

# Cycles and Girth

Instructor: Meng-Fen Chiang

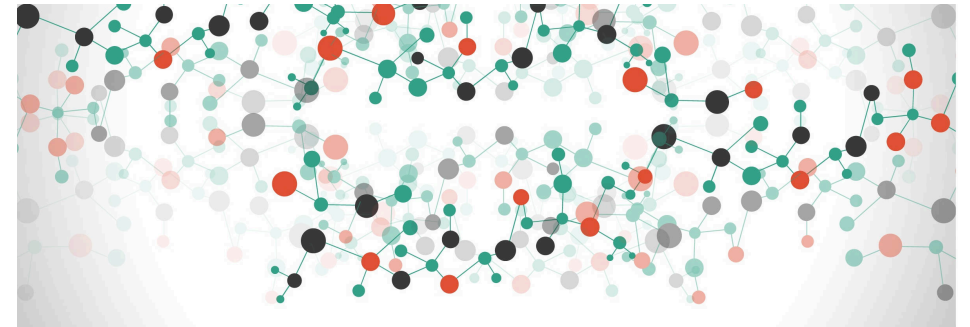
COMPCSI220: WEEK 10



Slides adapted from Mark Wilson, Georgy Gimel'farb, Simone Linz and Tanya Gvozdeva

# OUTLINE

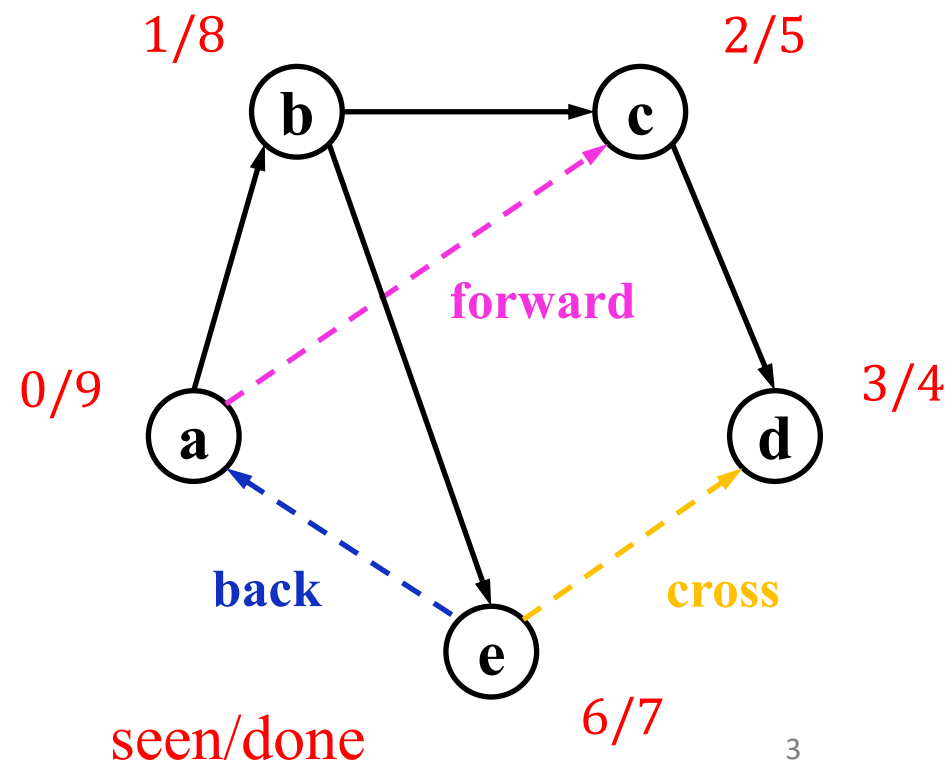
- Terminology
  - Cycle
  - Girth
- DFS Application: Cycle Detection
- BFS Application: Girth Detection
- Correctness of Girth Computation



# RECAP: Example 23.7

- Explain how to determine  $(u,v)$  in DFS algorithm whether it is a tree-, back-, forward- or cross-arc?

- If  $v$  is **white**, then  $(u,v)$  is a **tree** arc
- If  $v$  is **grey**, then  $(u,v)$  is a **back** arc
- If  $v$  is **black**, then  $(u,v)$  is
  - a **cross** arc( $seen[v] < seen[u]$ ,  $done[v] < seen[u]$ ), or
  - a **forward** arc( $seen[u] < seen[v] < done[v] < done[u]$ ).

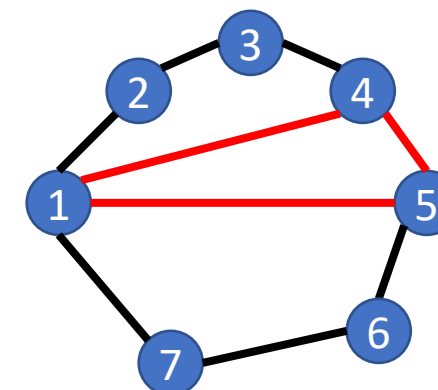
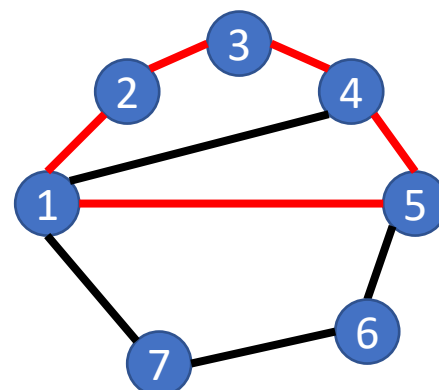
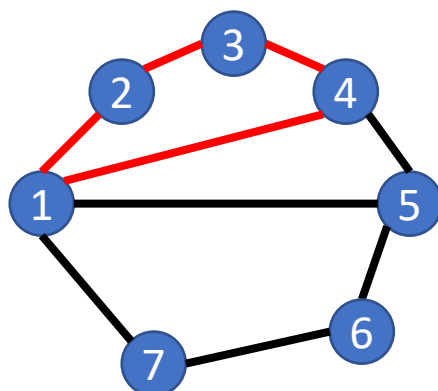
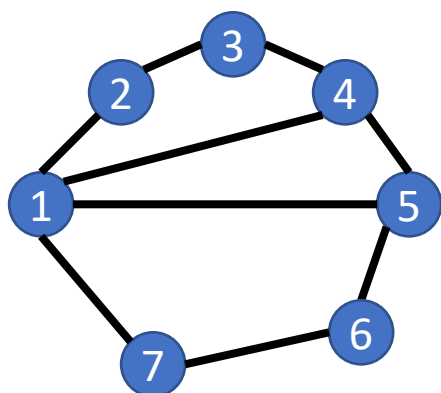


