Demystifying the most significant Java language features from 9 to 11

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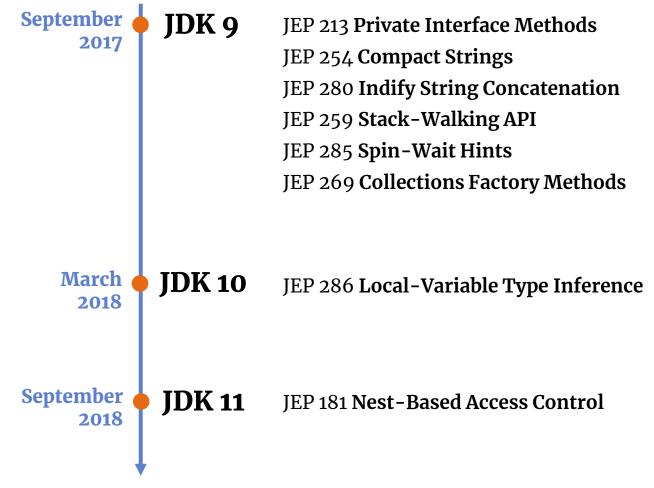






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- 01 Private Interface Methods
- 02 Compact Strings
- 03 Indify String Concatenation
- 04 Stack-Walking API
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Private Interface Methods



JEP 213: Milling Project Coin Author Joseph D. Darcy

Owner Joe Darcy Type Feature Scope SE Status Closed / Delivered Release 9

Summary

should be turned into an error in Java SE 9. It is also proposed that interfaces be allowed to have private methods. Non-goals

The small language changes included in Project Coin / JSR 334 as part of JDK 7 / Java SE 7 have been easy to use and have worked well in practice. However, a few amendments could address the rough edges of those changes. In addition, using underscore (" ") as an identifier, which generates a warning as of Java SE 8,

This JEP does **not** propose to run a "Coin 2.0" effort or to generally solicit new language proposals.

Description

Five small amendments to the Java Programming Language are proposed: 5. Support for private methods in interfaces was briefly in consideration for

tasks for Java SE 8. It is now proposed that support for private interface

inclusion in Java SE 8 as part of the effort to add support for Lambda Expressions, but was withdrawn to enable better focus on higher priority

6

```
public interface IPrivateMethod {
    default void foo() {
        print();
    }
        private interface method

private void print() {
        System.out.println("Private Method");
    }
}
```

```
public interface IPrivateMethod {
    default void foo() {
        print();
    }
        private interface method

private void print() {
        System.out.println("Private Method");
    }
}
```

```
public void foo();
    aload_0
    invokeinterface #1, 1 // InterfaceMethod print:()V

private void print();
    getstatic #2 // Field System.out:LPrintStream;
    ldc #3 // String "Private Method"
    invokevirtual #4 // Method PrintStream.println:(LString;)V
```

BYTECODE

private interface method is called using INVOKEINTERFACE, as any other interface method.

default is a keyword present only at the language level, it doesn't exist at bytecode level.

```
public static void baz();
   invokestatic #5  // InterfaceMethod staticPrint:()V

private static void staticPrint();
   getstatic #2  // Field System.out:LPrintStream;
   ldc  #6  // String "Private Static Method"
   invokevirtual #4  // Method PrintStream.println:(LString;)V
```

BYTECODE

private static interface method is called using INVOKESTATIC, as any static class method.

Compact Strings



Author Brent Christian

JEP 254: Compact Strings

Owner Xueming Shen Type Feature Scope Implementation Status Closed / Delivered Release 9

Summary

Adopt a more space-efficient internal representation for strings.

Goals

related Java and native interfaces.

Non-Goals

representation of strings. A subsequent JEP may explore that approach.

Motivation The current implementation of the String class stores characters in a char array, using two bytes (sixteen bits) for each character. Data gathered from many different applications indicates that strings are a major component of heap usage and, moreover, that most String objects contain only Latin-1 characters. Such

www.ionutbalosin.com characters require only one byte of storage, hence half of the space in the internal

Improve the space efficiency of the String class and related classes while

It is not a goal to use alternate encodings such as UTF-8 in the internal

maintaining performance in most scenarios and preserving full compatibility for all

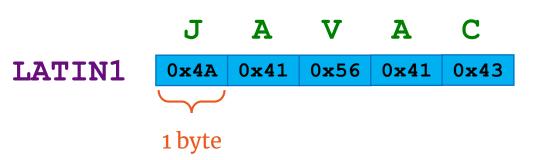
14

```
public final class String {
    private final byte[] value; <- JDK 8: private final char value[];</pre>
    private int hash;
    private final byte coder; // LATIN1 || UTF16
    @Native static final byte LATIN1 = 0;
    @Native static final byte UTF16 = 1;
    static final boolean COMPACT STRINGS;
         static {
              COMPACT STRINGS = true;
```

Note: -XX:+CompactStrings is enabled by default

String javaString = new String("JAVAC");





