Introduction

Outline

- 1. Introduction of PBFT
- 2. Mechanism of PBFT
- 3. Proof of PBFT
- 4. Detail of PBFT
- 5. Optimization of PBFT
- 6. Application: Hyperledger
- 7. Microsoft talking time

Introduction

- 1. Practical Byzantine Fault Tolerance
 - a. A consensus algorithm
 - i. solving Byzantine Fault Tolerance
 - ii. versus to proof of work
 - b. Replication System
 - i. Not just used for Blockchain, it can be seen as NFS
 - (Network file system)

PBFT vs POW

- 1. Node Scalability a. POW (no limited) > PBFT $(4\sim100 \text{ node})$
- 2. Fault Tolerance a. POW(50%) > PBFT (25~33%)
- 3. Transaction Throughputa. PBFT(1000/s Fabirc) >> POW(7/s Bitcoin)b. Make it practical !!!

PBFT Properties

- 1. No need mining
- 2. Fault tolerance: n = 3f + 1
- 3. Every node must connect with all nodes
- 4. View concept
 - a. In a view, one of the nodes is the primary node.
 - b. Others will be replica node.

Mechanism

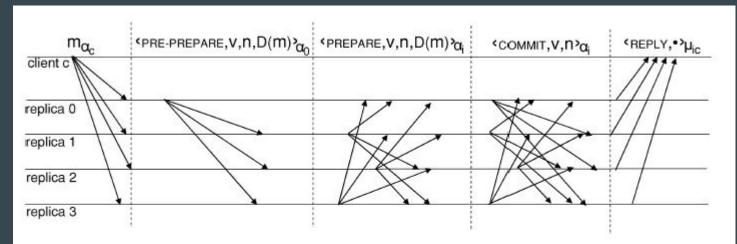
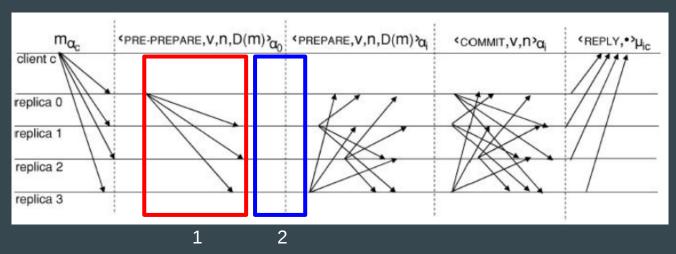
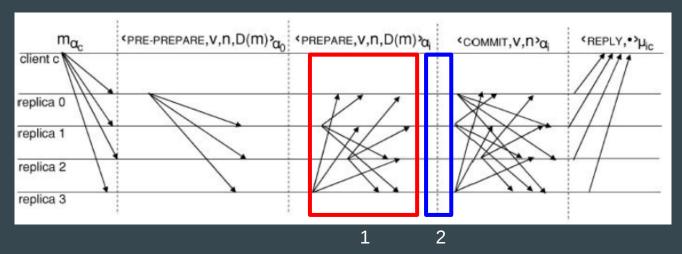


Fig. 1. Normal case operation: the primary (replica 0) assigns sequence number n to request m in its current view v and multicasts a pre-prepare message with the assignment. If a backup agrees with the assignment, it multicasts a matching prepare message. When a replica receives messages that agree with the assignment from a quorum, it sends a commit message. Replicas execute m after receiving commit messages from a quorum.



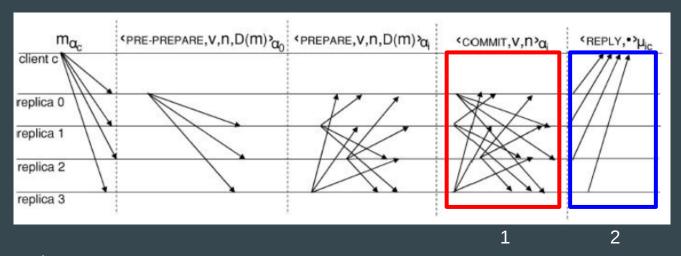
Pre-Prepare:

- 1. Primary node validate & Multicast the request got from client
- 2. Validate & Order requests got from primary node.
 - a. Distinguish the primary node is bad node when the digest of requests is different between primary node and replica node.



Prepare:

- 1. Multicast the requests the node got and compare from other node. Make sure all the node got the same requests at same time.
- 2. Execute these requests



Commit:

- 1. Multicast to other node the node have executed, then compare with other node's state. Make sure every node have the same current state.
- 2. Each node write in NFS and replies conclusion to client.

