

Cycles and Girth

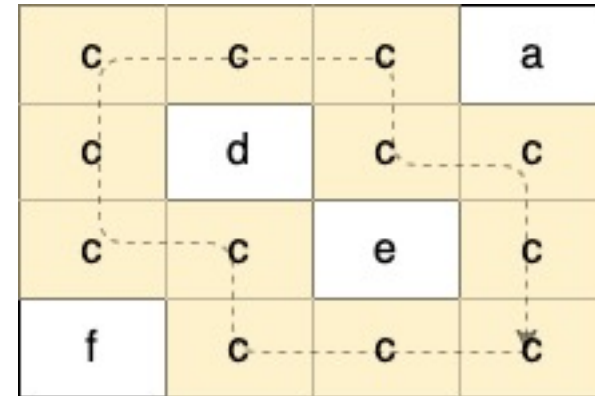
Instructor: Meng-Fen Chiang

COMPSCI: WEEK 10.1



OUTLINE

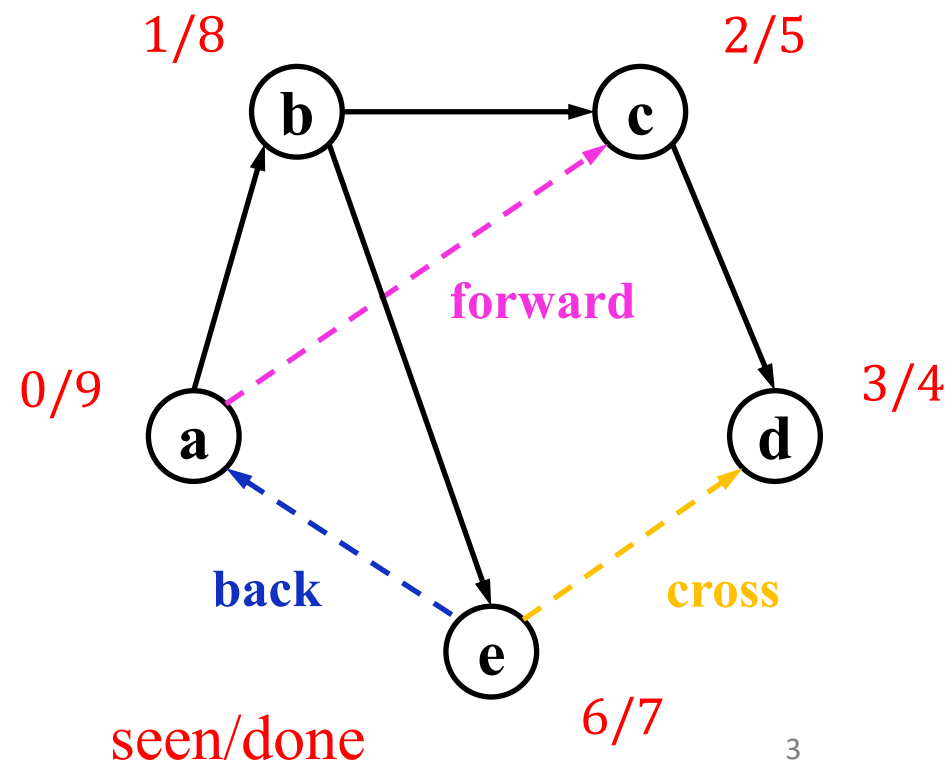
- Terminology
 - Cycle
 - Girth
- DFS Application: Cycle Detection
- BFS Application: Girth Computation
- Correctness of Girth Computation



RECAP: Example 23.7

- Explain how to determine (u,v) in DFS algorithm whether it is a tree-, back-, forward- or cross-arc?

- If v is **white**, then (u,v) is a **tree** arc
- If v is **grey**, then (u,v) is a **back** arc
- If v is **black**, then (u,v) is
 - a **cross** arc($seen[v] < seen[u]$, $done[v] < seen[u]$), or
 - a **forward** arc($seen[u] < seen[v] < done[v] < done[u]$).

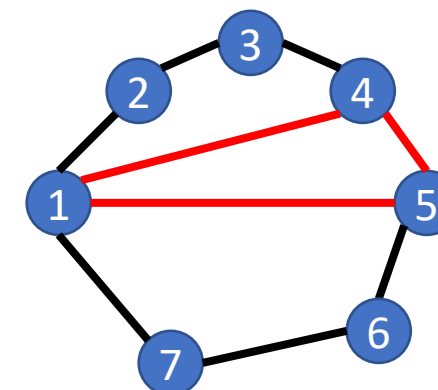
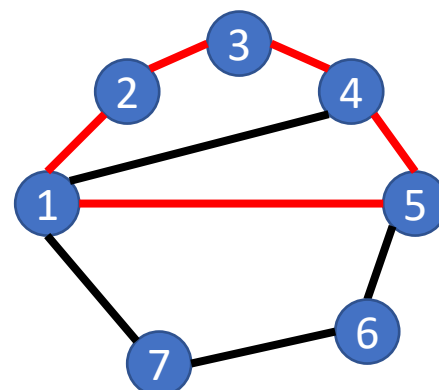
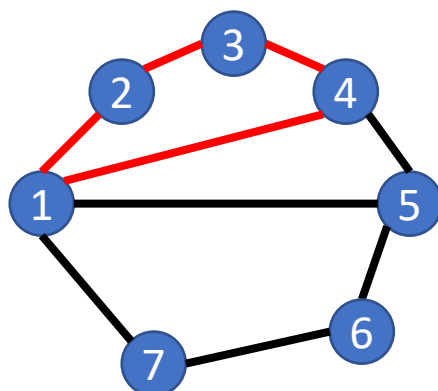
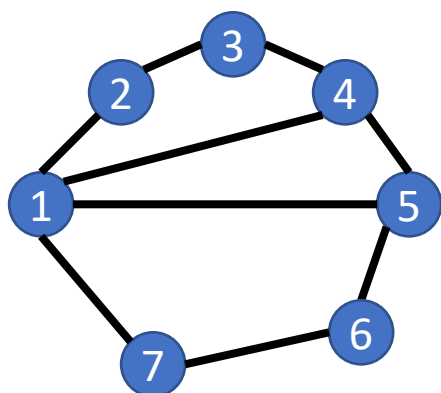


Cycle Detection

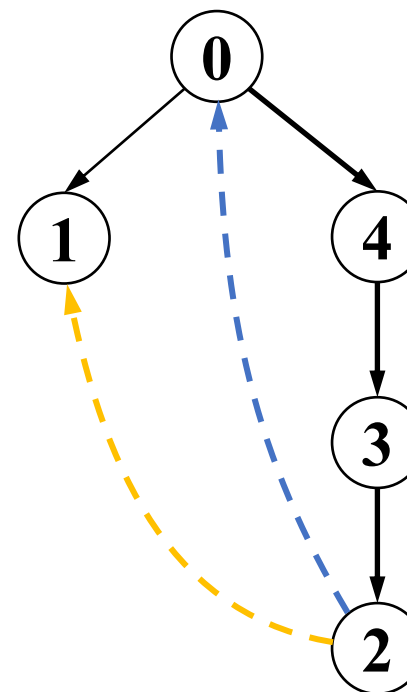
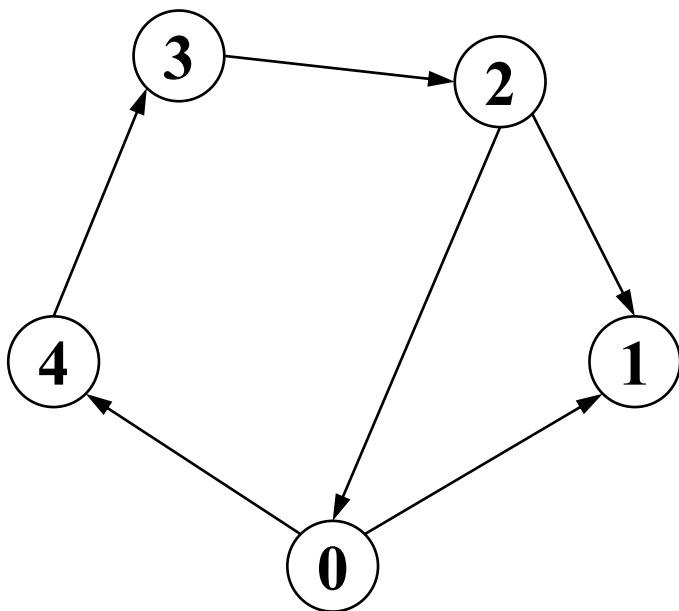
- Suppose that there is a cycle in G and let v be the node in the cycle visited first by **DFS**. If (u, v) is an arc in the cycle then it must be a back arc.
- Conversely if there is a **back arc**, we must have a **cycle**.
- Suppose that DFS is run on a digraph G . Then G is acyclic if and only if G does not contain a back arc.

Example: Cycles in Graphs

- A graph can have several cycles



Using DFS to Find Cycles in Digraphs



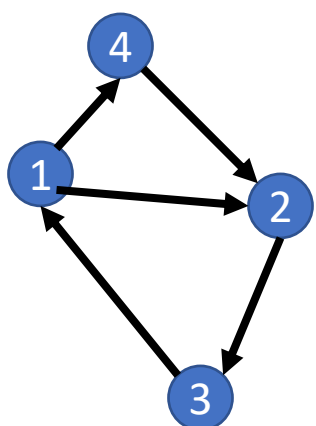
Back arc detected!

Once DFS finds a cycle, the stack contains the nodes that form the cycle.

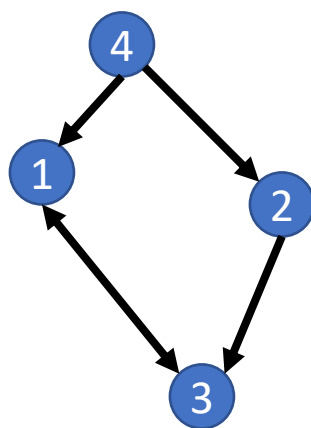
Girth and Digirth

- **Definition.** For a graph (with a cycle), the length of the shortest cycle is called the **girth** of the graph. If the graph has no cycle then the girth is undefined, can be viewed as $+\infty$.
- **Convention.** For a digraph we use the term **girth** for its **underlying graph** and the term **directed girth** for the length of the smallest directed cycle.

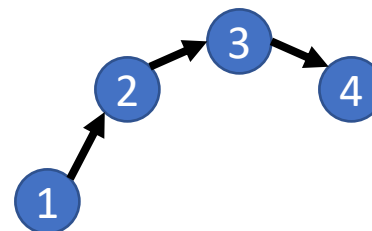
Examples: Girth of a Graph (Digraph)



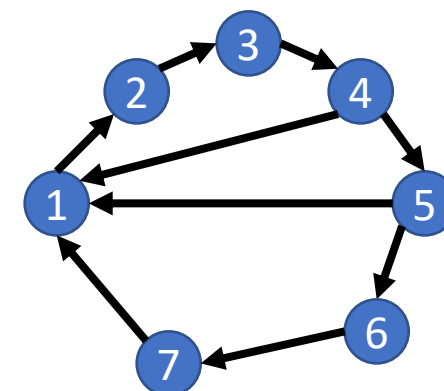
Girth: 3
Directed girth: 3



Girth: 4
Directed girth: 2

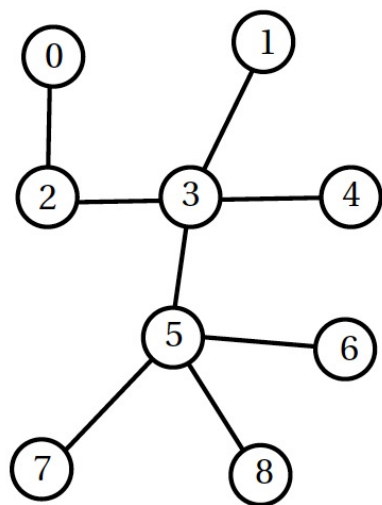


Girth: $+\infty$
Directed girth: $+\infty$

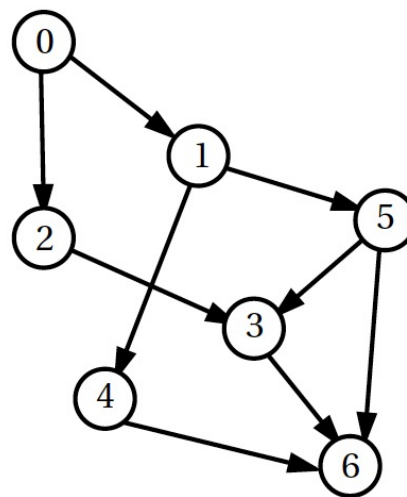


Girth: 3
Directed girth: 4

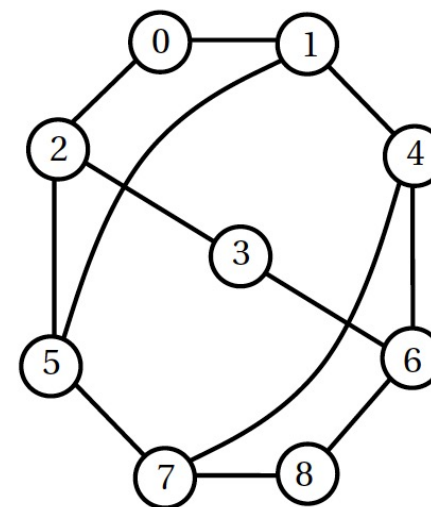
Examples: Girth of a Graph (Digraph)



