

# VE475 Introduction to Cryptography

## Homework 7

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### Homework 6

#### Ex. 5 - Merkle-Damgård construction

1. a) Since  $f(0) = 0$  and  $f(1) = 01$ ,  $f(x_i)$  is always start with 0. So  $y$  can be separated into several segments start from 0, except for the first two digits. Those segments are injective with  $x_i$ , so the map  $s$  from  $x$  to  $y$  is injective.  
b) If  $z$  is empty, from what previous proved, there's no such  $x'$ . If  $z$  is not empty, since we have 11 at the beginning of  $y_{i+1}$ , so no this no such  $x'$  and  $z$  such that  $s(x) = z||s(x')$ .  
2. Because the previous conditions guarantee the mapping is collision resistant.

### Homework 7

#### Ex. 1 - Cramer-Shoup cryptosystem

##### 1. Key generation:

- Alice generates a cyclic group  $G$  of order  $q$  with two distinct generators  $g_1, g_2$ .  $G$  could be  $U(\mathbb{Z}/p\mathbb{Z})$ .
- Alice chooses five random values  $(x_1, x_2, y_1, y_2, z)$  from  $\{0, 1, \dots, q-1\}$ .
- Alice computes  $c = g_1^{x_1} g_2^{x_2}$ ,  $d = g_1^{y_1} g_2^{y_2}$ , and  $h = g_1^z$ .
- Alice publishes  $(c, d, h)$  and  $G, q, g_1, g_2$  as her public key. She retains  $(x_1, x_2, y_1, y_2, z)$  as her private key.

##### Encryption:

- Bob converts plaintext into an element  $m$  in group  $G$ .
- Bob chooses a random  $k$  from  $\{0, 1, \dots, q-1\}$ , then calculates:
  - $u_1 = g_1^k, u_2 = g_2^k$ .
  - $e = h^k m$ .
  - $\alpha = H(u_1, u_2, e)$ , where  $H$  is a collision-resistant cryptographic hash function.
  - $v = c^k d^{k\alpha}$ .
- Bob sends the ciphertext  $(u_1, u_2, e, v)$  to Alice.

##### Decryption:

- Alice computes  $\alpha = H(u_1, u_2, e)$  and verifies that

2.

3.

### **Ex. 2 - Simple questions**

1.

2.

### **Ex. 3 - Birthday paradox**

1.

2.

3.

4.

### **Ex. 4 - Birthday attack**

1.

2.

3.

### **Ex. 5 - Faster multiple modular exponentiation**

1.

2.

3.

4.