Parallel and Distributed Computing (CS3006)

Course Instructor(s):

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Section(s): A,B,C,D,E,F,G,H,J,K

Sessional-II Exam

Total Time (Hrs): 1

Total Marks: 55

Total Questions: 4

Date: Apr 7, 2025

Roll No Course Section Student Signature

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Attempt all the questions.

[CLO 1: Demonstrate understanding of various concepts involved in parallel and distributed computer architectures.]

Q1: Multiple choice questions.

[5x1 & 10x2 = 25 marks]

1) Consider the machine file with the following configuration for MPI:

node1:4
node2:2

node3:3

Assume an MPI program is launched using the following command, choose the correct order of process spawning across these nodes?

mpirun -np 15 -machinefile machinefile myprogram.

- a) 4 processes on node1, 2 processes on node2, 9 processes on node3.
- b) 3 processes on node1, 2 processes on node2, 10 processes on node3.
- c) 6 processes on node1, 2 processes on node2, 7 processes on node3.
- **d)** 5 processes on node1, 3 processes on node2, 7 processes on node3.
- e) Any random order
- 2) What is the main advantage of OpenCL over CUDA?
 - a) OpenCL supports all programming languages
 - **b)** OpenCL is more efficient for multi-threaded CPU programming
 - c) OpenCL supports a wide range of devices, including CPUs, GPUs, and FPGAs
 - d) OpenCL has a simpler programming model
 - e) None of above
- 3) Source code in an OpenMP that is not covered by a pragma is executed by:
 - a) All threadsb) Single thread

c) Error, it must be specified with thread id

d) None of the above

- 4) Which of the following loop scheduling in OpenMP has the most overhead:
 - **a)** Static

d) #pragma omp for

b) Dynamic

- e) Guided
- c) It does not support loop scheduling

- 5) In OpenCL, which of the following is a key difference between buffers and images?
 - a) Buffers store 1D or 2D data, while images store 2D or 3D data.
 - b) Buffers are accessed directly via global IDs, while images are accessed using texture coordinates.
 - c) Buffers support only int and float types, while images support half, float, and double.
 - d) All of the above are true.
 - e) Only a) and b) are true.
- **6)** What will be the output of the following code:

```
int i;
int j;
int n=10;
int t=0;
omp set num threads(2);
#pragma omp parallel private(j)
     # pragma omp for
     for (i=0; i<n; i++) {
           printf("a\n");
     for(j=0; j< n; j++) {
          printf("b\n");
```

- a) 10 a's and 10 b's
- **b)** 20 a's and 20 b's
- c) 10 a's and 20 b's

- d) 20 a's and 10 b's
- e) Cannot be determined
- 7) Consider the following OpenMP code and determine the execution behavior of the loop:

```
omp_set_num_threads(10);
int i;
#pragma omp parallel for private(i)
for (i=1; i<99; i++)
      A[i+1] = A[i-1] + 20;
```

- a) Serial Execution
- **b)** Parallel Execution
- c) Runtime Error
- d) Master thread only
- e) Syntax Error

8) What will be the output of the following OpenMP code:

```
omp set num threads (4);
int n=1, x=10;
#pragma omp parallel for lastprivate(n)
For(i=1; i<x; i++)
     If (x/2 == i)
        n=i;
printf("%d", n);
```

a) 1

d) 5

b) 10

e) None of the above

- **c)** 9
- 9) What will be printed by process 0, considering total 4 processes?

```
int rank, data[5] = \{10, 20, 30, 40, 50\};
MPI Comm rank (MPI COMM WORLD, &rank);
if (rank == 0) data[0] = 5;
MPI Bcast(&data, 1, MPI INT, 2,
MPI COMM WORLD);
if (rank == 0) printf("Rank 0: %d\n", data);
```