# Cycle Detection

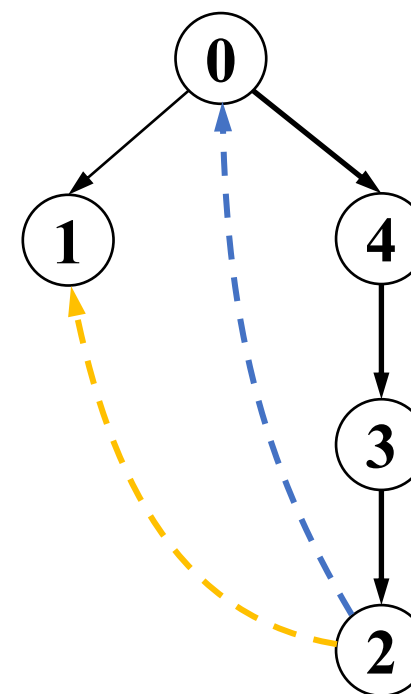
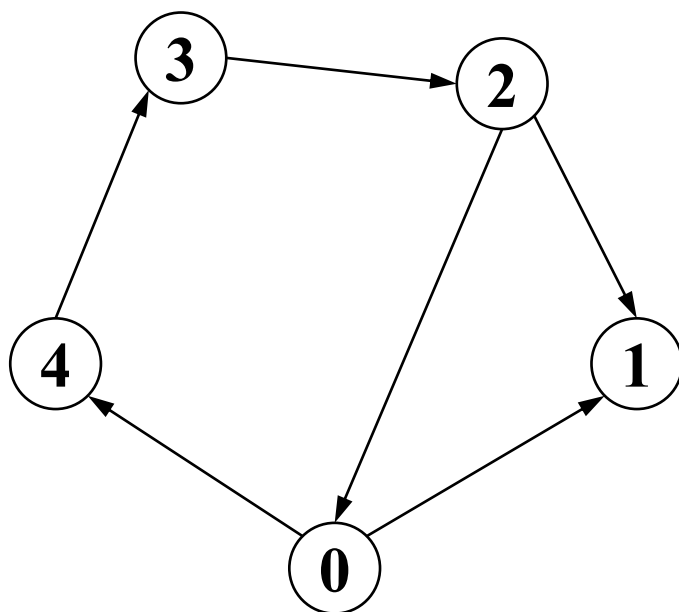
- Suppose that there is a cycle in  $G$  and let  $v$  be the node in the cycle visited first by **DFS**. If  $(u, v)$  is an arc in the cycle then it must be a back arc.
- Conversely if there is a **back arc**, we must have a **cycle**.
- Suppose that DFS is run on a digraph  $G$ . Then  $G$  is acyclic if and only if  $G$  does not contain a back arc.

# Example: Cycles in Graphs

- A graph can have several cycles



# Using DFS to Find Cycles in Digraphs



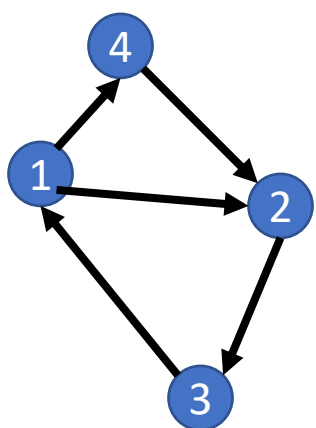
Back arc detected!

Once DFS finds a cycle, the stack contains the nodes that form the cycle.

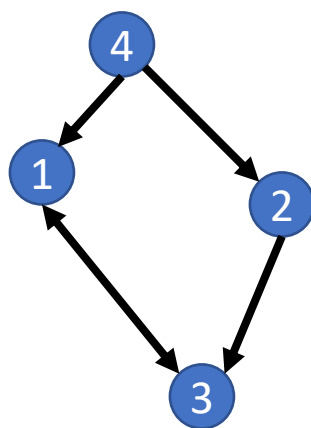
# Girth and Digirth

- **Definition.** For a graph (with a cycle), the length of the shortest cycle is called the **girth** of the graph. If the graph has no cycle then the girth is undefined but may be viewed as  $+\infty$ .
- **Convention.** For a digraph we use the term **girth** for its **underlying graph** and the term **directed girth** for the length of the smallest directed cycle.

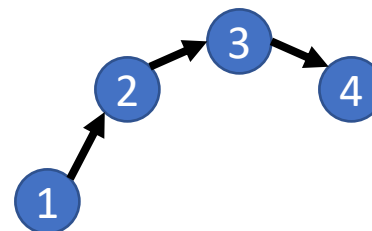
# Examples: Girth of a Graph (Digraph)



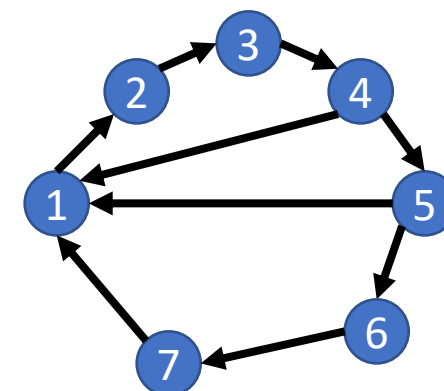
Girth: 3  
Directed girth: 3



Girth: 4  
Directed girth: 2



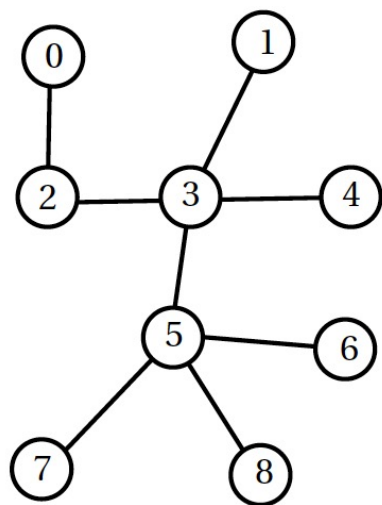
Girth:  $+\infty$   
Directed girth:  $+\infty$



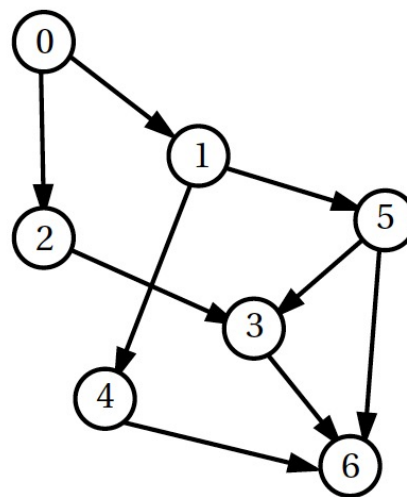
Girth: 3  
Directed girth: 4



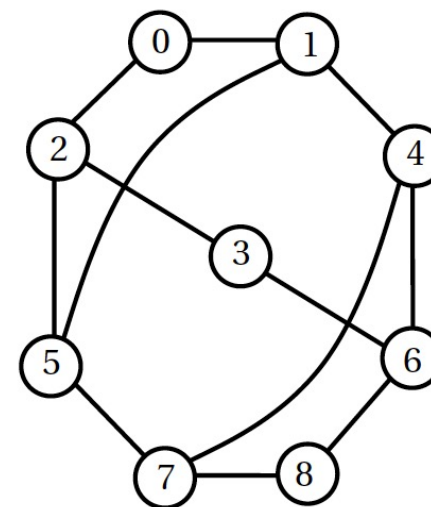
# Examples: Girth of a Graph (Digraph)



