



ESTRUCTURA DE DATOS

FORMAL SPECIFICATIONS OF AN ADT

Data logic = Cierto ||| !False * When the parameter is a generator you must omit the output

or, and :: Lógico → Lógico → Lógico (Role very important in Haskell because of new commands)

not :: Lógico → Lógico (Bool)

not Cierto = Falso // not Falso = Cierto

andor x y = z

orand

falso = not cierto.

UNIT 1 INTRODUCTION

1.1 Terminology:

- A type (sort) is a set of values → boolean < true
false
 - int by values 0, -1, 1, -2, 2 ... ; int [] composed by {1}, {1, 2, 3, 4}, {1, 4, 6}
 - Par (int, char) composed by (1, 'a'), (0, 'g')
- A datum is a value belonging to concrete type (true, 0, {1, 2})
- An Abstract Data Type (ADT) is a type together with a set of operations for managing its data, without taking into account any implementation detail.
Each operation is defined in terms of its inputs and outputs: Boolean ADT = {True, False, and, or, not}; Natural ADT = {0, succ, +, -, <, ...}
- A Data Structure is a physical implementation of one or more ADT's. Each operations defined into an ADT is associated to one or more functions/procedures methods written with a concrete programming language. Its data is usually stored in central memory. When it is stored into secondary memory we have file

1.2 Data Structures and Efficiency

- A computer program can be decomposed as: Program = Algorithm + Data Structure [N. Wirth]
" Logic + Control "
- The most important task of a program (VERY IMPORTANT) consists in storing and retrieving information. It is crucial to organize data in a way that information can be manipulated as more efficient as possible.
- Algorithms describe the sequence of actions to be applied on data and they have 2 components: definition of operations and implementation.
- A good choice of appropriate data structures to manipulate the data involved in a program, as well as a clever design might imply that such program is executed in a few seconds offering the pursued results.
- Some criteria which will help us to select a good Data Structure for a given problem:
 - To analyze the problem to determine the Data Structure
 - To decide the basic sets of operations to be applied on each Data Structure
 - Inserting, deleting, searching... data, are considered basic operations on a Data Structure
 - To select the Data Structure which minimizes the computational cost

13 Data Structures and Object Oriented

- O.O. Programming becomes a suitable paradigm for implementing a Data Structure thanks to its wide repertoire of resources and tools for performing this task:
 - The notion of class allows us to clearly define and create a Data Structure
 - Information hiding enables the possibility of hiding most implementation details of the Data Structures to the final user.
 - Inheritance is useful for creating new Data Structure from previous existing ones, as well as for building generic D.S.
 - Polymorphism allows to change the behaviour of a given operation depending on the D.S where it is applied

UNIT 2 FORMAL DESCRIPTION OF DATA TYPES

2.1 Introduction:

- The definition of an ADT is performed independently of its implementation. We can use formal languages which sometimes offer an implementation of the ADT.
- There exist a mathematical definition for ADT used for implementing this concept with a programming language. It is mandatory to know such mathematical description for using in JAVA. Something similar occurs with other types like boolean, real, array, arraylist... The more complex an ADT is, the more required is too its formal definition.
- It's important separate the specification from its implementation(s)

2.2 Formal Specification of an ADT.

- GOAL: To establish what a Data Type is by means of an algebraic model.

- We define the meaning or semantics of the operations by using a set of equations or axioms $(\forall n, m \in \text{Nat}) \text{ add}(\text{Suc}(m), n) = \text{Suc}(\text{add}(m, n))$
- In order to define the meaning of any ADT, the set of axioms must be sufficient (enough), any property of the ADT should be proven from the set of axioms, but it will not be possible to demonstrate false properties.
- Axioms describe in an abstract way a minimal set of constraints or conditions which must be verified by any concrete implementation of the ADT

ADVANTAGES

- Users share the same (unique) interpretation of the ADT
- Programs making use of the ADT admit formal (automatic) verification.
- Possibility of using tools for executing specifications → prototype generation.

ELEMENTS APPEARING IN A FORMAL SPECIFICATION OF AN ADT

- 1. Set of types required in the definition → Interested Type or Intended Type (IT) → Sorts or Domains.
- 2. Operations for manipulating the data of the IT → Provide names and prototypes or shapes.

We can distinguish the following kinds of operations:

- Generators or constructors: Used to build all data of the IT. Data are really the terms produced with these operations and they can't no seen as abstract which produce an output
- Selectors: Useful for extracting data of the IT from other operations which aren't generators.
- Consult operations: Return objects different from IT
 - Predicates: They provide a Boolean value
 - Selectors: They provide attributes of data collected in a IT.
- 3. Sufficient set of axioms defining operations and their properties. It is usual that some axioms be recursive.
 - Axioms are mandatory for proving that any free term is always equivalent to other one built only with generators → inform about equivalence

FREE AND NOT FREE GENERATORS

- Generators of an ADT are free if different terms always represent different data.

- Not free generators force us to establish a set of axioms for describing the equality among basic terms, called equivalence axioms required too for identifying non valid terms. They introduce an equivalence relation among terms, where there might exist a particular class collecting the set of all possible non valid terms.
- Equivalence → shape $\text{Suc}(\text{Suc}(\dots \text{Suc}(\text{Zero})))$ or $\text{Pre}(\text{Pre}(\dots \text{Pre}(\text{Zero})))$ called canonic representation → terms in normal form

HASKELL INSERT:

$\text{insert} :: \text{List } x \rightarrow \text{List } y \rightarrow \text{List } a$
 $\text{insert } x \quad y = x : y : \text{List } a$

30th October 2017

a) terms (specified) nf-Inte

$\text{data Nat} = \text{zero} \mid \text{Suc Nat}$

$\text{terms} :: \text{Nat} \rightarrow \text{Nat} \rightarrow \text{Nat}$

$\text{fact} :: \text{Nat} \rightarrow \text{Nat}$

$\text{fact} :: \text{Zero} \rightarrow \text{SucZero}$

$\text{fact} :: \text{Suc}(x) \rightarrow \text{terms}(\text{fact}(x)) (\text{Suc } x)$

$\text{data Inte} = \text{z} \mid \text{s Inte} \mid \text{p Inte}$

$\text{nf-Inte} :: \text{Inte} \rightarrow \text{Inte}$

$\text{abs I}, \text{aux-abs I} :: \text{Inte} \rightarrow \text{Nat}$

$\text{abs I } x = \text{aux-abs I}(\text{nf-Inte } x)$

$\text{aux-abs I} :: \text{Inte} \rightarrow \text{Nat}$

$\text{aux-abs I } \text{z} \rightarrow \text{Zero}$

$\text{aux-abs I } \text{s}(x) \rightarrow \text{Suc}(\text{aux-abs I } x)$

$\text{aux-abs I } \text{p}(x) \rightarrow \text{Suc}(\text{aux-abs I } x)$

$\text{factabs I} :: \text{Inte} \rightarrow \text{Inte}$

$\text{factabs I } \text{z} = \text{Zero}$

$\text{factabs I } x = \text{aux-fact}(\text{fact}(\text{abs-} \text{I } x))$

$\text{aux-fact} :: \text{Nat} \rightarrow \text{Inte}$

$\text{aux-fact } \text{Zero} = \text{z}$

$\text{aux-fact } \text{Suc}(x) = \text{s}(\text{aux-fact } x)$

EXERCISES

andor: $(x \wedge y) \vee z$

orand (false, a, b) = false

andor ($\text{true}, \text{true}, c$) = true

orand ($\text{false}, b, \text{true}$) = b

andor ($\text{true}, \text{false}, c$) = c

orand ($\text{true}, b, \text{true}$) = true .

andor (a, b, c) = c

EXAMPLE IN HASKELL: NATURAL NUMBERS:

`data Nat = Cero | Suc Nat deriving Show`

`suma :: Nat → Nat → Nat`. Only 2 formats.

`suma Suc(Cero) (Suc Cero) → Suc (Suc (Suc Cero))`

To equation
Suma Cero $x = x$
Suma ($\text{Suc } y$) $x = \text{Suc} (\text{suma } x \ y)$

Recursion parameter

Example exercise:

`inf Cero = Cero`

`fixed Cero = Suc (Suc Cero)`

`fixed Cero = Suc Cero`

`inf (Suc Cero) = Suc Cero`

`fixed Suc (Suc Cero) = Suc Cero`

`fixed Suc Cero = Suc (Suc Cero)`

`inf (Suc (Suc Cero)) = Suc (Suc (Suc Cero))`



`fixed Suc (Suc (Suc Cero)) = Cero`

nb

`fixed Suc (Suc (Suc Cero)) = Cero`

`fixed (Suc (Suc (Suc Cero))) = fixed x`

TEST IF THE INT IS IN NORMAL FORM

`test :: Int#(8) → Bool`

`test 2 = True`

`test (P(x)) = False`

`test (S(x)) = False`

`test Px = test x`

`test Sx = test x`

FACTORIAL

`Parityfactorial :: Nat → Nat`

`Parityfactorial :: Nat → Nat`

`Parityfactorial Cero = Suc Cero`

`Parityfactorial Cero = Cero`

`Parityfactorial Suc Cero = Suc Cero`

`Parityfactorial Suc Cero = Suc Cero`

`Parityfactorial (Suc (Suc x)) = times (Parityfactorial (x)) (Suc (Suc x))`

`Parityfactorial (Suc (Suc x)) = times (Parityfactorial (x)) (Suc (Suc x))`

We add here `Parityfactorial Suc (Suc Cero) = Suc (Suc Cero)`

JAVA LISTS

generators
 $[] \leftrightarrow$ empty list
 $(x:t) \leftrightarrow$ elements
 a

head :: $[a] \rightarrow a$

head $[] =$ ~~andor~~ "...."

head $(x:__)$ = x

tail :: $[a] \rightarrow$

Dynamic
LOW LEVEL ↗
Static



LIFO \Rightarrow last input is first output.

$(3:(2:(1:[])))$

$++ :: [a] \rightarrow [a] \rightarrow [a]$

$list \quad ++ \quad [] = list$

$list \quad ++ \quad (x:t) = (x:list) \quad ++ \quad t$ {WHAT IS IT DOING?}

$$\hookrightarrow l \quad ++ \quad (x:t) = (l \quad ++ \quad x \quad) \quad ++ \quad t$$

SUMMARY

IMPORTANT \Rightarrow LIFO (last input first output)

list | $[]$ | $[]$ | $(x:t)$ | head | tail Ex: $[1:2]$ tail output is the list containing 1.

Stack | Stack a | EST | pv | $(x:/s)$ | top | pop

Queue | Queue a | EQn | cv | $(q:x)$

qsg :: Queue a \rightarrow Stack a \rightarrow Queue a

qss :: Queue a \rightarrow Stack a \rightarrow Stack a } recursion needed.

sqq :: Stack a \rightarrow Queue a \rightarrow Queue a

sqs :: Stack a \rightarrow Queue a \rightarrow Stack a

cq :: Queue a \rightarrow Queue a \rightarrow Queue a

cq q $EQn = q$

{ cq q1 q2 = cq (q1 :. (first q2)) (rest q2)

 cq q1 (q2 : x) = cq (q1 : x) q2

$\hookrightarrow cq q1 q2 : x = cq (con : 1) (\overbrace{tQn : 2 : 3}^{q1}) : 4 = cq (\overbrace{tQn : 1}^{q1}) (\overbrace{tQn : 2}^{q2}) : 3 : 4 = cq (tQn : 1) (tQn : 2 : 3 : 4) =$

$\therefore con : 1 : 2 : 3 : 4$

sq :: Stack a \rightarrow Queue a

sq EST = EQn

sq $(x:/s) = (sq s) : x$

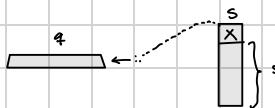
\rightarrow aux :: Stack a \rightarrow Queue a \rightarrow Queue a

sq2 s \in aux s EQn

aux $(x:/s) = aux s (q : x)$

aux EST $= q = q$

$q \in Q \rightarrow$ Queue a \rightarrow Stack a \rightarrow Queue a



STUDY 1st PROGRESS TEST

NOVEMBER 8TH, 2022

a) SPECIFICATION:

a) (NON-FREE generators) $\text{data Int} = \mathbb{Z} \mid s \text{ Int} \mid p \text{ Int}$ deriving $\text{subI} :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$ Input = Output.

