

Q1-a

Variables: x_e , where x_e is the flow in the edge, $\forall e \in E$

Objective Function: $\max \sum_{e \in in(t)} x_e$

Constraint:

$$1. \text{ capacity: } \begin{cases} x_e \geq 0 & \forall e \in E \\ x_e \leq c(e) & \forall e \in E \end{cases}$$

$$2. \text{ conservation : } \forall v \in V - \{s, t\}$$

$$l_u \sum_{e \in in(v)} x_e - \sum_{e \in out(v)} x_e = 0$$

Q1-b

Variables: x_e , where x_e is the flow in the edge, $\forall e \in E$

Objective Function: $\min \sum_{e \in E} x_e p(e)$

Constraint:

$$1. \text{ capacity: } \begin{cases} x_e \geq 0 & \forall e \in E \\ x_e \leq c(e) & \forall e \in E \end{cases}$$

$$2. \text{ conservation : } \forall v \in V - \{s, t\}$$

$$\sum_{e \in in(v)} x_e - \sum_{e \in out(v)} x_e = 0$$

$$3. \sum_{e \in in(t)} x_e = d$$

Q1-c

Variables: $x_i(e)$, where $x_i(e)$ is the number of product i go through e , $\forall e \in E, \forall i \in \{1, \dots, k\}$

Objective Function: $\max \sum_{i=1}^k \sum_{e \in out(s_i)} x_i(e)$

Constraint:

$$1. \text{ capacity: } \begin{cases} x_i(e) \geq 0 & \forall e \in E, \forall i \in \{1, \dots, k\} \\ \sum_{i=1}^k x_i(e) \leq c(e) & \forall e \in E \end{cases}$$

$$2. \text{ conservation : } \forall i \in \{1, \dots, k\}, \forall v \in V - \{s_i, t_i\}$$

$$\sum_{e \in in(v)} x_i(e) - \sum_{e \in out(v)} x_i(e) = 0$$

Q2-a

Variables: $x(u, v)$

where $x(u, v) = 1, \forall u, v \in V$ and $u \neq v$ **iff** there is a $G = (V, E)$ that realizes the given set of degree pairs and the edge $(u, v) \in E$.

Object Function: Nothing.

Any solution will be fine, as long as it satisfies the below constraints. (i.e. there is only feasible or infeasible solutions and no such thing as optimal solution)

Constraints:

1. $x(u, v) \in \{0, 1\}, \forall u, v \in V$
2. $\text{in}(u) - \sum_{v \in V - \{u\}} x(v, u) = 0, \forall u \in V$
3. $\text{out}(u) - \sum_{v \in V - \{u\}} x(u, v) = 0, \forall u \in V$

Construct the graph:

We are given V

Let X be the set of all variables defined above.

$E = \{\}$

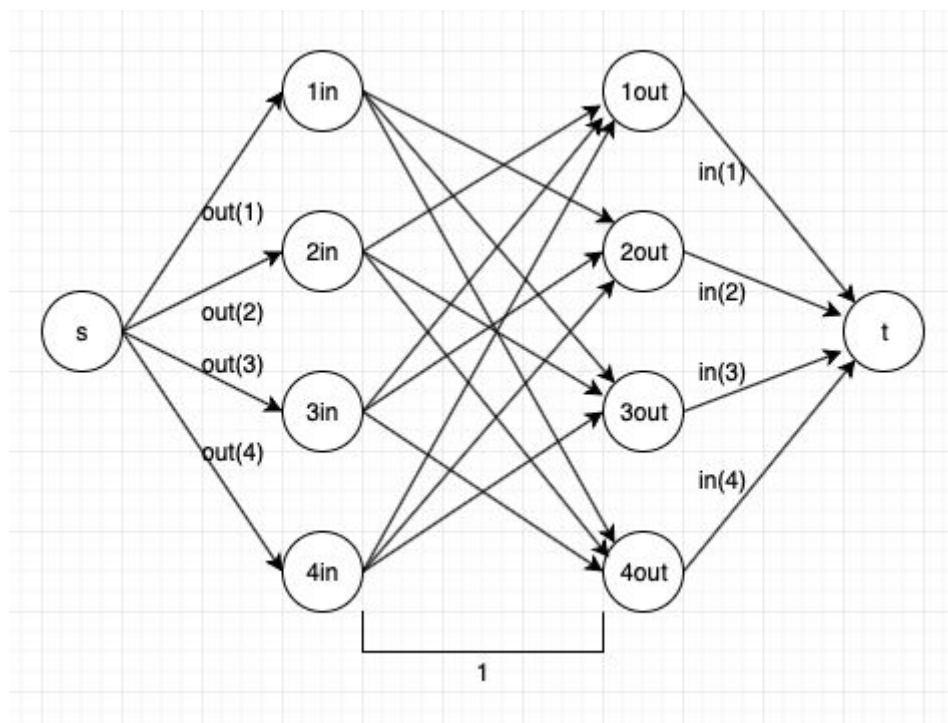
for $x(u, v)$ in X :

if $x(u, v) == 1$:

$E = E \cup \{(u, v)\}$

return $G = (V, E)$

Q2-b



Construction Instruction:

