Core of a Polytope and Rich Cells in Random Arrangement of Hyperplanes

Justin Lewis

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1 Introduction

These notes are the beginning of an analysis of the distribution of polytope facet complexity within a random hyperplane arrangement. The notion of the Core of a polytope is introduced as a new analysis tool. With this notion formalized, hopefully it can be used to show that the proportion of rich cells in a random hyperplane arrangement is low. NOTE: the region described below is actually not correctly characterized for polytopes for $d \ge 3$ (ask Justin for more details)

2 Notation

First, some standard definitions as well as an overview of notation:

- Hyperplane h is defined as $\{x \mid \langle c, x \rangle + d = 0\}$ for $x \in \mathbb{R}^d$
- Halfspace H is similarly defined as $\{x \mid \langle c, x \rangle + d \leq 0\}$
- Polytope P is the intersection of n halfspaces. $P = \bigcap_{i=1}^{n} H_i$ or equivalently $P = \{x | C^{\mathsf{T}} x + \vec{d} \leq 0 \}$
- Vertex set $\mathcal{V}(P)$ denotes the set of vertices of polytope P.
- Vertex neighborhood N(v) takes vertex v of a polytope and returns the 1-hop neighborhood of v on the 1-skeleton of the polytope.

3 Polytope Core

In words, the Core of a polytope C(P) is a region such that any hyperplane with nonempty intersection with C(P) is guaranteed to not increase the polytope's number of facets. Equivalently, such an intersecting hyperplane is guaranteed to not cleanly separate any vertex $v \in \mathcal{V}(P)$ of polytope P. A vertex v is considered cleanly separated by hyperplane h if v is strictly separated from every neighbor

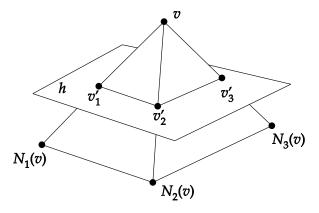


Figure 1: Example of a cleanly cut vertex in \mathbb{R}^3

 $N_j \in N(v)$. More specifically, hyperplane $h = \{x | \langle c, x \rangle + d = 0\}$ cleanly separates v if it satisfies:

$$sign(\langle c, v \rangle + d) \neq sign(\langle c, N_j(v) \rangle + d) \quad \forall j$$
$$\langle c, v \rangle + d \neq 0, \langle c, N_j(v) \rangle + d \neq 0 \quad \forall j$$

To mathematically define C(P), first consider a single vertex $v \in \mathcal{V}(P)$. One can define a region R(v) such that any hyperplane with nonempty intersection with R(v) is guaranteed to not cleanly separate v. To specify R(v) consider its complement $R(v)^C$ which is the region that can be reached by the space of hyperplanes cleanly separating v. In general position, each vertex v has d neighbors (i.e. the cardinality |N(v)| = d). Thus, the space of hyperplanes cleanly separating v is characterized by d independent choices of points $\{v'_1, v'_2, ..., v'_d\}$ s.t:

$$v_j' = \alpha_j v + (1 - \alpha_j) N_j(v) \quad \forall j$$
 (1)

$$\alpha_i \in (0,1) \tag{2}$$

There are two regions, $S_1(v)$ and $S_2(v)$, which cannot be reached by the space of cleanly separating hyperplanes (Proof \rightarrow section 5.1) (???). Let $\tilde{h}(v)_j \, \forall j \in \{1, 2, ..., d\}$ denote one of the d hyperplanes which pass through v. Then, let $\tilde{H}(v)_j$ denote the corresponding negative halfspace associated with each of these hyperplanes. Lastly, let $\hat{H}(v)$ denote the negative halfspace supported by the hyperplane which passes through the points N(v) and whose normal vector has positive inner product with v, then we have:

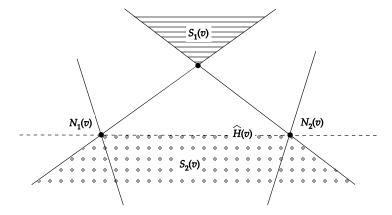


Figure 2: Example of R(v) for a vertex of a polytope in \mathbb{R}^2

$$S_1(v) = \bigcap_{j=1}^d \tilde{H}(v)_j^C$$

$$S_2(v) = \{\bigcap_{j=1}^d \tilde{H}(v)_j\} \cap \hat{H}(v)$$

$$(3)$$

$$S_2(v) = \{ \bigcap_{j=1}^d \tilde{H}(v)_j \} \bigcap \hat{H}(v)$$

$$\tag{4}$$

$$R(v) = S_1(v) \bigcup S_2(v) \tag{5}$$

Now, we have all of the components to define the Core of a polytope. For each vertex, we have defined R(v) which is the region which guarantees nonclean separation for a single vertex. The Core of P is the region which possesses this property for all vertices, therefore:

$$C(P) = \bigcap_{j=1}^{|\mathcal{V}(P)|} R(v_j)$$

$$= \{ \bigcap_{j=1}^{|\mathcal{V}(P)|} \hat{H}(v_j) \} \bigcap P$$
(6)

