Outlier Detection:

Density and Partition-Based

Mining Massive Datasets

Prof. Carlos Castillo — https://chato.cl/teach



Sources

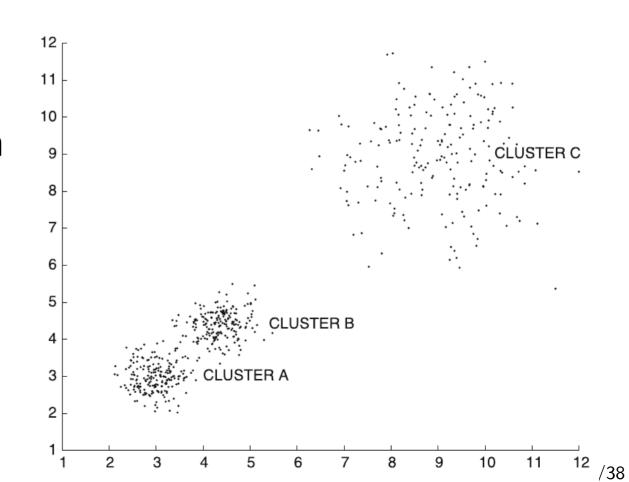
Liu, F. T., Ting, K. M., & Zhou, Z. H. Isolation forest. ICDM 2008.

- (1) Eryk Lewinson: Outlier detection with isolation forest (2018)
- (2) Tobias Sterbak: Detecting network attacks with isolation forests (2018)

Density-based methods

Density-based methods

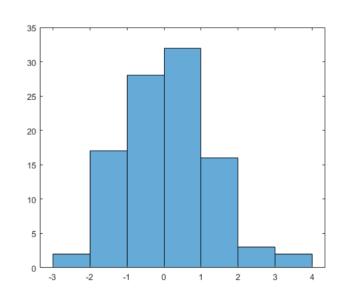
- Key idea: find sparse regions in the data
- Limitation: cannot handle variations of density



Histogram- and grid-based methods

Histogram-based method:

- 1. Put data into bins
- 2.Outlier score: num 1, where num is the number of items in the same **bin**

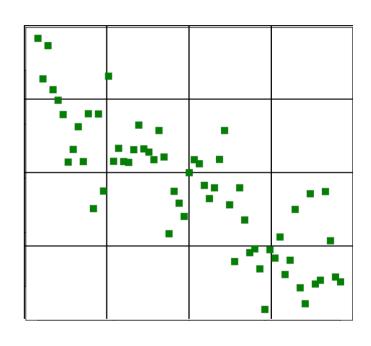


Clear outliers are alone or almost alone in a bin

Histogram- and grid-based methods

Grid-based method

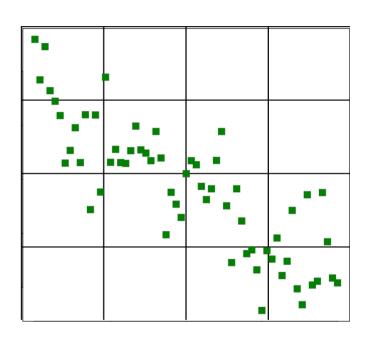
- 1. Put data into a grid
- 2.Outlier score: num 1, where num is the number of items in the same **cell**



Clear outliers are alone or almost alone in a cell

Problems with grid-based methods

- How to choose the grid size?
- Grid size should be chosen considering data density, but density might vary across regions
- If dimensionality is high, then most cells will be empty



Kernel-based methods

• Given n points $\overline{X_1}, \overline{X_2}, \dots, \overline{X_n}$

$$f(\overline{X}) = \frac{1}{n} \sum_{i=1}^{n} K_h(\overline{X} - \overline{X_i})$$

- K_h is a function peaking at \overline{X}_i with bandwidth h
- For instance, a Gaussian kernel:

$$K_h(\overline{X} - \overline{X_i}) = \left(\frac{1}{\sqrt{2\pi \cdot h}}\right)^d \cdot e^{-\|\overline{X} - \overline{X_i}\|^2/(2h^2)}$$