$\text{subI} :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$

$$\text{subI } x \quad z = x$$

$$\text{subI } (sx) \quad (sy) = \text{SubI } x \quad y$$

$$\text{subI } x \quad (sy) = P(\text{subI } x \quad y)$$

$$\text{subI } px \quad py = \text{subI } x \quad y$$

$$\text{subI } x \quad py = S(\text{subI } x \quad y)$$

b) Haskell:

$$op22 (\text{EQ} :: \mathbb{A} :: \mathbb{B} :: \mathbb{S}) \quad (4 :: \mathbb{E} :: \mathbb{S}) \rightarrow 3 :: / (\text{op22 } \text{EQ} :: \mathbb{A} :: \mathbb{B} :: \mathbb{S}) = 3 :: / (2 :: / (4 :: \mathbb{E} :: \mathbb{S})) \quad ? \text{ Y el } 4?$$

$$op22 \text{ EQ } 4 :: \mathbb{S} :: \mathbb{B} :: \mathbb{S} \rightarrow 6 :: / (\text{op22 } \text{EQ } 4 :: \mathbb{S} :: \mathbb{B} :: \mathbb{S}) = 6 :: / (5 :: / (\text{op22 } \text{EQ } 4 :: \mathbb{S} :: \mathbb{B} :: \mathbb{S})) = 6 :: / (5 :: / (4 :: \mathbb{S} :: \mathbb{B} :: \mathbb{S})) = 5 :: / (6 :: \mathbb{B} :: \mathbb{S})$$

$$\text{op22 } (\text{EQ} :: \mathbb{A} :: \mathbb{B} :: \mathbb{S}) \quad (4 :: \mathbb{S} :: \mathbb{B} :: \mathbb{S}) = 3 :: / 2 :: / 1 :: / 5 :: / 6 :: / \mathbb{S}$$

② a) HIGH LEVEL : $\text{public static } < E > \text{ Pila } (E) \text{ op22 } (\text{Coda } < E > q, \text{ Pila } < E > s)$

| Pila < E > s; Coda < E > q; ... |

b) LOW LEVEL : $\text{public void op22 } (\text{CD } < E > q)$

| Nodo < E > aux = null; ... |

JAVA

LOW STATIC → array
HIGH DYNAMIC → links, pointers



class Nodo < E>
| Info E; Nodo < E> Siguiente.

public class Nodo (x, p)
{ Info = x; Siguiente = p; }

Preserve the state of the parameter → Do a clone()

public void Apila (E e); public void APila (E x, Pila < E > p); ; public Pila < E > APila (E e); ; Pila < E > APila (E e, Pila < E > p);

(Codigo en la practica)

```

    | if (p. EsVacia())
      p. NodoCabeza = x. NodoCabeza
    else
      s = new PD ( p. topo(), p. Desapila())
      s. NodoCabeza = new Nodo (x, s. NodoCabeza)
    { return s; }
  
```

Ejercicio:

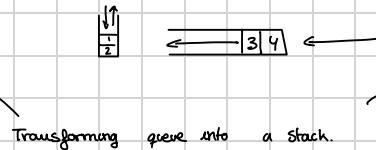
qsq:: Queue a → Stack a → Queue a

qsq q Est = q

qsq q. x / s = (qsq q s) : x ← NOT GOOD BECAUSE STACK INFO IS INSERTED IN REVERSE WAY.
(o) qsq (q. x) s

sqa:: Stack a → Queue a → Stack a

sqa Est q = q (qs q s)



qs :: Queue → Stack a

qs EQu = EST

qs q. x = x / (qs q) → example qs EQu : 3 : 4 → 4 : (qs EQu : 3) = 4 : 3 : (qs EQu) = 4 : 3 : EST.

public PD () Haskell CJ

| NodoCabeza = null;

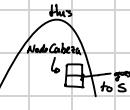
public PD (E x, PD < E > s)

| if (e != null)

New Nodo (x, s. NodoCabeza)

| PD < E > pd = new PD < E > ();

this. NodoCabeza = new Nodo (x, s. NodoCabeza);



Use this line:

infixl 5 : . Lo usaremos como encola (por comodidad)

ordenar :: (Ord a) \rightarrow Cola a \rightarrow Cola a

ordenar CV = CV

ordenar q = (ordenar (quitar g (m, or q))) \therefore (mayor q)

If we can compare objects its classes are comparable.

infixr 5 : / insertar.

polish notation:

$$2 \cdot 3 + 4 = 10 \rightarrow 2 \ 3 * \ 4 + = 10$$

$$2 * (3 + 4) = 14 \rightarrow [2 \ [3 \ 4 +] *]$$

IMPORTANT THINGS:

* Generators of an ADT are free if different terms always represent different data

EXERCISES: REVIEW UNIT 2

① Bool3 ADT

data Bool3 = True3 | False3 | Maybe3 deriving (Show)

Example for using: Function to combine 2 values value1 = True value2 = False value3 = Maybe

andBool3 :: Bool3 \rightarrow Bool3 \rightarrow Bool3.

andBool3 True3 True3 = True3

andBool3 False3 False3 = False3

andBool3 x False3 = False3

andBool3 x Maybe3 = Maybe3

② Bool10 ADT \rightarrow Not exist false, True is value < 10 but 0 = False.

Bool10 :: Bool10 \rightarrow Bool10 \rightarrow Bool10.

Bool10 (True10 x) (True10 y) = True10 (min x y)

Bool10 _ _ = False10

3) Nat ADT: greater, equal, even, odd, times, pow, gcd, sub, divide, mcd, mcm

greater:: Nat \rightarrow Nat \rightarrow Nat.

equal:: Nat \rightarrow Nat \rightarrow Bool

$$\text{greater zero } x = x$$

$$\text{greater } x \text{ zero} = x$$

$$\text{greater } (\text{suc } x) (\text{succ } y) = \text{suc } (\text{greater } x y)$$

times:: Nat \rightarrow Nat \rightarrow Nat

$$\text{times zero } x = \text{zero}$$

$$\text{times } (\text{suc } y) x = \text{add } x (\text{times } y x)$$

pow:: Nat \rightarrow Nat \rightarrow Nat

$$\text{pow } x \text{ zero} = \text{suc zero}$$

$$\text{pow } x (\text{suc } y) = \text{times } (\text{pow } x y) x$$

add:: Nat \rightarrow Nat \rightarrow Nat

$$\text{add } x \text{ zero} = x$$

$$\text{add } x (\text{suc } y) = \text{suc } (\text{add } x y)$$

factt:: Nat \rightarrow Nat

$$\text{factt zero} = \text{suc zero}$$

$$\text{factt } (\text{suc } x) = \text{times } \text{succ } x (\text{factt } x)$$

HACER DIAPOSITIVA 3.1 . QUEUES

REPASO:

Examen 8th November 2022.

④ Specification:

a) Sub I:: Inte \rightarrow Inte \rightarrow Inte

$$\text{subI } x z = P x$$

$$\text{subI } x S y = P (\text{sub } P x y)$$

$$\text{subI } x P y = S (\text{sub } S x y)$$

$$\text{subI } P x P y = \text{subI } (x y)$$

$$\text{subI } S x S y = \text{subI } (x y)$$

b) infixr 5:/ data stack a = ES | a:/ (stack a) deriving Show
infixl 5:/ data Queue a = EQ a | Queue : a deriving Show.

op22:: Queue a \rightarrow Stack a \rightarrow Stack a

$$\text{op22 } (q :: x) s = x:/ (\text{op22 } q s)$$

$$\text{op22 } q (x:y:s) = y:/ (\text{op22 } q x:s)$$

$$\text{op22 } _-_ _-_ = ES$$

b) Evaluate the following expressions

$$1) \text{op22 } (EQ :: 1:2:3) (4:/ES) = 3:/(\text{op22 } (EQ :: 1:2) (4:/ES)) = 3:/2:/(\text{op22 } EQ :: 1) 4:/ES$$
$$3:/2:/1:/(\text{op22 } EQ 4:/ES) = 3:/2:/1:/ES$$

$$2) \text{op22 } EQ (4:/5:/6:/ES) = 5:/6:/ES$$

$$3) \text{op22 } (EQ :: 1:2:3) (4:/S:/6:/ES) = 3:/2:/1:/(\text{op22 } EQ 4:/S:/6:/ES) = 321S6ES$$

2. IMPLEMENTATION:

- a) 2 efficient java methods
- 1) no more variables than the ones declared in the templates
 - 2) Try { ... } catch (TADVacioException ex) { ex.printStackTrace(); } must be used when necessary.
 - 3) Input in HIGH level method will remain unchanged after calling such method.
 - 4) LOW level methods will be able to alter the content of the input object thus inside the PD class

a) High level:

```
public static <E> Pila <E> op22 (Cola <E> q, Pila <E> s)
{
    Pila <E> sl = (Pila <E>) s. clone();
    Cola <E> ql = (Cola <E>) q. clone();

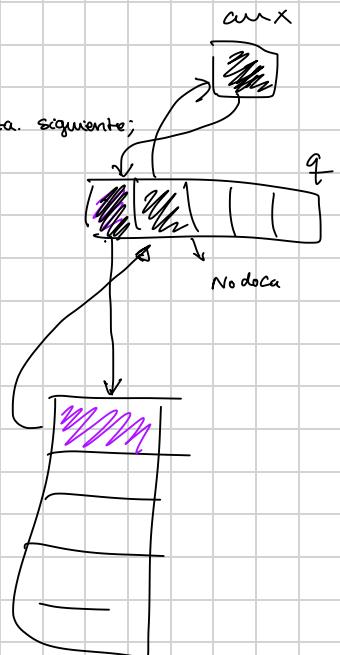
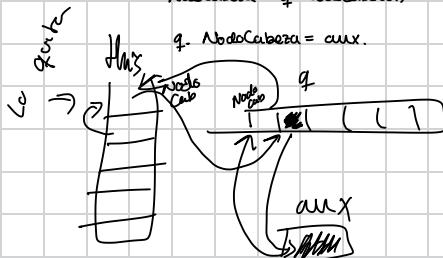
    try {
        if (!sl. EsVacia())
            sl = sl. Desapila();
        while (!ql. EsVacia())
            sl. Apila (ql. Cabeza());
            ql = ql. resto(); Me lo buego que saber.
    }
    catch (TADVacioException ex) { ex. PrintStackTrace(); }

    return sl;
}
```

b) public void op22 (CD <e> q)

```

    Nodo <E> aux = null;
    if (NodoCabeza != null)
        NodoCabeza = NodoCabeza. siguiente;
    while (q. NodoCabeza != null) {
        aux = q. NodoCabeza. siguiente;
        q. NodoCabeza. siguiente = NodoCabeza;
        NodoCabeza = q. NodoCabeza;
        q. NodoCabeza = aux;
    }
    q. NodoCabeza = aux.
}
```



```

Public void Desapilar () throws TADVacioException
{
    if (NodoCabeza == null) throws new TADVacioException();
    else { NodoCabeza = NodoCabeza.Siguiente;
    }
}

```

op22:: Queue a → Stack a → Stack a

op22 (q:.x) ES = x:/ES

op22 (q:.x) (y:/s) = x:/ (op22 q(y:/s))



op22 EQu (x:/y:/s) = y:/ (op22 EQu (x:/s))

op22 _____ = ES

Ex: op22 (EQu:.1:.2:.3) (4:/5:/6:/ES) = 3:/ (Op22 (EQu:.4:.2) (4:/5:/6:/ES)) = 3:/2:/1:/ (Op22 EQu (4:/5:/6:/ES)

(3:/2:/1:/5:/6:/ES) (Exam solution)

Ex: op22 (EQu:.1:.2:.3) (4:/5:/6:/ES) = 3:/2:/1:/5:/6:/ES

Luxan op22:: Queue a → Stack a → Stack a

op22 EQu ESt = ESt

op22 EQu (x:/s) = s

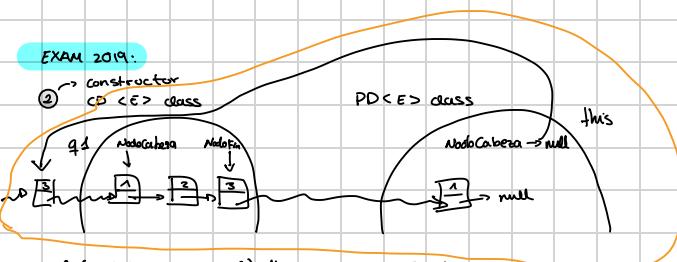
op22 (q:.x) s = x:/ op22 (q s)