Girth: $+\infty$
Directed girth: undefined



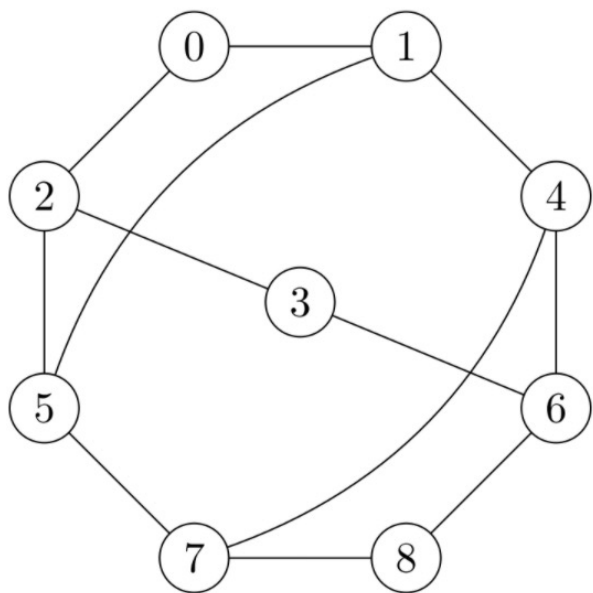
Girth: 3
Directed girth: $+\infty$



Girth: 4
Directed girth: undefined

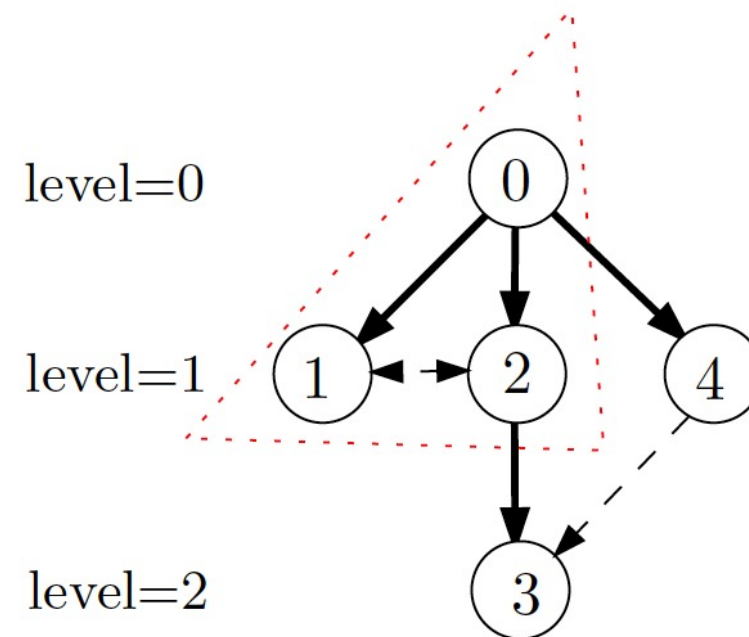
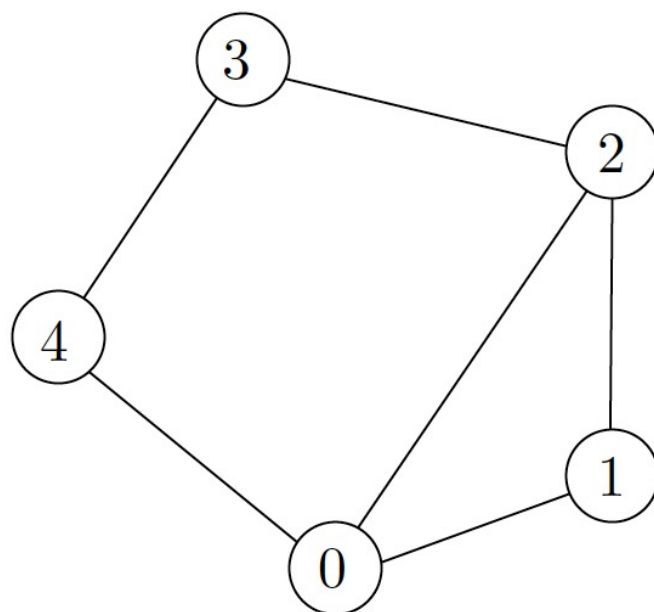
Finding the Girth of a Graph

- An easy-to-implement **DFS** idea may not work properly. Consider the **DFS** tree originating from vertex 0 of the graph below. Which is the smallest cycle it finds with the back edges? Which smaller cycle does it miss?



Finding the Girth of a Graph

- **Using BFS to find cycles in graphs.** Cycles can also be easily detected in a graph using BFS. Finding a cycle of **minimum** length in a graph is not difficult using BFS (better than DFS).

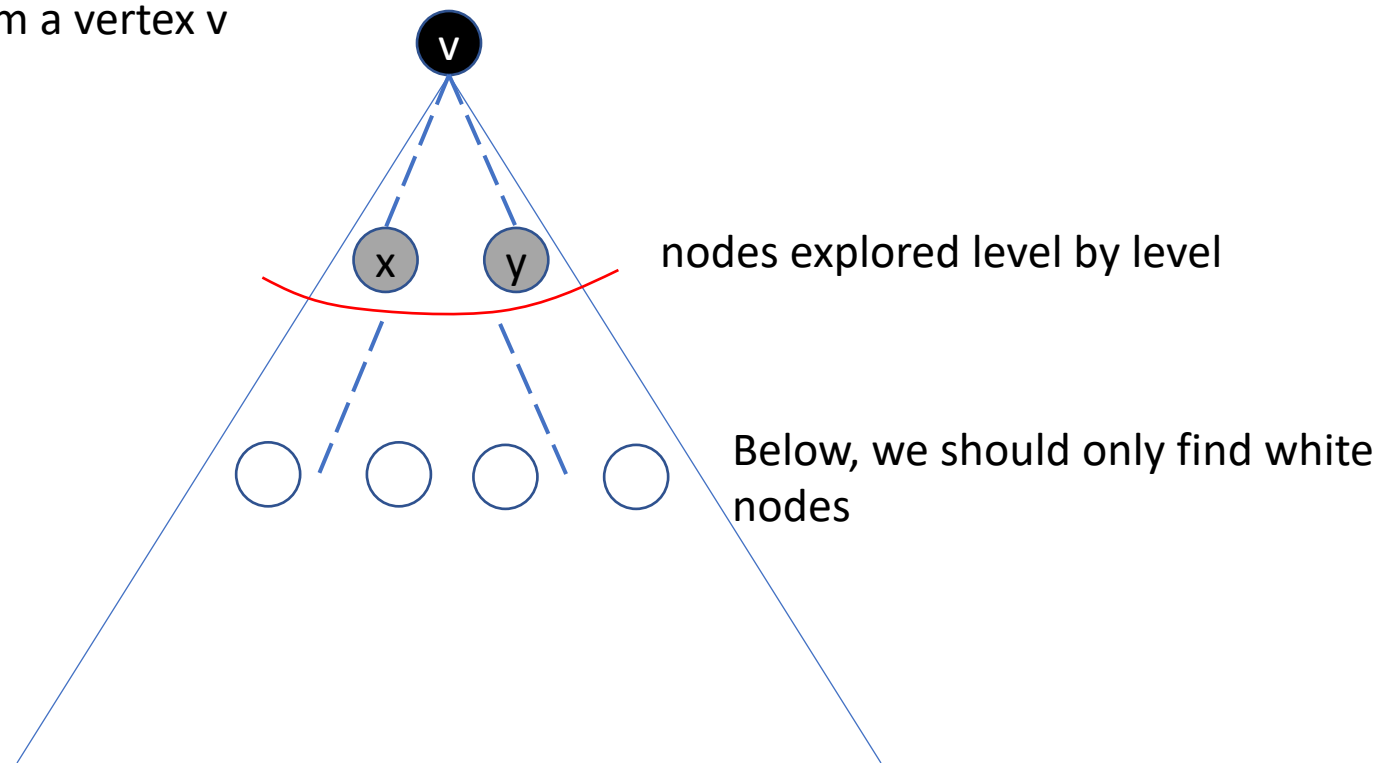


Finding the Girth of a Graph

- Perform `BFSVisit()` $|V|$ times, starting at each vertex $v' \in V$ in turn.
- If during a `BFSVisit`, we encounter a **grey neighbour** (rather than a white) we have found a cycle
 - The grey node was visited before following another path
- An important property of BFS is that if it runs from vertex v , then every vertex s , when it is first reached, then the path that was found from v to s is **minimal**. Thus, reaching v' from v with BFS finds the shortest path from v to v' , namely the **shortest cycle** that contains v .