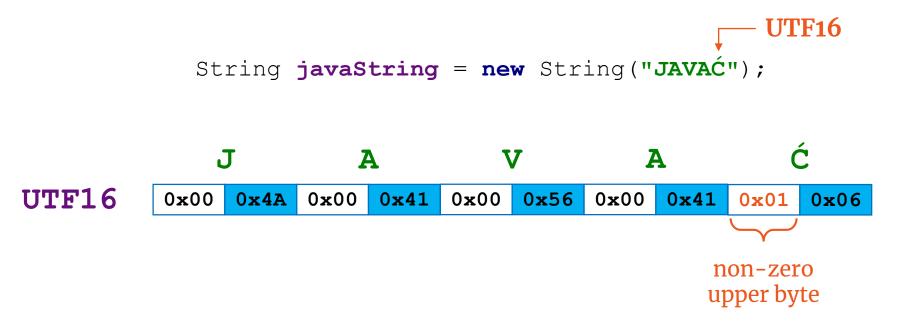
String javaString = new String("JAVAC");

UTF16

Class Name	Shallow Heap	Retained Heap
→ <regex></regex>	<numeric></numeric>	<numeric></numeric>
√ □ javaString java.lang.String @ 0x6d21347b0	24	56
> <class> class java.lang.String @ 0x6d2700980 System Class, Native Stack</class>	24	56
> ualue byte[10] @ 0x6d21347e0 J.A.V.A.C.	32	32
Σ Total: 2 entries		

LATIN1

Class Name	Shallow Heap	Retained Heap
→ <regex></regex>	<numeric></numeric>	<numeric></numeric>
√ □ javaString java.lang.String @ 0x6d2190ee8	24	48
> <class> class java.lang.String @ 0x6d2700980 System Class, Native Stack</class>	24	56
> ualue byte[5] @ 0x6d2190f18 JAVAC	24	24
Σ Total: 2 entries		



Strings containing characters with non-zero upper byte cannot be compressed. They are stored as 2 byte characters using UTF16 encoding.

String Construction

```
@see java.lang.String.java
String(char[] value, int off, int len, Void sig) {
    if (COMPACT STRINGS) {
        byte[] val = StringUTF16.compress(value, off, len);
        if (val != null) {
            this.value = val;
            this.coder = LATIN1;
            return:
    this.coder = UTF16;
    this.value = StringUTF16.toBytes(value, off, len);
```

- 1. first, it tries to compress the input char[] to Latin1 by stripping off upper bytes
- 2. if it fails, UTF16 encoding is used (i.e. each char spreads across 2 bytes)

```
\operatorname{QParam}(\{"\tilde{\mathbf{D}}\tilde{\mathbf{z}}\tilde{\mathbf{D}}^{1},\tilde{\mathbf{D}}^{2}\tilde{\mathbf{N}}\tilde{\mathbf{N}}^{1}\tilde{\mathbf{v}}\tilde{\mathbf{D}}"\})
                                                          @Param({"ΦȾ₩XĐ»ĐΛρ8Ë"})
public String utf 16 str1;
                                                          public String utf 16 str4;
@Param ({"ϑ; Ñ€Đ¾Đ; φаX"})
                                                          @Param ({"\ΩΔΘΞΨθςώς ηδο"})
                                                          public String utf 16 str5;
public String utf 16 str2;
                                                          @Param({"XĐςϠΨθϑ¿ĐžÈ¾"})
@Param({"Đςζ»Đ¾,Ñ^ĐμÑ,"})
public String utf 16 str3;
                                                          public String utf 16 str6;
@Benchmark
public String utf16 multiple concat() {
     return utf 16 str1 + utf 16 str2 + utf 16 str3 +
               utf 16 str4 + utf 16 str5 + utf 16 str6;
```

ult			
- defa		Average Time (ns/op)	Allocation Rate (B/op)
	-XX:+CompactStrings	54.234 ± 2.081	192.0
	-XX:-CompactStrings	45.957 ± 1.777	192.0

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

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	-XX:-CompactStrings	45.957 ± ^{1.777}	192.0

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

For applications that extensively use **UTF16** characters, consider disabling **Compact Strings** for a better performance.

String.equals()

```
// @see java.lang.String.java
public boolean equals(Object anObject) {
    // ...
    if (anObject instanceof String) {
        String aString = (String)anObject;
        if (coder() == aString.coder())
            return isLatin1() ? StringLatin1.equals(value, aString.value)
                              : StringUTF16.equals(value, aString.value);
    return false;
```

```
@Param({"123456789ABCDEFGHIJKLMNOPQRSTUVWXYZ"}) // length 35
public String latin 1;
@Param ({"123456789ABCDEFGHIJKLMNOPQRSTUVWXYZ"})
public String latin 1 identical;
                                            ____ UTF16
@Param({"123456789ABCDEFGHIJKLMNOPQRSTUVWX¾\"}) // length 35
public String utf 16;
@Param({"123456789ABCDEFGHIJKLMNOPQRSTUVWX¾\\"})
public String utf 16 identical;
```

```
@Benchmark
public boolean latin1 toLatin1 equals() {
    return latin 1.equals(latin 1 identical);
@Benchmark
public boolean latin1 toUtf16 equals() {
    return latin 1.equals(utf_16);
@Benchmark
public boolean utf16 toUtf16 equals() {
    return utf 16.equals(utf 16 identical);
```

default		Average Time (ns/op)
	-XX:+CompactStrings	-XX:-CompactStrings
latin1_toLatin1_equals	6.745 ± 0.418	7.793 ± 0.411
utf16_toUtf16_equals	7.910 ± 0.624	7.491 ± 0.372
latin1_toUtf16_equals	3.316 ± 0.201	7.114 ± 0.301

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

— default		Average Time (ns/op)
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Runtime: [OpenJDK 11.0.2 Linux 64-bit]

-XX:+CompactStrings:

- if coders differ, Strings comparison is really cheap.
- if the same coder, slightly better performance for Latin1.