Girth:  $+\infty$   
Directed girth: undefined



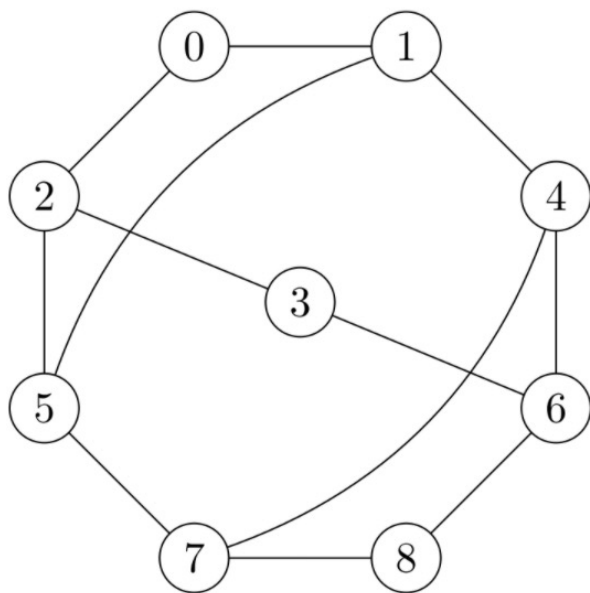
Girth: 3  
Directed girth:  $+\infty$



Girth: 4  
Directed girth: undefined

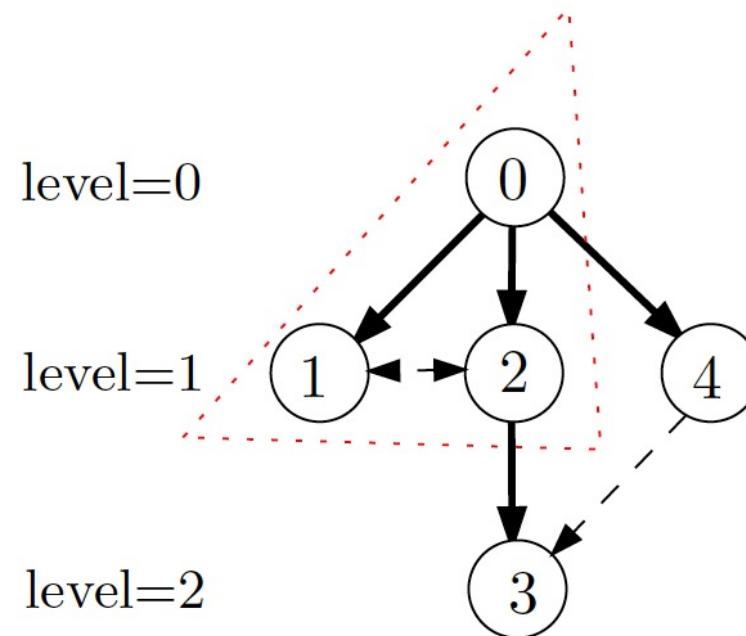
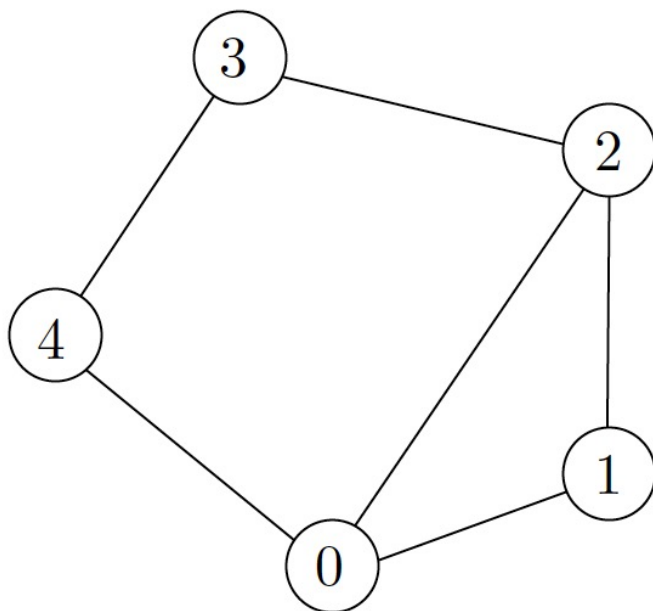
# Finding the Girth of a Graph

- An easy-to-implement **DFS** idea may not work properly. Consider the **DFS** tree originating from vertex 0 of the graph below. Which is the smallest cycle it finds with the back edges? Which smaller cycle does it miss?



# Finding the Girth of a Graph

- **Using BFS to find cycles in graphs.** Cycles can also be easily detected in a graph using BFS. Finding a cycle of **minimum** length in a graph is not difficult using BFS (better than DFS).

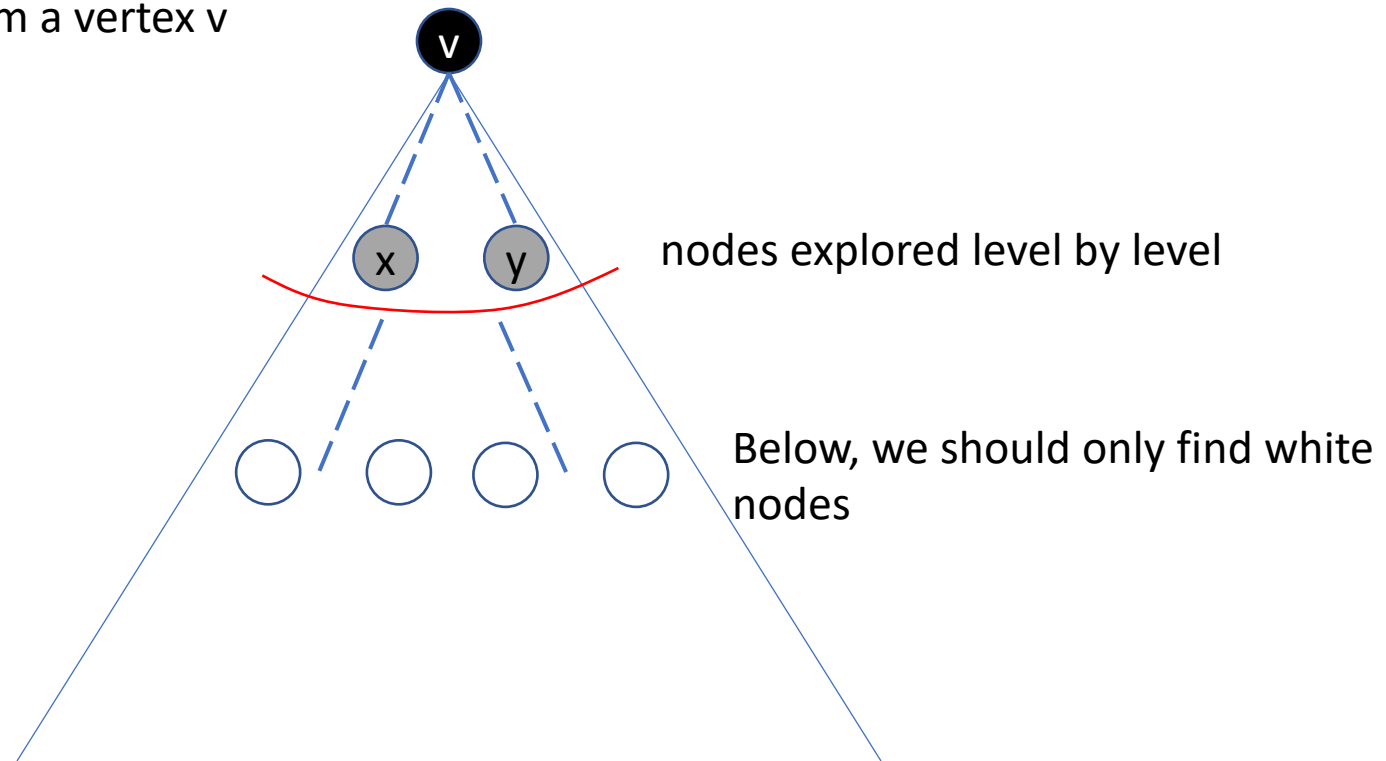


# Finding the Girth of a Graph

- Perform `BFSVisit()`  $|V|$  times, starting at each vertex  $v' \in V$  in turn.
- If during a `BFSVisit`, we encounter a **grey neighbour** (rather than a white) we have found a cycle
  - The grey node was visited before following another path
- An important property of BFS is that if it runs from vertex  $v$ , then every vertex  $s$ , when it is first reached, then the path that was found from  $v$  to  $s$  is **minimal**. Thus, reaching  $v'$  from  $v$  with BFS finds the shortest path from  $v$  to  $v'$ , namely the **shortest cycle** that contains  $v$ .

