VE475 Introduction to Cryptography
Homework 4

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Ex. 1 - Euler's totient

1. Notice that for a given prime p, we have $\varphi(p) = p - 1$. So, positive integers n that is smaller than p^k, so that $\gcd(n, p^k) \neq 1$, can be $1 \times p$, $2 \times p$, $3 \times p$, $\cdots$, $(p^{k-1} - 1) \times p$. So the amount of possible n is $p^{k-1} - 1$. Also, there are $p_k - 1$ positive integers are smaller than p^k, so for any prime p, $\varphi(p^k) = (p^k - 1) - (p^{k-1} - 1) = p^{k-1}(p - 1)$.

2. Since m and n are coprime integers, according to Chinese Remainder theorem, there exists a ring isomorphism between $\mathbb{Z}/mn\mathbb{Z}$ and $\mathbb{Z}/m\mathbb{Z} \times \mathbb{Z}/n\mathbb{Z}$. We have $\varphi(mn)$ is the order of $\mathbb{Z}/mn\mathbb{Z}$, $\varphi(m)$ is the order of $\mathbb{Z}/m\mathbb{Z}$, and $\varphi(n)$ is the order of $\mathbb{Z}/n\mathbb{Z}$. Since an isomorphism is a bijection that preserves algebraic structures, we would have $\varphi(mn) = \varphi(m) \times \varphi(n)$.

3. Assume $n = \prod_i p_i^{k_i}$, then applying the previous results to integer $n > 1$, we have

$$
\begin{aligned}
\varphi(n) =& \prod_i \varphi(p_i^{k_i}) \\
=& \prod_i p_i^{k_i - 1}(p_i - 1) \\
=& \prod_i p_i^{k_i}(1 - \frac{1}{p_i}) \\
=& n \prod_{p|n} (1 - \frac{1}{p})
\end{aligned}
$$

4. The three last digits of 7^{803} can be obtained by calculating $7^{803} \mod 1000$. We note that $1000 = 2^3 \times 5^3$, thus $\varphi(1000) = 1000 \times (1 - \frac{1}{2}) \times (1 - \frac{1}{5}) = 400$ according to the previous result. So we would have

$$
\begin{aligned}
7^{803} \equiv& (7^{400})^2 \times 7^3 \mod 1000 \\
\equiv& 7^3 \mod 1000 \\
\equiv& 7^3 \mod 1000 \\
\equiv& 343 \mod 1000
\end{aligned}
$$

So, the three last digits of 7^{803} are 343.

Ex. 2 - AES

1. The key used for round 1 is given by the columns $K(4), \cdots, K(7)$. Also, recall that for $i \not\equiv 0 \mod 4$, $K(i) = K(i-4) \oplus K(i-1)$, and for $i \equiv 0 \mod 4$, $K(i) = K(i-4) \oplus T(K(i-1))$.

Ex. 3 - Simple questions

1. In mode ECB, each block is encrypted independently with a function E and a key k, so corruption of one block wouldn't influence other blocks. So the number of plaintext decrypted incorrectly is one for the ECB mode.

   In mode CBC, after the second block, XOR operation between the previous ciphertext and the current plaintext is first done before the E function and key k. So if one block is corrupted, the next block will also be influenced, thus the number of plaintext decrypted incorrectly is two for the CBC mode.

2. Since the length of block is finite, so the IV would be repeated after 2^n trials, where n is the block length. The attacker then can use whatever plaintext and compare the ciphertext generated with the same IV to find the pattern. In this way, the schemes are not CPA secure.

3. Since $p - 1 = 29 - 1 = 28 = 2 \times 2 \times 7$, so $q \in \{2, 7\}$.

   - When $q = 2$, we have
   $$2^{(29-1)/2} \equiv 2^{14} \equiv 2^4 \cdot 32^2 \equiv 2^4 \cdot 3^2 \equiv 2^2 \cdot 7 \equiv 28 \not\equiv 1 \mod 29$$
   - When $q = 7$, we have
   $$2^{(29-1)/7} \equiv 2^4 \equiv 16 \not\equiv 1 \mod 29$$

   So, 2 is a generator of $\mathsf{U}(\mathbb{Z}/29\mathbb{Z})$.

4. Using proposition from Jacobi symbol, since 1801 and 8191 are odd prime positive integers, and $1801 \equiv 1 \mod 4$, we have
$$\begin{aligned}\left(\frac{1801}{8191}\right) &= +\left(\frac{8191}{1801}\right) = +\left(\frac{987}{1801}\right) \\ &= +\left(\frac{1801}{987}\right) = +\left(\frac{814}{987}\right) = +\left(\frac{2 \times 11 \times 37}{3 \times 7 \times 47}\right) = +\left(\frac{2}{987}\right)\left(\frac{11}{987}\right)\left(\frac{37}{987}\right) \\ &= +\left(\frac{987}{11}\right)\left(\frac{987}{37}\right) = +\left(\frac{8}{11}\right)\left(\frac{25}{37}\right) = +\left(\frac{2^3}{11}\right)\left(\frac{5^2}{37}\right) = +\left(\frac{2}{11}\right)^3\left(\frac{5}{37}\right)^2 \\ &= -\left(\frac{37}{5}\right)^2 = -\left(\frac{2}{5}\right) = -1\end{aligned}$$

5. If $\left(\frac{b^2-4ac}{p}\right) = 0$, meaning $b^2 - 4ac = 0$, then the equation have one solution $x = -\frac{b}{2a}$ (more technically speaking, two same solutions). Since $-\frac{b}{2a}$ can always mod p, it's true that the number of solutions is $1 + \left(\frac{b^2-4ac}{p}\right) = 1$.

