



Generating EVM bytecode from Yul in the new via-IR pipeline

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🎨 Yul is a thin, low-level abstraction over EVM opcodes.

🎨 But it retains more structure than a mere list of opcodes:

🎨 Control-flow Structures: `if, for, switch`

🎨 User-defined Functions: `function f(a) -> r { leave }`

🎨 Variables: `let x := 42`

🎨 No explicit jumps and no explicit stack manipulation (swaps, dups).

There is a trivial transformation from Yul to EVM assembly:

{	0x00
	calldataload
let x := calldataload(0)	
if gt(x, 0x15) {	0x15
revert(0,0)	dup2
}	gt
	iszero
}	tag_1
	jumpi
	0x00
	0x00
	revert
	tag_1:
	pop

But is this optimal?

```
solc --strict-assembly
```

```
0x00
```

```
calldataload
```

```
0x15
```

```
dup2
```

```
gt
```

```
iszero
```

```
tag_1
```

```
jumpi
```

```
0x00
```

```
0x00
```

```
revert
```

```
tag_1:
```

```
pop
```

```
solc --strict-assembly --optimize
```

```
0x15
```

```
0x00
```

```
calldataload
```

```
gt
```

```
tag_1
```

```
jumpi
```

```
stop
```

```
tag_1:
```

```
0x00
```

```
dup1
```

```
revert
```

Goals

- 🎨 Produce optimal code:
 - 🎨 Avoid unnecessary swaps and dups.
 - 🎨 Efficient arrangement of basic blocks.
- 🎨 Efficiently use the stack to avoid stack-too-deep errors.
- 🎨 Simple and robust:
 - Avoid introducing bugs due to complex transformations.

Goals

- 🎨 Produce optimal code:

- 🎨 Avoid unnecessary swaps and dups.


- 🎨 Efficient arrangement of basic blocks.

- 🎨 Efficiently use the stack to avoid stack-too-deep errors.

- 🎨 "Simple" and robust:

- Avoid introducing bugs due to complex transformations.


Approach

 Step 1 [simple]:

Generate a *Control Flow Graph*.

 Step 2 [complex]:

Algorithmically generate the full *stack layouts*
between blocks and between operations.

 Step 3 [simple]:

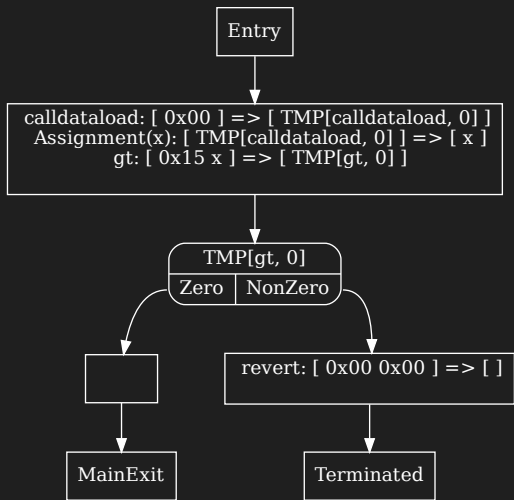
Generate concrete opcodes following the stack layouts of Step 2.
Can validate correctness relative to Step 1!

Step 1: Generate a Control Flow Graph

- Similar to naive direct opcode generation.
- Split into basic control flow blocks.
- Split basic blocks into sequences of operations with inputs and outputs.
- Simple and bug-resistant.

Step 1: Control Flow Graph

```
{  
    let x := calldataload(0)  
    if gt(x, 0x15) {  
        revert(0,0)  
    }  
}
```



Step 2: Algorithmically generate full Stack Layouts.

Generate the full stack layouts (relative to the base of the current Yul function) for entering and exiting basic control flow blocks and between operations.

- Generate stack requirements "backwards" and "bottom-up".

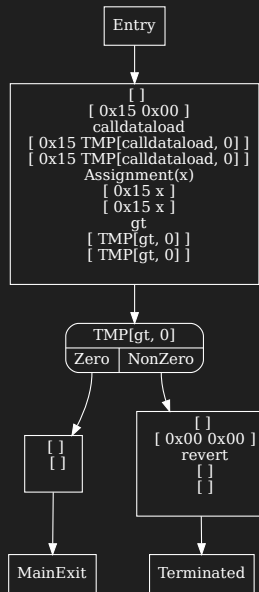
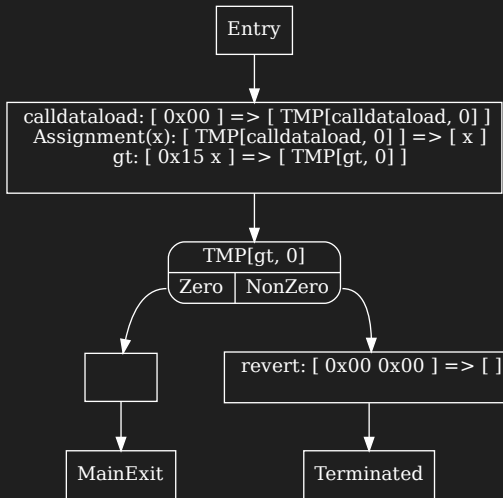
- Maintain optimal locations for variables on stack.