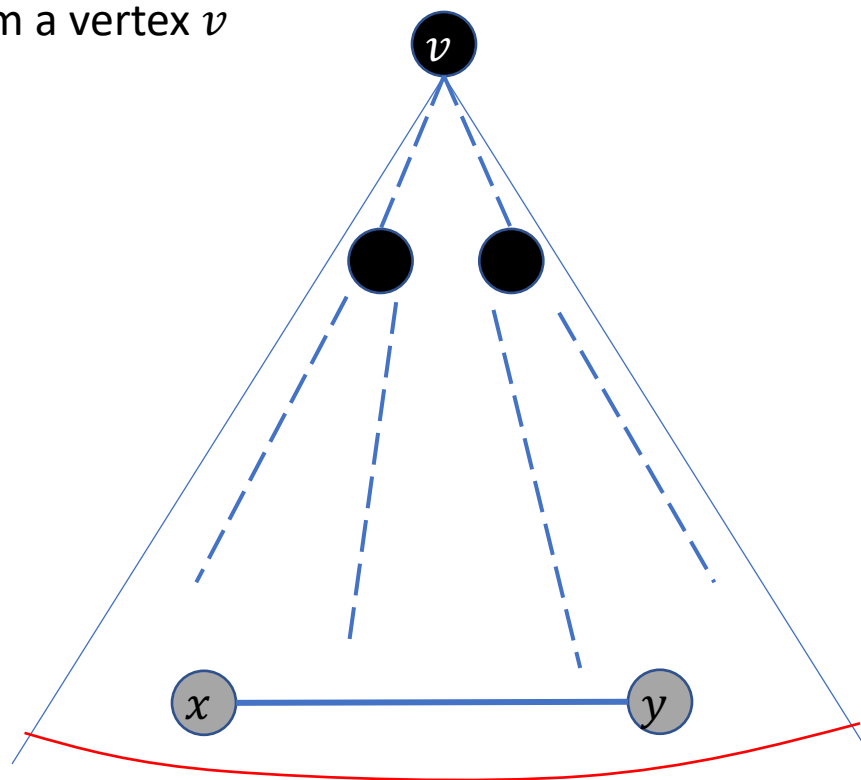
Example: Girth of a Graph

Executing BFSVisit from a vertex v



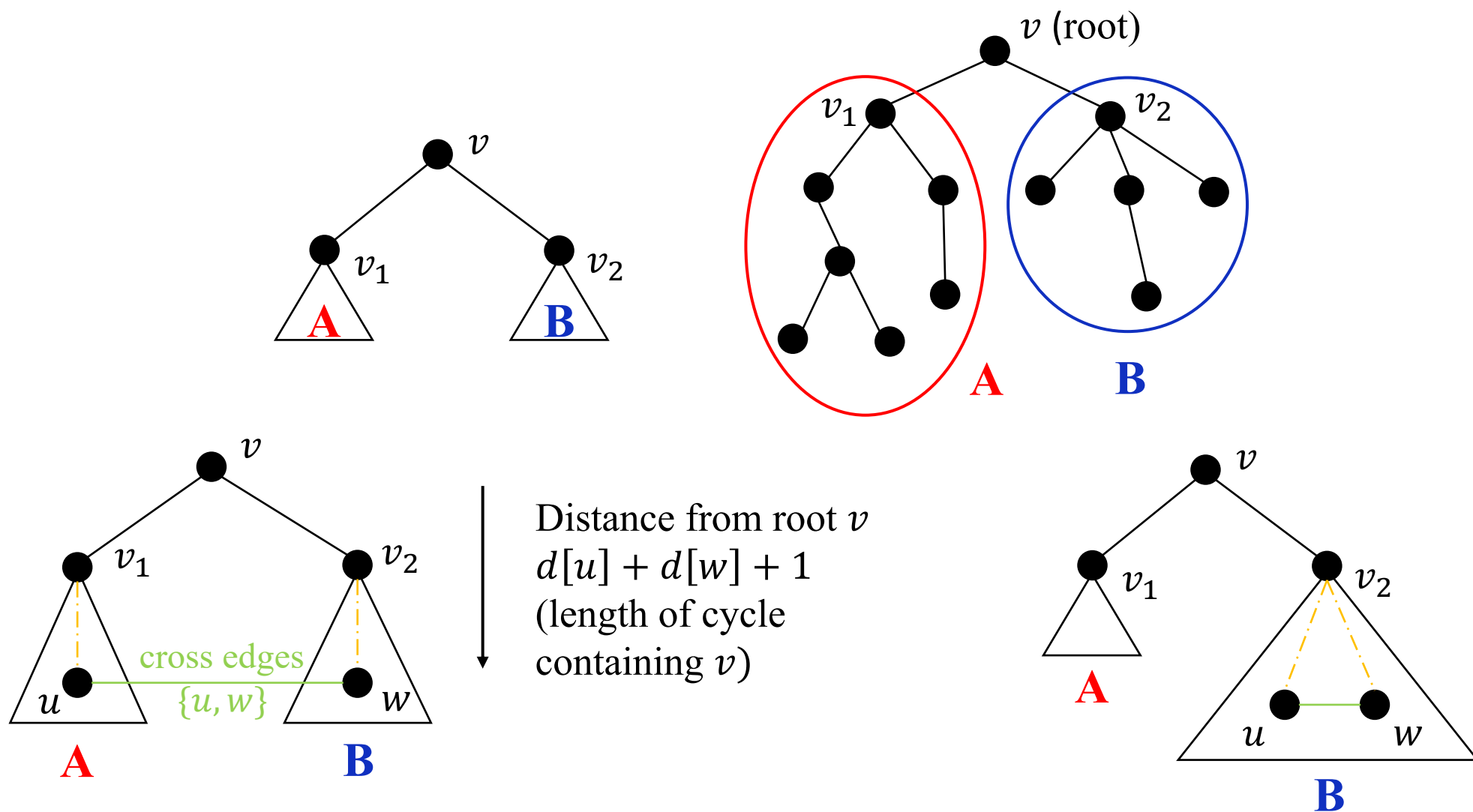
Example: Girth of a Graph

Executing BFSVisit from a vertex v



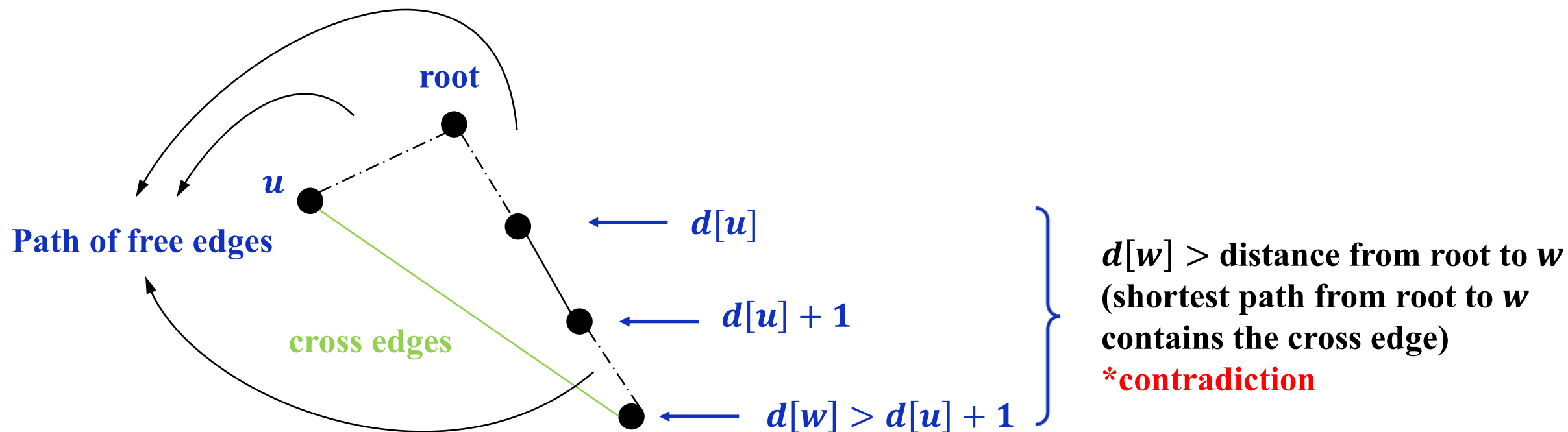
If a **cross** edge/arc exists then it will connect to a grey vertex
Both x and y are reachable from v therefore there is a cycle through the **cross** edge

Example: Girth of a Graph



BFS: Cross Edges in a Graph

- Let $\{u, w\}$ be a cross edge in a tree. Then either $d[u] = d[w]$ or $|d[u] - d[w]| = 1$
- Sketch of proof:** Suppose $|d[u] - d[w]| > 1$



Facts about Cycle Length

- If vertex v is on at least one cycle then BFS starting at v will find it.
- On detection of a cross edge between descendants u and w , determine whether u and w are in **different subtrees below v** (the root of the tree).
 - If yes, then a cycle of length $d[u]+d[w]+1$ is found.
 - If no, then a cycle of shorter length is found (but avoids v).
- If $d[u]=d[w]$ then **odd** length, where v is a common ancestor.
- Otherwise, $d[u]+1=d[w]$ and **even** length.

Facts about Cycle Length

- cross edge $\{u, w\}$

1. $d[u] = d[w]$

cycle length: $2d[u] + 1 \Rightarrow$ **odd** length
cross edge

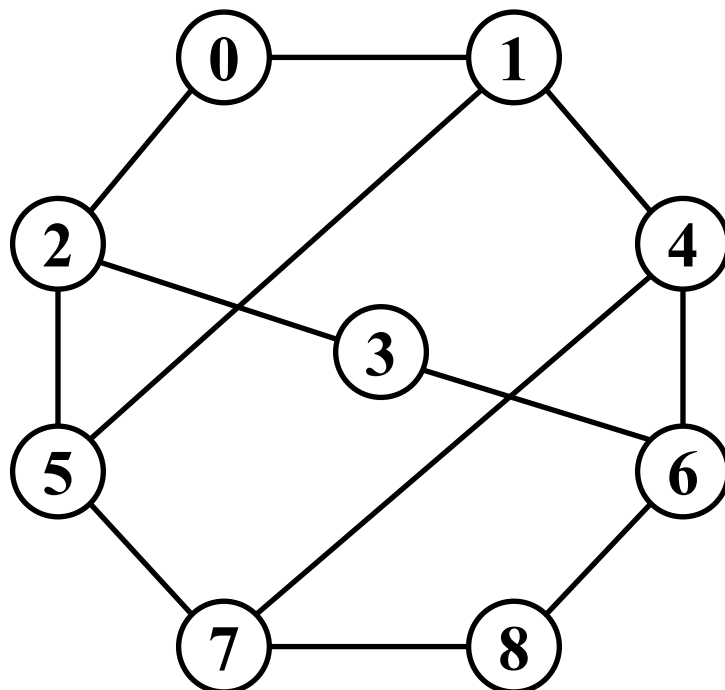
2. $d[u] + 1 = d[w]$

cycle length: $d[u] + d[w] + 1$
 $= d[u] + d[u] + 1 + 1$
 $= 2d[u] + 2 \Rightarrow$ **even** length

Finding the Girth of a Graph

- To compute girth, we perform $\text{BFSVisit}(v)$ procedure once for each $v \in V(G)$ and take **minimum**.
 - If a grey neighbour is met, e.g., an edge (x, y) is explored from x where y is grey, continue to the end of the current level and then stop.
 - For each edge (x, y) as above on this level, if v is the lowest common ancestor of x and y in the BFS tree, then there is a cycle containing x, y, v of length $l = d(x) + d(y) + 1$.
 - Report the minimum value of l obtained along the current level.
- The minimum of these lengths at the level is the smallest cycle that **involves v**
- The smallest cycle among all possible start vertices v is the girth.

Example (1): Finding the Girth of a Graph

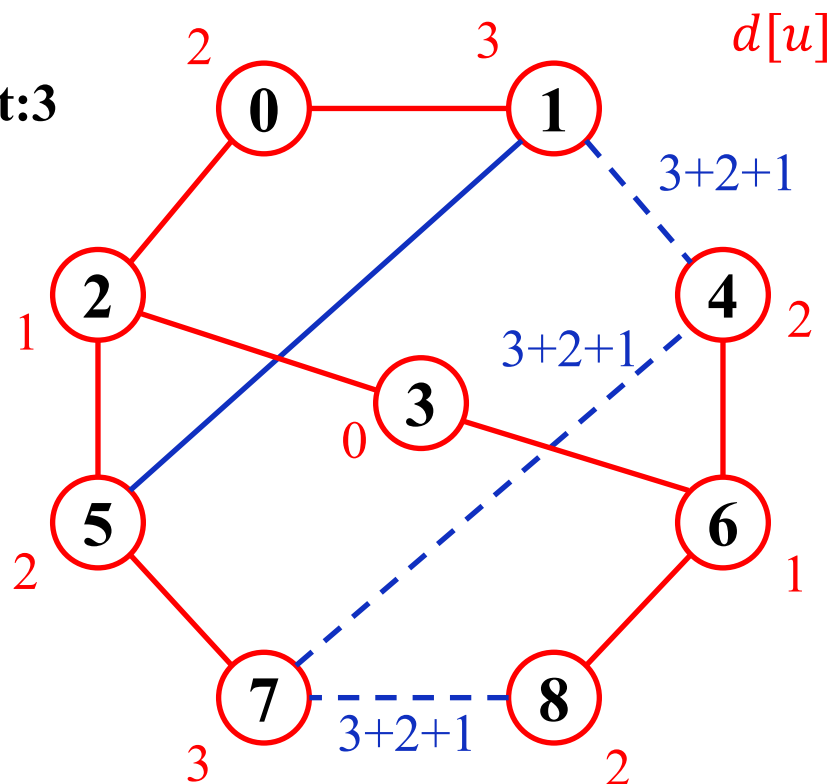


- Run BFS for 3, shortest cycle has length 6.
- Run BFS for any other vertex, shortest cycle has length 4.
- Girth of the graph is 4..

Example (1)

BFS

Start:3



Shortest cycle containing **3** has length 6

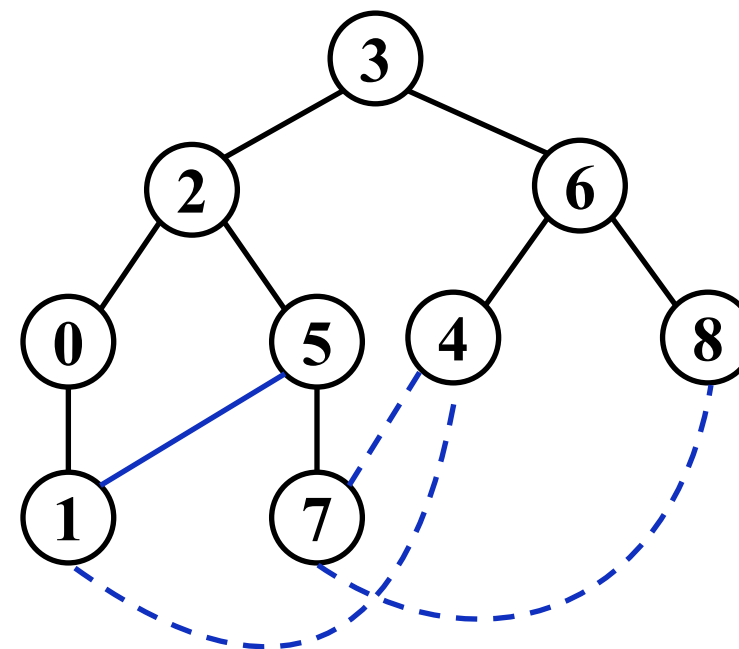
— tree edges
— cross edge same subtree
- - - tree edges different subtree
(relative to root **3**)

$d = 0$

$d = 1$

$d = 2$

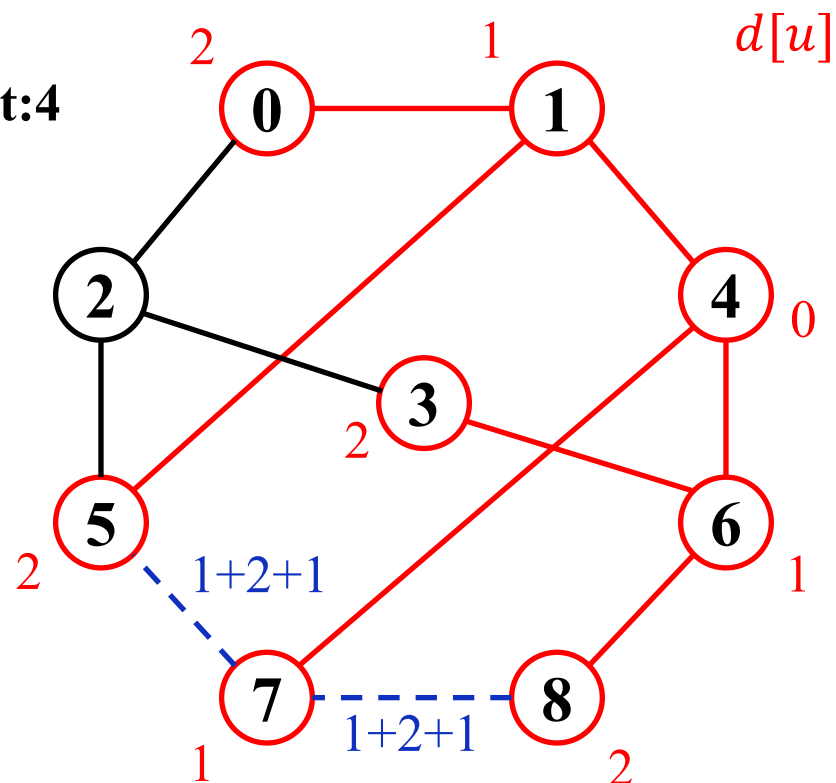
$d = 3$



Example (1)

BFS

Start:4



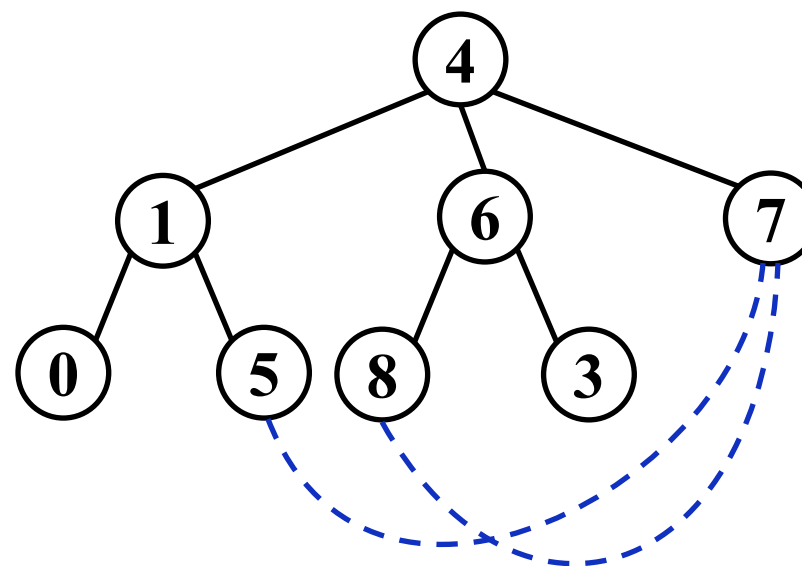
Shortest cycle containing **4** has length 4