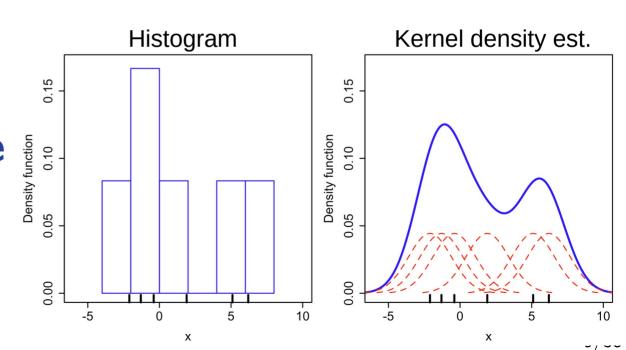
Kernel-based methods (cont.)

Example with a Gaussian kernel

$$\overline{X} = \langle -2.1, -1.3, -0.4, 1.9, 5.1, 6.2 \rangle$$

- Each K_k in red

• f = sum of K_h in blue
$$f(\overline{X}) = \frac{1}{n} \sum_{i=1}^{n} K_h(\overline{X} - \overline{X_i})$$



[Wikipedia: Kernel density estimation]

Partitioning-based method: isolation forest

Isolation forest method

- tree_build(X)
 - Pick a random dimension r of dataset X
 - Pick a random point p in $[\min_{r}(X), \max_{r}(X)]$
 - Divide the data into two pieces: $x_r < p$ and $x_r \ge p$
 - Recursively process each piece

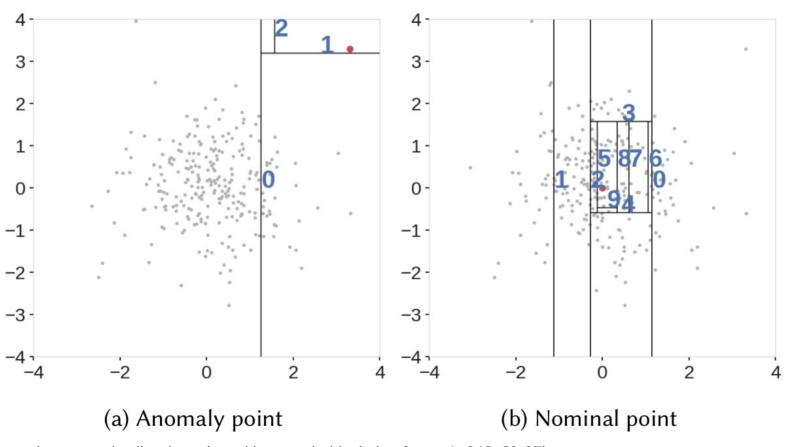
Stopping criteria for recursion

Stop when a maximum depth has been reached

-or-

Stop when each point is alone in one partition

Key: outliers lie at small depths



Outlier score

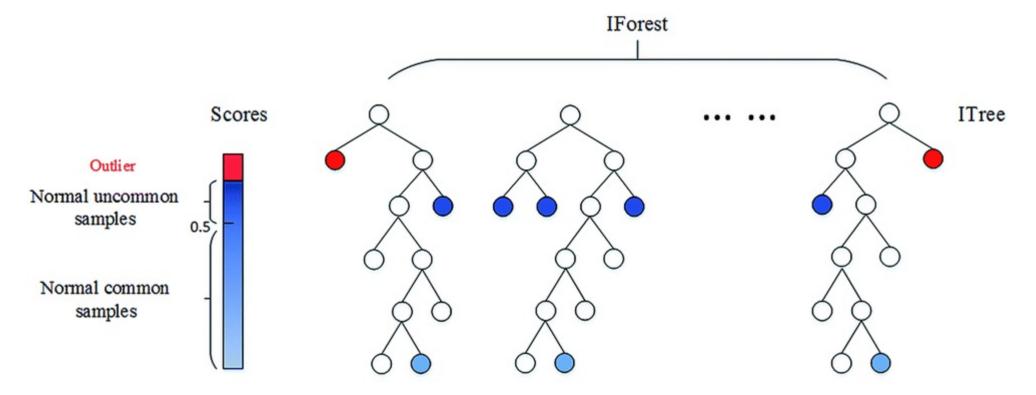
• Let c(n) be the average path length of an unsuccessful search in a binary tree of n items $c(n)=2H(n-1)-(2(n-1)/n) \qquad \qquad H(n)=\sum_{k=1}^n\frac{1}{k}$