-XX:-CompactStrings:

- almost the same response time in all cases.

Indify String Concatenation



```
private String aString;  // aString::length = 128
private int anInt;  // anInt = 2019
private float aFloat;  // aFloat = 42.0f

@Benchmark
public String concat() {
    return aString + anInt + aFloat;
}
```

```
public String concat();
                 #2 // class StringBuilder
    new
    dup
    invokespecial #12 // Method StringBuilder. "<init>":() V
    aload 0
    getfield #8 // Field aString:LString;
    invokevirtual #6 // Method StringBuilder.append: (LString;) LStringBuilder;
    aload 0
    getfield #9 // Field anInt:I
    invokevirtual #13 // Method StringBuilder.append: (I) LStringBuilder;
    aload 0
    getfield #11 // Field aFloat:F
    invokevirtual #14 // Method StringBuilder.append: (F) LStringBuilder;
    invokevirtual #7 // Method StringBuilder.toString:()LString;
    areturn
```

JEP 280: Indify String Concatenation

Owner Aleksey Shipilev
Type Feature
Scope SE
Status Closed / Delivered
Release 9

Summary

Change the static String-concatenation bytecode sequence generated by javac to use invokedynamic calls to JDK library functions. This will enable future optimizations of String concatenation without requiring further changes to the bytecode emitted by javac.

Goals

Lay the groundwork for building optimized String concatenation handlers, implementable without the need to change the Java-to-bytecode compiler. Multiple translation strategies should be tried as the motivational examples for this work. Producing (and possibly switching to) the optimized translation strategies is a stretch goal for this JEP.

Non-Goals

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It is not a goal to: Introduce any new String and/or StringBuilder APIs that might help to build better translation strategies, revisit the JIT compiler's support for optimizing String concatenation, support advanced String interpolation use

BootstrapMethods:

```
REF_invokeStatic StringConcatFactory.makeConcatWithConstants:() LCallSite;
Method arguments:
    \u0001\u0001\u0001
```

String Concat Strategies

-Djava.lang.invoke.stringConcat=<strategy>

BC_SB - spins up the bytecode that has the same StringBuilder chain javac would emit. It moves the javac generated bytecode to runtime. Baseline for current JDK behavior.

BC_SB_SIZED - similar to BC_SB, but it tries to guess the capacity required for StringBuilder to accept all arguments without resizing. It only guesses the space required for known types (e.g. primitives and Strings), but does not convert arguments. The capacity estimate may be wrong and StringBuilder may have to resize.

BC_SB_SIZED_EXACT - improves over BC_SB_SIZED, by first converting all arguments to String in order to get the exact capacity StringBuilder should have.

StringBuilder

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BC_SB_SIZED_EXACT - improves over **BC_SB_SIZED**, by first converting all arguments to String in order to get the exact capacity StringBuilder should have.

MH_SB_SIZED - builds the computation on public MethodHandle APIs which, at the end, ends up in calling the public StringBuilder API. This strategy does not use any private API at all.

MH_SB_SIZED_EXACT – improved over MH_SB_SIZED by first trying to get the exact capacity StringBuilder should have.

MH_INLINE_SIZED_EXACT - builds the char[] array on its own and passes that char[] array to String constructor. It opens the door for building a very optimal concatenation sequence, it does not use StringBuilder. **Default strategy**.

	Average Time (ns/op)	Allocation Rate (B/op)
JRE 1.8.0_211	134.996 ± 3.133	1,192.0
BC_SB	120.501 ± 5.621	656.0
BC_SB_SIZED	82.234 ± 3.297	384.0
BC_SB_SIZED_EXACT	59.217 ± 2.818	224.0
MH_SB_SIZED	88.952 ± 3.799	384.0
MH_SB_SIZED_EXACT	87.715 ± 3.124	376.0
MH_INLINE_SIZED_EXACT	54.773 ± 2.260	200.0

	Average Time (ns/op)	Allocation Rate (B/op)
JRE 1.8.0_211	134.996 ± 3.133	1,192.0
BC_SB	120.501 ± 5.621	656.0
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MH_SB_SIZED	88.952 ± 3.799	384.0
MH_SB_SIZED_EXACT	87.715 ± 3.124	376.0
MH_INLINE_SIZED_EXACT	54.773 ± 2.260	200.0

Advanced strategies (e.g. MH_INLINE_SIZED_EXACT) can improve performance when JIT compiler is not able to optimize the append StringBuilder chains (e.g. *_SB_* strategies).

Such advanced strategies can bring noticeable performance and allocation pressure improvements.

Stack-Walking API



JEP 259: Stack-Walking API Owner Mandy Chung

Type Feature Scope SE Status Closed / Delivered Release 9

Summary Define an efficient standard API for stack walking that allows easy filtering of, and lazy access to, the information in stack traces.

Non-Goal

this new API.

Motivation

There is no standard API to traverse selected frames on the execution stack

efficiently and access the Class instance of each frame.

There are existing APIs that provide access to a thread's stack:

frames on the stack. Both the Throwable::getStackTrace and

• Throwable::getStackTrace and Thread::getStackTrace return an array

It is not a goal to convert all existing stack walking code in the JDK to use

- of StackTraceElement objects, which contain the class name and method name of each stack-trace element.
- SecurityManager::getClassContext is a protected method, which allows

a SecurityManager subclass to access the class context. These APIs require the VM to eagerly capture a snapshot of the entire stack, and they return information representing the entire stack. There is no way to avoid the cost of examining all the frames if the caller is only interested in the top few

Stack-Walker - flexible mechanism to traverse and materialize the required stack-frame information allowing efficient lazy access to additional stack frames when required.