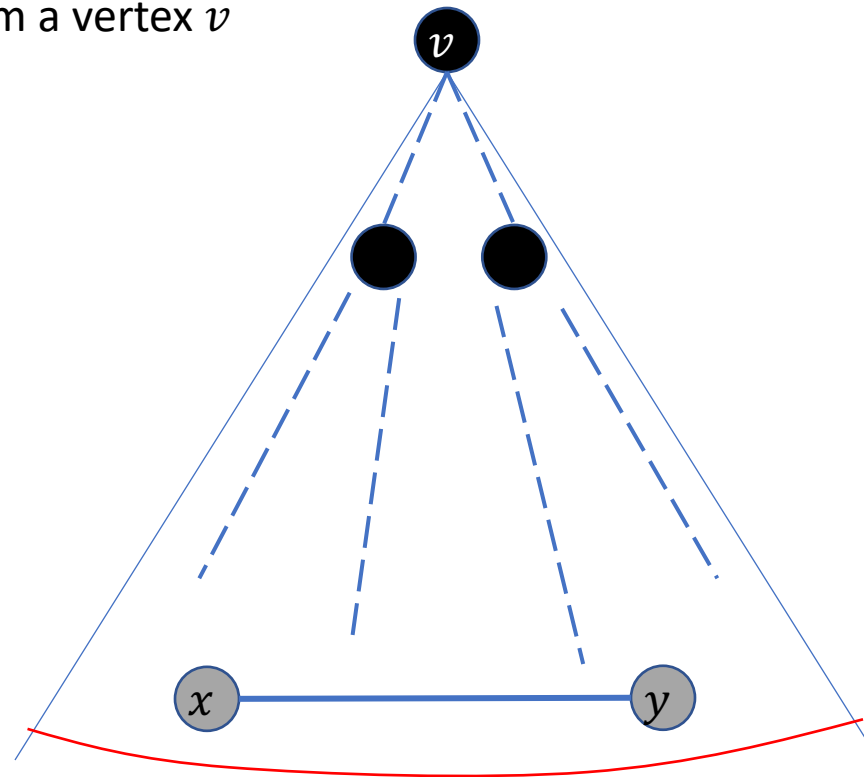
# Example: Girth of a Graph

Executing BFSVisit from a vertex  $v$



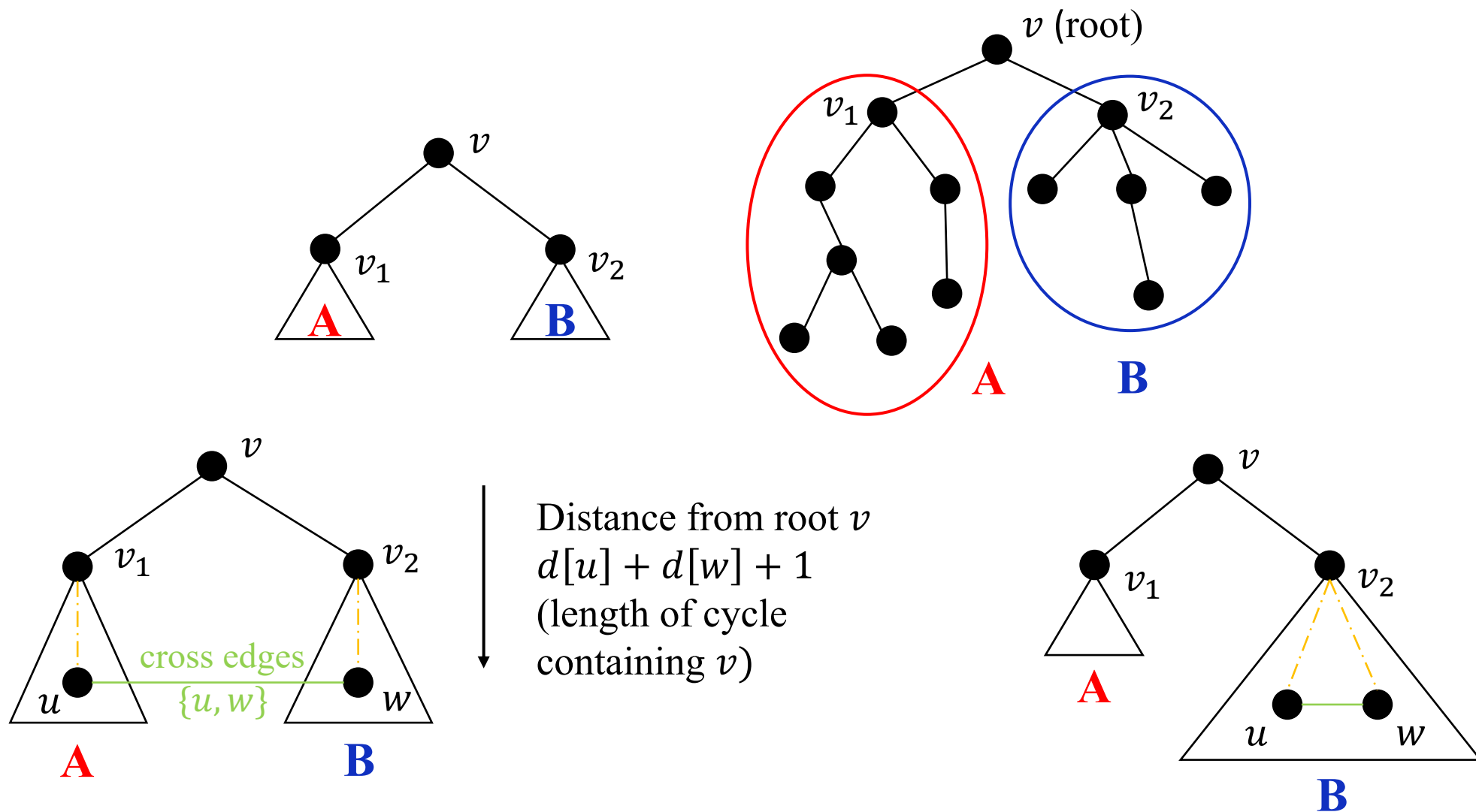
# Example: Girth of a Graph

Executing BFSVisit from a vertex  $v$



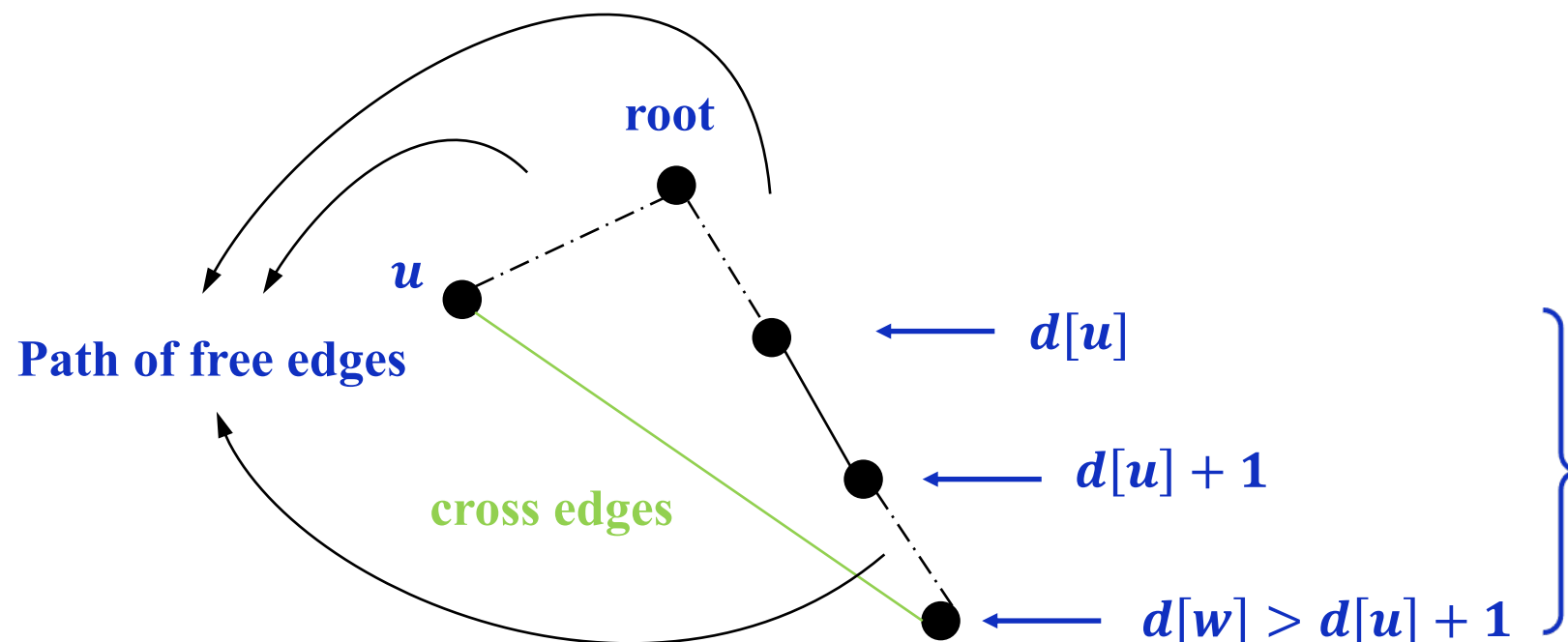
If a **cross** edge/arc exits then it will connect to a grey vertex  
Both  $x$  and  $y$  are reachable from  $v$  therefore there is a cycle through the **cross** edge

