

### CPSC 331 Assignment 3

Question 1: • Loop invariant: ~~Q1~~ IF  $T$  is a BST, then cursor is a BST as well

- if  $\text{key} > \text{cursor.key}$  then key is not in  $\text{cursor.left}$  (if key exists)
- if  $\text{key} < \text{cursor.key}$  then key is not in  $\text{cursor.right}$  (if key exists)

Before the first loop, it is a precondition that  $T$  is a BST, Cursor is assigned  $T.\text{root}$ , therefore cursor is a BST. If  $\text{key} < \text{cursor.key}$  then  $\text{cursor.right}$  is eliminated and we search  $\text{cursor.left}$  ~~at~~ in the next iteration. Since cursor was a BST  $\text{cursor.left}$  is a BST as well.

The loop ~~variant~~ variant is the height of cursor. Each iteration the height decreases by one since it searches the BST of one of its children. Thus it gives us an upper bounds of iterations of  $\text{cursor.height}$ . When  $\text{height} = 0 - 1$  the cursor will be pointing to null, thus exiting the loop. Therefore the loop terminates

## Question 2

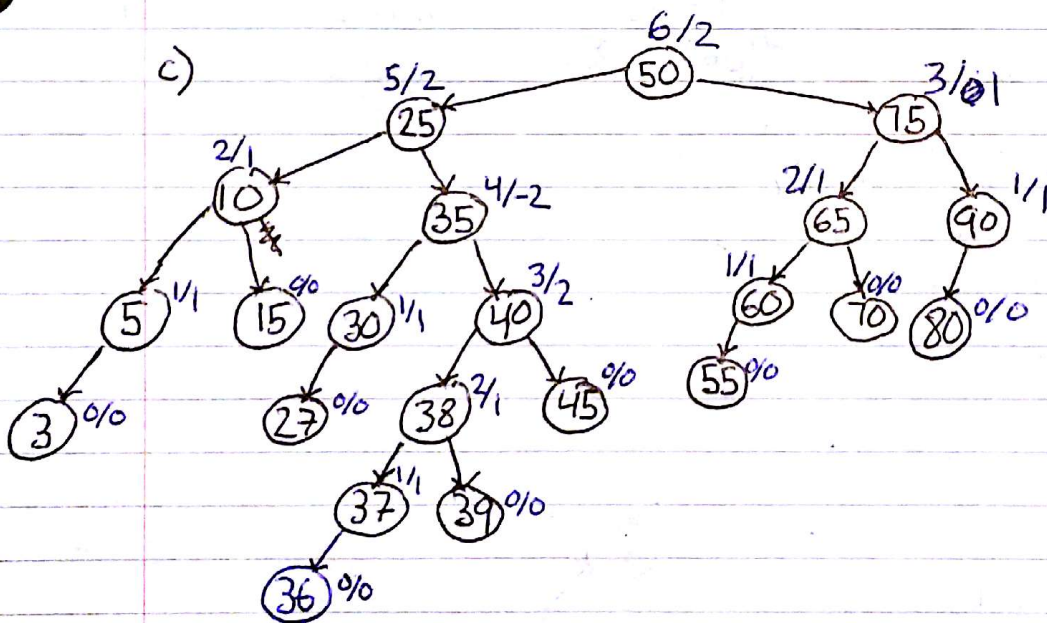
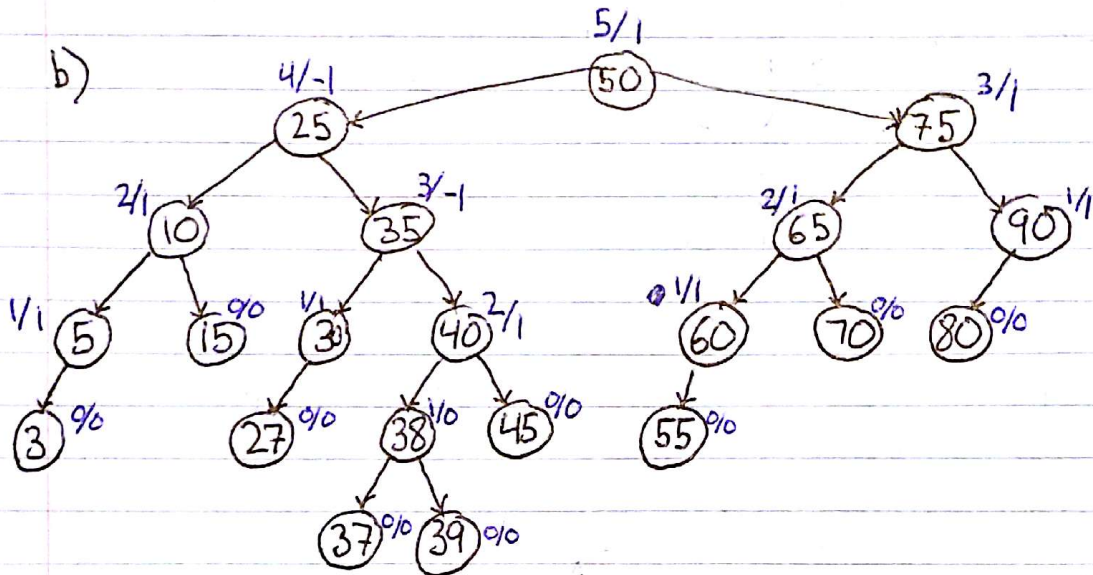
a) while (cursor != null)  
    if (element == cursor) throw FoundException;  
    if (element < cursor) then  
        | cursor = cursor.left;  
    else  
        | cursor = cursor.right;  
end  
if (cursor == null) then  
    | cursor = element;