It avoids the cost of examining all frames if caller is interested in a few.

Stack-Walker features:

- ✓ lazy frames construction
- ✓ limit stack depth
- ✓ filter frames

```
public static StackWalker getInstance()
public static StackWalker getInstance(Set<Option> options)
      Option:
             // supports getCallerClass() and getDeclaringClass()
             RETAIN CLASS REFERENCE
             // shows all reflection frames and other hidden frames
             // that are implementation-specific.
             SHOW REFLECT FRAMES
             // shows all hidden frames (including reflection frames).
             SHOW HIDDEN FRAMES
public static StackWalker getInstance(Set<Option> options, int estimateDepth)
```

retrieve current stack frame - top stack frame -

Current Stack Frame

•••

Stack Frame N + 3

Stack Frame N + 2

Stack Frame N + 1

•••

main()

stack depth

```
@Benchmark
public StackWalker.StackFrame get top stack frame() {
    return recTopStackFrame(stackDepth);
private StackWalker.StackFrame recTopStackFrame(int depth) {
    if (depth == 0) {
        return getTopStackFrame();
    return recTopStackFrame(depth - 1);
private StackWalker.StackFrame getTopStackFrame() {
    return StackWalker.getInstance()
            .walk(stream -> stream.findFirst())
            .orElseThrow(NoSuchElementException::new);
```

```
@Benchmark
public StackWalker.StackFrame get top stack frame() {
    return recTopStackFrame(stackDepth);
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    if (depth == 0) {
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private StackWalker.StackFrame getTopStackFrame() {
    return StackWalker.getInstance()
            .walk(stream -> stream.findFirst())
            .orElseThrow(NoSuchElementException::new);
```

		Average Time (us/op)
Stack Depth	StackWalker	getStackTrace()
1	1.397 ± 0.074	31.029 ± 1.812
10	1.347 ± 0.066	47.933 ± 7.189
100	1.482 ± 0.078	166.431 ± 9.379
1000	3.202 ± 0.157	1,446.025 ± 71.891

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

		Average Time (us/op)
Stack Depth	StackWalker	getStackTrace()
1	1.397 ± 0.074	31.029 ± 1.812
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100	1.482 ± 0.078	166.431 ± 9.379
1000	3.202 ± 0.157	1,446.025 ± 71.891

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

- **StackWalker** offers almost constant access to current/top frame for different StackDepth sizes.
- **getStackTrace()** cost depends on StackDepth size and is at least one order of magnitude slower.

Performance Scenario: process asynchronously throwable frames (in another thread) vs. [synchronously] stack walking

```
Thread #1
Throwable aThrowable = new Throwable();
// do other stuff ...
              Thread #2
aThrowable.getStackTrace();
// asynchronously process stack frames
```

```
Throwable aThrowable = new Throwable();
                                            VS.
aThrowable.getStackTrace();
```

Thread **Top Stack Frame** Stack Frame N + 3 Stack Frame N + 2 Stack Frame N + 1 main()

Average Time (us/op)		
Stack Depth	StackWalker [1]	new Throwable() [2]
1	1.411 ± 0.079	1.186 ± 0.064
10	4.647 ± 0.215	1.518 ± 0.075
100	26.909 ± 1.412	5.199 ± 0.230
1000	250.267 ± 13.402	40.944 ± 2.089

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

^{[1] -} includes the cost of traversing synchronously (backwards) stack frames.

^{[2] -} excludes the cost of generating stack trace elements (i.e. getStackTrace()).

Average Time (us/op)		
Stack Depth	StackWalker [1]	new Throwable() ^[2]
1	1.411 ± 0.079	1.186 ± 0.064
10	4.647 ± 0.215	1.518 ± 0.075
100	26.909 ± 1.412	5.199 ± 0.230
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Runtime: [OpenJDK 11.0.2 Linux 64-bit]

^{[1] -} includes the cost of traversing synchronously (backwards) stack frames.

^{[2] -} excludes the cost of generating stack trace elements (i.e. getStackTrace()).

Average Time
(us/op)

Stack Depth StackWalker [1] new Throwable() [2]

Processing asynchronously throwable frames (i.e. in another thread) is even faster than Stack Walker API.

Advice: leverage to this mechanism when affordable, for a better performance!

[1] - includes the cost of synchronously traversing (backwards) N stack frames [2] - excludes the cost of generating stack trace elements (i.e. getStackTrace()

Spin-Wait Hints



JEP 285: Spin-Wait Hints

Authors Gil Tene, Ivan Krylov
Owner Paul Sandoz
Type Feature
Scope SE
Status Closed / Delivered
Release 9

Summary

Define an API to allow Java code to hint that a spin loop is being executed.

Goals

Define an API that would allow Java code to hint to the run-time system that it is in a spin loop. The API will be a pure hint, and will carry no semantic behaviour requirements (for example, a no-op is a valid implementation). Allow the JVM to benefit from spin loop specific behaviours that may be useful on certain hardware platforms. Provide both a no-op implementation and an intrinsic implementation in the JDK, and demonstrate an execution benefit on at least one major hardware platform.