•
$$h(x)$$
 is the depth at which x is found in tree

• Score: outlier $(x, n) = 2^{-\frac{E(h(x))}{c(n)}}$

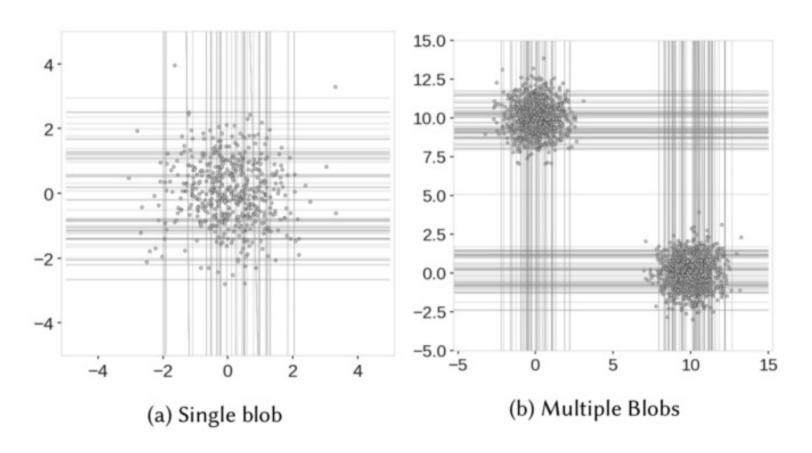
Outlier scores in isolation forests

(each tree is built from a sub-sample of original data)



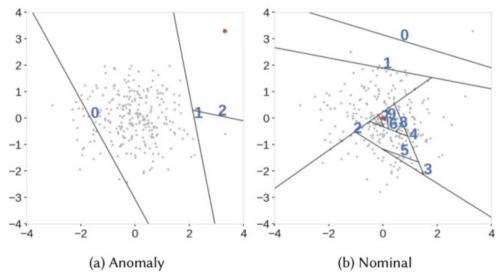
Example

(Note: here lines cross each other: we do not cross lines)



Extended Isolation Forest

 More freedom to partitioning by choosing a random slope and a random intercept

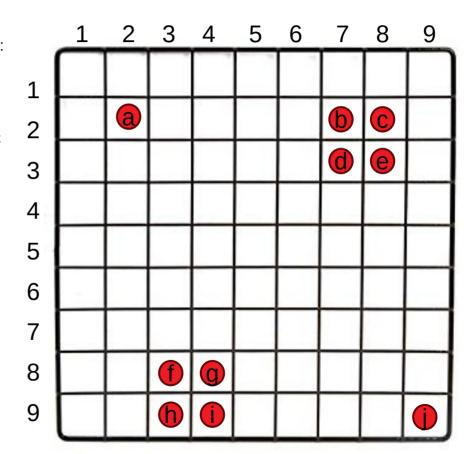


Exercise: isolation forest

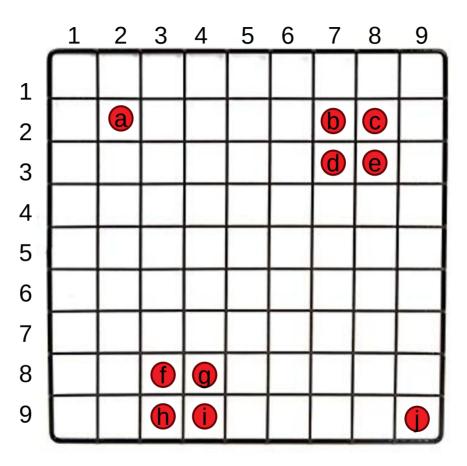
- Create one tree of the isolation forest by repeating 4 times:
 - Picking a sector containing >1 element
 - Picking a random dimension
 - Picking a random cut-off between min and max value along that dimension
 - Draw the line of your cut do not cross lines, and label each line with a number 0, 1, 2, ...
- Stop when each point is isolated
- Label each point with its depth h(x)