EXAM 2019:

② constructor
CD < E > class

PD < E > class

this



if (nodoFinal = null) throws TADVacioException
else { do { nodoFinal.Siguiente = new NodoCE> (); null)
nodoFinal = nodoFinal.Siguiente.

this (after calling constructor)

EXAM NOVEMBER 2020C

2) queueexam :: Queue Int → Queue Int

queueexam (q:.x:.y) = (q:.y:.x)

queueexam _____ = queueexam (EQu:.1:.2)

queueexam (EQu:.3) = queueexam (EQu:.1:.2) = (EQu:.2:.1)

queueexam (EQu:.3:.2:.1) = (EQu:.3:.1:.2)

OTHER VERSION:

queueexam :: Queue Int → Queue Int

queueexam EQu { EQu:.2:.1

queueexam EQu:.x

queue (EQu:.x:.y) = EQu:.y:.x
queueexam q = (first q) .. queueexam (rest q) Necesito más operaciones auxiliares

b) HIGH LEVEL:

```
public static Cola<Integer> quevexam (Cola<Integer> q)
{
    Cola_Dinamica<Integer> sol = new Cola_Dinamica<Integer>();
    try {
        if (q.EstaVacia () || q.Resta().EsVacio())
            sol.Encola (2);
        sol.Encola (1);
    } else {
        while (!q.Resta().Resto().EsVacio())
            { sol.Encola (q.Cabecera());
              q=q.Resta();
              Integer x;
              x=q.Cabecera();
              q=q.Resta();
              sol.Encola (x); }
    }
    catch (TADVacioException ex) { System.out.println ("TAD vacio"); }
    return sol;
}
```

DATA STRUCTURE

1ST PROGRESS TEST October 30th, 2019

① Haskell:

```
data N1 = One | Next N1  
data InteN1 = None | Pos N1 | Neg N1.  
legN1 :: N1 -> N1 -> Bool  
legN1 One x = True  
legN1 (Next x) One = False  
legN1 (Next x) (Next y) = legN1 x y
```

b) legInteN1 :: InteN1 -> InteN1 -> Bool.

```
legInteN1 None None Pos x = True  
legInteN1 None Neg x = False  
legInteN1 Pos x Pos y = legInteN1 x y  
legInteN1 Neg x Neg y = legInteN1 x y  
legInteN1 Pos x Neg y = False  
legInteN1 Neg x Pos y = True.
```

c) add N1 :: N1 -> N1 -> N1

```
addN1 One x = Next x  
addN1 Next x y = Neg(addN1 x y)
```

d) addmex :: InteN1 -> N1 -> InteN1.

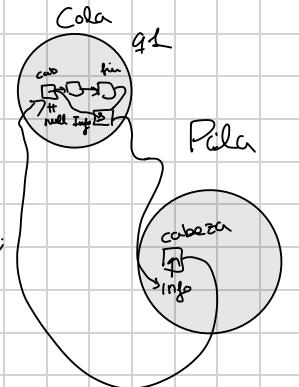
```
addmex None x = Pos x  
addmex (Pos x) y = Pos (addmex x y)  
addmex (Neg One) One = None  
addmex (Neg One) Next x = Pos x  
addmex (Neg(Next x)) One = Neg x  
addmex (Neg(Next x)) (Next y) = addmex (Neg x) y
```

② JAVA

```

public Pila_Dinamica(Cola_Dinamica < E > q)
{
    Cola_Dinamica < E > q1 = (Cola_Dinamica < E >) (q.clone());
    NodoCabeza = null;
    if (q1.NodoCabeza != null)
        { q1.NodoFinal.Siguiente = new Nodo < E > (q1.NodoCabeza.info, null);
        NodoCabeza = new Nodo < E > (q1.NodoFinal.info, q1.NodoCabeza); }
}

```

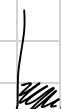


a) october 19 :: Cola a \rightarrow Pila a

october 19 CV = PV

12/12/1

october 19 q::x = x:/ (019 q::x ((cabeza q::-) :/ PV))



019 Cola a \rightarrow Pila a \rightarrow Pila a

019 CV p = p

019 q::x p = 019 q x:/ p

cabeza :: Cola a \rightarrow a

cabeza CV::a = a

cabeza q::- = cabeza q

october 19 (CV::1::3::2) = 2:/ 019 (CV::1::3::2)

(Cabeza q::-)/PV

1:/PV

2:/1:/3:/2:/1:/PV

c) public static < E > Pila < E > october 19 (Cola < E > q)

{ Cola < E > q1 = (Cola < E >) (q.clone());

Pila < E > p1, p2; p1 = new Pila_Dinamica < E >(); p2 = new Pila_Dinamica < E >();

if (!q1.EsVacia())

{

try { E first, aux; first = q1.Cabeza(); aux = first;

while (!q1.EsVacia())

{ aux = q1.Cabeza(); q1 = q1.resto(); p1.Apila (aux); }

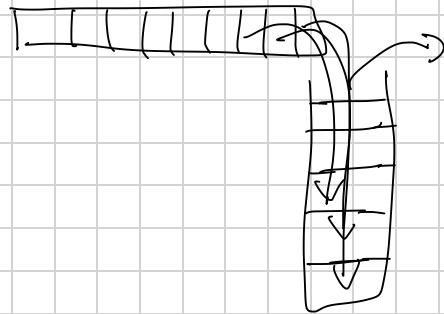
p2.Apila (first)

while (!p1.EsVacia()) { p2.Apila(p1.tope()); p1 = p1.Desopila(); } p2.Apila(first);

catch -- -- --

EXÁMENES

3rd November, 2022.



- ④ a) subI :: Inte → Inte → Inte .

$$\begin{array}{ll} \text{subI } x & z = x \\ \text{sub I } S(x) & S(y) = \text{subI } x \times y \\ \text{sub I } P(x) & P(y) = \text{subI } x \times y \\ \text{subI } x & P(y) = S(\text{subI } x \times y) \\ \text{subI } x & S(y) = P(\text{subI } x \times y) \end{array}$$

- b) op22 (EQu:..1:..2:..3) (4:/ES) → 3:/2:/1:/ES
 op22 (EQu) (4:/5:/6:/ES) → 5:/6:/ES.
 op22 (EQu:..1:..2:..3) (4:/5:/6:/ES) → 3:/2:/1:/5:/6:/ES

Queue stack stack

2

- a) public static < E > PCola < E > op22 (Cola < E > q, Pila < E > s)

↳ Pila < E > s2 = (Pila < E >) s. clone(); ↳ Para salvar lo que entra como parámetro)
 Cola < E > q1 = (Cola < E >) q. clone();

try {

```
if (!q1.EsVacia())
    s2 = s1. Desapila();
while (!q1.EsVacia())
    {
        int aux, first;
        s1.Apila(q1.Cabecera()); q1 = q1.Cabecera();
        {
            catch (TAD_Vacio_Exception ex) { ex.printStackTrace(); }
        }
    }
return s2;
```

- b) public void op22 (CD < E > q)

Nodo < E > aux = null;

if (nodoCabeza != null) NodoCabeza = NodoCabeza.Siguiente.

while (q.NodoCabeza != null) {

aux = q.NodoCabeza.Siguiente; q.NodoCabeza.Siguiente = NodoCabeza

NodoCabeza = q.NodoCabeza; q.NodoCabeza = aux; }

26 Octubre 2021:

1) NAND (negation of the conjunction)

$\text{nfBoolNAND} :: \text{BoolNAND} \rightarrow \text{BoolNAND}$

$\text{nfBoolNAND} \text{ Yes} = \text{Yes}$

$\text{nfBoolNAND} \text{ Nand Yes Yes} = \text{Nand} \text{ Yes Yes}$

$\text{nfBoolNAND} \text{ Nand Yes Yes} \text{ Nand Yes Yes} = \text{Nand} \text{ Yes Yes}$

$\text{nfBoolNAND} \text{ Nand} (\text{Nand} \text{ Yes Yes}) \text{ Yes} = \text{Yes}$

2) mix 2A :: Stack a \rightarrow Queue a \rightarrow Queue a

$\text{mix 2A} (x:/s) (q:y) = (\text{mix 2A} (y:/s) q) : x$

$\text{mix 2A} (x:/s) q = \text{mix 2A} s q$

$\text{mix 2A} \text{ ES } q = q$

a) $\text{mix 2A} \text{ ES } (\text{EQu}:3..4..5) \rightarrow (\text{EQu}:3..4..5)$

$\text{mix 2A} (1:/2:/\text{ES}) \text{ EQu} \rightarrow \text{EQu}$

$\text{mix 2A} (1:/2:/\text{ES}) (\text{EQu}:3) \rightarrow \text{EQu}:1$

$\text{mix 2A} (1:/2:/\text{ES}) (\text{EQu}:3..4..5) \rightarrow \text{EQu}:4..5 \cancel{..1}$

A) HIGH LEVEL

public static <E> Cola <E> mix2A (Pila <E> s, Cola <E> q)

{ Pila <E> sr = Pila <E> s.clone();

Cola <E> qr = Cola <E> q.clone();

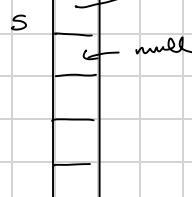
{ Scan Buckle ?

q

while /if (!sr.EsVacia() && !qr.EsVacia())

| qr = qr.Rastoc(); qr.Enqueue(sr.Topo()); }

return qr;



b) low level:

public void mix2A (PD <E> s)

if (NodoCabeza != null && s. NodoCabeza != null)

Public Nodo <E> = new Nodo<E> aux;

aux = NodoCabeza;

s. NodoCabeza.Siguiente = null;

NodoFinal.Siguiente = s. NodoCabeza

NodoFinal = NodoFinal.siguiente

NodoCabeza = NodoCabeza.siguiente;

30th October, 2017

1) times (specified) nf-Inte

data Nat = zero | Suc Nat

times:: Nat → Nat → Nat

fact:: Nat → Nat

fact: zero → SucZero.

fact: Suc(x) → times (fact x) (Suc x)

data Inte = z | s Inte | p Inte

nf-Inte:: Inte → Inte.

abs I, aux-abs I:: Inte → Nat

abs I x = aux-abs I (nf-Inte x)

aux-abs I:: Inte → Nat.

aux-abs I z → zero.

aux-abs I s(x) → Suc (aux-abs x)

aux-abs I p(x) → Suc (aux-abs x)

factabs I:: Inte → Inte.

factabs I z = zero

factabs I x = aux-fact (fact (abs-+ x))

aux-fact:: Nat → Inte

aux-fact zero = z

aux-fact Suc(x) = s (aux-fact x)

2) public static Pila < Integer > interRPN (Pila exp) (HIGH LEVEL)

```
{ Pila aux = new PD();  
Pila < Integer > sol = new PD< Integer >();  
Integer x, y, z;  
try { while (!exp.EsVacia())  
{ if (exp.Tope() instanceof Integer)  
    aux.Apila (exp.Tope());  
else { x = (Integer) aux.Tope();  
    aux = aux.Desapila();  
    y = (Integer).Tope();  
    aux = aux.Desapila();  
    z = (eval(y, (char) exp.Tope(), x));  
    aux.Apila (z); sol.Apila (z);  
}  
exp = exp.Desapila();  
sol.sol.DesApila();  
} catch (TAD... ) { sout (TAD... );  
return sol;
```

3) NO REPEAT.