- tree edges
- cross edge same subtree
- - - tree edges different subtree
(relative to root **4**)

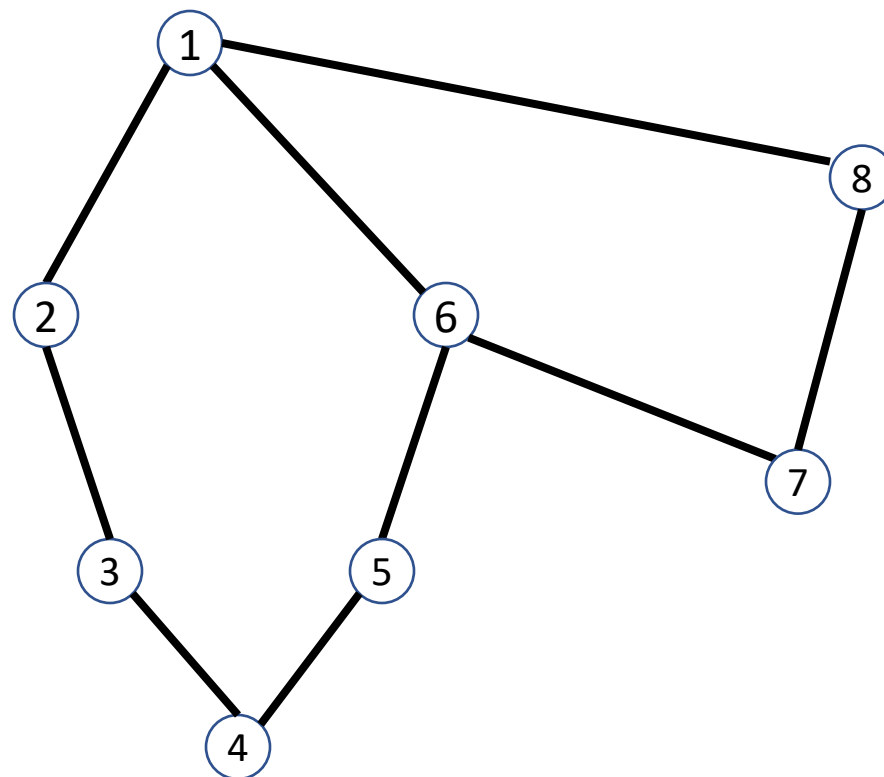
$d = 0$

$d = 1$

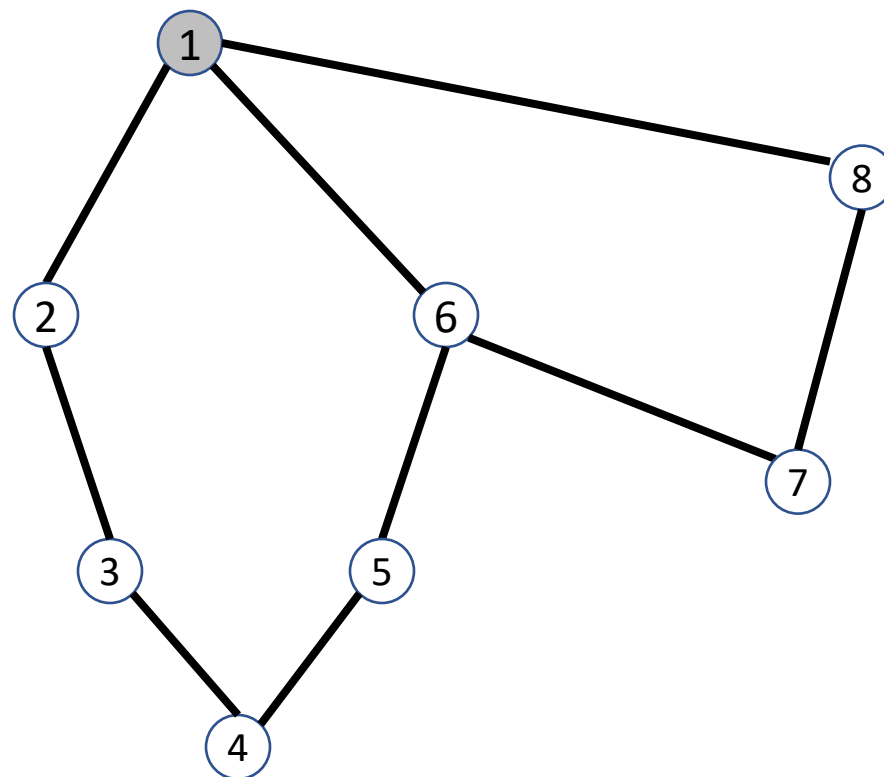
$d = 2$



Exercise: Finding the Smallest Cycle involving Vertex 1

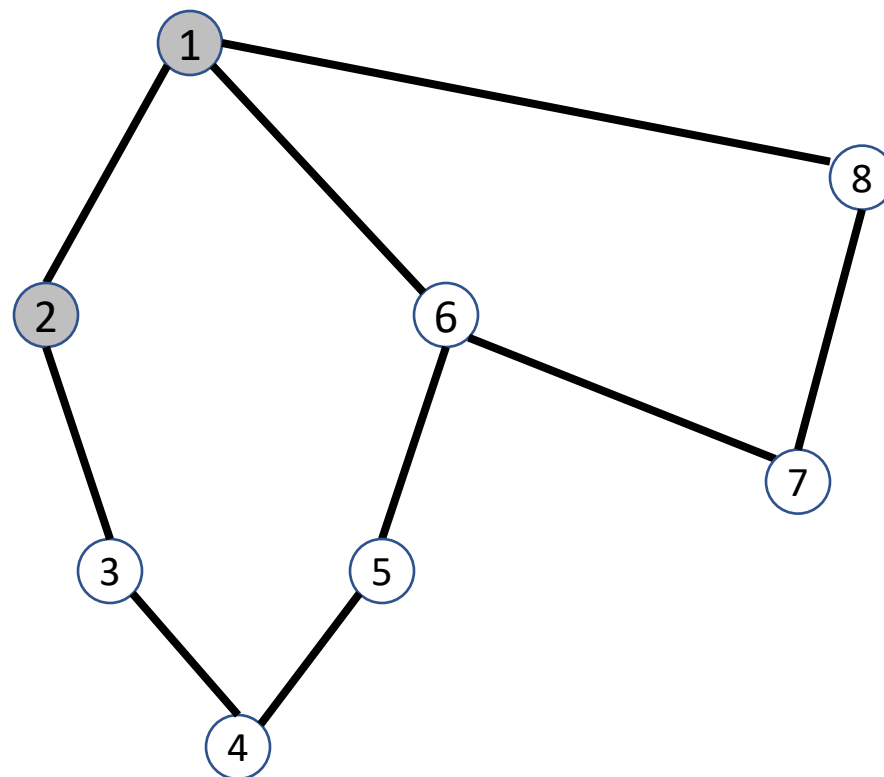


Exercise: Finding the Smallest Cycle involving Vertex 1



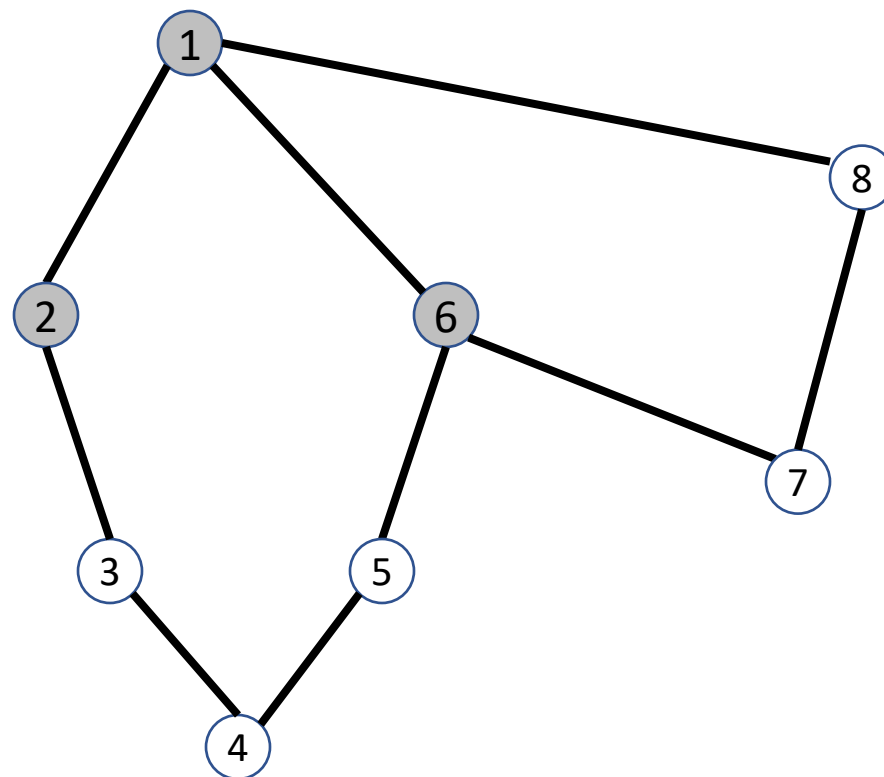
Queue: 1

Exercise : Finding the Smallest Cycle involving Vertex 1



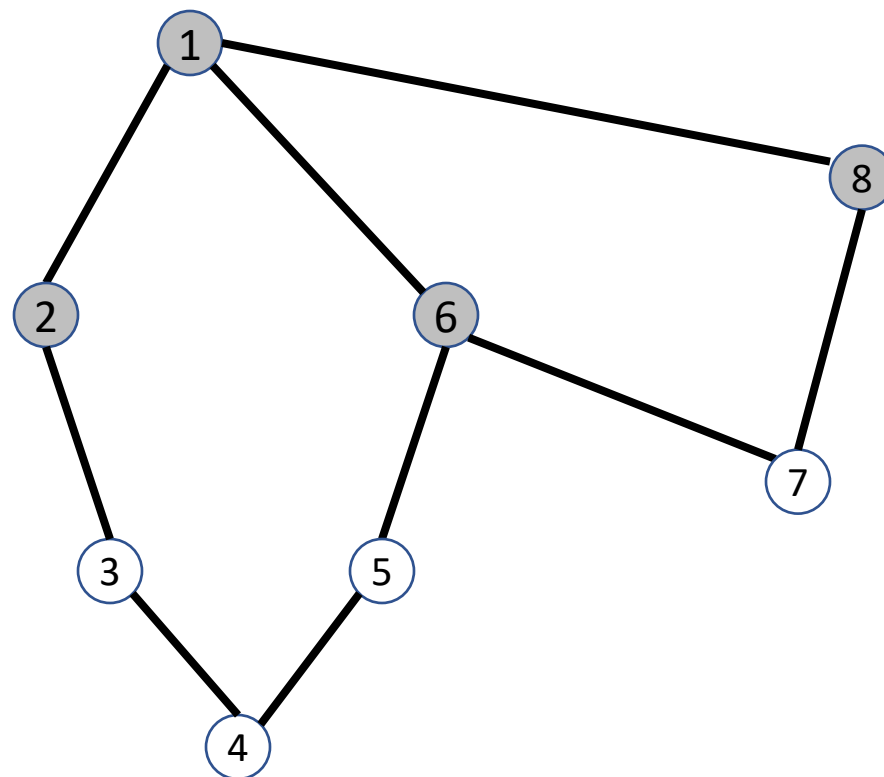
Queue: 1 2

Exercise: Finding the Smallest Cycle involving Vertex 1



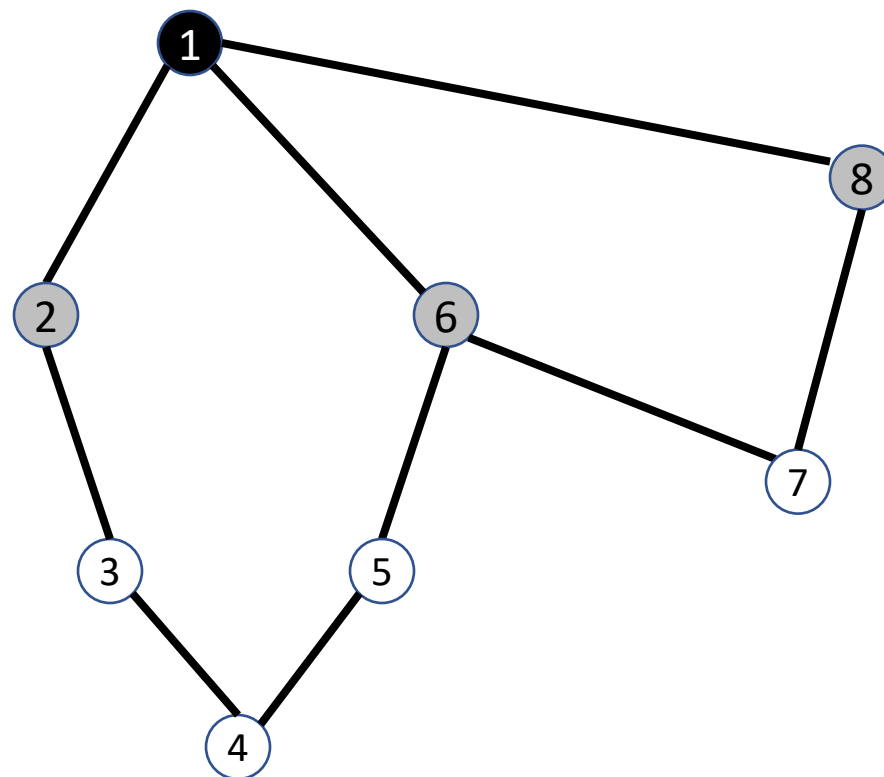
