

C Review

GWU CSCI 3411 - Fall 2019 - Lab 2
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Topics

In this lab, you will exercise some of the most basic principles of C in preparation for your work in this course. This lab is designed to remind you of C fundamentals and operations with a primary emphasis on pointers.

1 Basic Linked List

Recall the linked list data structure allows elements to be dynamically inserted or removed from a set. A basic *linked list* consists of a reference to the `head` of the list. Each element in the list may be called a *node* and each node minimally consists of a reference to the subsequent node in the list which is typically called `next`. Each node may also reference data associated with that node. The list can be traversed by iterating through the list through `next` references.

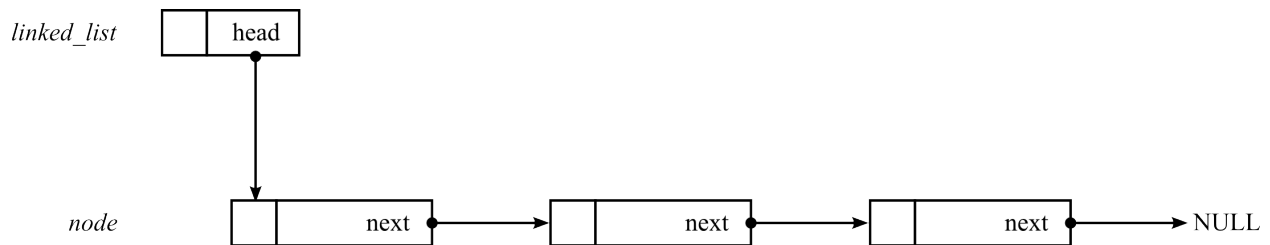


Figure 1: Basic Linked List

1.1 $O(c)$ Insertion List

The basic linked list is inefficient for insertion. In order to insert a new item into the list, the entire list must be traversed which means that in order to insert an element at the end of the list, n operations must be performed. The number of operations can be reduced by introducing a `tail` reference. The `tail` points to the last element of the list and allows an insertion to be performed in constant time by following the `next` reference on the current `tail` and then updating the `tail` to reference the item just inserted.

2 Doubly-Linked List

The basic linked list can only be traversed from front to back. A more sophisticated version of the linked list is the doubly linked list. The doubly-linked list introduces an additional pointer for each node in the list called `prev` which

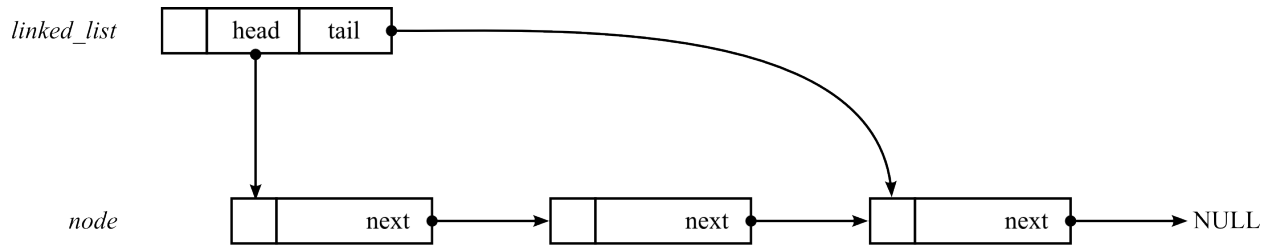


Figure 2: $O(1)$ Insertion Linked List

references the previous node. The doubly-linked list allows the list to be traversed from either the `head` or `tail` and references to all neighbors in the list can be followed or maintained from any node in the list. The drawback of the doubly-linked list is that the number of references in the list is doubled and requires more careful attention by the programmer when maintaining the list.

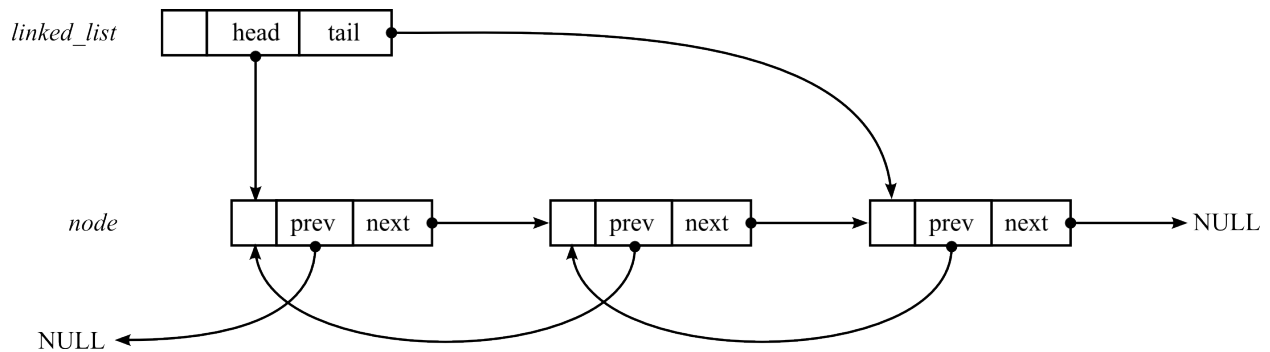


Figure 3: Doubly Linked List

3 Devices

In this lab, we are simulating a list of devices on a system. Each device will be identified by a name, a major number and a minor number. This list will be implemented as a doubly linked list. Each node of the doubly linked list will hold the device name - a char array, a major and a minor number - unsigned ints. The following methods must be implemented on the list:

1. `create` - to create the list. Return a newly created doubly linked list.
2. `insert` - insert to the end of the list. Ideally, insertion should happen by major number and then minor number - devices should be sorted by major number, all devices with the same major number should be together and within each major number devices should be sorted by minor numbers.
3. `at_index` - return -1 if the device is not found; the index otherwise. Two devices are considered to be the same if they have the same major and minor numbers.
4. `del` - delete a device from the tail of the list; alternatively build in support to delete specific devices. Return the device if successfully deleted; `NULL` otherwise. Note: Can you use `at_index` inside `del`?
5. `size` - should return the size of the list
6. `is_empty` - check if the list is empty

7. print - print the devices in the list
8. destroy - destroy the list. You can decide whether to delete only if list is completely empty or if to delete constituent data nodes while destroying. Return 0 on successfully destroying the list; 1 otherwise.