# Example: Girth of a Graph



# BFS: Cross Edges in a Graph

- Let  $\{u, w\}$  be a cross edge in a tree. Then either  $d[u] = d[w]$  or  $|d[u] - d[w]| = 1$
- Sketch of proof:** Suppose  $|d[u] - d[w]| > 1$



$d[w] > \text{distance from root to } w$   
(shortest path from root to  $w$   
contains the cross edge)

**\*contradiction**



# Facts about Cycle Length

- If vertex  $v$  is on at least one cycle then BFS starting at  $v$  will find it.
- On detection of a cross edge between descendants  $u$  and  $w$ , determine whether  $u$  and  $w$  are in **different subtrees below  $v$**  (the root of the tree).
  - If yes, then a cycle of length  $d[u]+d[w]+1$  is found.
  - If no, then a cycle of shorter length is found (but avoids  $v$ ).
- If  $d[u]=d[w]$  then **odd** length, where  $v$  is a common ancestor.
- Otherwise,  $d[u]+1=d[w]$  and **even** length.

# Facts about Cycle Length

- cross edge  $\{u, w\}$

1.  $d[u] = d[w]$

cycle length:  $2d[u] + 1 \Rightarrow$  **odd** length  
cross edge

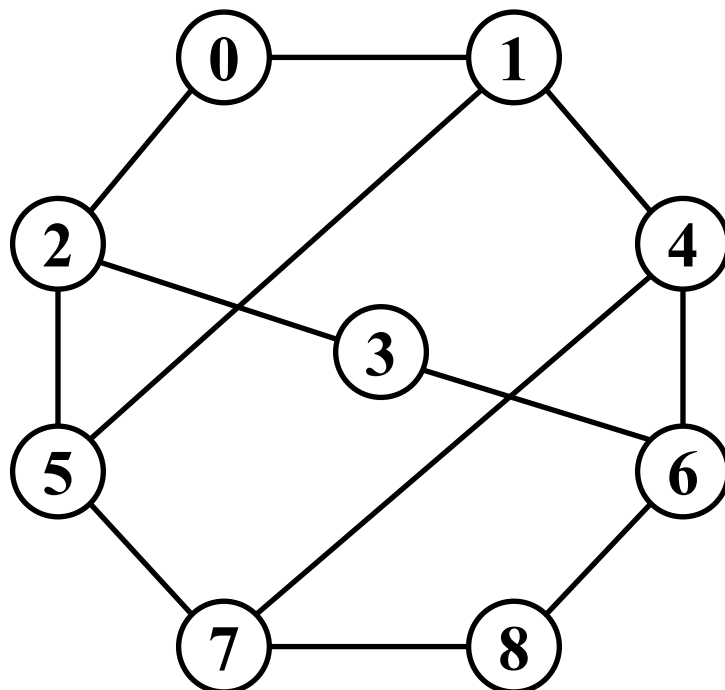
2.  $d[u] + 1 = d[w]$

cycle length:  $d[u] + d[w] + 1$   
 $= d[u] + d[u] + 1 + 1$   
 $= 2d[u] + 2 \Rightarrow$  **even** length

# Finding the Girth of a Graph

- To compute girth, we perform  $\text{BFSVisit}(v)$  procedure once for each  $v \in V(G)$  and take **minimum**.
  - If a grey neighbour is met, e.g., an edge  $(x, y)$  is explored from  $x$  where  $y$  is grey, continue to the end of the current level and then stop.
  - For each edge  $(x, y)$  as above on this level, if  $v$  is the lowest common ancestor of  $x$  and  $y$  in the BFS tree, then there is a cycle containing  $x, y, v$  of length  $l = d(x) + d(y) + 1$ .
  - Report the minimum value of  $l$  obtained along the current level.
- The minimum of these lengths at the level is the smallest cycle that **involves  $v$**
- The smallest cycle among all possible start vertices  $v$  is the girth.

# Example (1): Finding the Girth of a Graph

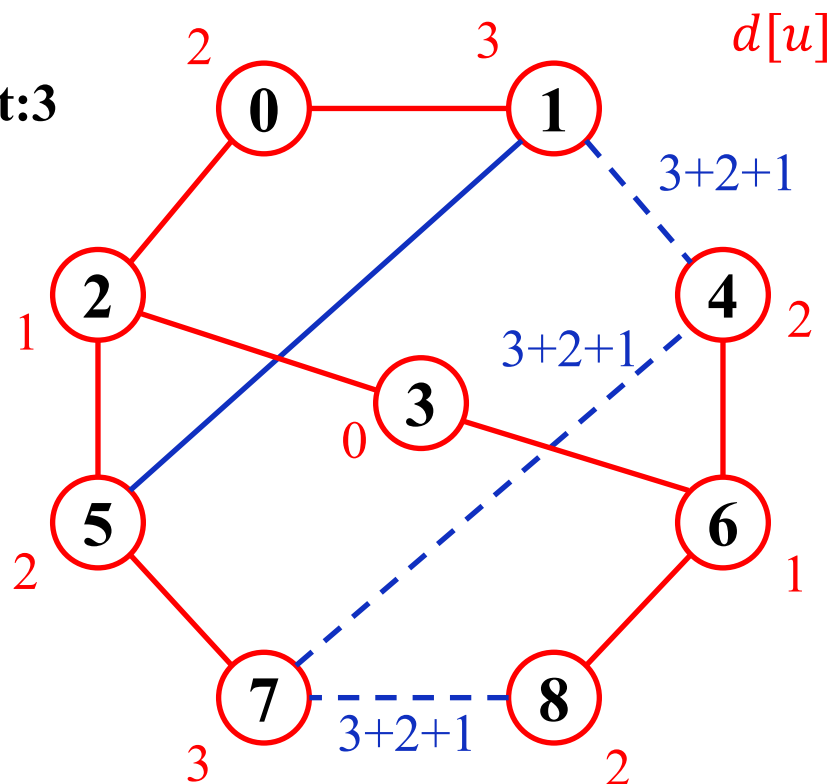


- Run BFS for 3, shortest cycle has length 6.
- Run BFS for any other vertex, shortest cycle has length 4.
- Girth of the graph is 4..