- Consolidate stack layouts on control flow joins.

- Avoid the necessity for access to deep stack slots.

- Complex transformation!

Step 2: Stack Layouts



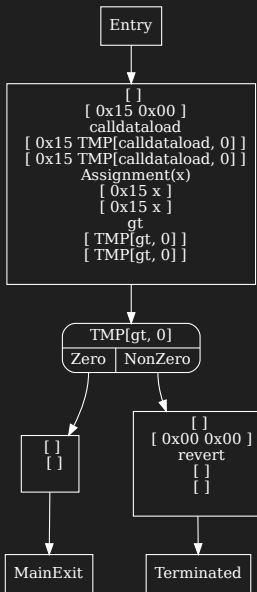
Step 2: Stack Layouts

calldataload:	[0x15 0x00]
[0x00] => [TMP[calldataload, 0]]	calldataload
Assignment(x):	[0x15 TMP[calldataload, 0]]
[TMP[calldataload, 0]] => [x]	[0x15 TMP[calldataload, 0]]
gt:	Assignment(x)
[0x15 x] => [TMP[gt, 0]]	[0x15 x]
	[0x15 x]
	gt
	[TMP[gt, 0]]
	[TMP[gt, 0]]

Step 3: Generate concrete opcodes following the stack layouts.

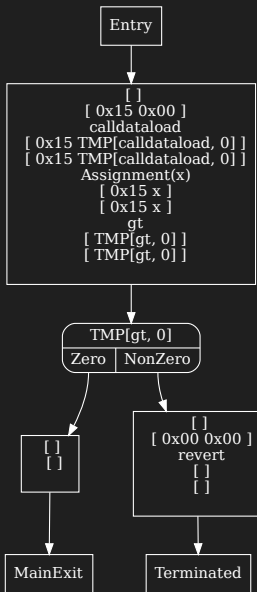
- 🎨 Maintain a symbolic copy of the current stack.
- 🎨 Traverse the Control Flow Graph.
- 🎨 Generate the stack layouts at each block transition and between all operations.
- 🎨 Validate the sanity of the current stack at each operation and each control flow transition.

Step 3: Generate EVM Assembly.

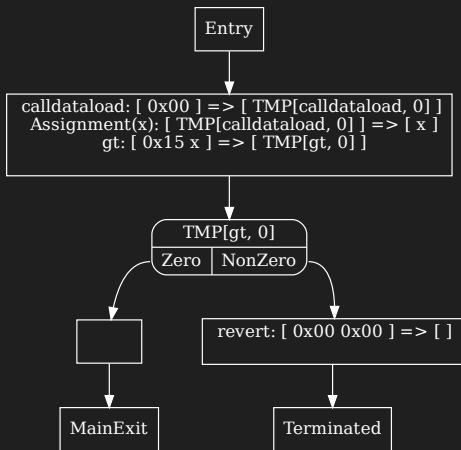


```
/* stack = [ ] */
0x15
0x00
/* stack = [ 0x15, 0x00 ]; Validate stack top */
calldataload
/* stack = [ 0x15, TMP[calldataload, 0] ] */
/* Assignment to x => stack = [ 0x15, x ] */
/* Create stack layout required for gt operation [no-op] */
gt
/* stack = [ TMP[gt, 0] ] */
/* Create block exit layout [no-op] */
/* Block exit is conditional jump; validate condition */
tag_1
jumpi
...
```

Step 3: Generate EVM Assembly.



```
...  
jumpi  
/* stack = [] */  
stop  
tag_1:  
/* stack = [] */  
/* Create stack layout for revert operation. */  
0x00  
dup1  
/* stack = [ 0x00, 0x00];  
 * Validate stack top for operation */  
revert
```



Step 3: Generate EVM Assembly.

```
/* stack = [ ] */
```

```
0x15
```

```
0x00
```

```
/* stack = [ 0x15, 0x00 ] */ ✓
```

```
calldataload
```

```
/* stack = [ 0x15, TMP[calldataload, 0] ] */ ✓
```

```
/* Assignment to x */
```

```
/* stack = [ 0x15, x ] */ ✓
```

```
gt
```

```
/* stack = [ TMP[gt, 0] ] */ ✓
```

```
tag_1
```

