# CO 351: Network Flow Theory A9

Term: Spring 2019

Instructor: M. Pei

### 1. min flow max cut

## (a) FF

Modified Ford Fulkerson Algorithm:

Start with any x that is a feasible flow for a directed graph G. And form a residual digraph D' where the value for residual is  $x_e - l_e$  for forward arcs and  $c_e - x_e$  for backward arcs.

While D' has an s-t dipath:

Find s,t-dipath P in D'

Push flow r along p (-r for forward arcs and +r for backwards arcs) where r is the minimum residual in P

Update D'

## (b) Theorem

The value of a minimum flow is equal to the maximum value of  $l(\delta(S)) - c(\delta(\bar{S}))$  among all s, t cuts  $\delta(S)$ 

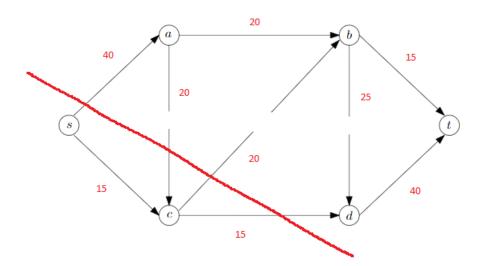
Proposition: Given a digraph G with sink and source t, s, capacities c and lower bounds l the value of any s,t flow is at least  $l(\delta(S)) - c(\delta(\bar{S}))$ 

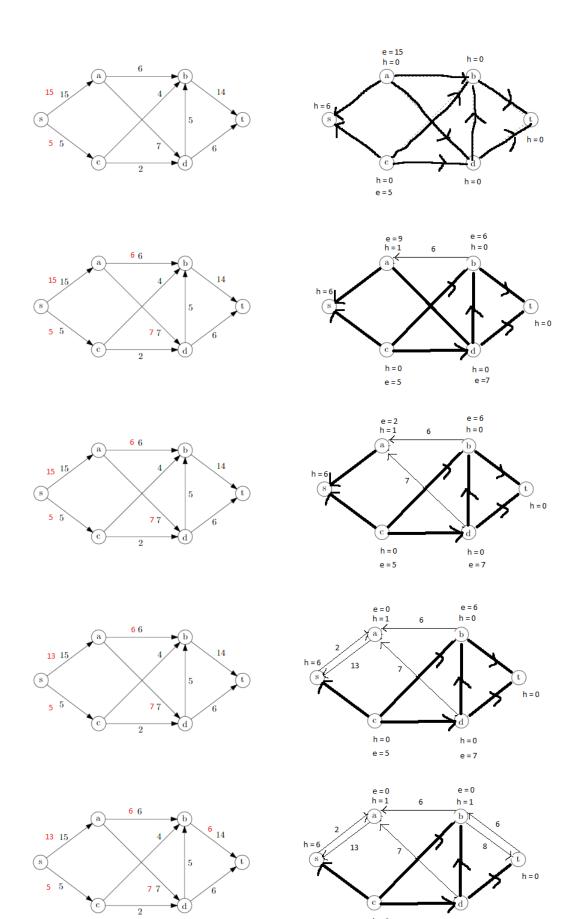
Proof: Let x be any s,t flow and  $\delta(S)$  be any s,t cut. We know that the s,t flow in an s,t cut is equal to the flow at s so  $x(\delta(s)) - x(\delta(\bar{s})) = x(\delta(S)) - x(\delta(\bar{s})) \geq l(\delta(S)) - c(\delta(\bar{s}))$ 

Proof of theorem above: Let x be a minimum flow for G. This means that in D' there is no s,t dipath remaining otherwise modified FF Alg can find a smaller flow. Then that means there exists an empty s,t cut  $\delta_{D'}(S)$ . If uv is an arc in  $\delta_D(S)$  then it is not an arc in D'(S) since  $\delta_{D'}(S)$  is empty. So then by the algorithm  $x_{uv} = l_{uv}$ . If uv is an arc in  $\delta_D(\bar{S})$  then vu is not an arc in D' and  $x_{uv} = c_{uv}$  So then the value of x is  $x(\delta(S)) - x(\delta(\bar{S})) = l(\delta(S)) - c(\delta(\bar{S}))$ . Then by our proposition we see that x is a minimu flow and  $\delta(S)$  is a cut with the maximum value of  $l(\delta(S)) - c(\delta(\bar{S}))$ 

#### (c) Sol

red numbers indicate flow on each arc. Optimal value is 55





h = 0 e = 5

h = 0 e = 7