a) Rank 0: undefined

d) Rank 0: 0

b) Rank 0: 10

e) Segmentation Fault

- c) Rank 0: 5
- 10) What will be the output of the following code, considering total 4 processes?

```
int rank, send = 10, recv = 0;
MPI Comm rank (MPI COMM WORLD, &rank);
send = rank;
MPI Reduce (&send, &recv, 1, MPI INT, MPI MAX,
1, MPI COMM WORLD);
if(rank == 0) printf("Recv: %d\n", recv);
```

a) Recv: 10

d) Recv: 3

b) Recv: 0

e) Segmentation Fault

- c) Recv: garbage
- 11) What will be the output of the following code, considering total 4 processes?

```
int rank;
int send = 1;
int recv;
MPI Comm rank (MPI COMM WORLD, &rank);
MPI Allgather (&send, 1, MPI INT, &recv, 1,
MPI INT, MPI COMM WORLD);
if(rank == 0) printf("Recv: %d\n", recv);
```

a) Recv: 3

d) Compile-time error

b) Recv: 1

e) Runtime error

c) Recv: garbage

12) What happens in this kernel execution, considering Global size = 2. Initial A[0] = 0.?

```
kernel void test( global int* A) {
   int id = get global id(0);
   if (id == 1)
       A[0] = 10;
   else
       A[0] = 5;
}
```

```
a) A[0] = 10
```

d) A[0] = non-deterministic

b) A[0] = 5

e) Compile-time error

- **c)** A[0] = 0
- 13) What will be the output of the following OpenCL kernel when executed with global size = 3 and local size = 2?

```
kernel void check() {
    int gid = get global id(0);
    int lid = get local id(0);
    int group = get group id(0);
    int temp = gid * 2 + lid + group;
    if (1id == 2)
       printf("Temp: %d\n", temp);
    // This will help us confirm thread count
    // printf("GID: %d, LID: %d, Group:
%d\n", gid, lid, group);
```

```
a) Temp: 6
```

d) Temp: 10

b) Temp: 4

e) No output

- c) Temp: 5
- 14) In the given MPI program what is the value of the global product printed by process 0 when the program is executed with 3 processes?

```
int rank, size, n = 12, buffer[12], result[4];
MPI Init(&argc, &argv);
MPI Comm rank (MPI COMM WORLD, &rank);
MPI Comm size (MPI COMM WORLD, &size);
if (rank == 0) {
    for (int i = 0; i < n; i++) {
        buffer[i] = 2 * (i % 2) + 1;
}
```

```
MPI Scatter(buffer, n / size, MPI INT, result, n / size,
MPI INT, 0, MPI COMM WORLD);
   int local product = 1;
    for (int i = 0; i < n / size; i++) {
        local product *= result[i];
int global product;
MPI Reduce (&local product, &global product, 1, MPI INT,
MPI PROD, 0, MPI COMM WORLD);
if (rank == 0) {
   printf("Global product: %d\n", global product);
```

```
a) Global product: 24
                                        c) Global product: 64
b) Global product: 32
                                        d) Global product: 96
```

15) What will be the output of the following code snippet?

```
int N=4, s data = 10, p data = 0, M=0, S=0, L=-1;
omp_set_num_threads(5);
#pragma omp parallel num threads(3) firstprivate(s data) shared(M, S)
for (int i = 0; i < N; i++) {
  int thread id = omp get thread num();
  #pragma omp master
    M += s data + i;
  #pragma omp single
    S++;
  }
  s_data += (thread_id + i);
  L = thread id;
}
printf("M: , ", M);
printf("S: , ", S);
printf("s_data: %d\n", s_data);
```

```
a) M: 50, S: 4, S data: 10
b) M: 50, S: 12, S_data: 10
c) M: 50, S: 4, S data: 16
d) M: 46, S: 4, S_data: 10
e) None of above
```

[CLO 2: Implement different parallel and distributed programming paradigms and algorithms using Message-Passing Interface (MPI) and OpenMP]

Q2: Write output of the following program.