HAS KELL:

create:: (Eq a) \rightarrow Queue a \rightarrow Quaselist a

create Eq = EqL

remove:: (Eq a) \rightarrow a \rightarrow Quaselist a

create q := x = IgL x (remove x (create q)) remove — EqL = EqL

remove x (IgL y q1) = < if x == y q1
else IgL y (remove x q1)

Eq: 1: 2: 1: 3

IgL 3 (remove 3 (IgL 1: (remove 1 (IgL 2 (remove 2 (IgL 1 remove 1 EqL)

IgL 3 IgL 1 IgL 2 EqL

JAVA:

```
public class Quaselist <E> {  
    public int longitud;  Public Nodo <E> NodoCabeza;  
    public Quaselist (ColaDinamica <E> q)  
    {  
        longitud = 0; nodoCabeza = null; Nodo <E> aux 1 = q. NodoCabeza;  
        while ( aux 1 != null)  
        {  
            Nodo <E> aux 2 = aux 1. siguiente;  
            while ( aux 2 != null && !aux 2. info. equals (aux 1. info))  
                aux 2 = aux 2. siguiente;  
            if (aux 2 == null) longitud ++;  
            NodoCabeza = New Nodo <E> (aux 1. Info, NodoCabeza);  
            aux1= aux1. siguiente;  
        }  
    }  
}
```

5th November 2018

4. Nat ADT $\text{ReqN } x \ y = x + y \text{ // addN } \& \text{ subN (sum negatives)}$

$\text{IN} = \text{Pos Nat} \ | \ \text{Neg Nat}$

$\text{abs} :: \text{IN} \rightarrow \text{IN}$

$\text{abs } x = x$

$\text{abs Neg}(x) = \text{Pos}(x)$

$\text{ReqN} :: \text{IN} \rightarrow \text{IN} \rightarrow \text{Bool}$

$\text{ReqN pos}(x) \ pos(y) = \text{ReqN } x \ y$

$\text{Req pos}(0) \ neg(0) = \text{True}$

$\text{Req pos}(x) \ neg(y) = \text{False}$

$\text{Req neg}(x) \ pos(y) = \text{True}$

$\text{Req neg}(x) \ pos(x) = \text{Req } x \ y$

$\text{Req} :: \text{Nat} \rightarrow \text{Nat} \rightarrow \text{Bool}$

$\text{Req } 0 \ 0 = \text{True}$

$\text{Req } 0 \ \text{Suc}(x) = \text{True}$

$\text{Req Suc}(x) \ 0 = \text{False}$

$\text{Req Suc}(x) \ Suc(y) = \text{ReqN } x \ y$

$\text{addIN} :: \text{IN} \rightarrow \text{IN} \rightarrow \text{IN}$

$\text{addIN } x \ 0 = x$

$\text{addIN pos}(x) \ pos(y) = \text{addN } x \ y$

$\text{addIN pos}(x) \ Neg(y) = \begin{cases} \text{Req } y \ x : pos(\text{sub } x \ y) \\ \text{else neg } (\text{sub } y \ x) \end{cases}$

$\text{addIN Neg}(x) \ Neg(y) = \text{neg } (\text{add } x \ y)$

$\text{addIN neg}(x) \ pos(y) = \begin{cases} \text{Req } y \ x : neg(\text{sub } x \ y) \\ \text{else pos } (\text{sub } y \ x) \end{cases}$

2. Evaluate arithmetic expression

public static Pila plusPN(Pila exp)

{ Pila aux = new Pila(); Integer x, y;

try { while (!exp.EsVacia()) {

if (exp.Toper() instanceof Integer && !aux.EsVacia() && aux.Toper() instanceof Integer && aux.DesApila().equals(Toper().toString()))

x = (Integer) aux.Toper(); aux = aux.DesApila();

Toper.equals(newPila());

y = (Integer) exp.Toper(); exp = exp.DesApila();

aux = aux.DesApila(); exp.Apila(x+y);

else { aux.Apila(exp.Toper()); exp = exp.DesApila(); }

while (!aux.EsVacia()) {

if (aux.Toper instanceof Character) exp.Apila((char) aux.Toper());

else { exp.Apila(Aux.Toper()); aux = aux.DesApila(); }

return exp; catch (TAD vacio exc e) { Sout("Empty "+e); return exp; }

$$3) a) \text{lqueuel } [(CV..2..3..4) : CV : (CV..3..4)] = \text{lqueuel } [(CV..2..3..4) (CV..3..4)] = [1:3:2:4;3]$$

b) HIGH LEVELS

```
public static <E> Lista <E> lqueuel (Lista <Cola <E>> lg) throws TADVacioException
{
    Lista <E> l = new Lista <E>();
    try { if (!lg.esVacia())
        { Cola <E> q1 = lg.Cabeza(); l = lqueuel(lg.Cola());}
        while (!q1.esVacia())
            { l.Atrape(q1.Cabeza()); q1 = q1.Rostro();}
    }
    return l;
}
```

c) Haskell

```
lstack l :: [Polar a] → [a]
lstack [] = []
lstack l [es:t] = lstack t
lstack l [x:g:t] = (lstack l [g]) ++ (x:(lstack t))
```

30th October 2019

$$a) \text{legN1} : \text{N1} \rightarrow \text{Bool}$$

$$\text{leg N1 One Next x} = \text{True}$$

$$\text{leg Next x One} = \text{False}$$

$$\text{leg Next x Next y} = \text{legN1 x y}$$

$$\text{legIntN1} :: \text{IntN1} \rightarrow \text{IntN1} \rightarrow \text{Bool}$$

$$\text{legIntN1 None None} = \text{True}$$

$$\text{legIntN1 None Neg x} = \text{False}$$

$$\text{legIntN1 Neg x None} = \text{True}$$

$$\text{legIntN1 Pos x None} = \text{False}$$

$$\text{legIntN1 Pos x Pos y} = \text{legN1 x y}$$

$$\text{legIntN1 Neg x Neg y} = \text{legN1 y x}$$

$$\text{addN1} :: \text{N1} \rightarrow \text{N1} \rightarrow \text{N1}$$

$$\text{addN1 x One} = \text{Next x}$$

$$\text{addN1 x Next y} = \text{add Next x y}$$

$$\text{addmex} :: \text{oneN1} \rightarrow \text{N1} \rightarrow \text{IntN1}$$

$$\text{addmex None x} = \text{Pos x}$$

$$\text{addmex Pos x y} = \text{Pos addN1(x y)}$$

$$\text{addmex Neg(One) One} = \text{None}$$

$$\text{addmex Neg(One) Next x} = \text{Pos x}$$

$$\text{addmex Neg(Next x) One} = \text{Neg x}$$

$$\text{addmex Neg(Next x) (Next y)} = \text{Addmex (neg x y)}$$

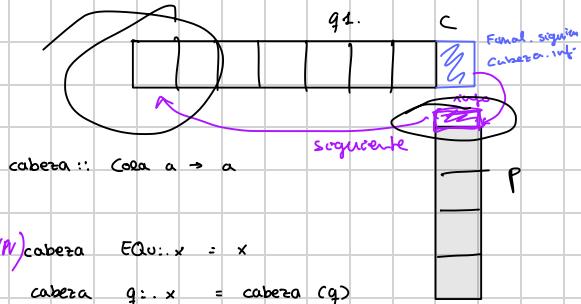
2)

a) Haskell rules:

october 19:: Cola a → Pila a

october 19 EQu = ES

october 19 q:: x = x:/ Cola a (q:: x) ((cabera (q:: r)):/ M) cabera EQu:: x = x

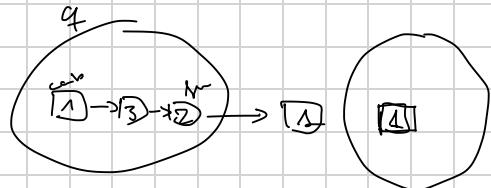


019:: Cola a → Pila a → Pila u

019 EQu p = p

019 q:: x p = 019 q x:/ p

october 19 (CV, 1; 3; 2) = 2 / 1:/ 3; 2 / 1:/ ES



public static <E> Pila<E> october19 (Cola <E> q)

{ Cola <E> q1 = new Cola <E> (q.Clone()); Pila <E> p1 = new Pila-Dinamica <E> ();
try { if (!q1.EsVacia()) { p1.push(q1.Pop()); } } catch (Exception e) { e.printStackTrace(); }

} Pr. Apila (q1.Cabera());

while (!q1.EsVacia())

} Pr. Apila (q1.Cabera()); Aux = q1.Cabera();

q1 = q1.Restos();

{

while (!p1.EsVacia())

} p1.Apila (p1.Top());

p1 = p1.DesApila();

{

p1.Apila (Aux);

{ catch

return p1;

1) $\text{nfBN} :: \text{BoolThreeNat} \rightarrow \text{BoolThreeNat}$

$$\text{nfBN} (\text{BN zero zero}) = \text{BN zero zero}$$

$$\text{nfBN} (\text{Suc } x \text{ zero}) = \text{BN Suc zero zero}$$

$$\text{nfBN} (\text{Zero Suc } x) = \text{BN zero Suc zero}$$

$$\text{nfBN} (\text{BN Suc } x \text{ Suc } y) = \text{nfBN} (x \text{ } y)$$

$\text{andBN} :: \text{Bool} \rightarrow \text{Bool} \rightarrow \text{Bool} \rightarrow \text{Bool}$

$$\text{andBN} (\text{B zero zero}) = \text{B zero} (\text{Suc zero}) = \text{Zero Suc zero}$$

$$\text{andBN} (\text{B zero} (\text{Suc zero})) = \text{B zero} (\text{Suc zero})$$

$$\text{andBN} (\text{---}) = \text{Zero zero}$$

$$\text{andBN} (\text{Suc zero zero}) = X$$

$$\text{andBN} (\text{BN } x \text{ BN } y) = \text{andBN} (\text{nfBN BN } x) (\text{nfBN BN } y)$$

2) a) $\text{quevexam} (\text{EQu}: 3) = \text{EQu}: 2 : 1$

$$\text{quevexam} (\text{EQu}: 3 - 2 : 1) = \text{EQu}: 3 : 1 : 2$$

b) HIGH LEVEL

```
public static Cola<Integer> quevexam (Cola<Integer> q)
{
    CD<Integer> sol = new CD<Integer>();
    try { if (q.EsVacia() || q.Restos().EsVacia())
        | sol.Encola(2);
        | sol.Encola(1);
    }
    else { int x, y;
            while (!q.SiRestoC().RestoC().EsVacia())
                sol.encola(q.Cabecera()); q=q.Restos();
            x=q.Cabecera(); q=q.Restos();
            sol.Encola(q.Cabecera()); q=q.Restos();
            sol.Encola(x);
        }
    catch (TAD_VacioException exp) { exp.sout("empty queue"); }
    return sol;
}
```

LOW LEVEL

```
public void preexam()
{
    if (NodoCabeza == null || NodoCabeza.Siguiente == null)
        NodoFinal = new Nodo(1, null);
    NodoCabeza = new Nodo(2, NodoFinal); }

else {
    Nodo <E> aux = NodoCabeza;
    while (aux.Siguiente != NodoFinal)
        aux = aux.Siguiente();
    E x = aux.Info(); aux.Info = NodoFinal.Info; NodoFinal.Info = x; }
```

October (6th 2021)

nfBoolNand :: BoolNand → BoolNand

nfBoolNand yes = yes

nfBool Nand Nand yes yes = Nand yes yes

nf BoolNand Nand yes (Nand yes yes) = yes

nfBoolNand (Nand (Nand yes yes) Yes = yes

nfBoolNand Nand (Nand yes yes)(Nand yes yes) = Yes

nf Bool Nand (Nand x y) = nfBoolNand (Nand (nfBoolNand x) (nfBoolNand y))

(2) a) $1 = \text{EQv} \vdash 3 \vdash 4 : S$ $3 = (\text{num}21 \quad (3; 2; / \text{ES} \quad \text{EQv})) \therefore 1 = \text{EQv} \vdash 1$
 $2 = \text{EQv}$ $4 = (\text{EQv} \quad) \therefore 4 : S \vdash 1$

b) HIGH LEVEL

```
public static <E> Cola <E> max21 (Pila <E> s, Cola <E> q)
{
```

```
    Cola q2 = new Cola <E> (q.clone());
    if (!s.EsVacia && !q2.EsVacia())
        q2 = q2.Pisto(); q2.Encola (s.Tope()); }

    return q2;
```

LOW LEVEL

```
public void mux2A (POC E> s)
if ( NodoCabeza != null && s. NodoCabeza != null)
    { NodoCabeza = NodoCabeza. Siguiente.
      NodoFinal. Siguiente = new Nodo <E> ( s. NodoCabeza. Info, null);
```

CONSTRUCTORS

- No returns
- Can combine LOW and HIGH level.
- Domains = LOW, interfaces = HIGH

HASKELL:

exam:: [a] → [a] → [a] (from constructor of the exam)