The above equality in (6) holds as every non-redundant hyperplane of the polytope passes through at least one vertex and $S_2(v) \notin C(P) \ \forall v \text{ of } P$ (Proof \rightarrow section 5.2) (????). Thus, there is a simple mathematical expression for C(P), it is the intersection of $|\mathcal{V}(P)|$ neighbor-defined halfspaces in turn intersected with the polytope P. Note, this rearrangement shows that $C(P) \subset P$.

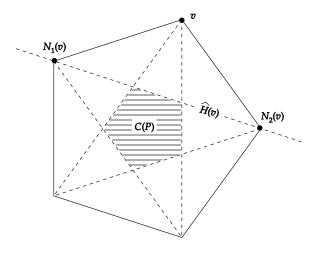


Figure 3: Example of C(P) for a polytope in \mathbb{R}^2

4 Polytope Core Expansion

In this section, we begin to consider the evolution of C(P) as the number of facets of P increases. Let P' represent the new polytope after increasing the facets of P by one. With Lemmas ?? and ?? we will prove Theorem 4.2, that is, the volume of the Core always increases with each added facet. To understand the change from C(P) to C(P'), first we understand how the polytope changes. Note, the new facet added via new hyperplane h' cleanly separates v from P, thus:

$$\mathcal{V}(P') = \{ \mathcal{V}(P) \setminus v \} \bigcup \{ v'_1, v'_2, ..., v'_d \}$$
 (7)

That is, the vertex v is removed and d new vertices are added. Note that these new vertices satisfy equations (1) and (2) (i.e. each is a strict convex combination of v and $N_j(v)$ for some j). Now, we consider the changes to the Core due to the vertex set change. The changes from C(P) to C(P') are as follows:

- 1) The halfspace $\hat{H}(v)$ is removed
- 2) d new halfspaces $\{\hat{H}(v_1'),...,\hat{H}(v_d')\}$ are added
- 3) d old halfspaces $\{\hat{H}(N_1(v)),...,\hat{H}(N_d(v))\}$ are modified

To show that $C(P) \subset C(P')$ first we show that $\hat{H}(v_i')$ is redundant to $R(v) \forall i$. This takes care of changes 1) and 2). Then we show that $\hat{H}(N_i(v))'$

is redundant to $\hat{H}(N_j(v)) \, \forall j$. This takes care of change 3). We start with the first redundancy claim, thus we wish to prove:

Lemma 4.1. Given hyperplane h' which cleanly separates v, d new vertices are added to P, and thus d new hyperplanes $\hat{h}_i = \{x | \langle \hat{c}_i, x \rangle + \hat{d}_i = 0\} \ \forall i \ are \ added$ to C(P). Furthermore, \hat{h}_i is redundant to $R(v) \ \forall i$:

$$\langle \hat{c}_i, x \rangle + \hat{d}_i \le 0 \quad \forall x \in R(v)$$
 (8)
 $\forall i$ (9)

$$\forall i$$
 (9)

Proof.

Theorem 4.2. Given P', the polytope which results from adding one new facet to an arbitrary polytope P, we have that: $C(P) \subset C(P')$

Appendix 5

- Proof of Expression of R(v)5.1
- 5.2Proof of Simplified Expression of C(P)

6 Random Notes

6.1Minor Details

- perhaps somewhere it should be mentioned that the beginning version of the proofs assume general position everywhere and later we will consider the case outside of general position.
- Much of the discussion about the Core of a polytope runs into confusion when the polytope is unbounded. (E.g. for unbounded polytope not every vertex has d neighbors, thus how is H(v) defined?)
- I think that the boundaries of the definition of R(v) might be improperly defined. (E.g. $S_1(v)$ is the complement plus the boundary?)
- Does "negative halfspace" include the boundary of the halfspace
- Figure 2 and Figure 3 should be changed from dashed lines to solid lines