[10 Marks]

```
#define NUM THREADS 4
#define ITERATIONS 2
float *create nums(int num elements, int rank) {
  float *nums array = (float *)malloc(sizeof(float) *
num elements);
  for (int i = 0; i < num elements; <math>i++) {
   nums array[i] = rank + i \star 0.1;
 return nums array;
int main(int argc, char** argv) {
 MPI Init(NULL, NULL);
 int rank, size;
 MPI Comm rank (MPI COMM WORLD, &rank);
 MPI Comm size (MPI COMM WORLD, &size);
 float *nums array = create nums(num elements per proc, rank);
 printf("Process %d - nums array: ", rank);
  for (int i = 0; i < num elements per proc; i++) {</pre>
    printf("%f ", nums array[i]);
  }
 printf("\n");
 float sum = 0;
 for (int i = 0; i < num elements per proc; i++) {
    sum += nums array[i];
 printf("Process %d - sum: %f\n", rank, sum);
 float total sum;
 MPI Allreduce (&sum, &total sum, 1, MPI FLOAT, MPI SUM,
MPI COMM WORLD);
 float mean = total sum / (num elements per proc * size);
 printf("Process %d - mean: %f\n", rank, mean);
 float sq diff = 0;
 for (int i = 0; i < num elements per proc; i++) {</pre>
    sq_diff += (nums_array[i] - mean) * (nums array[i] - mean);
 printf("Process %d - sq diff: %f\n", rank, sq diff);
```

```
float total sq diff;
  MPI Reduce (&sq diff, &total sq diff, 1, MPI FLOAT, MPI SUM, 0,
MPI COMM WORLD);
  if (rank == 0) {
    float stddev = sqrt(total sq diff / (num elements per proc *
    printf("Final Mean = %f, Standard Deviation = %f\n", mean,
stddev);
  }
  free (nums array);
  MPI Barrier (MPI COMM WORLD);
  MPI Finalize();
```

If we consider, num_elements_per_proc=3 (for other assumptions you can calculate likewise), then following will be output:

```
Process 0 - nums_array: 0.0 0.1 0.2
Process 1 - nums_array: 1.0 1.1 1.2
Process 2 - nums array: 2.0 2.1 2.2
Process 3 - nums_array: 3.0 3.1 3.2
Process 0 - sum: 0.30
Process 1 - sum: 3.30
Process 2 - sum: 6.30
Process 3 - sum: 9.30
Process 0 - mean: 1.60
Process 1 - mean: 1.60
Process 2 - mean: 1.60
Process 3 - mean: 1.60
Process 0 - sq_diff: 3.48
Process 1 - sq_diff: 0.48
Process 2 - sq diff: 0.48
Process 3 - sq_diff: 3.48
```

Final Mean = 1.600000, Standard Deviation = 0.836660

[CLO 1: Demonstrate understanding of various concepts involved in parallel and distributed computer architectures.]

Q3: Below is an OpenCL kernel that computes the degree of each node in an undirected graph using adjacency matrix. This kernel is invoked using following command

```
clEnqueueNDRangeKernel(queue, kernel, 1, NULL,
&global work size, NULL, 0, NULL, NULL);
```

[2+3+5=10 Marks]

```
kernel void computeDegree ( global int *adjMatrix, global int
*degree, int numNodes) {
    int id = get global id(0); // Get the index of the current
node
    if (id < numNodes) {</pre>
        for (int j = 0; j < numNodes; j++) {
            degree[id] += adjMatrix[id * numNodes + j];
        }
    }
}
```

a) Identify potential performance issue in this kernel. Discuss briefly inefficiency related to memory

Global Memory Access Latency:

Each work-item reads data[id] from global memory, computes the square, and writes it back. Global memory is slow, and frequent accesses create bottlenecks.