This is normally repeated several times, in the end:

outlier
$$(x,n) = 2^{-\frac{E(h(x))}{c(n)}}$$



In this case $c(10) = 2xH(9) - (2x9/10) \approx 3.857 \approx 4$



Summary

Things to remember

- Density-based methods
- Isolation forest

Exercises for TT19-TT21

- Data Mining, The Textbook (2015) by Charu Aggarwal
 - Exercises $8.11 \rightarrow \text{all except } 10, 15, 16, 17$

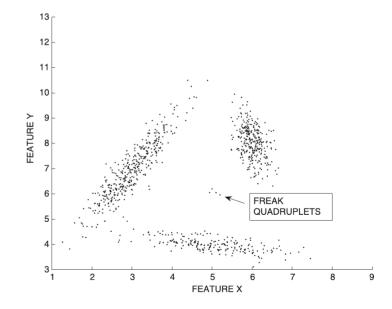
Additional contents (not included in exams)



Distance-based methods

Instance-specific definition

- The distance-based outlier score of an object x is its distance to its kth nearest neighbor
- In this example of a small group of 4 outliers, we can set k > 3

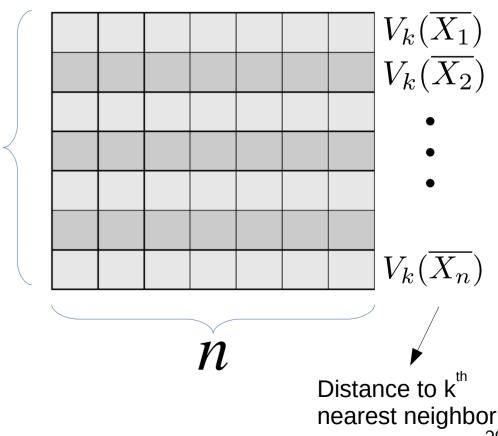


Problem: computational cost

- The distance-based outlier score of an object x is its distance to its k^{th} nearest neighbor
- In principle this requires O(n²) computations!
 - Index structure:
 useful only for cases of low data dimensionality
 - Pruning tricks:
 useful when only top-r outliers are needed

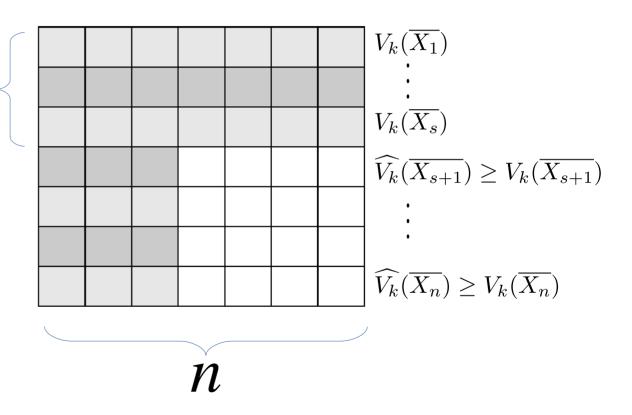
Problem: computational cost

- The distance-based outlier score of an item x is its distance to its kth nearest neighbor
- In principle this requires:
 - $O(n^2)$ computations for evaluating the n x n distance matrix
 - $O(n^2)$ computations for finding the r smallest values on each row



Pruning method: sampling

- Evaluate s x n distances
- For points1...s we are OK
- For points (s+1)...n
 we know only upper
 bounds



Pruning method: sampling (cont.)