Queue: 1 2 6

Exercise: Finding the Smallest Cycle involving Vertex 1



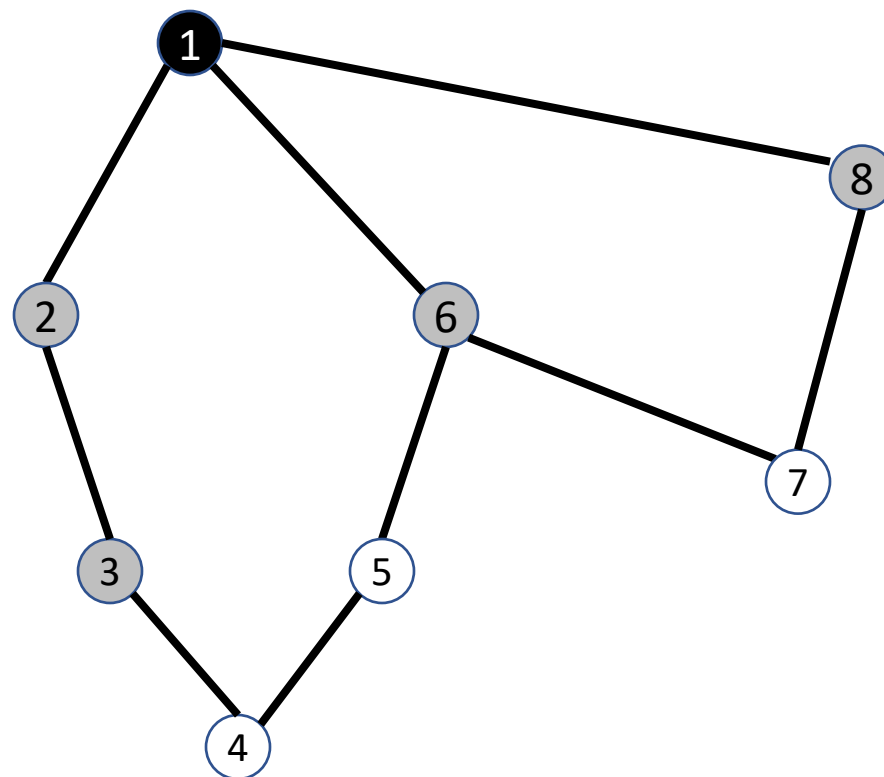
Queue: 1 2 6 8

Exercise: Finding the Smallest Cycle involving Vertex 1



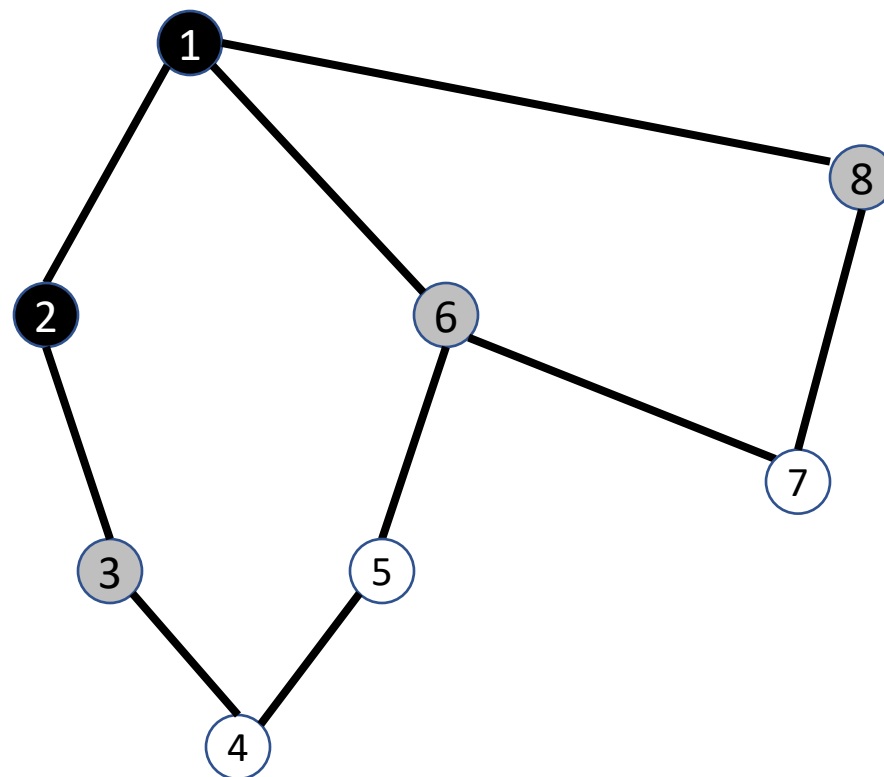
Queue: 2 6 8

Exercise: Finding the Smallest Cycle involving Vertex 1



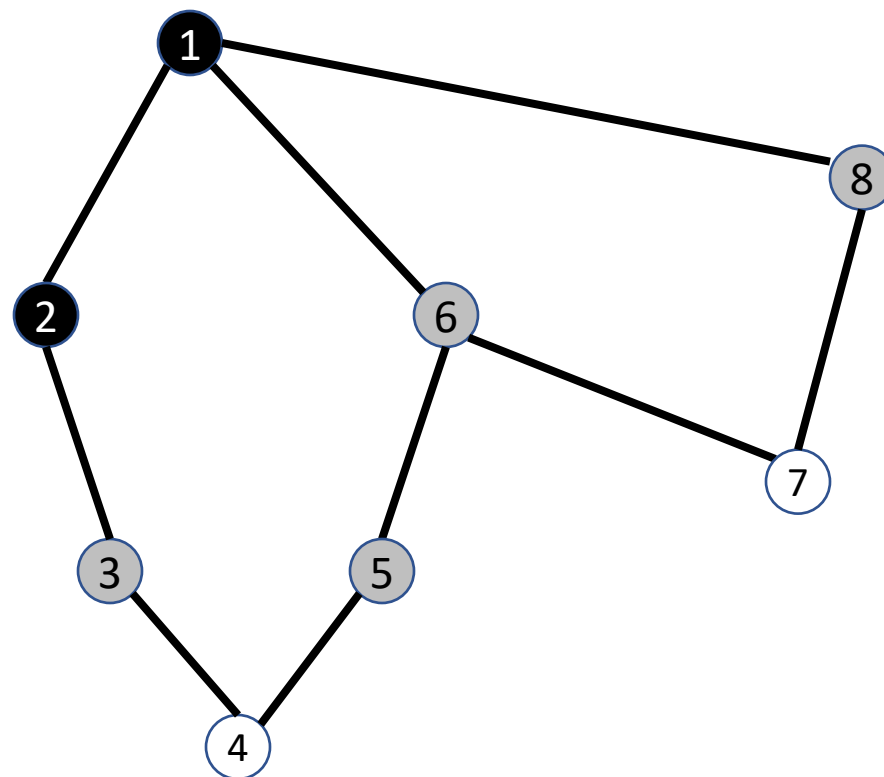
Queue: 2 6 8 3

Exercise: Finding the Smallest Cycle involving Vertex 1



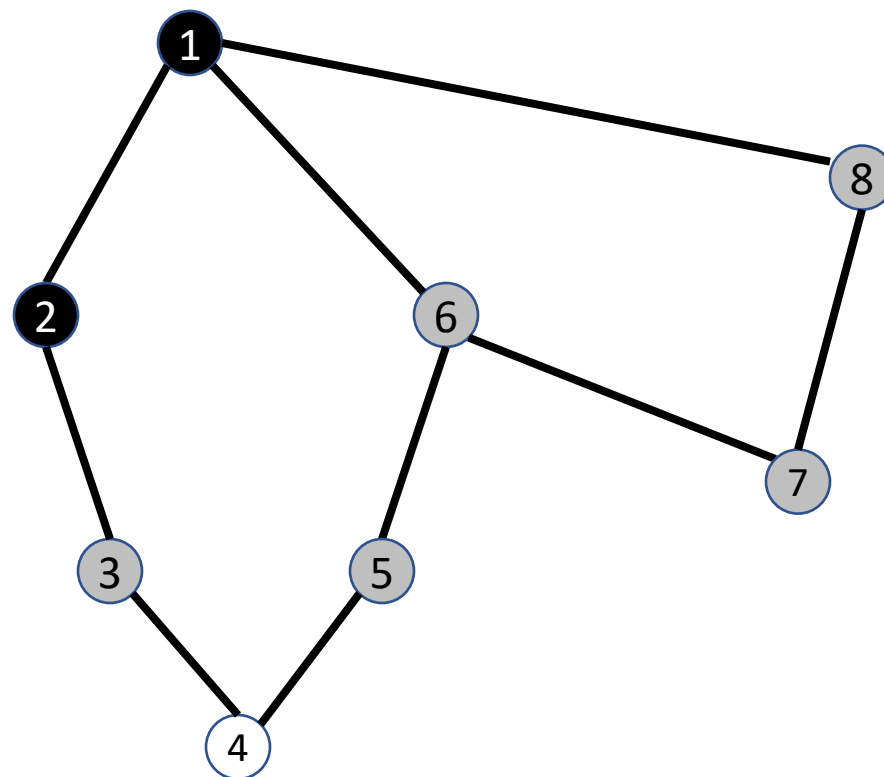
Queue: 6 8 3

Exercise: Finding the Smallest Cycle involving Vertex 1



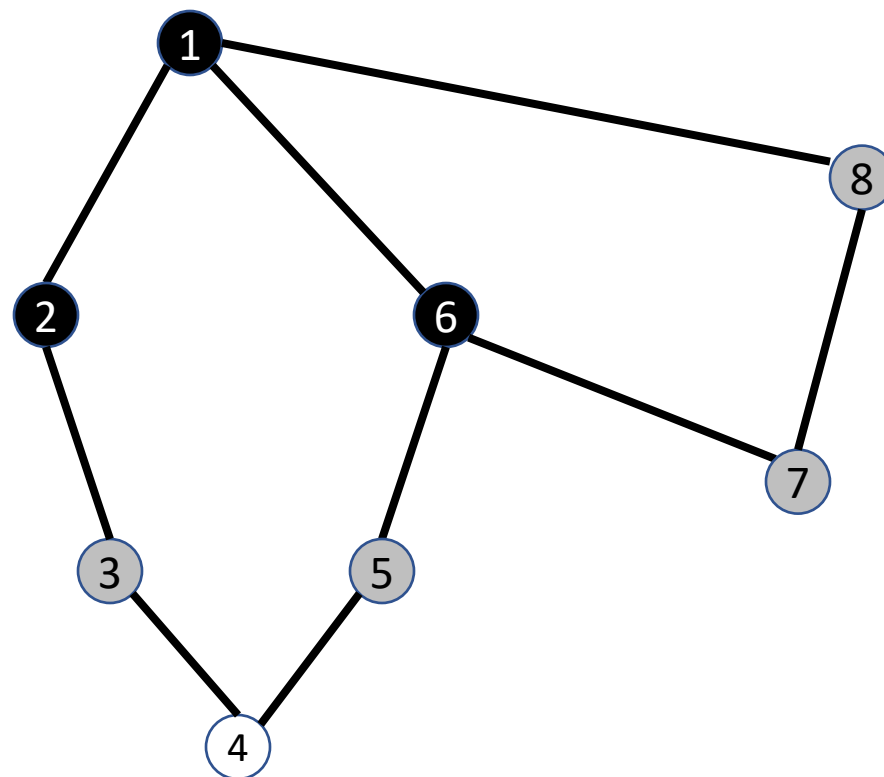
Queue: 6 8 3 5