Haskell specification for JAVA CODE of trees:

$$\begin{aligned} \text{part} &:: \text{Arbol a} \rightarrow \text{Arbol a} \rightarrow \text{Arbol a} \\ \text{part} &\quad \text{AV} \quad \text{AV} = \text{AV} \\ \text{part} &\quad (\text{AB} \ r \ --) \text{ AV} = \text{AB} \ r \ \text{AV} \ \text{AV} \\ \text{part} &\quad - \ (\text{AB} \ r \ --) = \text{AB} \ r \ \text{AV} \ \text{AV} \end{aligned}$$

exam:: [a] → [a] → [a]

exam [] l₂ = l₂

exam x:l₁ l₂ = exam (l₁) x:l₂)

exam:: [a] → [a] → [a]

exam2 EL EL = EL

exam2 EL L = L

exam2 L — = (reverse L) ++ —

Class:

aux:: [a] → [a] → [a]

aux [] l = l

aux x:l₁ l₂ = aux (l₁) (x:l₂)

reverse:: [a] → [a]

reverse EL = EL

reverse (x:y:l) = (reverse (x:l)) : y

Haskell specification for JAVA code of trees:

part:: Arbol a \rightarrow Arbol a \rightarrow Arbol a

part AV AV = AV

part (AB r --) AV = AB r AV AV

part _ (ABr --) = AB r AV AV

exam:: [a] \rightarrow [a] \rightarrow [a]

exam [] l₂ = l₂

exam x:l₂ l₂ = exam (l₂) x:l₂)

Class:

aux:: [a] \rightarrow [a] \rightarrow [a]

aux [] l = l

aux x:l₁ l₂ = aux (l₁) (x:l₂)

exam2:: [a] \rightarrow [a] \rightarrow [a]

exam2 [] EL = EL

exam2 EL l = l

exam2 l _ = (reverse l) ++ _

reverse:: [a] \rightarrow [a]

reverse EL = EL

reverse (x:y:l) = (reverse (x:l)) : y

```
public ABD(ArbolBinario <E> a1, ABD<E> a2) {
    if (!a1.EsVacio() || a2.raiz != null)
        if (a2.raiz == null && !a1.EsVacio())
            this.raiz = new NodoArbol <E> (a1.Raiz(), null, null);
        else if (a1.EsVacio && a2.raiz != null)
            this.raiz = new NodoArbol <E> (a2.raiz.info, null, null);
        else
            if (a2.raiz.info.compareTo(a1.Raiz()) > 0)
                this.raiz = new NodoArbol <E> (a1.Raiz(), null, a2.raiz.info);
            else if (a2.raiz.info.compareTo(a1.Raiz()) < 0)
                this.raiz = new NodoArbol <E> (a1.Raiz(), a2.raiz.info, null); ↓
            else ↓ this.raiz = null;
    }
    catch (TAD_Vacio_Exception exp) {exp.printStackTrace();}
}
```

Que el compareTo() de menor que 0

Significa que el objeto this utilizado es menor que el encontrado en los parentesis.

JAVA EXERCISE

If both have a root: the `ArbolBinario = new root`; and if $AB < r$ it will be left else right

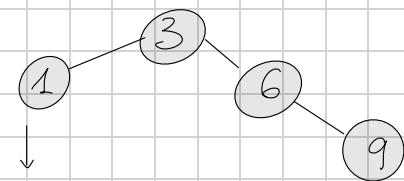
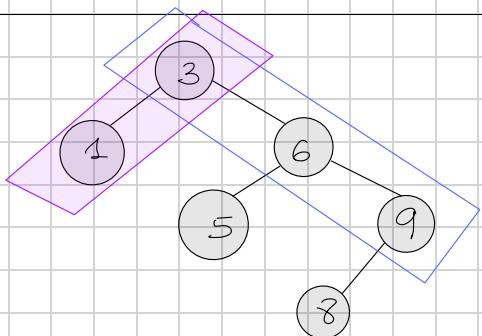
HASKELL

`data Arbol a = AV | AB a (Arbol a) (Arbol a)`

`belong :: (Ord a) → a → Arbol a → Bool.`

`belong x AV = False`

`belong x (AB r i d) = if $x == r$ = True
 $| x < r = \text{belongs } x \text{ i}$
 $| x > r = \text{belongs } x \text{ d}$`



$(AB\ 3\ (AB\ 1\ AV\ AV))\ (AB\ 6\ AV\ (AB\ 1\ AV\ AV))$

`mimax :: (Ord a) → Arbol a → Arbol a`

`mimax AV = AV`

`mimax (AB r i d) = AB r (aux True i) (aux False d)`

`aux :: Bool → Arbol a → Arbol a`

`aux — AV = AV`

`aux True (AB r i d) = AB r (aux True i) AV`

`aux False (AB r i d) = AB r AV (aux False d)`

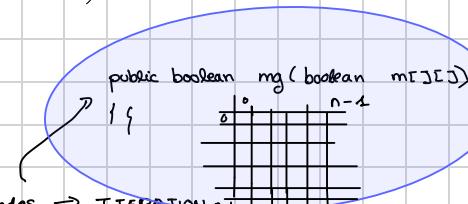
EXAM

GRÁFICOS: Nodos, Adyacencias, numnodos → ITERACIÓN $n-1$

GRAPHICS JAVA: \rightarrow LOW LEVEL ADYACENCIA.

\rightarrow HIGH LEVEL \rightarrow Nodos(); lista(E) adyacentes

METHOD PROPOSED WITH GRAPH will be something like \rightarrow (Constructor) `public GNODE (Lista<E> q) {}`



SECOND PROGRESS TEST

1. SPECIFICATION:

a) dec22:: Grafo A a \rightarrow Grafo A a

$$i) \text{dec22} (\text{Nodo } 1 \text{ (Arco } 1 \ 1 \text{ GVaux)}) \Rightarrow \text{dec22} (\text{Arco } 1 \ 1 \ (\text{Nodo } 1 \text{ GVaux})) = \text{dec22} (\text{Nodo } 1 \text{ GV})$$

$$= \text{Nodo } 1 \text{ GVaux}$$

$$ii) \text{dec22} (\text{Nodo } 1 \text{ (Arco } 2 \ 1 \ (\text{Nodo } 2 \text{ GVaux}))) \Rightarrow \text{dec22} (\text{Arco } 2 \ 1 \ (\text{Nodo } 1 \ (\text{Nodo } 2 \text{ GV}))) =$$

$$\text{dec22 } \text{Nodo } 1 \ (\text{Nodo } 2 \text{ GV}) = \text{Arco } 1 \ 2 \ (\text{Nodo } 2 \ (\text{dec22} (\text{Nodo } 1 \text{ GV})))$$

$$\text{Arco } 1 \ 2 \ (\text{Nodo } 2 \ (\text{Nodo } 1 \text{ GV}))$$

$$iii) \text{dec22} (\text{Nodo } 1 \text{ (Arco } 1 \ 2 \ (\text{Arco } 2 \ 3 \ (\text{Arco } 3 \ 2 \ (\text{Nodo } 3 \text{ GV})))))) \Rightarrow$$

$$\text{dec22} (\text{Arco } 1 \ 2 \ (\underbrace{\text{Nodo } 1 \ (\text{Arco } \dots)}_{g})) \Rightarrow$$

$$\text{dec22} (\text{Nodo } 1 \ (\text{Arco } g))$$

$$\text{dec22} (\text{Arco } 2 \ 3 \ (\text{Nodo } 1 \ (\text{Nodo } 2 (\text{Arco } 3 \ 2 \ (\text{Nodo } 3 \text{ GV}))))))$$

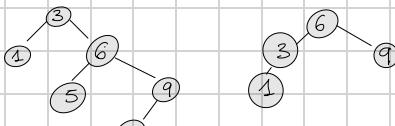
$$\text{dec22} (\text{Nodo } \underbrace{1}_{x} \ (\text{Nodo } \underbrace{2}_{y} (\text{Arco } \underbrace{3 \ 2}_{g} (\text{Nodo } 3 \text{ GV})))) =$$

$$\text{Arco } 1 \ 2 \ (\text{Nodo } 2 \ (\text{dec22} (\text{Nodo } 1 \ (\text{Arco } 3 \ 2 \ (\text{Nodo } 3 \text{ GV}))))))$$

$$\text{Arco } 1 \ 2 \ (\text{Nodo } 2 \ (\text{dec22} (\text{Nodo } 1 \ (\text{Nodo } 3 \text{ GV}))))$$

$$\text{Arco } 1 \ 2 \ (\text{Nodo } 2 \ (\text{Arco } 1 \ 3 \ (\text{Nodo } 3 \ (\text{Nodo } 1 \text{ GV}))))$$

b) mimax:: Arbol a \rightarrow Grafo A a



$$\text{mimax AV} = \text{GV}$$

$$\text{mimax } (AB \times i \ d) = \text{aux False } \times d \quad (\text{mimax True } \times i \ (\text{Nodo } \times 6V))$$

aux:: Bool \rightarrow a \rightarrow Arbol a \rightarrow Grafo a \rightarrow Grafo c

$$\text{aux } _ _ \text{ AV } g = g$$

$$\text{aux True } \times (AB \ y \ i \ -) \ g = \text{mimax True } \times i \ (\text{Arco } \times g \ (\text{Nodo } y \ g))$$

$$\text{aux False } \times (AB \ y \ - \ d) \ g = \text{mimax False } \times d \ (\text{Arco } \times g \ (\text{Nodo } y \ g))$$

no deberia ser y?

DUDA

$\text{memor}(\text{AB } 3 (\text{AB } 1 \text{ AV } \text{AV}) (\text{AB } 6 (\text{AB } 5 \text{ AV } \text{AV}) (\text{AB } 9 (\text{AB } 8 \text{ AV } \text{AV}) \text{AV})) \rightarrow q$
 aux False 3 \times $(\text{AB } 6 (\text{AB } 5 \text{ AV } \text{AV}) (\text{AB } 9 (\text{AB } 8 \text{ AV } \text{AV}) \text{AV}))$ $(\text{aux } 3 (\text{AB } 1 \text{ AV } \text{AV} (\text{Node } 3 \text{ GV}))$
 aux False 3 $(\text{AB } 9 (\text{AB } 8 \text{ AV } \text{AV}) \text{AV})$ Arco(3 6 (Node 6 (aux 3 (AB 1 AV AV (Node 3 GV)))))
 aux False 3 AV (Arco 3 9 (Node 9 (Arco(3 6 (Node 6 (aux 3 (AB 1 AV AV (Node 3 GV)))))))
 Arco(3 9 (Node 9 (Arco 3 6 (Node 6 (aux True 3 AV (Arco 3 1 (Node 1 (Node 3 GV)))))))
 Arco(3 9 (Node 9 (Arco 3 6 (Node 6 (Arco 3 1 (Node 1 (Node 3 GV))))))))

EXAMEN HASKEEL → NO GRAPHS

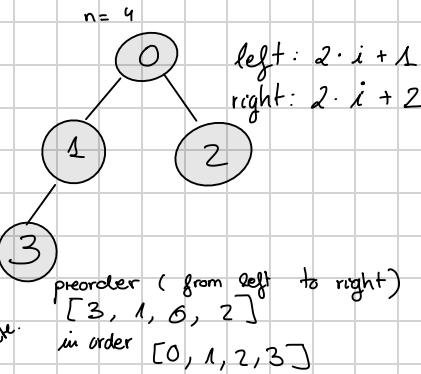
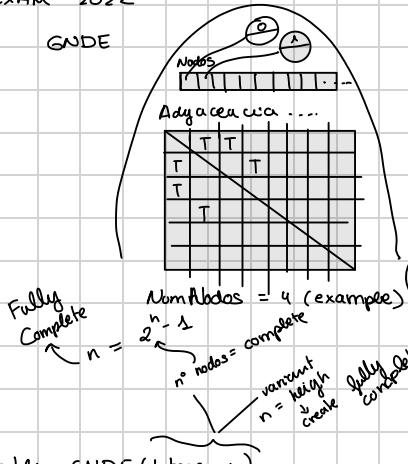
↳ NO GENERAL TREES

DIFFERENCE BETWEEN COMPLETE AND FULLY COMPLETE

COMPLETE: Holes in last level
FULLY COMPLETE: No holes.