# Example (1)

**BFS**

**Start:3**



Shortest cycle containing **3** has length 6

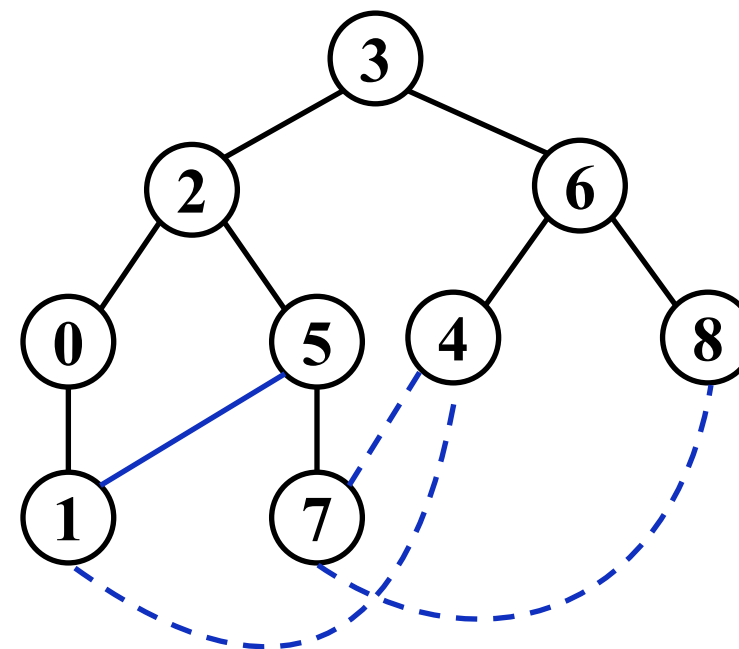
— tree edges  
— cross edge same subtree  
- - - tree edges different subtree  
(relative to root **3**)

$d = 0$

$d = 1$

$d = 2$

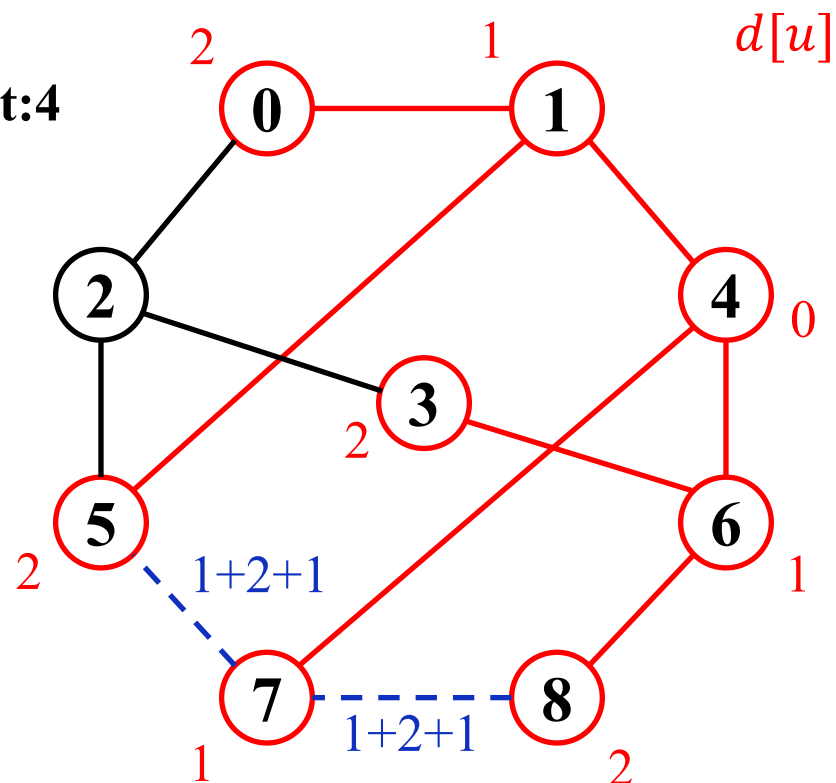
$d = 3$



# Example (1)

BFS

Start:4



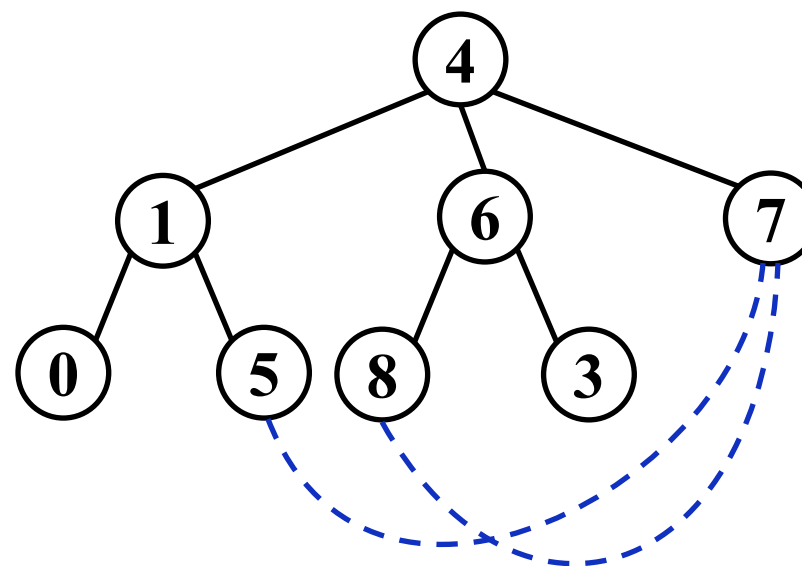
Shortest cycle containing **4** has length 4

- tree edges
- cross edge same subtree
- - - tree edges different subtree  
(relative to root **4**)