1. $V_F = \{v_{in}, v_{out} \mid \forall v \in V\} \cup \{s, t\}$
2. $E_F = \{(s, v_{in}) \mid \forall v \in V\} \cup \{(u_{in}, v_{out}) \mid \forall u, v \in V \& u \neq v\} \cup \{(v_{out}, t) \mid \forall v \in V\}$

$$3. c(e) = \begin{cases} out(v) & \text{if } e = (s, v_{in}) \\ in(v) & \text{if } e = (v_{out}, t) \\ 1 & O/W \end{cases}$$

Let $F = (G_F = (V_F, E_F), s, t, c)$

Algorithm:

Realize($V, (in(u), out(u))$):

construct flow net work F from V . $\Leftarrow O(n+m)$

$f = \text{MAXFLOW}(F)$ (integral) $\Leftarrow O(nm)$

if $V(f) \neq \sum_{u \in V} out(u)$:

return non-realizable

else:

$E = \{(u, v) \in E_F \mid u \neq s \text{ and } v \neq t \text{ and } f(u, v) = 1\}$ $\Leftarrow O(m)$

return $G = (V, E)$

Running Time:

Note: from the construction we can see that $m = n^2$.

$$O(n+m) + O(nm) + O(m) = O(n+n^2) + O(n^3) + O(n^2) \in O(n^3)$$

Correctness:

realizable set of degree (a) \Rightarrow Integral max flow of

pairs $V, (in(u), out(u)) \Leftarrow (b)$ value $\sum_{u \in V} out(u)$ in F

(a) Given a realizable set of degree pairs $V, (in(u), out(u))$

There is a $G = (V, E)$ that realizes the given set of degree pairs.

And construct $F = (G_F = (V_F, E_F), s, t, c)$ from G

f is the max flow in F by "MAXFLOW(F)"

WTS, $V(f)$ is indeed $\sum_{u \in V} out(u)$:

$$f(u_{in}, v_{out}) = \begin{cases} = 1 & \text{if there is a edge } (u, v) \in E \text{ of } G \\ = 0 & o/w \end{cases}$$

$$\forall u \in V, \quad \text{total flow out of } u = \sum_{v \in V - \{u\}} f(u_{in}, v_{out})$$

$= out(u)$ (i.e out degree of u) [G realizes the given set of degree pairs]

$=$ total flow in u [by conservation of a flow]

$= f(s, u_{in})$ (*)

$$V(f) = \sum_{u \in V} f(s, u_{\text{in}}) \quad [\text{by definition of value of flow}]$$

$$= \sum_{u \in V} \text{out}(u) \quad [\text{by *}]$$

(b) Given an integral max flow of value $\sum_{u \in V} \text{out}(u)$ in F

Let $E = \{(u, v) \in E_F \mid u \neq s \text{ and } v \neq t \text{ and } f(u, v) = 1\}$

Let $G = (V, E)$

WST: $\forall u \in V, |\{(u, v) \mid (u, v) \in E\}| = \text{out}(u)$

and $\forall u \in V, |\{(v, u) \mid (v, u) \in E\}| = \text{in}(u)$

Fact: $\forall u \in V, f(s, u_{\text{in}}) = \text{out}(u)$

$$\text{we know } \sum_{u \in V} f(s, u_{\text{in}}) = \sum_{u \in V} \text{out}(u) \quad [\text{given}]$$

$$\text{we know } \sum_{u \in V} c(s, u_{\text{in}}) = \sum_{u \in V} \text{out}(u) \quad [\text{since } \forall u \in V, c(s, u_{\text{in}}) = \sum_{u \in V} \text{out}(u)]$$

we can say $\forall u \in V, f(s, u_{\text{in}}) = c(s, u_{\text{in}}) = \text{out}(u)$

Proof by contradiction:

$$\forall u \in V, f(s, u_{\text{in}}) \leq c(s, u_{\text{in}}) \quad (**)$$

If $\exists u \in V$ s.t $f(s, u_{\text{in}}) \neq c(s, u_{\text{in}})$,

$$\Rightarrow f(s, u_{\text{in}}) < c(s, u_{\text{in}}) \quad (***)$$

$$\Rightarrow \sum_{u \in V} f(s, u_{\text{in}}) < \sum_{u \in V} c(s, u_{\text{in}}) = \sum_{u \in V} \text{out}(u) \quad [\text{by ** \& ***}]$$

$$\text{Which contradict to } \sum_{u \in V} f(s, u_{\text{in}}) = \sum_{u \in V} \text{out}(u)$$

Recall def of f : $\forall u \in V$,

$$f(u_{\text{in}}, v_{\text{out}}) \begin{cases} = 1 & \text{if there is a edge } (u, v) \in E \text{ of } G \\ = 0 & \text{o/w} \end{cases}$$

$$|\{(u, v) \mid (u, v) \in E\}|$$

$$= \sum_{v \in V - \{u\}} f(u_{\text{in}}, v_{\text{out}}) \quad [\text{by def of } f \text{ and } f \in \{0, 1\}]$$

$$= f(s, u_{\text{in}}) \quad [\text{by conservation}]$$

$$= c(s, u_{\text{in}}) \quad [\text{by fact}]$$

$$= \text{out}(u) \quad [\text{by construction of the } F]$$

Similar for proof $|\{(v, u) \mid (v, u) \in E\}| = \text{in}(u)$

$\Rightarrow G = (V, E)$ realize the set of degree pairs.

Therefore the set of degree pairs is realizable.