MANDATORY ITERATIVE

EXAM 2022



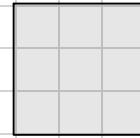
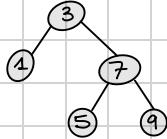
```

public GNDE (Integer n)
{
    Adyacencias = new boolean [MAX_NODOS] [MAX_NODOS];
    Nodos = new NodoGrafo [MAX_NODOS]; numnodos = 0;
    //Complete.
}

```

OTHER

a) [1, 3, 5, 7, 9]



testAVL::(Arbol a) → Bool

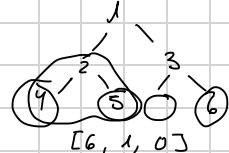
testAVL AV = True

testAVL (AB r x a) = if abs ((height i) - (height d)) > 1 then False.
else (test AVL i) && (test AVL d)EKAN December 13th 2022.

a) ArbolBinario.

```

public static <E> Lista<Integer> hole (ArbolBinario<E> t)
{
    Integer s, r; Lista<Integer> sol, left, right; sol = new LD();
    if (t. EsVacio()) sol. Arriade(0);
    else try {left = hole(t.subArbolIzquierdo()); right = hole(t.subArbolDerecho());
    l = left. Cabeza(); r = right. Cabeza();
    if (l < r) {sol = left. Cola(); sol. Arriade(0);}
    else if (l > r) {sol = right. Cola(); sol. Arriade(4);}
    sol. Arriade(l + r + 1)}
  
```

4 < 5
sol = 4Input → Complete tree. 1st hole
left = 0 Right = 1.

PREGUNTA

b) ITERATIVE CONSTRUCTOR

```

public GNODE (Arbol Binario<E> t)
{
  
```

```

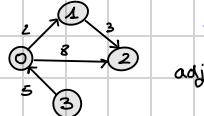
try{while (!t. Es Vacio())
  t. Raiz()=
```

PRACTICES:

~~NO SIDE~~

(We only work with adjacency matrix \rightarrow FLOYDS ALGORITHM

↑
No weighted
(Adjacency boolean)
No directed



adj

Weighted
(Directed
(symmetrical adj.)
(Integers in adjacency)

path weight

\rightarrow To go from 0 to 2 : $[0, 2] \rightarrow 8$

$[0, 1, 2] \rightarrow 2+3=5$

best path \rightarrow LOWEST WEIGHT

HIGHEST WEIGHT

	0	1	2	3	4	5	6
0	I	2	8	I			
1	I	I	3	I			
2	I	I	I	I			
3	5	I	I	I			
4					I		
5					I		
6						I	

Major poner 0 en la diagonal.

* always arrow with a node and itself.

Integer I = Integer. MAX_VALUE $\sim \infty$

adj

WEIGHT

WEIGHT

WEIGHT

WEIGHT

WEIGHT

WEIGHT

	0	1	2
0			5
1			
2			
3			
4			
5			
6			

path

	0	1	2
0			
1			
2			
3			
4			
5			
6			

best path
LOWEST
WEIGHT.

lists of
paths.

lpc::: Arbol a \rightarrow Int.

lpc AV = 0

lpc AB - AV AV = 1

lpc (AB - i AV) = -1

lpc (AB - AV d) = -1

lpc CAB - i d) = if (lpc i == -1) || (lpc d == -1) then = -1
else lpc i + lpc d

lpc::: Arbol a \rightarrow Int

lpc AV = 0

lpc (AB - i d) = aux(lpc i) (lpc d)

aux::: Int \rightarrow Int \rightarrow Int

aux li ld = if (li == -1) || (ld == -1)
else li + ld.

b) ITERATIVE \rightarrow GNDE

SEARCH
TREE
maxmax

public GNODE (ArbolBinario <E> t) ITERATIVE \rightarrow methods of Arbol Binario.

adyacencias

| Adyacencias = new boolean [MAX_NODOS][MAX_NODOS];

Nodos = new NodoGrafo [MAX_NODOS];

numnodos = 0; ArbolBinario <E> aux;

if (!t. EsVacio())

try { Nodos [0] = new NodoGrafo (t. Raiz(), false);

numnodos++; aux = t. SubArbolIZquierdo();

while (!aux. EsVacio())

{ Nodos [numnodos] = new NodoGrafo (aux. Raiz(), false);

Adyacencias [numnodos][numnodos - 1] = Adyacencias [numnodos - 1][numnodos] = true;

numnodos++; aux = aux. SubArbolIZquierdo();

if (!t. SubArbolDcho(). EsVacio ())

{ Nodos [numnodos] = new NodoGrafo (t. SubArbolDcho(). Raiz(), false);

Adyacencias [numnodos][0] = Adyacencias [0][numnodos] = true;

numnodos++; aux = t. SubArbolDcho();

while (!aux. EsVacio())

{ Nodos [numnodos] = new NodoGrafo (aux. Raiz(), false);

Adyacencias [numnodos][numnodos - 1] = Adyacencias [numnodos - 1][numnodos] = true;

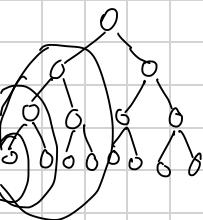
numnodos++; aux = aux. SubArbolDcho();

}

}

catch (TAD_VacioException ex) {ex. printStack (trace());}

0	1	2	3	4	5
0					
1					
2					
3					
4					
5					



IMPORTANT \rightarrow

30th November 2021.

1) No Graphs and no general trees.

2) public GNODE (Integer n)

```
{ Adyacencias = new boolean [ MAX_NODOS ][ MAX_NODOS ];
  Nodos = new NodoGrafo [ MAX_NODOS ]; numnodos = 0;
  for ( i = 0; i < n; i++ )
    { Nodos [ i ] = new NodoGrafo ( i, false );
      for ( i = 0; i < n; i++ )
        { if ( i * 2 + 1 < n ) Adyacencias [ i ][ i * 2 + 1 ] = Adyacencias [ 2 * i + 1 ][ i ] = true;
          else if ( i * 2 + 2 < n ) Adyacencias [ i ][ i * 2 + 2 ] = Adyacencias [ 2 * i + 2 ][ i ] = true;
          else break;
        }
    }
}
```

18th December 2020

My way

ITERATIVE

```
public Lista_Dinamica <E> FromLevel ( int level )
{ int counter_level = 1; aux = new ArbolDinamico(); aux2 = new ArbolDinamico(); sol = new ListaDinamica();
  if ( counter_level >= level )
    { sol . Añade ( thus . Raiz () );
      try { while ( ! thus . EsVacio () );
        { counter_level ++; aux = thus . SubArbolIzqdo (); aux2 = thus . SubArbolDcho ();
          if ( counter_level >= level )
            { sol . Añade ( aux . Raiz () );
              sol . Añade ( aux2 . Raiz () );
            }
        }
      }
    }
  return sol }
```

public Lista_Dinamica <E> FromLevel (int level)

Solucion

RECURSIVE

```
{ lista_Dinamica <E> sol = new Lista_Dinamica ();
  if ( ! EsVacio () )
    { try { lista_Dinamica aux = ((ArbolBinarioDinamico <E>) SubArbolDcho ()). FromLevel ( level - 1 );
      if ( level <= 1 ) sol . Añade ( Raiz ());
      sol = ((Arbol Binario Dinamico <E>) SubArbolIzqdo ()). FromLevel ( level - 1 );
      sol . Concatena ( aux );
    } catch ( TAD_VacioException exp ) { exp . printStackTrace ();
    }
  }
  return sol }
```

```

public GNODE (ArbolBinario < E > t)
{
    Adyacencias = new boolean [MAX_NODOS][MAX_NODOS]; Nodos= new NodoGrafo [MAX_NODOS]; numnodos = 0;
    if (!t.EsVacio())
        try { Nodos [0] = new NodoGrafo (t.Raiz(), false); numnodos = 1;
               aux20 (t.SubArbolIzqdo, 0); aux2 (t.SubArbolDcho, 0);
            } catch (TADVacioException exp) { System....(exp) }
    }
}

```

```

private void aux20 (ArbolBinario < E > t, int fat)
{
    if (!t.EsVacio())
        try { Nodos [numnodos] = new NodoGrafo (t.Raiz(), false);
               Adyacencias [numnodos][fat] = Adyacencias [fat][numnodos] = true;
               fat = numnodos; numnodos++;
               aux20 (t.SubArbolIzqdo, fat); aux (SubArbolDcho, fat);
            } catch (TAD Vacio Exception exp) { System....(exp); }
}

```

27th November 2019.

lemaleaf :: ArbolG a → Int

lemaleaf AGV = 0

lemaleaf AG - [] = 1

lemaleaf AG - [i:] = 1 + lemaleaf t

remoleaf :: ArbolG a → Int

remoleaf AGV = 0

remoleaf AG r [] = 1

remoleaf AG - [g] = 1 + remoleaf g

remoleaf AG r [- : t] = remoleaf (AG r t)

alleaves :: ArbolG → Bool

alleaves AGV = True

alleaves AG r [] = True

alleaves g = auxal (lemaleaf g) g

auxal Int → ArbolG a → Bool

auxal n AG - [] = if n == 1 → True else False

auxal n Ag r [g] = auxal (n-1) g

auxal n Ag r (g:t) = [auxal (n-1) g] && [auxal n (Ag r t)]

```

public static Integer minimum (ArbolBinario < Integer > t) throws TADValueException
{
    if (!t.esVacio())
        try { if (t.SubArbolIzquierdo().esVacio()) return t.raiz();
               else return minimum(t.SubArbolIzquierdo());
            }
    }

public void replaceRoot (NodoArbol < Integer > t, Integer y) throws IllegalArgumentException
{
    try { if (n.izq != null && (y <= (Integer) maximum (new ArbolBinarioDinamico (n.izq))))
          throw new IllegalArgumentException();
          if (n.dre != null && (y >= (Integer) minimum (new ArbolBinarioDinamico (n.dre))))
              throw new IllegalArgumentException();
          else n.valor = y;
      } catch (TADValueException exp) { Sys... (exp); }
}

```

```

public void replace (Integer x, Integer y) throws IllegalArgumentException
{
    if (raiz == null) throw new IllegalArgumentException();
    else { NodoArbol < E > n; n = raiz;
           while (n != null && x != (Integer) (n.valor))
               if (x < (Integer) (n.valor))
                   if (y >= (Integer) (n.valor))
                       throw new IllegalArgumentException();
                   else n = n.izq;
               else if (y <= (Integer) (n.valor))
                   throw new IllegalArgumentException();
               else
                   n = n.dre;
           }
    if (n == null) throw new IllegalArgumentException();
    else replaceRoot (NodoArbol < Integer > n, y);
}

```

REPASO TODOS EXAMENES

HASSUELL

$m\text{max} :: (\text{Ord } a) \rightarrow \text{Arbol } a \rightarrow \text{Arbol } a$

$m\text{max } AV = AV$

$m\text{max } (AB \ r \ i \ d) = AB \ r \ (\text{aux True } i) (\text{aux False } d)$

$\text{aux} :: \text{Bool} \rightarrow \text{Arbol } a \rightarrow \text{Arbol } a$

$\text{aux } — \quad AV = AV$

$\text{aux True } (AB \ r \ i \ d) = AB \ r \ (\text{aux True } i) AV$

$\text{aux False } (AB \ r \ i \ d) = AB \ r \ AV \ (\text{aux False } d)$

$m\text{max} :: (\text{Ord } a) \text{ Arbol } a \rightarrow \text{Arbol } a$

$m\text{max } AV = AV$

$m\text{max } AB \ r \ AV \ AV = AB \ r$

$m\text{max } AB \ r \ i \ d = AB \ r \ (\text{aux True } i) (\text{aux False } d)$

$\text{aux} :: \text{Bool} \rightarrow \text{Arbol } a \rightarrow \text{Arbol } a$