From points

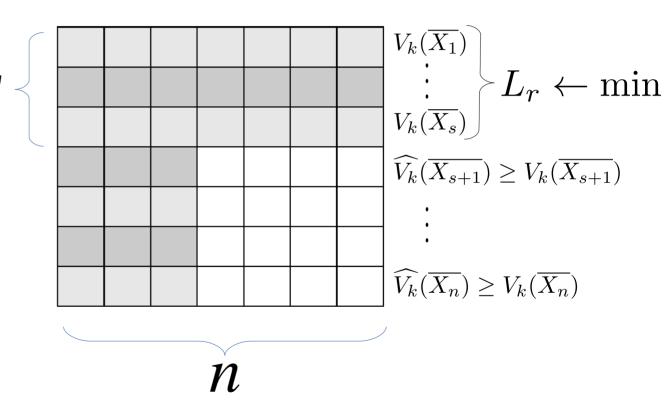
1...s we already know the \mathbf{S} "winners"

($r \le s$ nodes with the larger distance to their k^{th} nearest neighbor)

Any point having

 $V_{k} < L_{s}$ cannot be among

the top r outliers



Pruning method: sampling (cont.)

From points

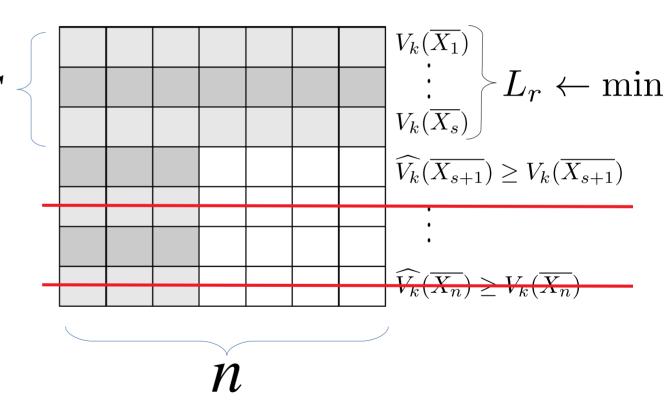
1...s we already know the ${\bf 5}$ "winners"

($r \le s$ nodes with the larger distance to their k^{th} nearest neighbor)

Any point having

 $V_{k} < L_{c}$ cannot be among

the top r outliers



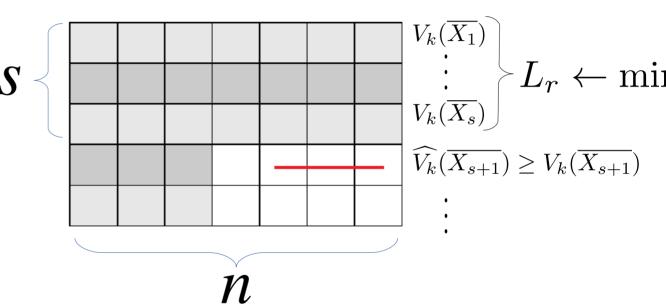
Pruning method: sampling (cont.)

Remove points

having
$$\widehat{V_k} \leq L_r$$

Update L_{keeping} r largest values, and stop computing for a row if one already finds k nearest neighbors in that row that

are all below distance L



Local outlier factor

Local Outlier Factor (LOF)

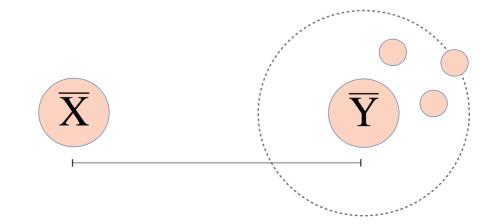
- Let $V_k(\overline{X})$ be the distance of X to its k-nearest neighbor
- Reachability distance

$$R_k(\overline{X}, \overline{Y}) = \max\{\text{Dist}(\overline{X}, \overline{Y}), V_k(\overline{Y})\}$$

- $V_k(\overline{X})$: distance of \overline{X} to its k-nearest neighbor
- Reachability distance

$$R_k(\overline{X}, \overline{Y}) = \max\{\text{Dist}(\overline{X}, \overline{Y}), V_k(\overline{Y})\}$$

- Not symmetric
- Equal to simple distance for long distances
- Smoothed by $V_k(\overline{X})$ for short distances



Reachability distance

$$R_k(\overline{X}, \overline{Y}) = \max\{\text{Dist}(\overline{X}, \overline{Y}), V_k(\overline{Y})\}$$