Non-Goals

onSpinWait() - tells the CPU there is busy-waiting loop that may burn few CPU-cycles waiting for something to happen.

CPU can assign more resources to other threads, without actually invoking the OS scheduler to dequeue the thread (which may be expensive).

```
while (!condition_not_satisfied) {
     Thread.onSpinWait();
}
```

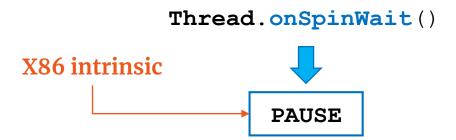
onSpinWait() - suitable for events which might happen very frequently
and finish very quickly (with a higher rate).

```
Producer |
```

```
for (long i = 0; i < items; i++) {
    while (!ready_to_produce) {
        Thread.onSpinWait();
    }
    produce_item();
}
// signal is_running = false
signal_finish();</pre>
```

Consumer

```
while (is_running) {
    while (!ready_to_consume) {
        Thread.onSpinWait();
    }
    consume_item();
}
```



PAUSE instruction delays the next instruction's execution for a finite period of time, hence parts of the pipeline are no longer being used → *reduces power consumption*!

Source: [https://software.intel.com/en-us/articles/benefitting-power-and-performance-sleep-loops]

	Average Time (us/op)	Context Switches [1]
onSpinWait()	39.419 ± 0.215	2,815
yield()	321.749 ± 5.175	2,826
parkNanos(1)	431.874 ± 16.510	2,821
sleep(1)	1,361,715.726 ± 31,258.272	18,504

^{[1] -} difficult to report an accurate number of context switches!

	Average Time (us/op)	Context Switches [1]
onSpinWait()	39.419 ± 0.215	2,815
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JDK / JDK-8159532

Investigate SPARC intrinsic for Thread.onSpinWait

Details

Priority: 4 P4 Resolution: Unresolved

Affects Version/s: 9, 10 Fix Version/s: tbd

Component/s: hotspot

Labels: bleaklow-interest c1 c2 c2-intrinsic performance

Subcomponent: compiler

CPU: sparc

Description

JDK 8147844 implements Thread.onSpinWait and the correspond x86 intrinsic.

Investigate whether it's possible to support an intrinsic method on SPARC.

Collections Factory Methods



JEP 269: Convenience Factory Methods for Collections

Owner Stuart Marks Type Feature Scope SE Status Closed / Delivered Release 9

Summary

Define library APIs to make it convenient to create instances of collections and maps with small numbers of elements, so as to ease the pain of not having collection literals in the Java programming language.

Goals

Provide static factory methods on the collection interfaces that will create compact, unmodifiable collection instances. The API is deliberately kept minimal.

Non-Goals

It is not a goal to provide a fully-general "collection builder" facility that, for example, lets the user control the collection implementation or various characteristics such as mutability, expected size, loading factor, concurrency level, and so forth.

It is not a goal to support high-performance, scalable collections with arbitrary numbers of elements. The focus is on small collections.

It is not a goal to provide "importable persistent" or "functional" collections

It is not a goal to provide unmodifiable collection types. That is, this proposal does not expose the characteristic of unmodifiability in the type system, even though the proposed implementations are actually unmodifiable.

Factory Methods – static factory methods on the collection interfaces that creates **compact**, **unmodifiable** collection instances.

```
List.of("a", "b", "c")

Set.of("a", "b", "c")

Map.of("a", "v1", "b", "v2")

Map.ofEntries(entry("a", "v1"), entry("b", "v2"))
```

List.of()

```
explicit methods for ≤
```

```
<E> List<E> of()
<E> List<E> of (E e1)
\langle E \rangle List\langle E \rangle of (E e1, E e2)
\langle E \rangle List\langle E \rangle of (E e1, E e2, E e3)
\langle E \rangle List\langle E \rangle of (E e1, E e2, E e3, E e4)
<E> List<E> of (E e1, E e2, E e3, E e4, E e5)
<E> List<E> of (E e1, E e2, E e3, E e4, E e5, E e6)
<E> List<E> of(E e1, E e2, E e3, E e4, E e5, E e6, E e7)
<E> List<E> of(E e1, E e2, E e3, E e4, E e5, E e6, E e7, E e8)
<E> List<E> of (E e1, E e2, E e3, E e4, E e5, E e6, E e7, E e8, E e9)
<E> List<E> of (E e1, E e2, E e3, E e4, E e5, E e6, E e7, E e8, E e9, E e10)
<E> List<E> of (E... elements) // var args method
```

```
var args(int ... args)
                                         // 4 args method call
   mov DWORD PTR [rdx+0x8], 0xf800016d
                                        // {metadata({array int})}
   mov DWORD PTR [rdx+0xc], 0x4
                                         // Collect stack params
                                        //*getfield param2
   mov r10d, DWORD PTR [rbp+0x10]
   mov r8d, DWORD PTR [rbp+0x14]
                                       //*getfield param3
   mov r11d, DWORD PTR [rbp+0x18]
                                       //*getfield param4
   mov ecx, DWORD PTR [rbp+0xc]
                                        //*getfield param1
                                         // Add params to array
                                     //*iastore
   mov DWORD PTR [rdx+0x10],ecx
                                      //*iastore
   mov DWORD PTR [rdx+0x14], r10d
   mov DWORD PTR [rdx+0x18], r8d
                                      //*iastore
   mov DWORD PTR [rdx+0x1c], r11d
                                       //*newarray
   call 0x00007f68e963e8c0
                                        //*invokespecial var args
                                         // {optimized virtual call}
```

```
var args(int ... args)
                                         // 4 args method call
   mov DWORD PTR [rdx+0x8], 0xf800016d
                                        // {metadata({array int})}
   mov DWORD PTR [rdx+0xc], 0x4
                                         // Collect stack params
                                        //*getfield param2
   mov r10d, DWORD PTR [rbp+0x10]
   mov r8d, DWORD PTR [rbp+0x14]
                                       //*getfield param3
   mov r11d, DWORD PTR [rbp+0x18]
                                       //*getfield param4
   mov ecx, DWORD PTR [rbp+0xc]
                                        //*getfield param1
                                         // Add params to array
                                     //*iastore
   mov DWORD PTR [rdx+0x10],ecx
                                      //*iastore
   mov DWORD PTR [rdx+0x14], r10d
                                      //*iastore
   mov DWORD PTR [rdx+0x18], r8d
   mov DWORD PTR [rdx+0x1c], r11d
                                       //*newarrav
   call 0x00007f68e963e8c0
                                         //*invokespecial var args
                                         // {optimized virtual call}
```

```
prior calling var_args
```

```
var_args method calls deal with extra stack parameter
          manipulation and array allocation
          (i.e. impacting the performance)
```