$\text{aux } — \quad AV = AV$

$\text{aux True } AB \ r \ i \ d = AB \ r \ (\text{aux True } i) AV$

$\text{aux False } AB \ r \ i \ d = AB \ r \ AV \ (\text{aux False } d)$

NOVEMBER 30TH 2024.

```
2) public static <E> ArbolBinario <E> prune (Lista ag)
{ ArbolBinario <E> sol = new ABD <E> ();
  if (!ag.EsVacia())
    try { ArbolBinario <E> left = new ABD <E> ();
           right = new ABD <E> (); E root = (E) ag.ag.Cabecera (); ag = ag.ColaC(); ag = ag.ColaC();
           while (!ag.EsVacia() && !ag.ColaC().EsVacia() && !ag.ColaC().Colac().EsVacia ())
             { ag = ag.ColaC(); if (!ag.EsVacia())
               { left = prune ((Lista) ag.Cabecera()); ag = ag.ColaC();
                 if (!ag.EsVacia()) right = prune ((Lista) ag.Cabecera()); }
               sol = new ABD <E> (root, left, right);
             } catch (TADVacioException exp) { Solv(exp); }
           return sol;
         }
```



b) public GNODE (Integer n).

```
{ Adyacentes = new boolean [MAX_NODOS][MAX_NODOS];
  Nodos = new NodoGrafo [MAX_NODOS]; numnodos = 0; Integer i; numnodos = n;
  for (i=0; i<n; i++) Nodos[i] = new NodoGrafo (i, false);
  for (i=0; i<n; i++)
    if (i*2+1 < n) Adyacentes [i][i*2+1] = Adyacentes [i*2+1][i] = True; else break;
    if (i*2+2 < n) Adyacentes [i][i*2+2] = Adyacentes [i*2+2][i] = True; else break;
  }
```

December 13th, 2022

memax:: Arbol a → Arbol a

memax AV = AV

memax AB r i d = AB r (aux True i) (aux False d)

aux :: Bool → Arbol a → Arbol a

aux — AV = AV

aux True (AB r i d) = AB r (aux True i) AV

aux False (AB r i d) = AB r AV (aux False d)

2. IMPLEMENTATION

```
public static < E > Lista < Integer > hole (ArbolBinario < E > t)
{
    Integer l, r; Lista < Integer > sol, left, right; sol = new LD< > ();
    if (t. EsVacio ()) sol. Añade (0);
    else
    {
        try { left = hole (t. SubArbolIzqdo ());
            right = hole (t. SubArbolDcho ());
            l = left. Cabeza (); r = right. Cabeza ();
            if (l < r) { sol = left. Colar (); sol. Añade (0); }
            if (l > r) { sol = right. Colar (); sol. Añade (1); }
            sol. Añade (l + r + 1);
        } catch (TAD...){}
        Return sol;
    }
}
```

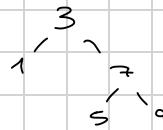
```
public GND (ArbolBinario < E > t)
{
    Adyacencias = new boolean [MAX_NODOS] [MAX_NODOS];
    Nodos = new NodoGrafo [MAX_NODOS]; numnodos = 0;
    ArbolBinario < E > aux;
    if (!t. EsVacio ())
    {
        Nodos [0] = new NodoGrafo (t. Raiz (), false); numnodos++;
        aux = t. SubArbolIzqdo ();
        while (!aux. EsVacio ())
        {
            Nodos [numnodos] = new NodoGrafo (aux. Raiz (), false);
            Adyacencias [numnodos] [numnodos - 1] = Adyacencias [numnodos - 1] [numnodos] = true;
            aux = aux. SubArbolIzqdo ();
        }
    }
}
```

```
| y (! t. SubArbolDcho . EsVacio ())
| Nodos [numnodos] = new NodoGrafo (t. SubArbolDcho, false);
| Adyacencias [0] [numnodos] = Adyacencias [numnodos - 1] [0] = true;
| numnodos++; aux = t. SubArbolDcho;
| while (! aux. EsVacio ())
| {
|   Nodos [numnodos] = new NodoGrafo (aux. Raiz (), false);
|   Adyacencias [numnodos] [numnodos - 1] = Adyacencias [numnodos - 1] [numnodos] = true;
|   numnodos++;
|   aux = aux. SubArbolDcho;
| }
| {
|   catch (VA.)
| }
```

December 18th 2020

2)

```
public Lista_Dinamica < E > FromLevel ( int level )
{
    Lista_Dinamica < E > sol = new Lista_Dinamica < E > ();
    if (!EsVacio)
        try { Lista_Dinamica < E > aux = ((ArbolBinario) (SubArbolDch.FromLevel (level - 1));
            if (level <= 1) aux.Arriade (Raiz ());
            sol = (ArbolBinario) SubArbolIzqdo.FromLevel (level - 1);
            sol.Concatena (aux);
        } catch (TADVacioException) { e.printStackTrace (); }
    return sol;
}
```



b) public GNODE (ArbolBinario < E > t)

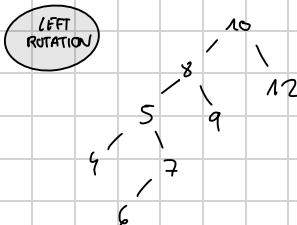
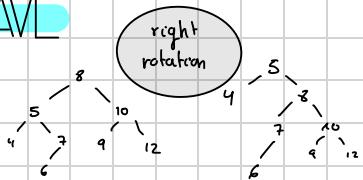
```
{ Adyacencias = new boolean [ MAX_Nodos ] [ MAX_Nodos ];
Nodos = new NodoGrafo [ MAX_Nodos ]; numnodos = 0;
if (!t. EsVacio)
    try { Nodos [ 0 ] = new NodoGrafo ( t. Raiz (), false ); numnodos++;
    while (!t. EsVacio )
        aux ( t. SubArbolIzqdo, 0 ); aux ( t. SubArbolDcho, 0 );
    }
catch ( TAD_VacioException exp ) ( sout ( exp ) )
{
```

private void aux (ArbolBinario < E > t, int fat)

{ while (!t. EsVacio)

```
{ try { Nodos [ numnodos ] = new NodoGrafo ( t. Raiz (), false );
Adyacencias [ fat ] [ numnodos ] = Adyacencias [ numnodos, fat ];
fat = numnodos; numnodos++;
aux20 ( t. SubArbolIzqdo (), fat ); aux20 ( SubArbolDcho (), fat );
}
catch ( TAD_VacioException exp ) ( sout ( exp ) )
{
```

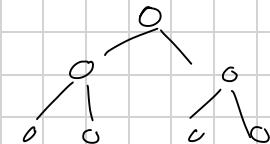
AVL



```
public static Integer miniumum (ArbolBinario < Integer > t) throws TADValueException  
{ try { while (!t.Esvacio())  
        { t. SubArbolIzquierdo(); }  
    } catch (TADValueException exp) { Sout(exp); }  
    return t.Raiz(); }
```

Iterante
Recorrido

```
public static Integer miniumum (ArbolBinario < Integer > t) throws TADValueException  
{ if (t. SubArbolIzquierdo(). Esvacio()) return t. Raiz();  
else return miniumum (t. SubArbolIzquierdo()); }
```



```
public void replaceRoot (NodoArbol < Integer > n, Integer y) throws IllegalArgumentException  
{ if (n == null) throw IllegalArgumentException;  
else if (n.izq >= y || n.dre <= y || y == n) throw IllegalArgumentException.  
else  
    { comproba maximo y <= (Integer) (maximum (new ArbolBinarioDinamico (n.izq)));  
     comproba minimo y >= (Integer) (minimum (new ArbolBinarioDinamico (n.dre)));  
     n.valor = y; } catch TAD... }
```

```
public void replace (Integer x, Integer y) throws IllegalArgumentException  
{ if (raiz == null) throw Ill...Arg...;  
else if (nodoArbol < E > n; n = raiz;  
while (n != null && x != (Integer) n. valor)  
    { if (x < (Integer) n. valor)  
        if (y >= (Integer) n. valor) throw Illega...;  
        else n = n. izq;  
    else  
        if (y <= (Integer) n. valor) throw Illega...;  
        else n = n. dre;  
    if (n == null) throw ...;  
else replace root ( n, y ) }
```

```

public GNDG (ArbolBinario<E> t) {
    Adyacentes = new boolean [M-N][M-N];
    nodos = new NodoGrafo[M-N];
    numnodos = 0;
    if (!t.Esvacio())
        try {
            nodos[0] = new NodoGrafo(t.Raiz(), false);
            numnodos = 1;
            aux (t.SabArbolIzqdo, 0); aux (t.SabArbolDcho, 0);
        } catch ...

```

```

private void aux (ArbolBinario<E> t, Integer i)
{
    if (!t.Esvacio())
        try {
            nodos [numnodos] = new NodoGrafo (t.Raiz(), false);
            Adyacentes [numnodos][fat] = Ady [fat][numnodos] = true;
            fat = numnodos; numnodos++;
            aux (t.SabArbolIzqdo, i); aux (t.SabArbolDcho, i);
        } catch ...
}

```

ITERATIVE

```

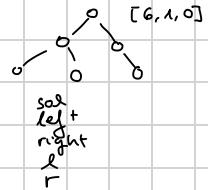
public GNDG (Integer n)
{
    Ady = new bool [N][N]; nodos = new NodoGrafo [N]; numnodos=0; Integer i;
    for (i=0; i < n; i++) nodos [i] = new Nodo (i, false); t
    for (i=0; i < n; i++)
        A

```

```

public static <E> Lista < Integer > hole (ArbolBinario < E > t)
{
    Integer l, r; Lista < Integer > sol, left, right; sol = new LD <>();
    if (!t.vacio) {
        left = hole( t. SubArbolIzquierdo ); right = ( t. SubArbolDcho );
        l = ( Integer ) left. Cabeza (); r = ( Integer ) right. Cabeza ();
        if ( l < r ) { sol = left. Cola (); sol. Anade ( o ); }
        if ( l > r ) { sol = right. Cola (); sol. Anade ( o ); }
        sol. Anade ( l + r + 1 );
    } catch ( TAD ... ) { }
    else if ( t. EsVacio ) sol. anade ( o );
}
return sol;

```



```

public GODE (ArbolBinario < E > t)
{
    Nodos = new NodoGrafo [ MN ]; Adyacencias = new boolean [ MN ] [ MN ]; numnodos = 0; ArbolBinario < E > aux;
    if ( !t. vacio )
    { try {
        Nodos [ numnodos ] = new NodoGrafo ( t. Raiz (), false );
        numnodos ++; aux = t. SubArbolIzquierdo ();
        while ( !aux. vacio () )
        { Nodos [ numnodos ] = new NodoGrafo ( aux. Raiz (), false );
            Adyacencias [ numnodos ] [ numnodos - 1 ] = Adyacencias [ numnodos - 1 ] [ numnodos ] = true;
            numnodos ++; aux = aux. SubArbolIzquierdo ();
        }
    } if ( !t. SubArbolDcho. vacio () )
    {

```

hola:: Arbol a \rightarrow Bool

hola AV = Polar

hola AB - AV AV = True

hola AB - AV d = True && hola d

hola AB - i AV = True && hola i

otro:: Arbol a \rightarrow AVL

otro AV = AV

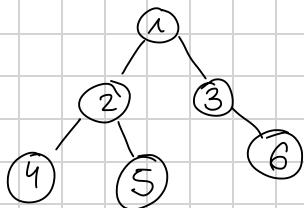
otro AB - AV AV =

aux:: Arbol a \rightarrow Arbol a \rightarrow int

aux AV AV = 0

aux AB r

aux AB



PRACTICE 3

```
public static Lista < ABD < Integer >> insert (ABD < Integer > x, Lista < ABD < Integer >> y) throws TAOVacuoException  
ArbolBinario aux = new ABD < Integer > ();  
{ try{ if( y.EsVacio() ) y.Arriade(x);  
else if( x.Raiz() <= y.Cabeza().Raiz() ) y.Arriade(x);  
else { ABD < Integer > c;  
c = y.Cabeza();  
y = y.Cola();  
y = insert(x,y); y.Arriade(c); }  
return y;  
}  
}
```

```
public static ABD < Integer > huffmanTree (Lista < ABD < Integer >> a)  
sol = new ABD();  
{ try{ if (a. EsVacio()) return sol; }  
else if (a. Cabeza(). EsVacio()) sol = a. Cabeza(); a = a. Cola();  
else { while (!a. Cola(). EsVacio())  
{  
ABD aux = new ABD ( a. Cabeza(). Raiz() ^ a. Cola(). Cabeza(). Raiz(), a. Cabeza(), a. Cola(). Cabeza());  
a = a. Cola(). Cola();  
a = insert (aux , a );  
}  
sol = a. Cabeza();  
}  
catch ...  
return sol;  
}
```