Average reachability distance

$$AR_k(\overline{X}) = \underset{\overline{Y} \in L_k(\overline{X})}{E} \left[R_k(\overline{X}, \overline{Y}) \right]$$

 $L_{\mathbf{k}}(\overline{X})$ is the set of points within distance $V_{\mathbf{k}}(\overline{X})$ of \overline{X} (might be more than k due to ties)

$$R_{k}(\overline{X}, \overline{Y}) = \max\{\operatorname{Dist}(\overline{X}, \overline{Y}), V_{k}(\overline{Y})\}$$

$$AR_{k}(\overline{X}) = \underset{\overline{Y} \in L_{k}(\overline{X})}{E} \left[R_{k}(\overline{X}, \overline{Y})\right]$$

Local outlier factor

$$LOF_k(\overline{X}) = \mathop{E}_{\overline{Y} \in L_k(\overline{X})} \frac{AR_k(\overline{X})}{AR_k(\overline{Y})}$$

Outlier score

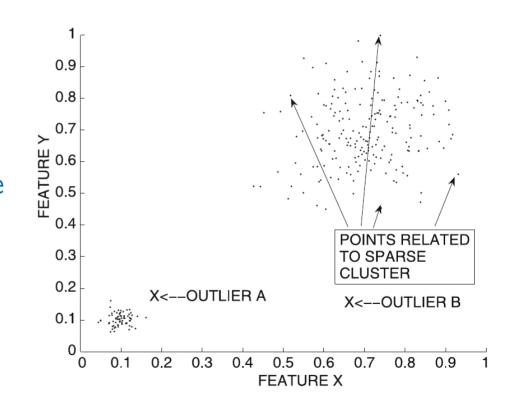
$$\max_k \mathrm{LOF}_k(\overline{X})$$

Large for outliers, close to 1 for others

Local outlier factor

$$LOF_k(\overline{X}) = \mathop{E}_{\overline{Y} \in L_k(\overline{X})} \frac{AR_k(X)}{AR_k(\overline{Y})}$$

- LOF values for points inside a cluster are close to one if cluster is homogeneous
- LOF values much higher for outliers: they are computed in terms of average distances of near-by clusters



Exercise

compare outlier score LOF(u), LOF(v)

- Let k=2
- LOF₂(u) = E[{AR₂(u) /AR₂(a), AR₂(u)/AR₂(b)}] = ______
- LOF₂(v) = E[{AR₂(v) /AR₂(b), AR₂(v)/AR₂(u)}] = ______
- $AR_2(u) = E[\{R_k(u,a), R_k(u,b) \}] = \underline{\hspace{1cm}}$
- $AR_2(v) = E[\{R_k(v,b), R_k(v,u) \}] = \underline{\hspace{1cm}}$
- $AR_2(a) = E[\{R_k(a,u), R_k(a,b) \}] = \underline{\hspace{1cm}}$
- $AR_{2}(b) = E[\{R_{k}(b,u), R_{k}(b,a)\}] = \underline{\hspace{1cm}}$
- $R_k(a,u) = \underline{\hspace{1cm}}; R_k(a,b) = \underline{\hspace{1cm}}; R_k(b,u) = \underline{\hspace{1cm}}; R_k(b,a) = \underline{\hspace{1cm}}$
- $R_k(u,a) = \underline{\hspace{1cm}}; R_k(u,b) = \underline{\hspace{1cm}}; R_k(v,b) = \underline{\hspace{1cm}}; R_k(v,u) = \underline{\hspace{1cm}}$
- $V_2 = distance to 2^{nd} nearest neighbor: V_2(u) = _____; V_2(v) = _____; V_2(a) = _____; V_2(b) = _____;$

$$AR_k(\overline{X}) = \underset{\overline{Y} \in L_k(\overline{X})}{E} \left[R_k(\overline{X}, \overline{Y}) \right]$$

 $R_k(\overline{X}, \overline{Y}) = \max{\{\mathrm{Dist}(\overline{X}, \overline{Y}), V_k(\overline{Y})\}}$