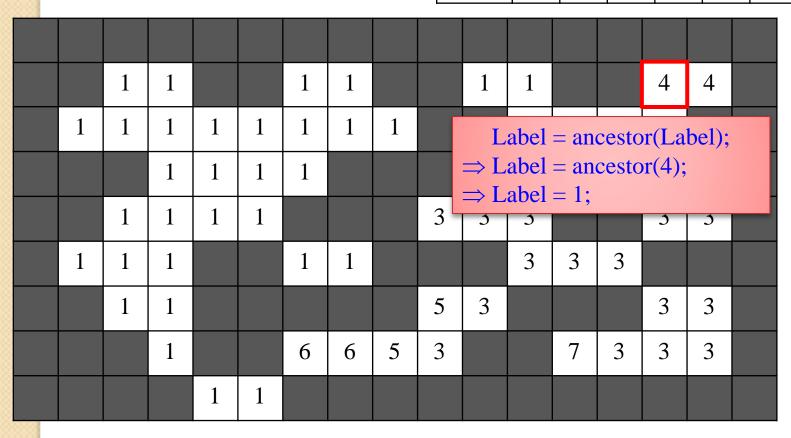
Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3

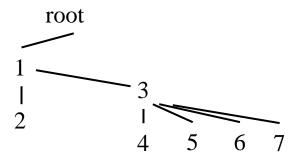
	1	1			1	1			1	3			4	4	
1	1	1	1	1	1	1	1		L	abel	= an	cesto	or(La	bel);	
		1	1	1	1				$\Rightarrow L$ $\Rightarrow L$			cesto	or(3);		
	1	1	1	1				3	3	3	– 1,		3	5	
1	1	1			1	1				3	3	3			
	1	1						5	3				3	3	
		1			6	6	5	3			7	3	3	3	
			1	1											



step7 (Second pass)

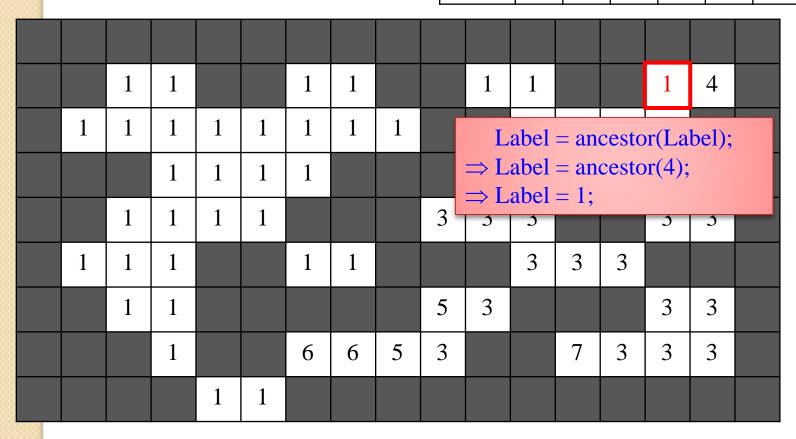
Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3

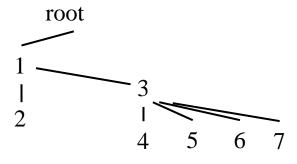




step7 (Second pass)

Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3





(Second pass)

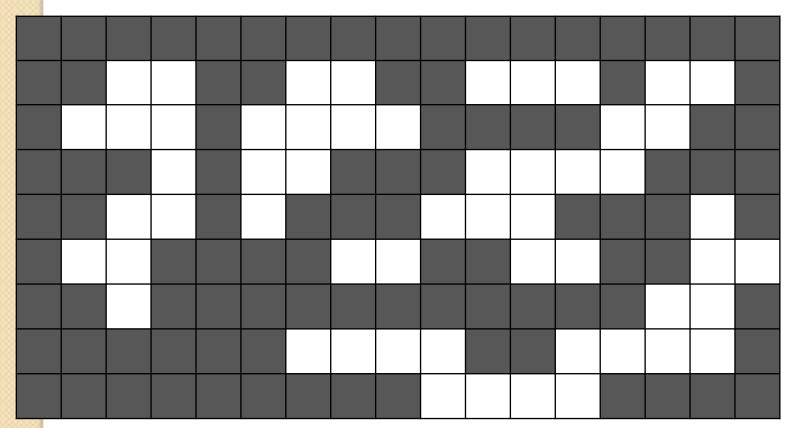
Label	1	2	3	4	5	6	7
parent	0	1	1	3	3	3	3

	1	1			1	1			1	1			1	1	
1	1	1	1	1	1	1	1			1	1	1	1		
		1	1	1	1				1	1	1	1			
	1	1	1	1				1	1	1			1	1	
1	1	1			1	1				1	1	1			
	1	1						1	1				1	1	
		1			1	1	1	1			1	1	1	1	
			1	1											

隨堂練習

• 手寫: first-pass, (8連通)

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隨堂練習

• 手寫: first-pass,

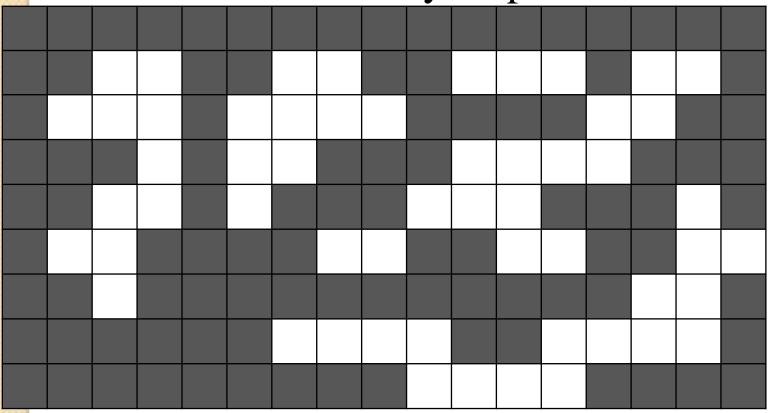
Label	1	2	3	4	5	6	7	8	9
parent	0	0	0	3	3	0	3	6	6

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	1	1		2	2			3	3	3		4	4	
1	1	1	2	2	2	2					3	3		
		1	2	2				5	5	3	3			
	1	1	2				5	5	3				6	
1	1				7	5			3	3			6	6
	1											6	6	
				8	8	8	8			9	6	6	6	
							8	8	8	6				

Homework #6.3

- 1. 以練習題目實做出Two-pass (8連通)
- 2. 讀取一張binary map並show出結果



```
algorithm TwoPass(data)
  linked = []
  labels = structure with dimensions of data, initialized with the value of Background
  First pass
  for row in data:
      for column in row:
           if data[row][column] is not Background
              neighbors = connected elements with the current element's value
              if neighbors is empty
                  linked[NextLabel] = set containing NextLabel
                  labels[row][column] = NextLabel
                  NextLabel += 1
              else
                  Find the smallest label
                  L = neighbors labels
                  labels[row][column] = min(L)
                  for label in L
                      linked[label] = union(linked[label], L)
  Second pass
  for row in data
      for column in row
          if data[row][column] is not Background
              labels[row][column] = find(labels[row][column])
  return labels
```