PRACTICE 4

FLOYDS ALGORITHM:

```
public static Lista < Integer > [ ] Floyd( int [ ] [ ] adj )
{ Lista < Integer > [ ] path = new LD< Integer> [adj.length] [adj.length];
if( !adj.length != 0 )
    for( i = 0; i < path.length; i++ )
        for( j = 0; j < path.length; j++ )
            path[ i ][ j ] = new LD< Integer>;
    if( i != j ) adj[ i ][ j ] = Integer.MAX_VALUE
    { path[ i ][ j ].Arrode( i ); path[ i ][ j ].Arrode( j ); }
// FLOYD ALGORITM
for( h = 0; h < path.length; h++ )
    for( k = 0; k < path.length; k++ )
        for( j = 0; j < path.length; j++ )
            if( adj[ i ][ h ] == Integer.MAX_VALUE ) || adj[ k ][ j ] == Integer.MAX_VALUE ) && ( adj[ i ][ k ] + adj[ k ][ j ] < adj[ i ][ j ] )
                try{
                    adj[ i ][ j ] = adj[ i ][ h ] + adj[ h ][ j ];
                    path[ i ][ j ] = ( LD< Integer > ) path[ i ][ k ].clone();
                    path[ i ][ j ].Concatena( path[ i ][ k ].clone() );
                } catch( TADU... e ) { e.printStackTrace() }
```

```
public static Lista < Integer > cycle( int [ ] [ ] adj )
int inf = Integer.MAX_VALUE, minweight = Integer.MAX_VALUE, n = adj.length, i = 0, j = 0, h = 0;
Lista < Integer > sol = new LD< Integer > ();
Lista < Integer > [ ] path = Floyd( adj );
for( x = 0; x < n; x++ )
    for( j = 0; j < n; j++ )
        if( i == j && adj[ i ][ j ] == inf && adj[ j ][ i ] == inf && ( adj[ i ][ j ] + adj[ j ][ i ] ) < minweight )
            minweight = adj[ i ][ j ] + adj[ j ][ i ];
            sol = ( Lista < Integer > ) ( path[ i ][ j ].clone() );
            sol.Concatena( path[ j ][ i ].clone() );
            sol.Concatena( path[ i ][ j ].clone() );
return sol;
```

PRACTICE EXAM

20th December 2023

$2 + ((12 \% 5) * 3)$ in polish notation: $+ 2 (* (% 12 5) 3) = + 2 * \% 12 5 3$

```

Pisa s = new PDC(); s.Aplica(3); s.Aplica(5); s.Aplica(12); s.Aplica("%"); s.Aplica("*"); s.Aplica("2"); s.Aplica("+");
public static Double PNTvalue(PDC exp)
{
    PDC aux = new PDC(); Double x, y; String op;
    try{ while(!exp.EsVacio())
    {
        if(exp.Tope() instanceof Char) Aux.Aplica(exp.Tope());
        else if(exp.Tope() instanceof Double) aux.Tope() instanceof Double)
            { x = aux.Tope(); y = exp.Tope(); aux=aux.DesAplica(); (op = aux.Tope());)
            aux.Aplica(eval(x,(op),y));
            (String)
            aux.Tope();
        }
        else { aux.Aplica((Exp.Tope());}
        aux.K.Aplica(exp.Tope());
        exp=exp.DesAplica();
    }
    catch(TADVacioException) {System.out.println("Error");}
    return (Double) exp.Tope();
}

```

2													
*	*												
%	%	%											
12	12	12	*	12	*	*	*	*	*	*			
5	5	2	5	2	5	2	5	2	2	2	2	2	
3	+	3	+	3	+	3	+	3	+	3	+	6	+
exp	aux												

2. [Huffman trees]

```

public static int percent(ArbolBinario<Integer> t, Lista<Integer> l)
{
    Integer p = -1; ArbolBinario<Integer> aux = new ABD<?>();
    if(!t.EsVacio())
        try{ if(l.EsVacio())
        {
            if(t.SubArbolDcho().EsVacio()) p = t.Raiz();
            else if(l.Cabeza() == 0) p = percent(t.SubArbolIzq(), l.Cola());
            else p = percent(t.SubArbolDcho(), l.Cola());
        }
        catch...-
        return p;
    }
}

```

```

public static List<Integer> cycle (int[][] adj)
{
    int inf = Integer.MAX_VALUE, minweight = Integer.MAX_VALUE, n = adj.length, i=0, j=0, k=0;
    List<Integer> sol = new ArrayList();
    List<Integer> path = Floyd (adj);
    for (int i=0; i < n; i++)
    {
        for (j=0; j < n; j++)
        {
            if (i != j && adj[i][j] == inf && adj[j][i] == inf && adj[i][j] + adj[j][i] < minweight)
                minweight = adj[i][j] + adj[j][i];
            sol = (List<Integer>) (path[i][j].clone());
            sol.add (path[j][i].clone());
        }
    }
    return sol;
}

```

FLOYD'S

```

public static int[][] transform (int[][] c)
{
    int i, j, k; int inf = Integer.MAX_VALUE, minweight = Integer.MAX_VALUE, n = adj.length;
    int[][] adj = new int[c.length][c.length]; minweight = 1;
    for (i=0; i < c.length; i++)
    {
        for (j=0; j < c.length; j++)
        {
            k = c[i][j];
            if (i == j || k == j && c[i][k] == k && c[k][j] == j)
                adj[i][j] = -1;
            else adj[i][j] = weight++;
        }
    }
}

```

```

public static int penultimate (int[][] c, int i, int j)
{
    int sol;
    if (i == j) sol = -1;
    else if (c[i][j] == -1) sol = -1;
    else if (c[i][j] == -j) sol = i;
    else return penultimate (c, c[i][j], j);
}

```

15th December 2021:

1) [Reverse Polish Notation]

```
public static Pila partialRPN (Pila exp, int i)
{
    Pila aux = new PD<>(); Double x, y; int counter;
    try { while (!exp.EsVacia()) && (i > 0))
        if (exp.Tope() instance of Double) aux.Apila(exp.Tope());
        else { x = (Double) aux.Tope(); aux = aux.DesApila();
            y = (Double) aux.Tope(); aux = aux.DesApila();
            aux.Apila (eval(x, (String) exp.Tope()), y); i--;
            exp = exp.DesApila();
        }
    } while (!aux.EsVacia())
    if exp.Apila(aux.Tope()); aux = aux.DesApila();
    return exp;
}
```

2) comhuff :: int → Arbol Flot

int i < h < 16
height

```
public static ArbolBinario<Float> comhuff (int h)
```

```
{ return auxch(h, new float(100), 1);
    }
```

```
public static ArbolBinario<Float> auxchr (int level, float perc, int ascii)
```

```
{ ABD<Float> out = new ABD<Float>();
```

```
if (level > 2) { out = new ABD<Float> (perc, auxch(level - 1, perc/2, ascii*2), auxch(level - 1, perc/2, ascii*2+1))
```

```
else out = new ABD<floats> (perc, new ABD<float> (new float(ascii), new ABD<float>(), new ABD<float>(), new ABD<float>()));
    } return out.
```

3) public static int[][] transform (int[][] c)

new ABD<float>()
new ABD<float>()

HUFFMAN TREE

fx21:: ArbolFloat → ArbolFloat

e19 AV = AV

fx21 AB - n AV = AB 100 n AV

fx21 AB p (AB - n AV) (AB - m AV) = AB p (P/2 n AV) (AB (P/2) m AV)

fx21 (AB p (AB - n AV)) (AB r i d) = AB p (p - r n AV) fx21 (AB r i d)

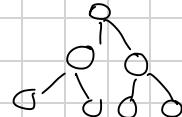
fx21 (AB p (AB r i d)) (AB - n AV) = AB p fx21 (AB r i d) (AB p - r n AV)

fx21 AB r i d = AB r fx21 i fx21 d

e19:: ArbolInt

e19 AB.

```
public static Integer RPNmaxbase(Pila exp) { PD<> curr = new PD(); Double x, y;
    if (!exp.EsVacio())
        try { while (!exp.Vacío())
            x = exp.Tope(); exp = exp.Desapilar();
            if (x instanceof Double && !exp.EsVacio() && exp.Desapilar().instanceof Character)
                z = eval(x, exp.Desapilar().Tope(), exp.Tope());
            exp = exp.Desapilar().Desapilar(); if (z > max) max = z
        }
    }
}
```



HUFFMAN

public static ArbolBinario<Integer> morrar(ArbolBinario<Integer> h)

{ ArbolBinario<Integer> s;

try { if (h.SubArbolDcha().EsVacio())

s = (ArbolBinario<Integer>) h.clone();

else s = new Arbol<Integer>(h.Raiz(), morrar(h.SubArbolDcha().derecho()));

e19:: Arbol a → Arbol a

e19 AV = AV

e19 AB r i d = morrar AB

morrar AB(r x AV) = AB r x AV

morrar AB(r i d) = AB r morrar d

	1	2	3	4	5	
1	-1					4 arc = δ
2		-1				2 arcs = n internode
3			-1			
4				-1		
5					-1	

3) public static $\text{Pila} < \text{Integer} \rangle$ internodes (int c[3][3] c)

{ Pila < Integer > sal = new Pila < Integer >();

int i, j, k, n = c.length;

for (i = 0; i < n; i++)

for (j = 0; j < n; j++)

{ k = c[i][j];

if (k == -1 || k == j || !Pertenece(sal, k))) sal.apilar(k);

return sal

public static boolean Pertenece (Pila < Integer > p, int k)

{ try { if (p.esVacia()) return false;

else return k == p.tope() || Pertenece(p.desapilar(), k);

FLOYD PATH

1) public static int permutante (int c[3][3] c, int i, int j) int sal;

try { if (!c.esVacia())

{ if (c[i][j] == -1) return -1;

else if (c[i][j] == j) return 1;

else return permutante(c, c[i][j], j);



-1 → no path
j → direct
k → connector



2) fix 21: ArbolFloat → ArbolFloat

fix 21 AV = AV

fix AB - n AV = AB.lod n A

public static int c[3][3] transform (int c[3][3] c) ⁻¹

int adj[] = new int[c.length][c.length]; int weight = 0;

for (int i = 0; i < c.length; i++)

for (int j = 0; j < c.length; j++)

{ if (c[i][j] == -1) adj[i][j] = -1;

else if (c[i][j] == j) adj[i][j] = weight++;

else if (i == j || c[i][j] == n || c[n][j] == j || w == -1 || w == j)

adj[i][j] = -1;

return adj.

HUFFMAN TREE:

2) comhuff:: Int → Arbol <Float>
 comhuff O = AV
 comhuff h = auxhuff h 100 1
 auxhuff Int → Float → Float → Arbol <Float>
 auxhuff level perc ascii = if (level > 2) then AB perc
 auxhuff level-1 p/2 ascii * 2)
 auxhuff level-1 p/2 ascii * 2+1
 else AB perc (AB ascii AV AV) AV
 auxop(h, 100, 1);
{

public static ArbolBinario <Float> comhuff (int h)

out = new ABD <Float> (c);

if (level > 2)

out = new ABD (perc, auxop(level - 1, perc/2, ascii * 2), auxop(level - 1, $\frac{perc}{2}$, ascii * 2))

else out = new ABD (perc, new ABD (ascii, new ABD <Float>, new ABD <Float>))

public static Lista <Integer> cycle (int adj[][]) adj)

| int iinf = Integer. MAX_VALUE, minweight = Integer. MAX_VALUE, n = adj.length, i = 0, j = 0, k = 0;

Lista <Integer> sol = new (0 < Integer>()); Lista <Integer> [] path = Floyd (adj);

for (int i = 0; i < n; i++)

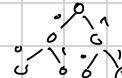
for (int j = 0; j < n; j++)

if (i != j && adj[i][j] != inf && adj[j][i] == inf && adj[i][j] + adj[j][i] < minweight)

minweight = adj[i][j] + adj[j][i];

sol (Lista <Integer>) path[i][j]. Cola(). clone());

sol. Concatenar (path[j][i]. Cola(). clone());



public static int percent (ArbolBinario <Integer> t, Lista <Integer> l)

| Integer p = -1; ArbolBinario <Integer> aux = new ABD <C>;

if (!t. EsVacio())

try { if (!l. EsVacio())

{ if (t. SubArbolDcho(). EsVacio ()) p = t. Raiz(); }

else if (l. Colera () == 0) p = percent (t. SubArbolIzqdo (), l. Colar());

else p. pe