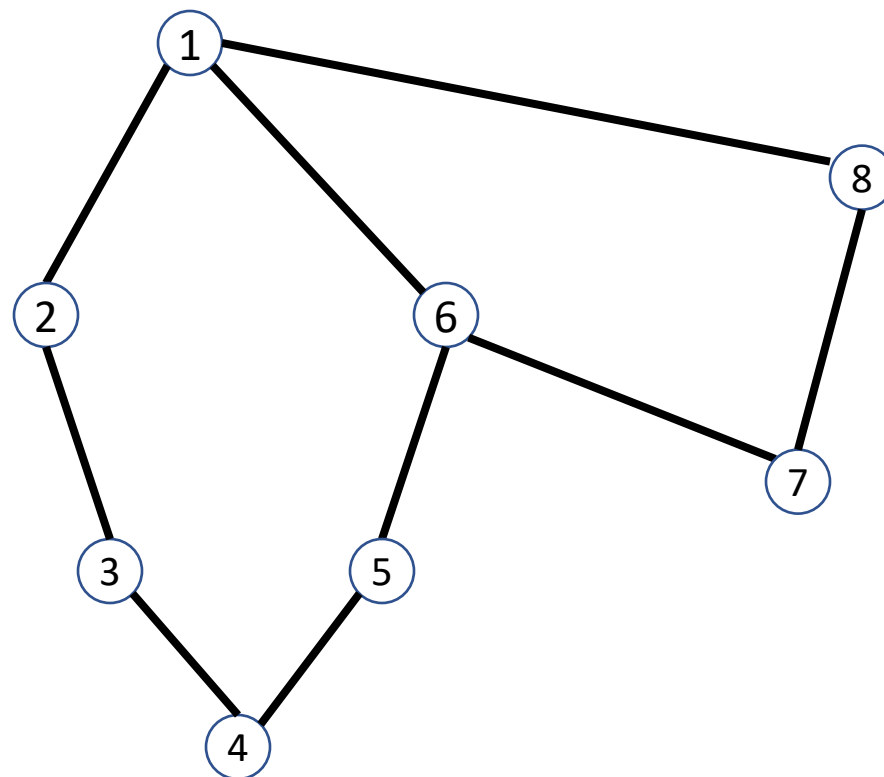
$d = 0$

$d = 1$

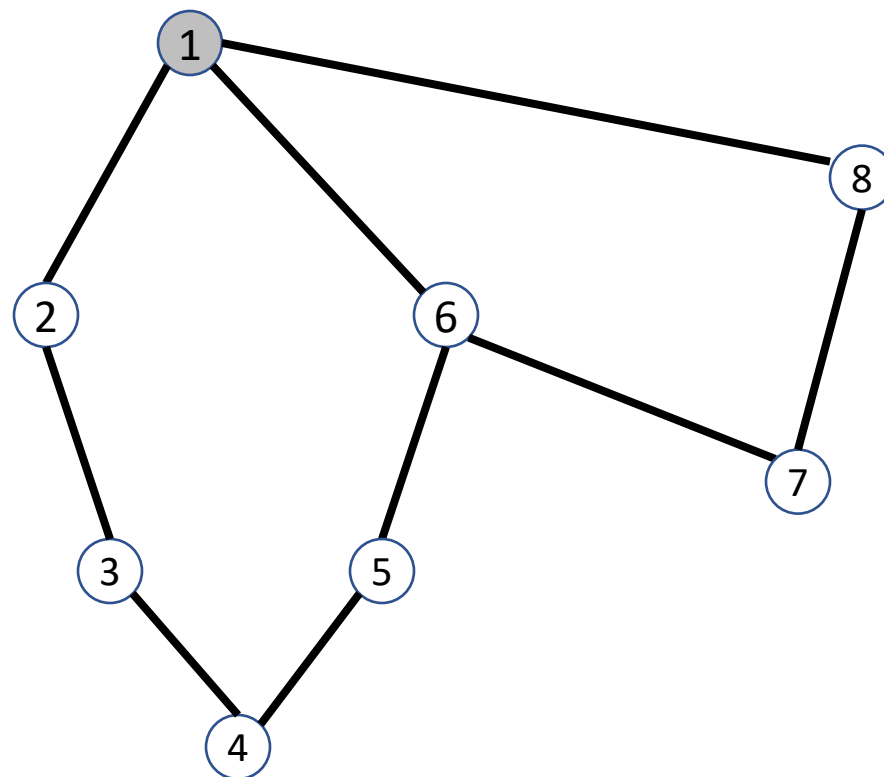
$d = 2$



# Exercise: Finding the Smallest Cycle involving Vertex 1



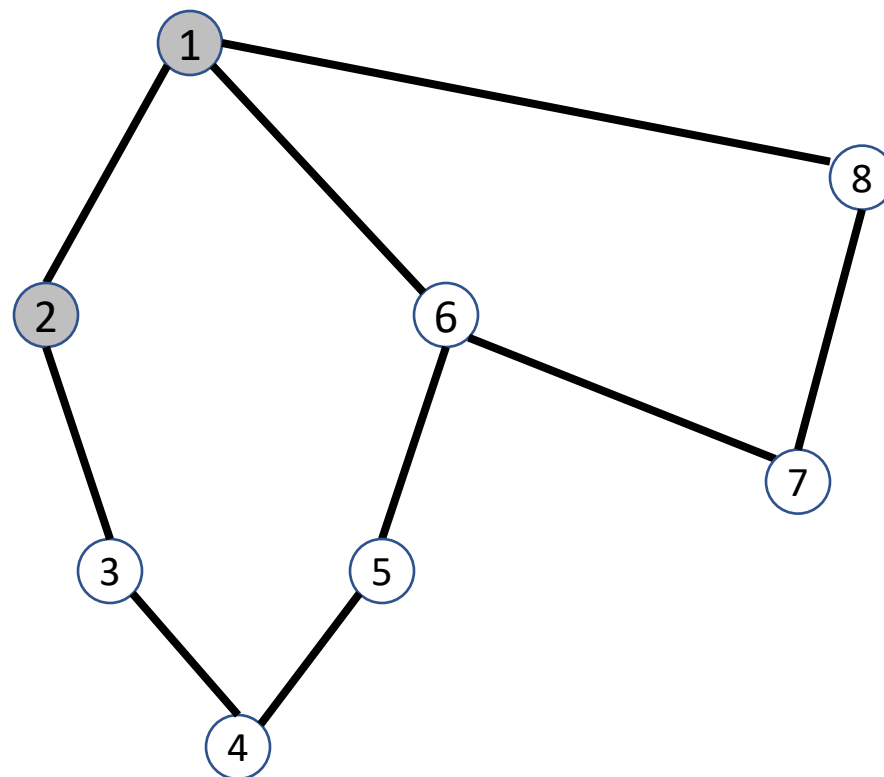
# Exercise: Finding the Smallest Cycle involving Vertex 1