Average Time (us/op)		
No. of args	explicit method params	var_args method
2	2.304 ± 0.097	4.367 ± 0.364
4	2.575 ± 0.022	4.958 ± 0.067
6	3.015 ± 0.206	7.068 ± 0.812
8	3.342 ± 0.062	7.950 ± 1.151
10	3.748 ± 0.110	8.676 ± 0.606

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

Average Time (us/op)		
No. of args	explicit method params	var_args method
2	2.304 ± 0.097	4.367 ± 0.364
4	2.575 ± 0.022	4.958 ± 0.067
6	3.015 ± 0.206	7.068 ± 0.812
8	3.342 ± 0.062	7.950 ± 1.151
10	3.748 ± 0.110	8.676 ± 0.606

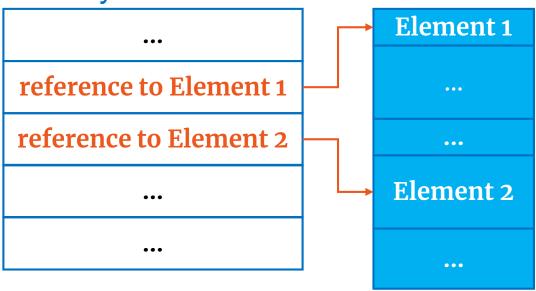
Runtime: [OpenJDK 11.0.2 Linux 64-bit]

```
static <E> List<E> of() {
    return ImmutableCollections.emptyList();
                                                       custom wrappers for
                                                           1 and 2 args
static <E> List<E> of (E e1) {
    return new ImmutableCollections.List12<>(e1);
static <E> List<E> of (E e1, E e2) {
    return new ImmutableCollections.List12<> (e1, e2);
static <E> List<E> of (E e1, E e2, E e3) {
    return new ImmutableCollections.ListN<>(e1, e2, e3);
                                                       custom wrapper for
                                                            N > 2 args
```

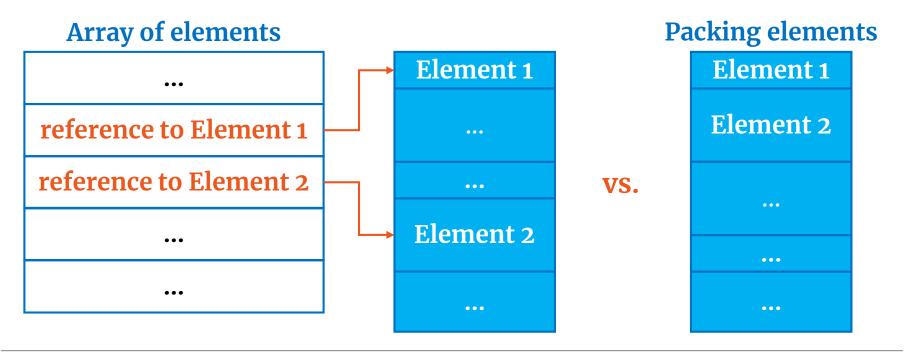
```
static final class List12<E> extends AbstractImmutableList<E> {
    @Stable private final E e0;
    @Stable private final E e1;
   // ...
                                             packed elements
                                             implementation
static final class ListN<E> extends AbstractImmutableList<E> {
    @Stable private final E[] elements;
                                               array based
                                             implementation
```

Data Locality

Array of elements



Data Locality



Packing elements reduces the number of indirections and improves data locality (i.e. minimize CPU D-Cache misses)

In a nutshell:

reduced memory footprint

improved data locality

Set.of() and **Map.of()** are implemented based on similar design principles:

{ custom wrappers, explicit method parameters }

Local-Variable Type Inference



JEP 286: Local-Variable Type Inference

Author Brian Goetz
Owner Dan Smith
Type Feature
Scope SE
Status Closed / Delivered
Release 10

Summary

Enhance the Java Language to extend type inference to declarations of local variables with initializers.

