Exercise 2

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1 SELECT Queries in SQL

We solve the task by using DB Browser.

1.1 a)

We run the query

 $\begin{array}{ll} \textbf{SELECT} \ \operatorname{songID} \ , \ \operatorname{name} \ , \ \operatorname{duration} \ , \ \ \textbf{year} \ , \ \ \operatorname{artistID} \\ \textbf{FROM} \ \operatorname{song} \ ; \end{array}$

and get the (cropped) output

	songID	name	duration	year	artistID	
1	1	Saved	178	2015	1	
2	2	Oops! I Did It Again	221	2000	2	
3	3	Don't Start Now	183	2019	3	
4	4	Strangers	233	2017	4	
5	5	I Went Too Far	294	2016	5	
6	6	Blasé	286	2015	1	
7	7	Hot Girl Summer	199	2019	a	

1.2 b)

We run the query

SELECT name, year FROM album WHERE year < 2017;

and get the output

	name	year
1	Free TC	2015
2	Oops! I Did It Again	2000
3	All My Demons Greeting Me as a	2016
4	SremmLife 2	2016
5	ANTI	2016
6	I Am Not a Human Being II	2013

1.3 c)

I interpret the logical condition as 2018 < year < 2021. The query is then

SELECT name, year FROM album WHERE year > 2017 AND year < 2021 ORDER BY year

which returns

	name	year
1	Scorpion	2018
2	Sucker Punch	2019
3	Fine Line	2019
4	thank u, next	2019
5	Future Nostalgia	2020
6	Positions	2020
7	Good News	2020
8	folklore	2020

1.4 d)

We need to merge such that the integrity of the relations are intact.

which returns the (cropped) output

songName	featuredName	^
Savage Remix	Beyoncé	
I'm the One	Chance the Rapper	
Love Me	Drake	
Work	Drake	
Saved	E-40	
Blasé	Future	
Don't Judge Me	Future	
Love Me	Future	
Plack Postles	Cuggi Mano	V

1.5 e)

We merge merge and extract

 $\textbf{SELECT} \ \, \textbf{song.name} \ \, \textbf{AS} \ \, \textbf{songName} \,, \ \, \textbf{album.name} \ \, \textbf{AS} \ \, \textbf{albumName} \,, \ \, \textbf{song.year} \ \, \textbf{AS} \ \, \textbf{releaseYear}$

FROM ((artist INNER JOIN song USING (artistID))

INNER JOIN songOnAlbum USING (songID))

INNER JOIN album USING(albumID)

WHERE artist.name = "Ariana-Grande"

ORDER BY song.year, album.name, song.name

the query returns

	songName	albumName	releaseYear
1	thank u, next	thank u, next	2018
2	7 rings	thank u, next	2019
3	positions	Positions	2020

1.6 f)

We treat the cases of main artist and features artist separately, join by a union clause and then order. In the second part artist takes two different roles.

SELECT artist.name AS artistName, song.name AS songName

FROM song INNER JOIN artist USING (artistID)

WHERE artist.name = "Ty-Dolla-Sign"

UNION

SELECT mainArtist.name AS artistName, song.name AS songName

FROM featuredOn

INNER JOIN song USING (songID)

INNER JOIN artist AS featuredArtist ON (featuredArtist.artistID = featuredOn.artistID)

INNER JOIN artist AS mainArtist ON (mainArtist.artistID = song.artistID)

WHERE featuredArtist.name = "Ty-Dolla-Sign"

ORDER BY artistName, songName

the query returns

	artistName	songName
1	Megan Thee Stallion	Hot Girl Summer
2	Ty Dolla Sign	Blasé
3	Ty Dolla Sign	Don't Judge Me
4	Ty Dolla Sign	Love U Better
5	Ty Dolla Sign	Saved

1.7 g)

Extract the info and check for "the" as part of the song name.