b) modify (T, key, value)  
    cursor = T.root;  
    while ((cursor != null) || (key != cursor.key & value != cursor.value))  
        | if (key == cursor.key) then  
            | cursor.value = value;  
        | else if (key > cursor.key) then  
            | cursor = cursor.right;  
        | else  
            | cursor = cursor.left;  
    end  
    if (cursor == null) then  
        | throw KeyNotFound Exception;  
end

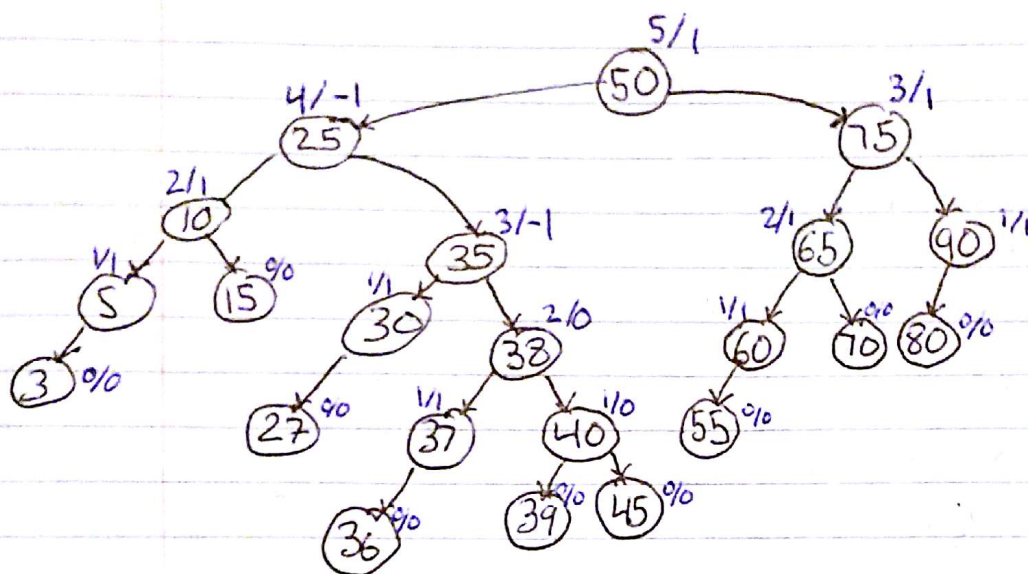


# Question 5

a) Student ID: 10017168  
Group: 0

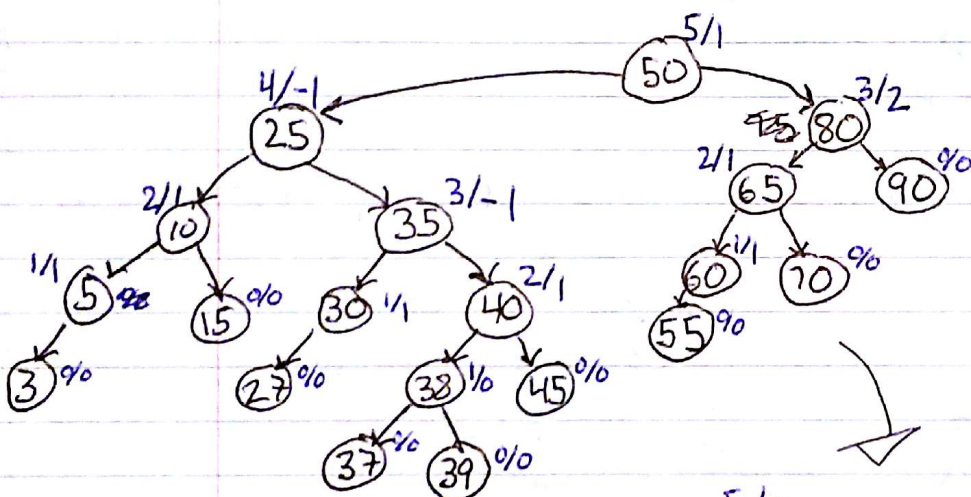


Insert 36

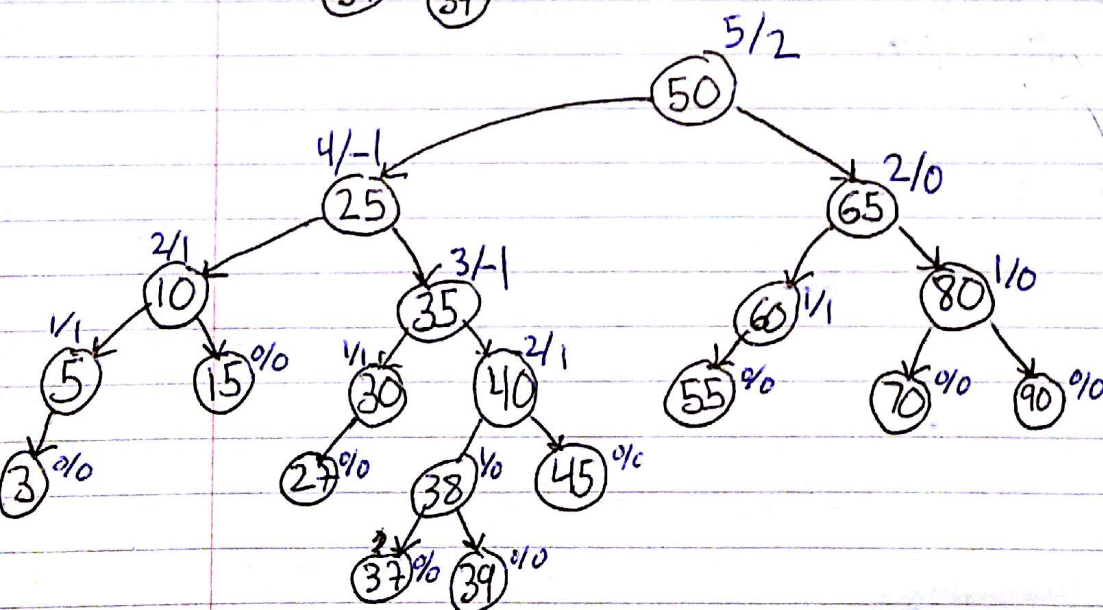


performed a right rotation on 40

d) delete 75 we take the MIN of the right subtree which is 80 to replace it.