Goals

We seek to improve the developer experience by reducing the ceremony associated with writing Java code, while maintaining Java's commitment to static type safety, by allowing developers to elide the often-unnecessary manifest declaration of local variable types. This feature would allow, for example, declarations such as:

```
var list = new ArrayList<String>(); // infers ArrayList<String>
var stream = list.stream(); // infers Stream<String>
```

This treatment would be restricted to local variables with initializers, indexes in the enhanced for-loop, and locals declared in a traditional for-loop; it would not be available for method formals, constructor formals, method return types, fields,

```
var aBool = true;
var aChar = '\ufffd';
var aLong = 0L;
var aString = "conference";
var aFloat = 1.0f;
var aDouble = 2.0;
var listOf = List.of("a", "b");
```

```
var aBool = true;
var aChar = '\ufffd';
var aLong = 0L;
var aString = "conference";
var aFloat = 1.0f;
var aDouble = 2.0;
var listOf = List.of("a", "b");
```

```
boolean aBool = true;
char aChar = 65533;
long aLong = 0L;
String aString = "conference";
float aFloat = 1.0F;
double aDouble = 2.0D;
List<String> listOf = List.of("a", "b");
www.jonutbalosin.com
```

GENERATED CODE

var is only syntactic sugar, it does not exist at bytecode level.

```
(!)
```

```
var aByte = 0;
var aShort = 0x7fff; // intended as short
var anotherLong = 17; // intended as long
```

GENERATED CODI

```
int aByte = 0;
int aShort = 32767;  // intended as short but widened to int!
int anotherLong = 17; // intended as long but narrowed to int!
```



When the initializer is a numeric value (especially for integer literals) without an explicit compiler hint, that numeric value may be silently widened or narrowed (e.g. inferred implicitly as an int).

```
(!)
```

```
var aList = new ArrayList<>(); // diamond operator

var anotherList = arrayToList(anArray); // generic method call
public static <T> List<T> arrayToList(String[] a) {...}
```

GENERATED CODE

```
ArrayList<Object> aList = new ArrayList();
List<Object> anotherList = arrayToList(anArray);
// the inferred type cannot be deduced, it falls back to Object!
```



In case of diamond operator or generic methods where there is no explicit target type neither the actual constructor/method arguments to provide sufficient type information, the value will be inferred as <Object>.

Non-denotable types

{ anonymous class type, intersection type, capture variable type, NULL type }

Non-denotable types are types internally used by the compiler, which cannot be explicitly written out in the program source code.

```
var obj = new Object() { // obj has a non-denotable type
    public void foo() {
        System.out.println("Java Conference!");
    }
}:
```

```
var obj = new Object() { // obj has a non-denotable type
    public void foo() {
        System.out.println("Java Conference!");
    }
};
```

```
BYTECODE
```

```
new NonDenotableTypeMain$1 // new type generated by javac
dup
invokespecial NonDenotableTypeMain$1.<init> () V
astore 1
```



obj.foo(); // call not possible without type inference!

BYTECODE

invokevirtual NonDenotableTypeMain\$1.foo () V

aload 1

Type inference becomes very handy in case of non-denotable types, making the language more flexible.

Nest-Based Access Control



```
public class Outer {
    private void prv foo() {
        System.out.println("private foo method");
    public class Inner {
        public void foo() {
            prv foo(); // private method call
```

```
// Outer.class
public class Outer {
    public Outer();
    static void access$000(Outer) { // bridge method - NEW!
        aload 0
        invokespecial #1 // calls prv foo:()V
    private void prv foo();
  Outer$Inner.class
public class Inner {
    final Outer this$0;
    public Inner(Outer);
    public void foo() {
        aload 0
        getfield #1 // field this$0:LOuter;
        invokestatic #3 // calls Outer.access$000:(LOuter;)V
```

```
// Outer.class
public class Outer {
    public Outer();
  ▶ static void access$000 (Outer) { // bridge method - NEW!
        aload 0
        invokespecial #1 // calls prv foo:()V
    private void prv foo();
// Outer$Inner.class
public class Inner {
    final Outer this$0;
    public Inner(Outer);
    public void foo() {
        aload 0
        getfield
                            / field this$0:LOuter;
       invokestatic
                           // calls Outer.access$000:(LOuter;)V
```

```
// Outer class
public class Outer {
    public Outer();
  ▶ static void access$000 (Outer) { // bridge method - NEW!
        aload 0
        invokespecial #1 // calls prv foo:()V
 private void prv foo();
  Outer$Inner.class
public class Inner {
    final Outer this$0;
    public Inner(Outer);
    public void foo() {
        aload 0
        getfield
                            field this$0:LOuter;
       invokestatic
                             calls Outer.access$000:(LOuter;)V
```

Potential drawbacks in case of bridges methods:

subvert encapsulation

slightly increase generated code size

can confuse users and tools

additional layer of indirection

JEP 181: Nest-Based Access Control

Author John Rose
Owner David Holmes
Type Feature
Scope SE
Status Closed / Delivered
Release 11

Summary

Introduce *nests*, an access-control context that aligns with the existing notion of nested types in the Java programming language. Nests allow classes that are logically part of the same code entity, but which are compiled to distinct class files, to access each other's private members without the need for compilers to insert accessibility-broadening bridge methods.

Non-Goals

This IEP is not concerned with large scales of access control, such as modules.