Exercise: Finding the Smallest Cycle involving Vertex 1



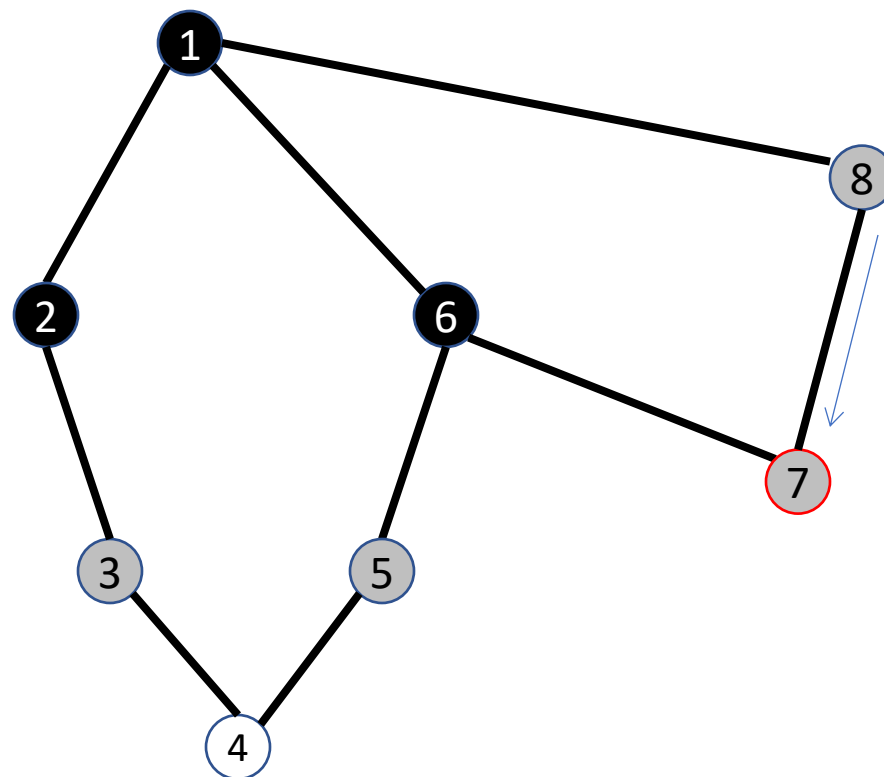
Queue: 6 8 3 5 7

Exercise: Finding the Smallest Cycle involving Vertex 1



Queue: 8 3 5 7

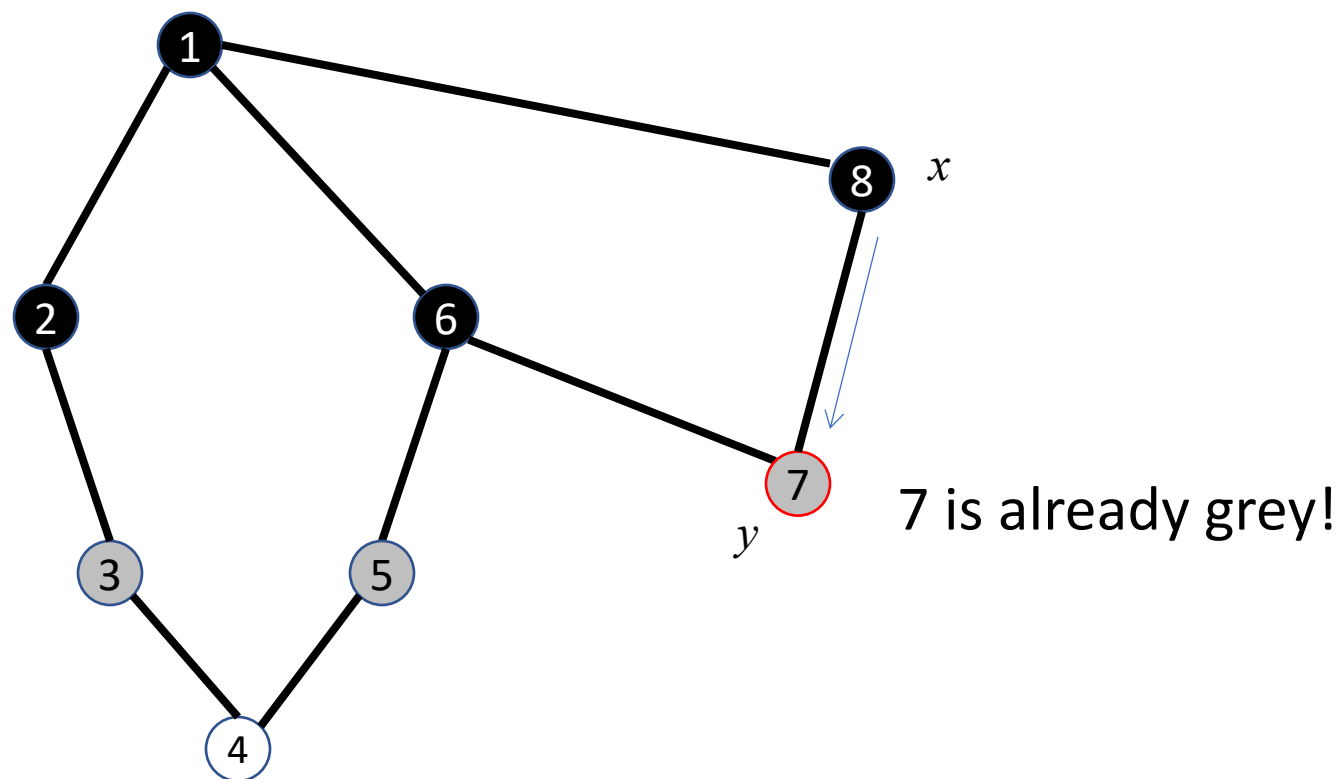
Exercise: Finding the Smallest Cycle involving Vertex 1



7 is already grey!

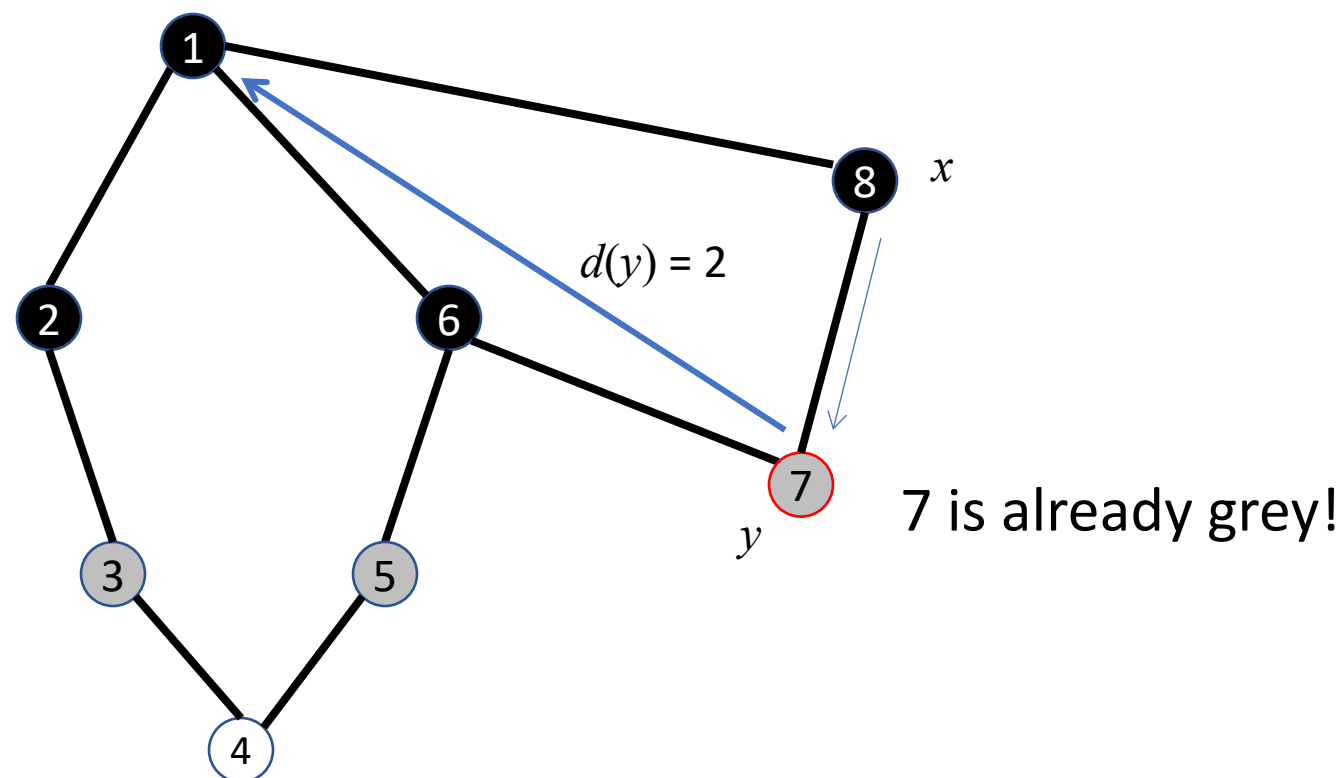
Queue: 8 3 5 7

Exercise: Finding the Smallest Cycle involving Vertex 1



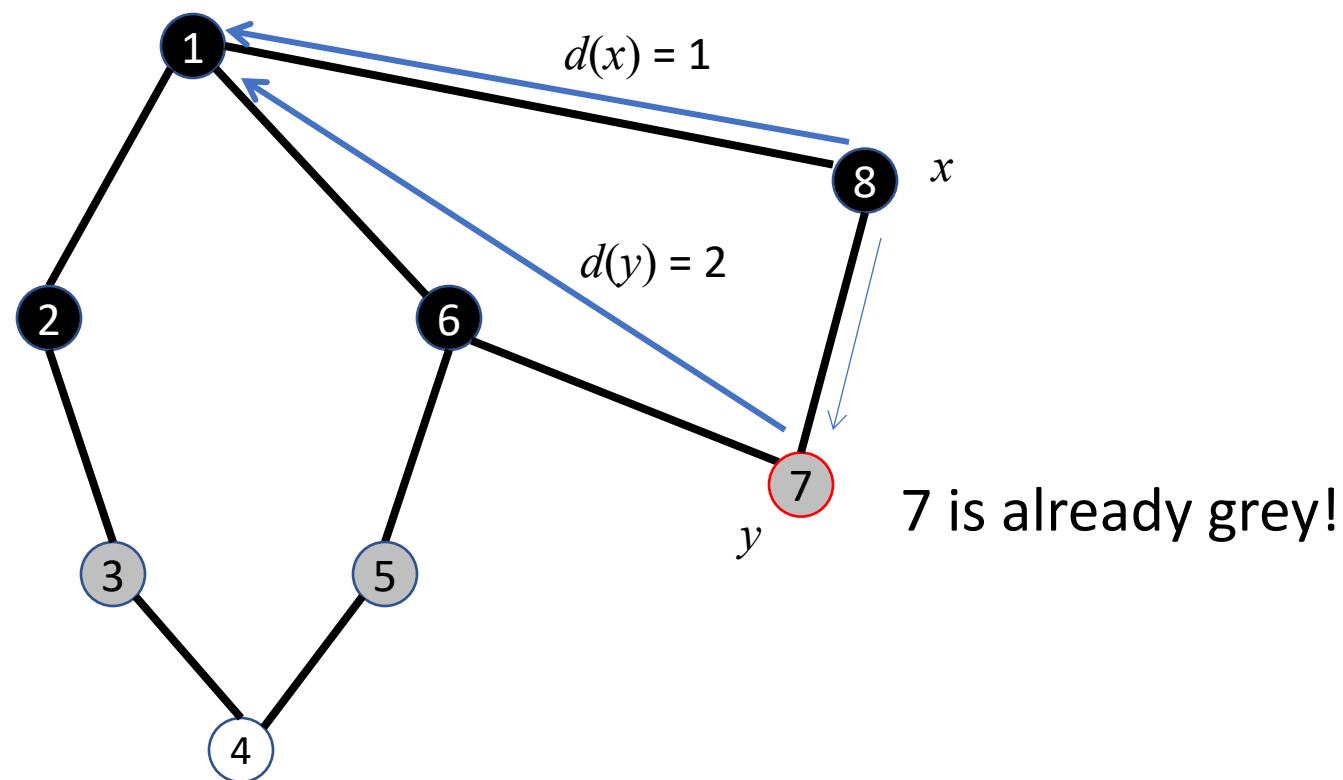
Queue: 3 5 7 (stopped)

Exercise: Finding the Smallest Cycle involving Vertex 1



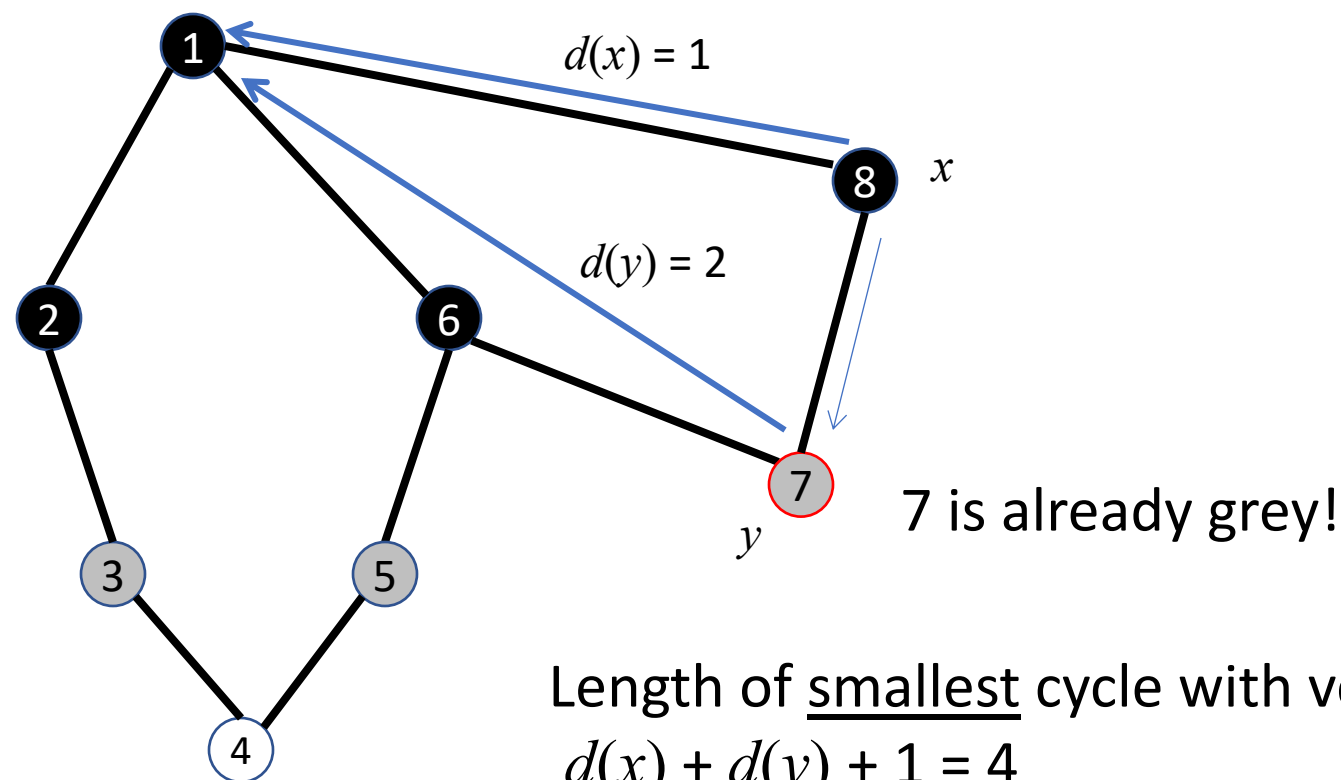
Queue: 3 5 7 (stopped)