No Work-Group Optimization:

All computations occur at the global memory level, missing potential optimizations with local memory.

b) Modify the kernel to utilize private memory for computing degree of each node for improving efficiency.

```
int sum = 0;
for (int j = 0; j < numNodes; j++) {
    sum += adjMatrix[id * numNodes + j];
}
degree[id] = sum;
```

c) Provide kernel code that utilizes local memory for buffering the given matrix with the aim of improving performance by reducing memory access time. The new kernel will be invoked using following command

```
clEnqueueNDRangeKernel(queue, kernel, 1, NULL, &global work size,
 &local work size, 0, NULL, NULL);
  kernel void computeDegreeOptimized(
 global int *adjMatrix,
__global int *degree,
int numNodes
int group_size = get_local_size(0); // Size of the work-group
__local int localRow[256]; // Adjust size based on hardware limits
```

// Load the row in chunks of `group_size`

for (int j = 0; j < numNodes; j += group_size) {</pre>

int sum = 0;

```
int index = gid * numNodes + (j + lid);  // Global index

// Bounds check to avoid out-of-range access
if ((j + lid) < numNodes) {
        localRow[lid] = adjMatrix[index];
} else {
        localRow[lid] = 0;
}

// Synchronize to make sure all work-items have written to localRow barrier(CLK_LOCAL_MEM_FENCE);

// Accumulate the values from the shared chunk for (int k = 0; k < group_size && (j + k) < numNodes; k++) {
        sum += localRow[k];
}

barrier(CLK_LOCAL_MEM_FENCE);
}

// Write the final degree to global memory degree[gid] = sum;</pre>
```

[CLO 2: Implement different parallel and distributed programming paradigms and algorithms using Message-Passing Interface (MPI) and OpenMP.]

Q4: A company requires a simple console-based utility to find the brightest pixel in a grayscale image. The grayscale image is already formatted as a 2D array of pixels, where each pixel has an intensity value ranging from 0 (black) to 255 (white). When the image is provided, the utility must output two things:

- 1. The brightest pixel in the entire image.
- **2.** The brightest pixel on the border of the image (i.e., the first row, last row, and the first and last columns of the image).

The company has some multicore machines having different architectures where this utility will be executed. Utility must process the image using multithreading with OpenMP, adhering to the following guidelines:

- The number of threads should be decided by the user.
- Each thread will process a row-wise subset of the image.
- Each thread will find the brightest pixel in its assigned region and the brightest pixel among the border pixels of its assigned region.

The basic skeleton code is already provided. Your task is to understand the code and complete it by adding the necessary code to implement the functionality.

Note: Please only write the missing code on the answer sheet with proper labels as provided here. Do not write the whole code again. Assume all required libraries are included. [10 marks]

```
int brightest pixel = brightest border pixel = local max = 0;
 int local border max = 0, num threads;
 printf("Enter the number of threads: ");
 scanf("%d", &num_threads);
 // A. Write the missing code here
 // B. Write the missing clauses
 #pragma omp parallel
     // C. Write the best suited scheduling scheme
     #pragma omp for
     for (int i = 0; i < IMAGE ROWS; i++) {
         for (int j = 0; j < IMAGE COLS; j++) {
             // Update the brightest pixel in the entire image
             if (image[i][j] > local_max) {
                 local max = image[i][j];
             // D. Update the brightest pixel on the border
                 if (image[i][j] > local border max) {
                    local border max = image[i][j];
                 }
             }
         // E. Update the global maximum values
 printf("Brightest pixel: %d\n", brightest pixel);
 printf("Brightest border pixel: %d\n", brightest border pixel);
 return 0;
A. omp set num threads(num threads);
B. shared(brightest pixel, brightest border pixel) private(local max,
 local border max)
C. schedule(dynamic)
D. if (i == 0 || i == IMAGE ROWS - 1 || j == 0 || j == IMAGE COLS - 1)
E. #pragma omp critical
         if (local max > brightest pixel) {
```

```
brightest pixel = local max;
}
if (local border max > brightest border pixel) {
    brightest border pixel = local border max;
```