Motivation

Many JVM languages support multiple classes in a single source file (such as Java's nested classes), or translate non-class source artifacts to class files. From a user perspective, however, these are generally considered to be all in "the same class", and therefore users expect them to share a common access control regime. To preserve these expectations, compilers frequently have to broaden the access of private members to package, through the addition of access bridges: an invocation of a private member is compiled into an invocation of a compiler-generated package-private method in the target class, which in turn accesses the intended private member. These bridges subvert encapsulation, slightly increase the size of a deployed application, and can confuse users and tools. A formal notion of a group of class files forming a *nest*, where *nest mates* share a common

access control mechanism, allows the desired result to be directly achieved in a

NestMates = {Outer.class, Inner.class}

```
class Outer {
    class Inner {
        //...
        NestHost:Outer // class attribute
    NestMembers: Inner // class attribute
```

Nests allow to access each other's private members without the need for compilers to insert accessibility broadening bridge methods.

```
// Outer class
public class Outer {
    public Outer();
    private void prv foo();
    NestMembers: Inner
// Outer$Inner.class
public class Inner {
    final Outer this$0;
    public Inner(Outer);
    public void foo() {
        aload 0
        getfield #1 // field this$0:LOuter;
        invokevirtual #3 // calls Outer.prv foo:()V
        return
    NestHost: Outer
```

```
Outer class
public class Outer {
    public Outer();
  private void prv foo();
                                        direct private member access,
    NestMembers: Inner
                                     no bridge method needed anymore!
  Outer$Inner.class
public class Inner {
    final Outer this$0;
    public Inner(Outer);
    public void foo() {
        aload 0
        getfield
                      #1 // field this$0:LOuter;
        invokevirtual #3 |// calls Outer.prv_foo:() V
        return
    NestHost: Outer
```

Thanks Gracias Danke Merci

Resources

Slides

https://IonutBalosin.com/talks

Further Readings

https://openjdk.java.net/projects/amber/LVTIstyle.html

https://www.youtube.com/watch?v=wIyeOaitmWM - Compact Strings

IonutBalosin.com/training **y**@IonutBalosin



UTF16 String Size Constraints

new String(new char[<length>]); // new String(<utfl6String>.toCharArray());

```
new String(new char[<length>]); // new String(<utf16String>.toCharArray());
// @see java.lang.StringUTF16.java
public static byte[] newBytesFor(int length) {
    // ...
    if (length > MAX LENGTH) { // MAX LENGTH = Integer.MAX VALUE >> 1
        throw new OutOfMemoryError("UTF16 String size is " + length +
                                    ", should be less than " + MAX LENGTH);
    return new byte[length << 1];</pre>
```

This constraint is not applicable for (1) LATIN1 Strings or (2) -XX:-CompactStrings.

List.of()

JDK 8

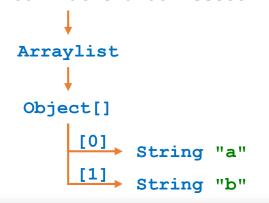
```
List<String> arrayList = new ArrayList<>();
    arrayList.add("a");
    arrayList.add("b");
    List<String> unmodifiableList = Collections.unmodifiableList(arrayList);

vs.

JDK 11
List<String> listOf = List.of("a", "b");
```

JDK 8

unmodifiableList --- Collections\$UnmodifiableRandomAccessList



Object / Stack Frame	Shallow Heap	Retained Heap
√ □ <local> java.util.ArrayList @ 0x74bcd4620</local>	24	80
> 🕝 < class> class java.util.ArrayList @ 0x74bf23c08 System Class	32	104
> 🗓 elementData java.lang.Object[10] @ 0x74bcd4668	56	56
$^{\Sigma}$ Total: 2 entries		
√ □ <local> java.util.Collections\$UnmodifiableRandomAccessList @ 0x74bcd47f8</local>	24	24
> 🕰 < class > class java.util.Collections\$UnmodifiableRandomAccessList @ 0x74be73428 System Class	8	8
> ☐ list, c java.util.ArrayList @ 0x74bcd4620	24	80
$^{\Sigma}$ Total: 2 entries		

Total Retained Heap: 80 + 24 bytes

JDK 11

e1	String	"a"
e2	String	"b"

Object / Stack Frame	Shallow Heap	Retained Heap
√	24	24
> <a> class > class java.util.ImmutableCollections\$List12 @ 0x74bf2afb8 System Class	8	32
> • e0 java.lang.String @ 0x74bcd4638 a	24	48
> □ e1 java.lang.String @ 0x74bcd46a0 b	24	48
$^{\Sigma}$ Total: 3 entries		

Total Retained Heap: 24 bytes

Average Time (us/op)		
No. of elements	of()	unmodifiableList()
empty	2.709 ± 0.162	8.349 ± 0.308
1	5.527 ± 0.283	19.051 ± 0.873
2	5.957 ± 0.258	23.931 ± 1.884
3	12.698 ± 0.549	27.078 ± 0.956
10	28.876 ± 1.675	46.579 ± 2.382
11	33.551 ± 1.849	67.406 ± 5.324

Runtime: [OpenJDK 11.0.2 Linux 64-bit]

Average Time (us/op)		
No. of elements	of()	unmodifiableList()
empty	2.709 ± 0.162	8.349 ± 0.308
1	5.527 ± 0.283	19.051 ± 0.873
2	5.957 ± 0.258	23.931 ± 1.884
3	12.698 ± 0.549	27.078 ± 0.956
10	28.876 ± 1.675	46.579 ± 2.382
11	33.551 ± 1.849	67.406 ± 5.324

Runtime: [OpenJDK 11.0.2 Linux 64-bit]