Exercise: Finding the Smallest Cycle involving Vertex 1



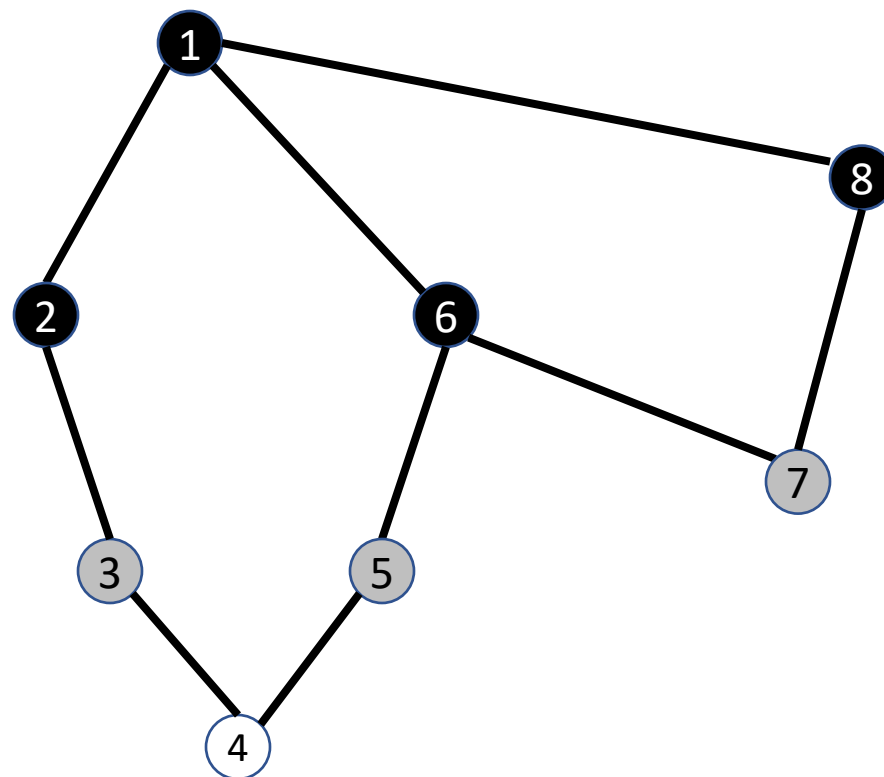
Queue: 3 5 7 (stopped)

Exercise: Finding the Smallest Cycle involving Vertex 1



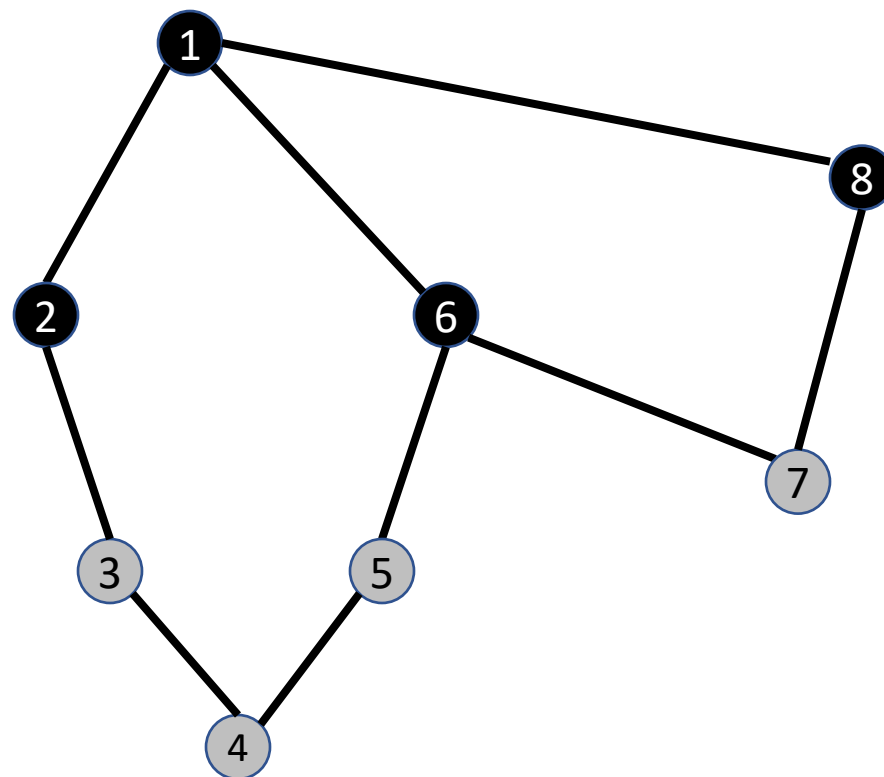
Queue: 3 5 7 (stopped)