   If $\left(\frac{b^2-4ac}{p}\right) \neq 0$, meaning $b^2 - 4ac \neq 0$, then the equation have two different solutions $x_1 = \frac{-b+\sqrt{b^2-4ac}}{2a}$ and $x_2 = \frac{-b-\sqrt{b^2-4ac}}{2a}$. So we would have
$$\begin{aligned}\frac{-b \pm \sqrt{b^2-4ac}}{2a} &\equiv x \mod p \\ \sqrt{b^2-4ac} &\equiv \pm(2ax+b) \mod p\end{aligned}$$

Then, if $\left(\frac{b^2-4ac}{p}\right) = 1$, meaning $b^2 - 4ac$ is a square mod p, then it's true that the number of solutions mod p is $1 + \left(\frac{b^2-4ac}{p}\right) = 2$.

Otherwise, if $\left(\frac{b^2-4ac}{p}\right) = -1$, meaning $b^2 - 4ac$ is not a square mod p, it's true that the number of solutions mod p is $1 + \left(\frac{b^2-4ac}{p}\right) = 0$.

6. Since $\gcd(n, pq) = 1$, we have $\gcd(n, p) = 1$ and $\gcd(n, q) = 1$. Also, since $q - 1$ divides $p - 1$, we have $(p - 1) = k(q - 1)$, where k is a positive integer. So, according to Euler's theorem, we have

$$\begin{aligned} n^{p-1} \equiv& 1 \mod p \\ (n^{q-1})^k \equiv n^{p-1} \equiv& 1 \mod q \end{aligned}$$

Since $\gcd(n^{p-1}, p) = 1$ and $\gcd(n^{p-1}, q) = 1$, we can conclude that $\gcd(n^{p-1}, pq) = 1$, which is

$$n^{p-1} \equiv 1 \mod pq$$

7. - Sufficiency: if $p \equiv 1 \mod 3$, then we can obtain

$$\left(\frac{p}{3}\right) = 1$$

Also, note that p is an odd prime. If $p \equiv 1 \mod 4$

$$\left(\frac{-3}{p}\right) = \left(\frac{-1}{p}\right)\left(\frac{3}{p}\right) = 1 \cdot 1 = 1$$

If $p \equiv 3 \mod 4$

$$\left(\frac{-3}{p}\right) = \left(\frac{-1}{p}\right)\left(\frac{3}{p}\right) = (-1) \cdot (-1) = 1$$

- Necessity: We already known $\left(\frac{-3}{p}\right) = \left(\frac{-1}{p}\right)\left(\frac{3}{p}\right) = 1$, since p is an odd prime, $p \equiv 1 \mod 4$ or $p \equiv 3 \mod 4$.
If $p \equiv 1 \mod 4$, $\left(\frac{-1}{p}\right) = 1$, thus $\left(\frac{3}{p}\right) = \left(\frac{p}{3}\right) = 1$.
If $p \equiv 3 \mod 4$, $\left(\frac{-1}{p}\right) = -1$, thus $\left(\frac{3}{p}\right) = -\left(\frac{p}{3}\right) = -1$.
So, in both case $\left(\frac{p}{3}\right) = 1$, which gives $p^{(3-1)/2} \equiv p \equiv 1 \mod 3$.

8. If $\left(\frac{a}{p}\right) = 1$, we would have $a^{(p-1)/2} \equiv 1 \mod p$. However 2 is a prime factor of $p - 1$, meaning $2|(p - 1)$. So conflict with the requirement to be a generator, thus a is not a generator of $\mathbb{F}_p^*$.

Ex. 4 - Prime vs. irreducible

1.

Ex. 5 - Primitive root mod 65537

1. Using proposition of Jacobi symbol, since $65537 \equiv 1 \mod 4$ and $\gcd(3, 65537) = 1$, we have

$$\begin{aligned}\left(\frac{3}{65537}\right) &= +\left(\frac{65537}{3}\right)\\ &= +\left(\frac{2}{3}\right)\\ &= -1\end{aligned}$$

   Meaning 3 is not a square mod 65537.

2. Since 65537 is a prime integer, $\frac{65537-1}{2} = 32768$, and 3 is not a square mod 65537, we can conclude that $3^{32768} \not\equiv 1 \mod 65537$.

3. First note that $p - 1 = 65537 - 1 = 2^{16}$, thus $q = 2$ is the only prime such that $q|(p-1)$. Also, since $3^{32768} \equiv \alpha^{(p-1)/q} \not\equiv 1 \mod p$, and according to the theorem, we can conclude that $\alpha = 3$ is a primitive root mod 65537.