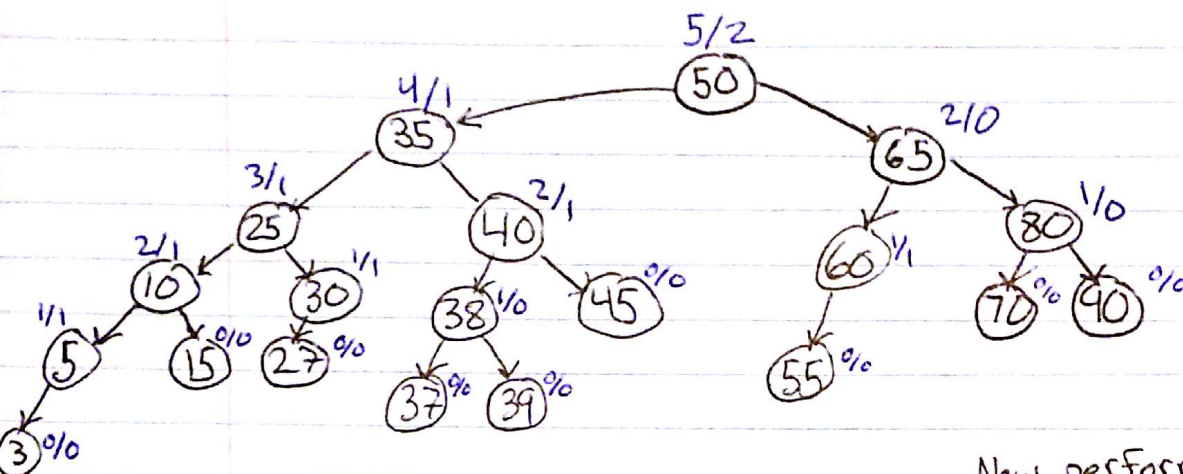
It is a left outer case so we will perform a right rotation



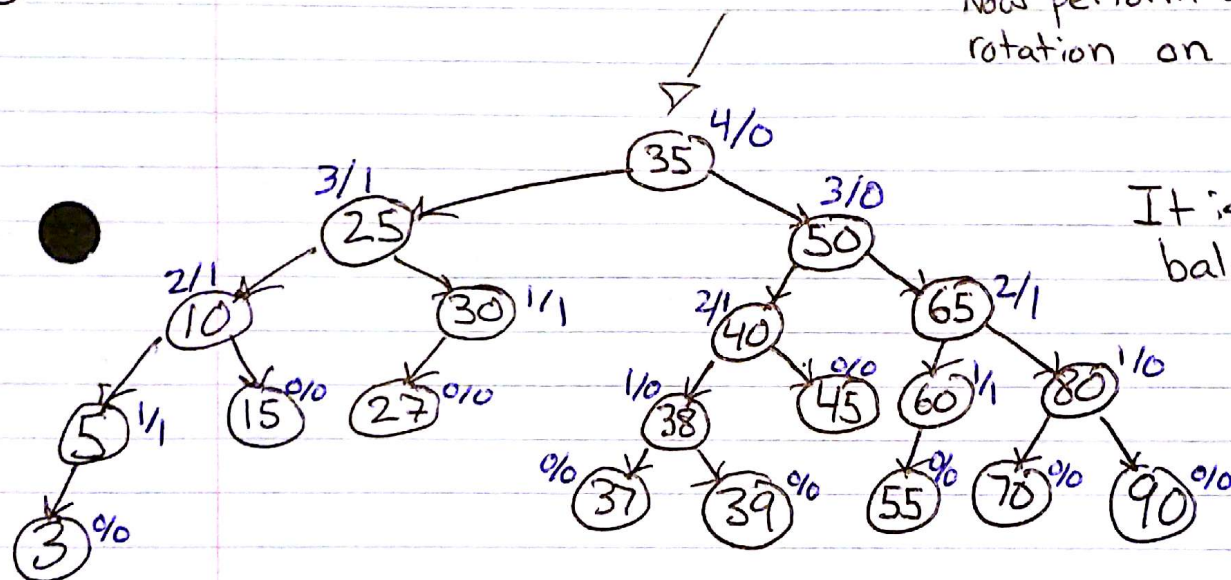
We now have an unbalance at the root. It is an inner left case, so we must perform a left right double rotation.



We will perform a left rotation on v.left, which is node 25.



Now perform a right rotation on v (50)



It is now a balanced AVL