Exercise: Finding the Smallest Cycle involving Vertex 1



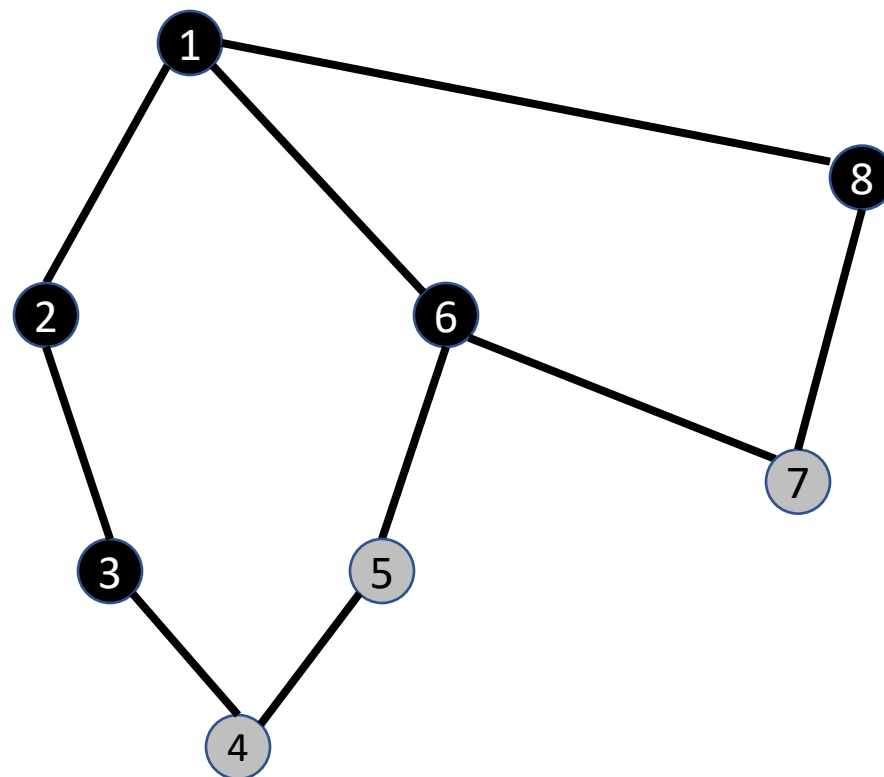
Queue: 3 5 7

Exercise: Finding the Smallest Cycle involving Vertex 1



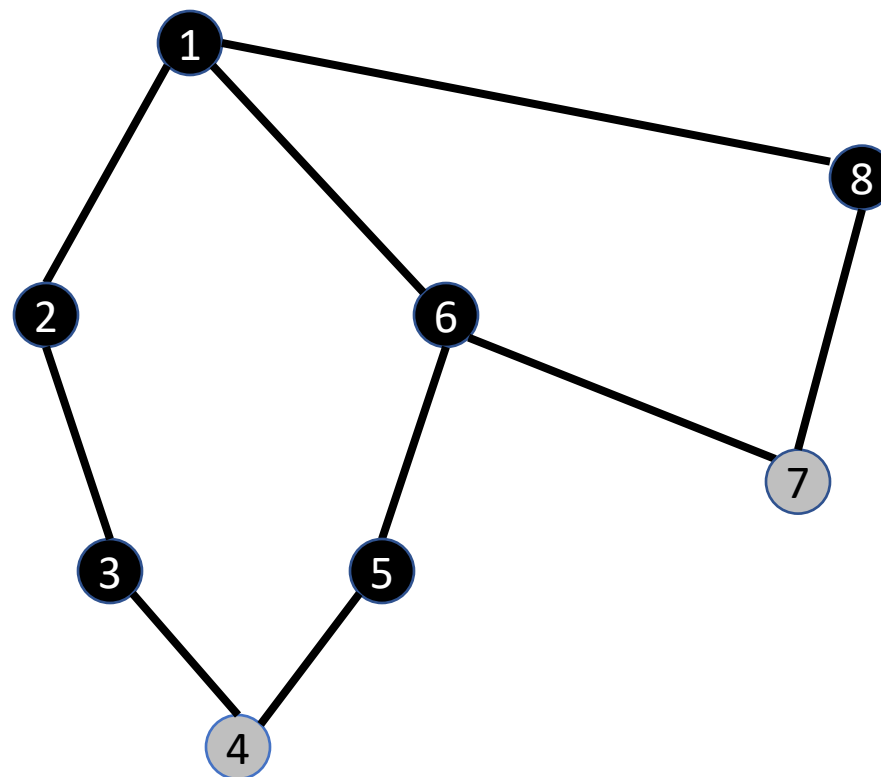
Queue: 3 5 7 4

Exercise: Finding the Smallest Cycle involving Vertex 1



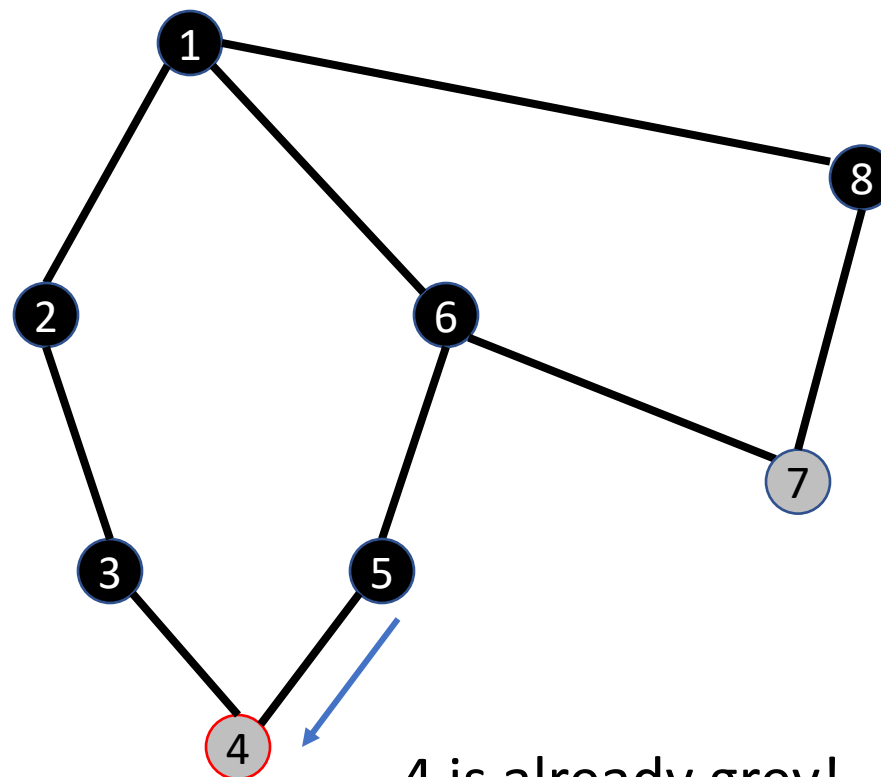
Queue: 5 7 4

Exercise: Finding the Smallest Cycle involving Vertex 1



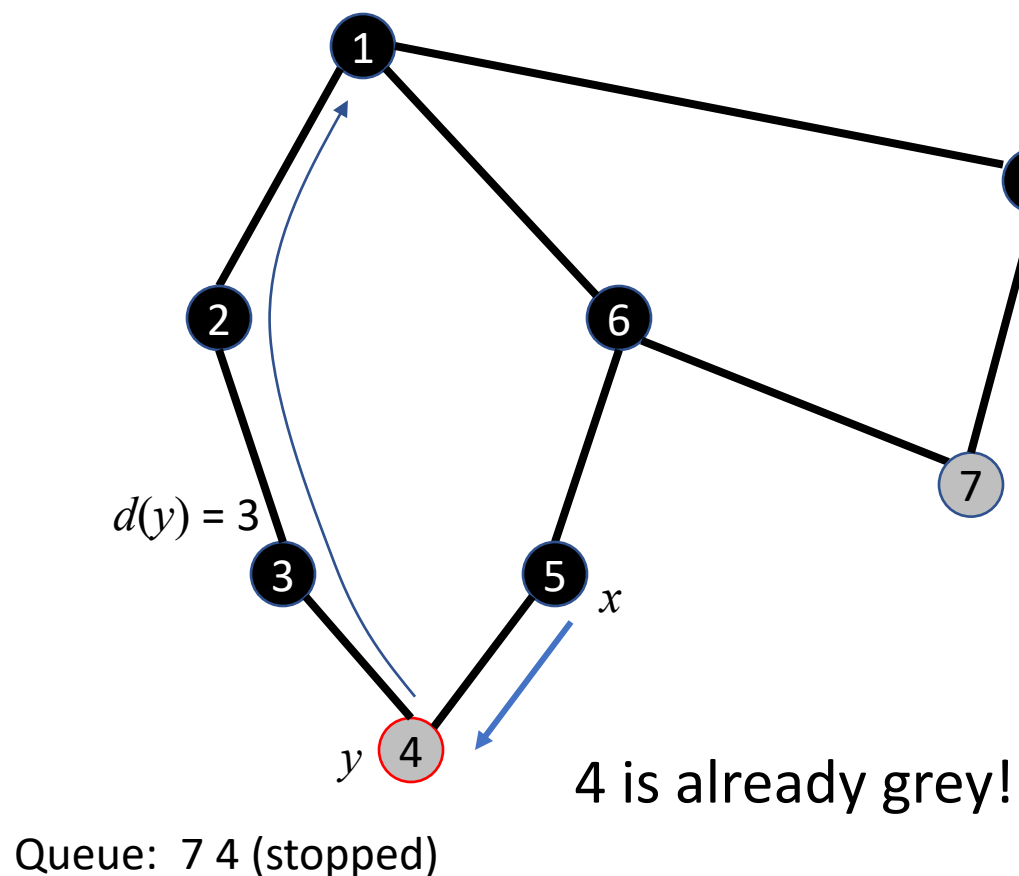
Queue: 7 4

Exercise: Finding the Smallest Cycle involving Vertex 1

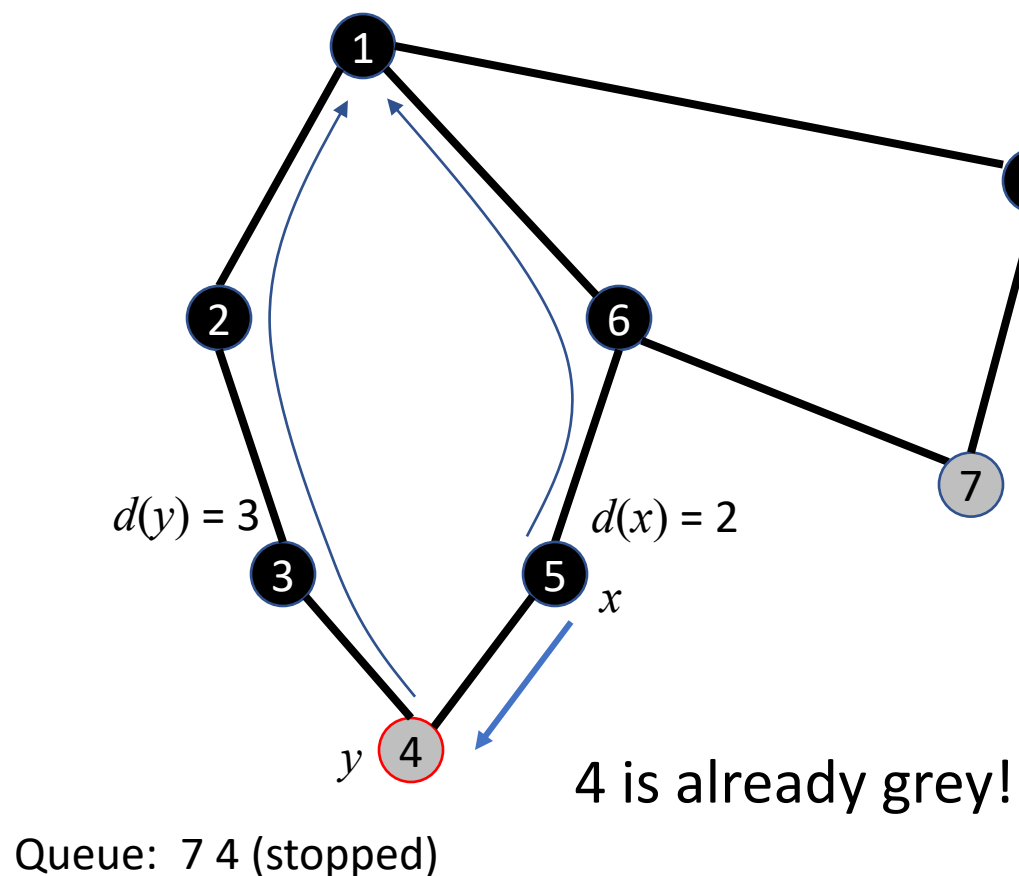


Queue: 7 4 (stopped)

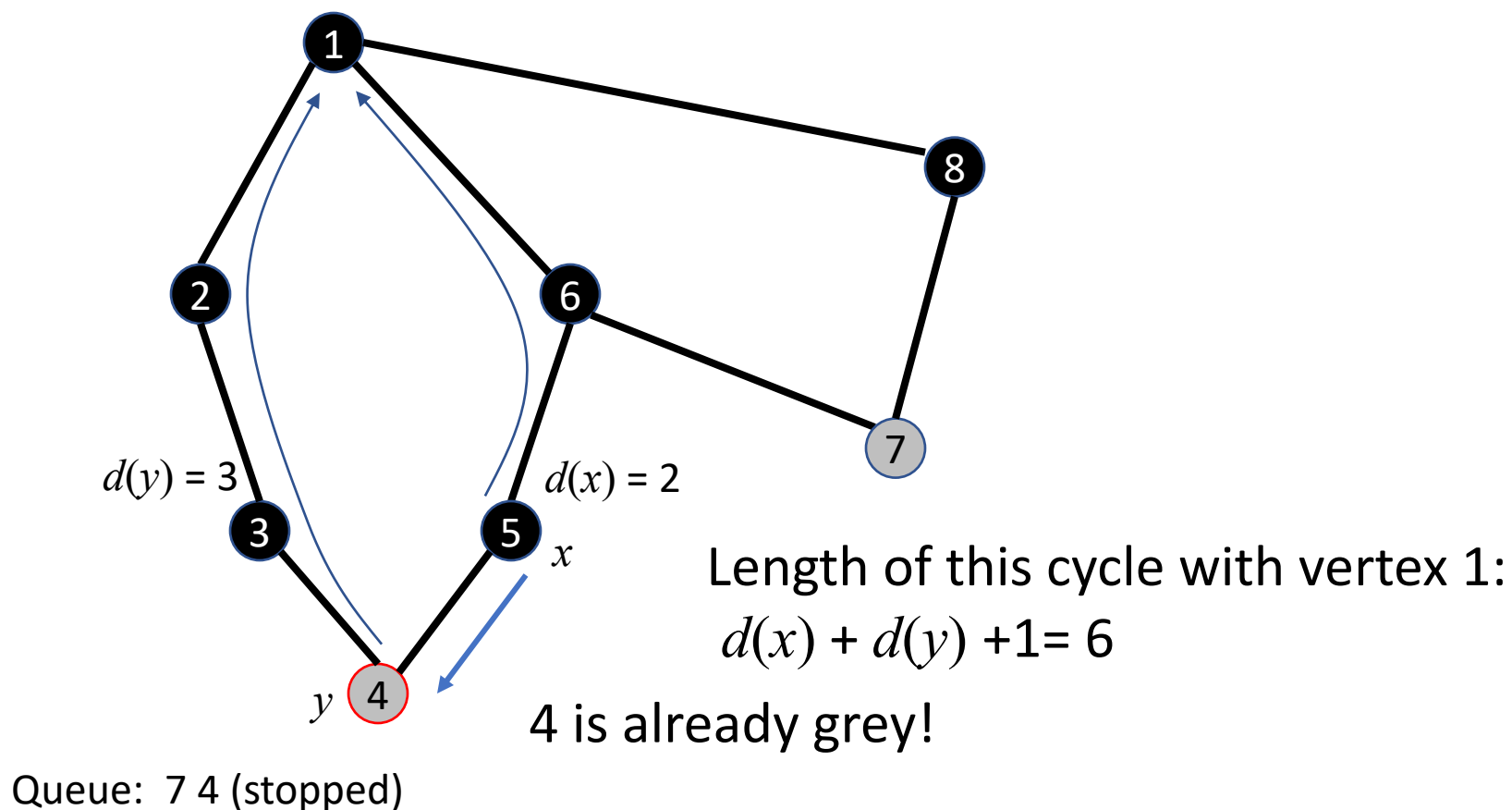
Exercise: Finding the Smallest Cycle involving Vertex 1



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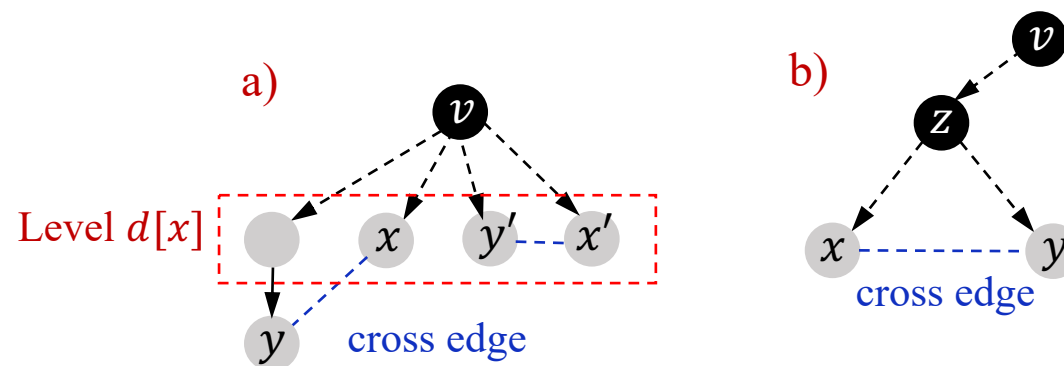


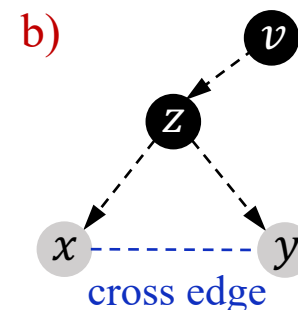
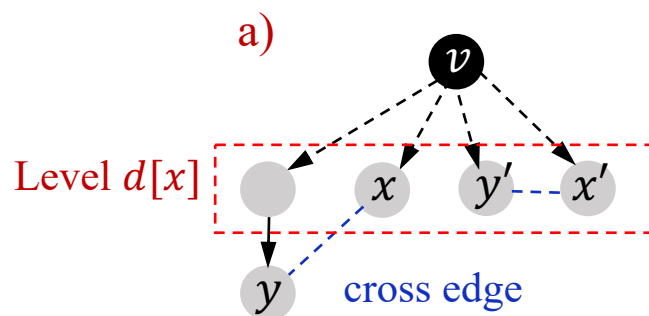
Correctness of the Algorithm

1. Meeting a grey neighbour means that we have found a cycle;
2. The cycle found in this way will be the shortest.

Running $\text{BFSVisit}(v)$ above, we have two cases when we first meet a cross edge (x, y) and let $d[x] \leq d[y]$:

- a) x and y are in different subtrees of v
- b) x and y are in the same subtree of v





Running $\text{BFSVisit}(v)$ above gives us two cases when we first meet a cross edge (x, y) and let $d[x] \leq d[y]$:

- a) x and y are in different subtrees of v : v, x, y are in the same cycle with length $d[x] + d[y] + 1 \leq 2d[x] + 2$. All cross edges detected at the same level have the same upper bound of cycle length. Then, the algorithm finds the smallest cycle containing v .
- b) x and y are in the same subtree of v : a smaller cycle (not including v) exists. Let z be the least common ancestor of x and y . When BFSVisit starts from z , we should find the smallest cycle containing z based on a).

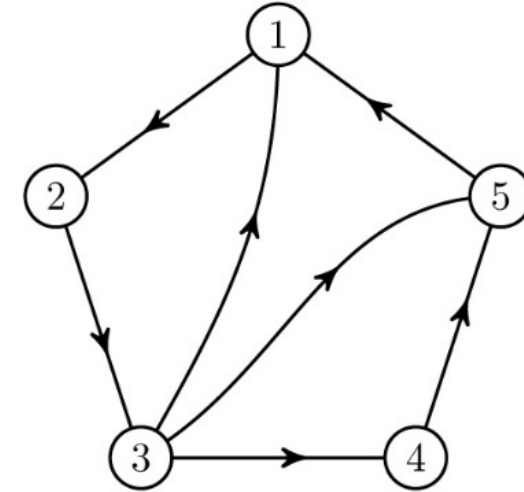
Finding Directed Girth of Digraph

Following procedure finds the (length of the) shortest directed cycle through node v in a digraph.

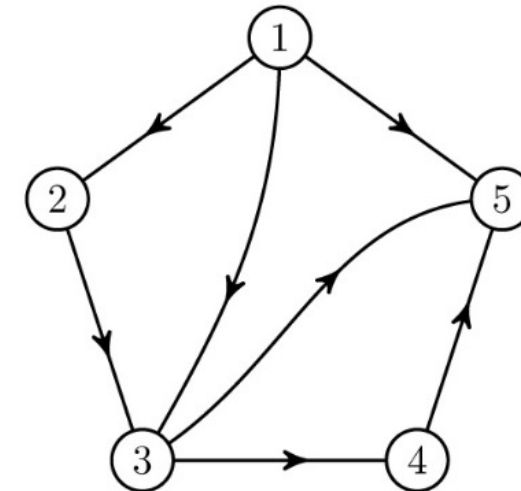
1. Run BFSVisit(v) starting at node v .
 2. The first time a **back-arc** of the form (x, v) is found, we have found a cycle of length equal to $(1 + \text{the depth of } x \text{ in the search tree})$.
 3. Stop and Return the length found.
- Run the above procedure on every node and pick the length of the shortest cycle

SUMMARY

- Terminology
 - Cycle
 - Girth
- DFS Application: Cycle Detection
- BFS Application: Girth Detection
- Correctness of Girth Computation
- Illustrative Examples



(a) A directed cyclic graph.



(b) A directed acyclic graph.