FROM song INNER JOIN artist ON (song.artistID = artist.artistID)

WHERE song.name LIKE "%the%"

the query returns

```
artistName songName

1 DJ Khaled I'm the One
```

1.8 h)

```
We group and count, and then check against the maximum of the same group and count.
```

1 Future 3

2 More Queries in Relational Algebra

2.1 a)

First we need to find out what song don't appear on an album. This will be the set difference between the songID's in song and the songID's in songOnAlbum

```
songIDnotOnAlbum = \Pi_{songID}(song) - \Pi_{songID}(songOnAlbum).
```

We can now use these ID's to filter for the desired songs with a selection-operator

```
songNotOnAlbum = \sigma_{songID \in \Pi_{songID}(songIDnotOnAlbum)}(song).
```

Now we join with artist

 $songNotOnAlbumWithArtist = artist \bowtie_{artistID} songNotOnAlbum.$

Now we project and rename

```
 artistAndSongName = \Pi_{artist.name,songNotOnAlbum.name} (songNotOnAlbumWithArtist) \\ songsWithoutAlbums = \rho_{artistAndSongName(artist.name, songNotOnAlbum.name)} (artistAndSongName).
```

2.2 b)

We first pick out the artists with songs from the oughts,

```
oughts = \Pi_{\text{artist.name, song.name}} (\sigma_{2000 \leq \text{year} \leq 2009}(\text{song}) \bowtie_{\text{artistID}} \text{artist}).
```

We select the songs from the current decade (which i guess is supposed to be the 2010s), merge with the merged sets of artist that start with B which feature on a song

```
currentArtistsB = \Pi_{artist.name, song.name} (\sigma_{2010} < \text{year} < 2019 (song) \bowtie_{artistID} featuredOn \bowtie_{songID} \sigma_{name} = \text{"B%"} (artist)).
```

The finall query will then be the union of the two

 $final = oughts \cup currentArtistsB$

2.3 c)

We first join song and artist

 $songArtist = song \bowtie_{artistID} artist.$

Then we aggregate with the count function on the grouping artist.name.

$$songArtistCount = \gamma_{artist.name, COUNT(*) \rightarrow numberOfSongs}(songArtist).$$

Then we can order according to numberOfSongs

$$songArtistCountOrderd = \tau_{numberOfSongs\ DESC}(songArtistCount).$$

Finally we project the desired information

 $final = \Pi_{artist.name, song.name}(songArtistCountOrderd).$

3 Introduction to Database Normalization

3.1 a)

We would have to update 4 birth year cells and 4 director name cells. This would mean a total of 8 updates for the given changes.

3.2 b)

We should split the table, a logical way to do it is to split into a film and a director table. This would look like

```
Film(<u>FilmID</u>, Name, Year, DirectorID)
Director(<u>DirectorID</u>, DirectorName, DirectorBirthYear)
```

The functional relations are preserved in each table, and we store the information of the director through a relationship instead of directly.

4 Functional Dependencies, Keys and Closures

4.1 a)

Task	Validity	Explanation
1	True	Trivial relation
2	False	a_3 has two different values b_3 and b_4
3	Unknown	No conflicting values
4	Unknown	No conflicting values
5	False	c_1 corresponds to both d_1 and d_2
6	False	Lots of conflicting values
7	True	Trivial superkey
8	False	The first two rows show that ABC dose not determine D
9	False	D dose not determine C, so it cannot be a candidate key
10	Unknown	All ABD rows are unique, so no way to disprove

Table 1: Caption

4.2 b)

We find the atomic relations of F. We have $D \to A$ and $B \to D$, and thus we also have $B \to A$. We can also use these relations to simplify $ABD \to C$. Since B determines both A and D, we simply have $B \to C$. With these relations in mind, $D^+ = \{D, A\}$. For BC^+ we at least know they determine B and C. Since B also determines D and A, we know that BC^+ determines all letters, and thus the entire table, i.e. $BC^+ = ABCD^+$. The same reasoning leads to $AB^+ = BD^+ = ABCD^+$. We see from the atomic relations that B is a candidate key for the table. No other letters can determine the whole table alone, so this is the only one. It is also trivially minimal.