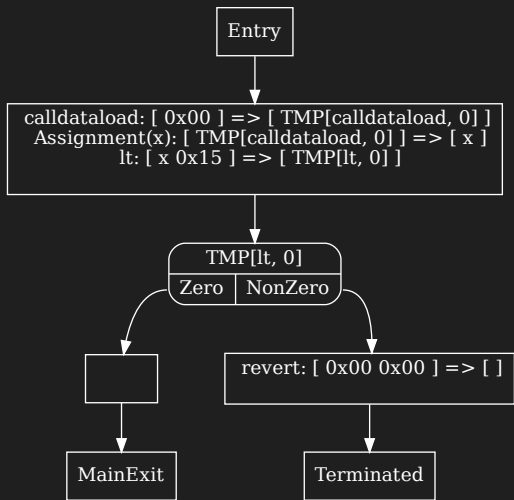
```
jumpi
```

```
...
```

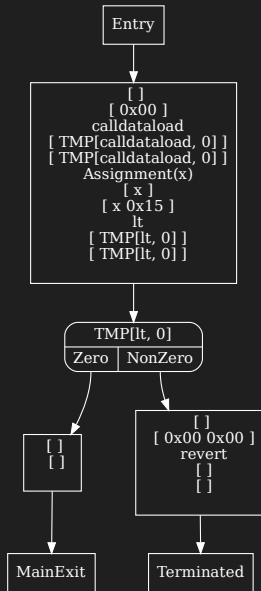
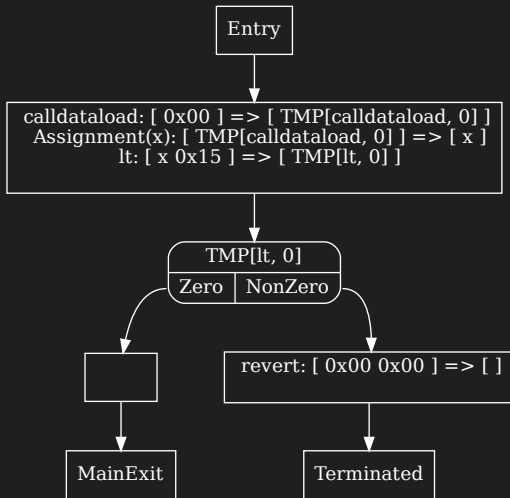

Minor changes can have non-local effects!

Step 1: Control Flow Graph

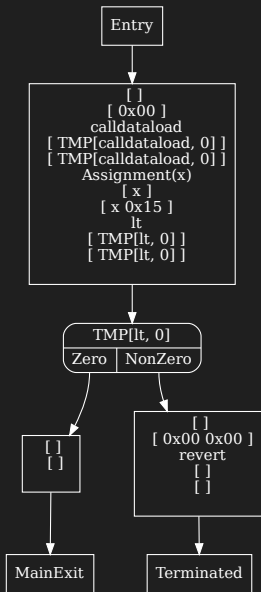
```
{  
    let x := calldataload(0)  
    // if gt(x, 0x15) {  
    if lt(0x15, x) {  
        revert(0,0)  
    }  
}
```



Step 2: Stack Layouts



Step 3: Generate EVM Assembly.



```
/* stack = [ ] */
0x00
/* stack = [ 0x00 ]; Validate stack top for operation */
calldataload
/* stack = [ TMP[calldataload, 0] ] */
/* Assignment to x => stack = [ x ] */
/* Create stack layout required for lt operation */
0x15
/* stack = [x, 0x15] */
lt
/* stack = [ TMP[lt, 0] ] */
/* Create block exit layout [no-op] */
/* Block exit is conditional jump;
   validate condition is on stack top */
tag_1
jumpi
...
```

Stack Shuffling


- 🎨 The approach requires a general stack shuffling algorithm.
- 🎨 Any stack layout may contain duplicates and arbitrary slots (JUNK).
- 🎨 Greedy algorithm that iteratively chooses the best operation.
- 🎨 Performs the shuffling on a symbolic representation of the stack.
- 🎨 Can verify correctness.
- 🎨 Can precisely report unreachable slots.


Advantages

- 🎨 Validity of the Stack Layouts can be verified.
- 🎨 Stack Layout Generation has full freedom for complex transformations.
- 🎨 The Stack Layout Generator can react to Stack-Too-Deep situations and try to resolve them.
- 🎨 If unsuccessful, the Stack Layout Generator can report a precise set of candidate variables for moving to memory.

Advantages

 Gas reduction by several percentage points.

 Reduction in code size.

 `solc --via-ir --optimize`

should (almost) never suffer from stack-too-deep errors anymore.

Disadvantages

🎨 Less predictable byte code.

🎨 But we can output the full stack layout at each proper operation as debug data!

We just need to coordinate defining a format for that.

@Tooling and @Debuggers

🌈 DO NOT try to reproduce this by hand.

🌈 DO NOT try to employ heuristics on this.

🌈 Whatever you do, it WILL break.

🌈 Let's really define a proper debugging format
instead.

Thank you!