Proof of PBFT

Define a good node

- 1. Always processing request correctly a. Deterministic
- 2. Time limited Network
 - a. the Interval sending message and getting reply always
 - limited
- 3. Otherwise, the node is seen as bad node

Many kinds of attack

- Control the node
 a. Tempering the message
- 2. Delay the Network a. Denial of service
- 3. Pretend as good node

Tempering the message

1. Voting

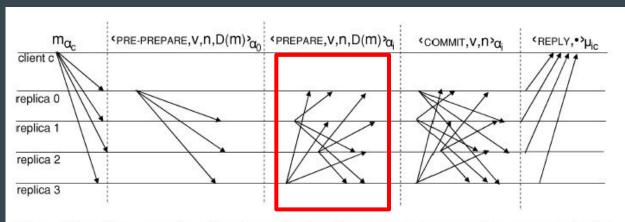


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Delay the network

- Once receive 2f + 1 of reply
 a. Make the decision and keep process.
- 2. All the bad node give bad reply a. Can not change the decision.
- 3. All the bad node delay to replay a. Can not delay the decision.

Pretend as good node

- 1. Change private key periodically
- 2. Turning to recover node detect which the node loss the control.
- 3. Other node reject message sign by old key.

Something detail

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View Change

1. If the primary node is bad node

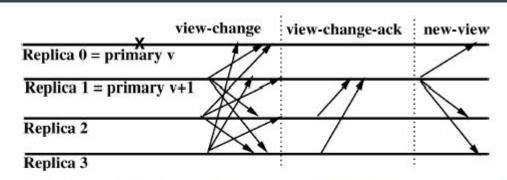


Fig. 2. View-change protocol: the primary for view v (replica 0) fails causing a view change to view v+1.

Garbage collection

- 1. Any message received or sent to other node will log in nodes.
- 2. The log will growth infinitly.
- 3. Use checkpoint consider which state is absolutly correct.
- 4. Then discard the logs recording prior to that state
- 5. Once the k requests committed, multicast checkpoint message
- 6. Get 2f+1 replay, then discard the logs

Optimalize

Chunk Processing

- 1. Seen as numbers of request a chunk.
- 2. Sending message in unit of chunk.

3. Decrease network usage, amortize protocol

overhead.

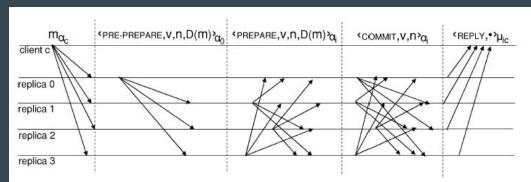


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Digest

- 1. Instead of sending whole requests content
- 2. Just send the encode of requests
- 3. Decrease network usage
- 4. MD5

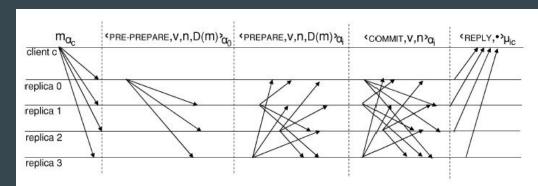


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Others

- 1. Using symmetric cryptography authenticate message.
 - a. Every node node have private key, use non symmetric cryptography exchange public key.
- 2. According to paper, the speed of BFT NFS just faster than normal NFS 2% to slower it 24%

Application

Hyperledger Fabric

- 1. Website
- 2. Using PBFT as consensus algorithm
- 3. Using Nodejs / Java as client interface
- 4. Using golang as chaincode smart contract
 - a. No cryptocurrency, UTXO, just a ledger you can record everything.
 - b. No fee for transaction

1. Require & version

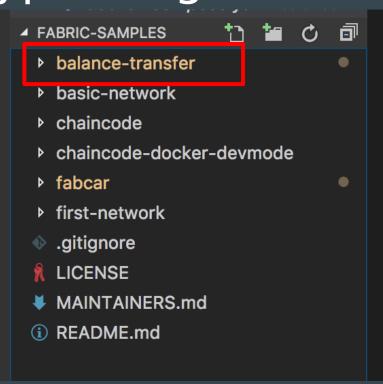
- a. docker 17.03.0-ce or greater
- b. docker-compose 1.8 or greater
- c. Go language 1.7.x
- d. Nodejs 6.9.x or greater but < 7.x
- e. npm 3.10.10

1. Download the docker Image

- a. sudo curl -sSL https://goo.gl/5ftp2f | bash
- b. sudo docker images

2. Download the example

- a. git clone https://github.com/hyperledger/fabric-samples.git
- b. cd fabric-samples



- 1. sudo bash runApp.sh
- 2. run sudo docker ps see im

NAMES

```
peer1.org2.example.com
peer1.org1.example.com
peer0.org1.example.com
peer0.org2.example.com
ca_peer0rg1
orderer.example.com
ca_peer0rg2
```

- 1. See testAPI.sh to know the node client rest api format.
- 2. artifacts/src/github.com/example.cc folder a. put the example code of chaincode smart contract.
- 3. app.js show how the node rest api working
- 4. ChainCode package shim