Queue: 1

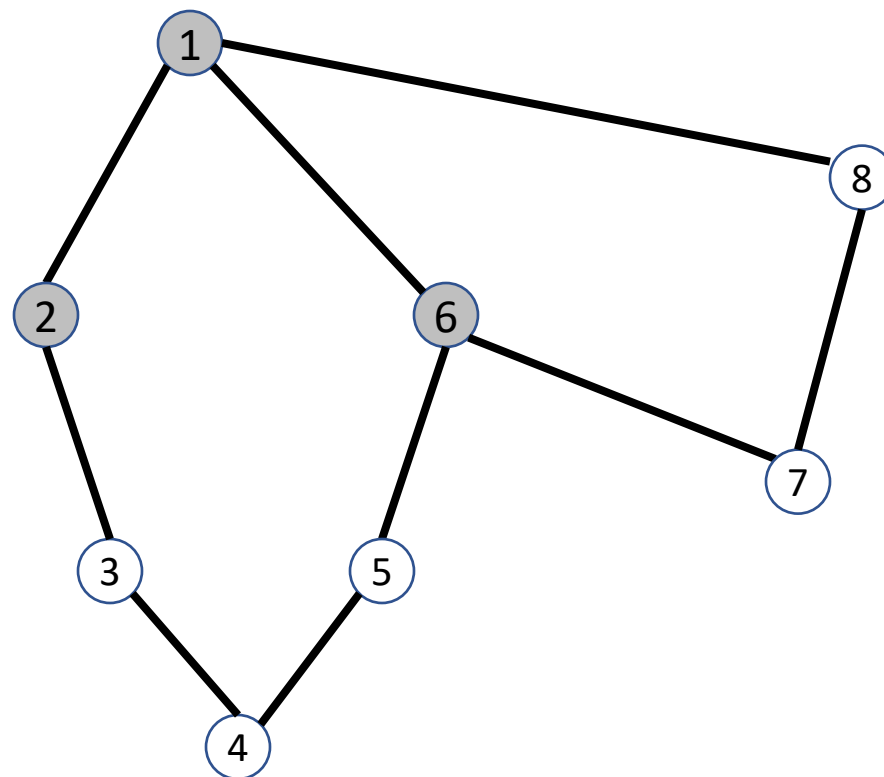


# Exercise : Finding the Smallest Cycle involving Vertex 1



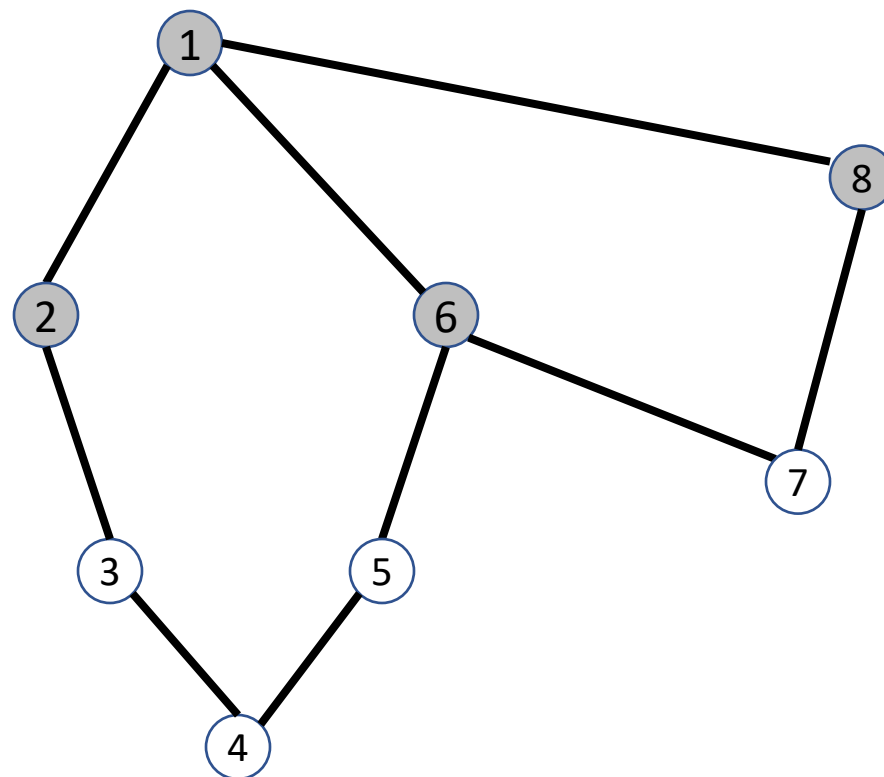
Queue: 1 2

# Exercise: Finding the Smallest Cycle involving Vertex 1



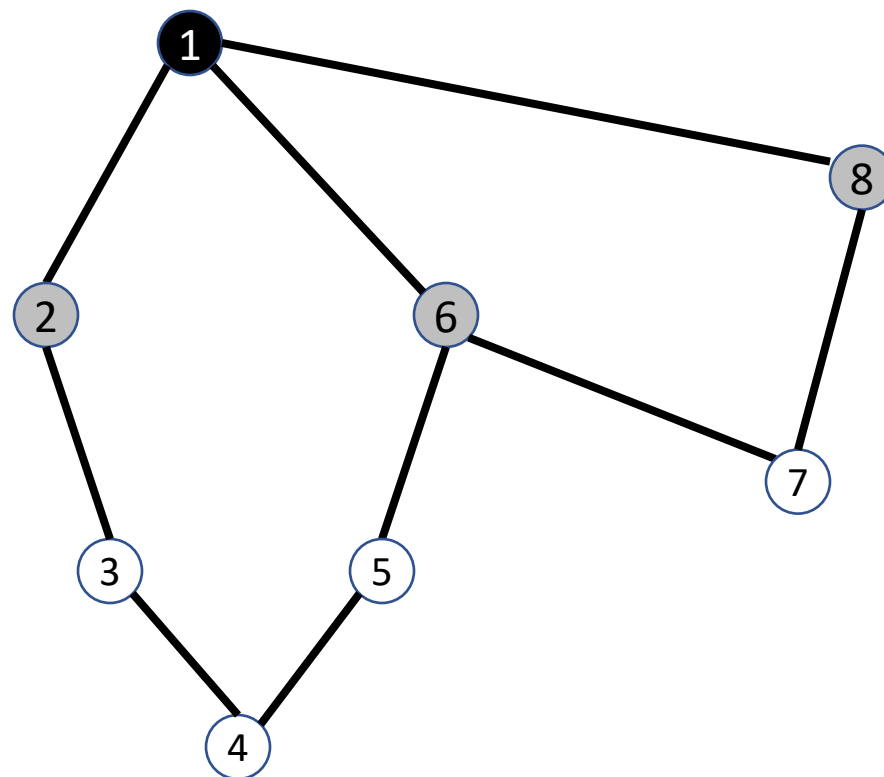
Queue: 1 2 6

# Exercise: Finding the Smallest Cycle involving Vertex 1



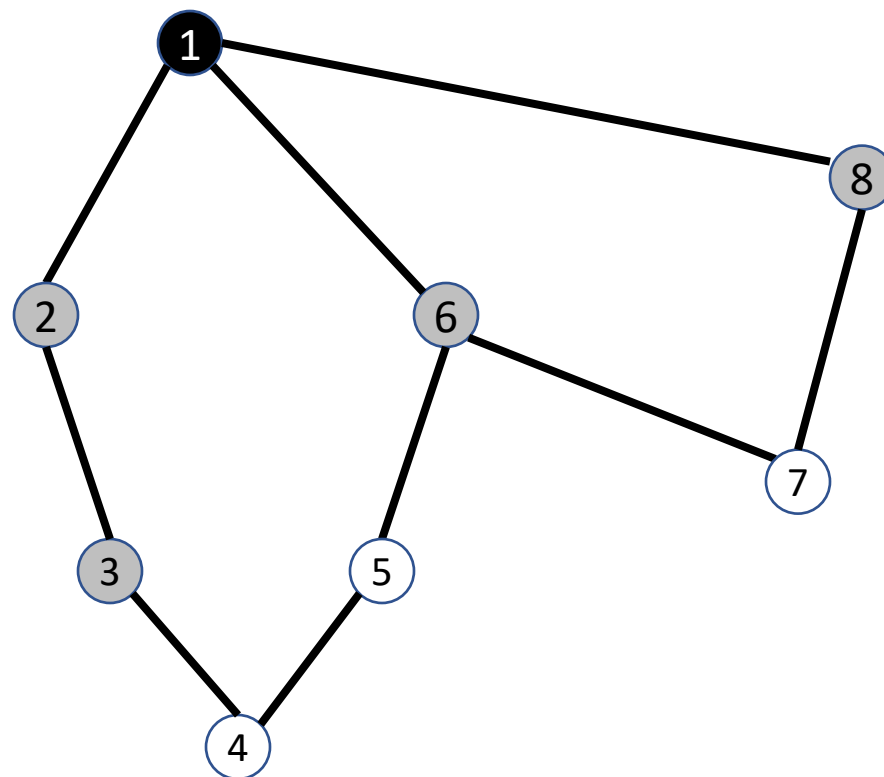
Queue: 1 2 6 8

# Exercise: Finding the Smallest Cycle involving Vertex 1



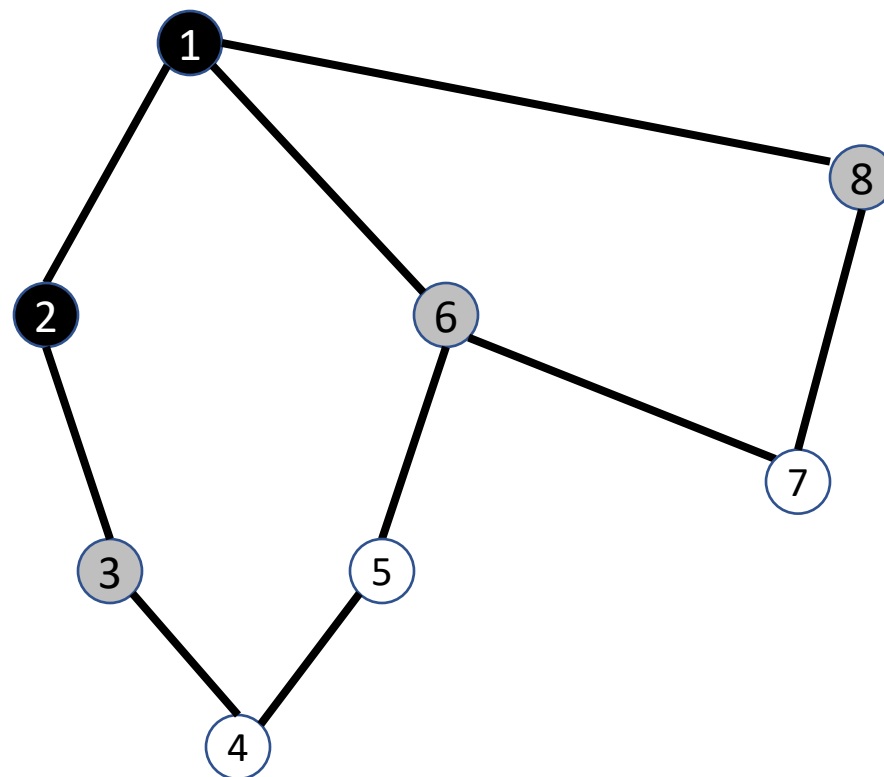
Queue: 2 6 8

# Exercise: Finding the Smallest Cycle involving Vertex 1



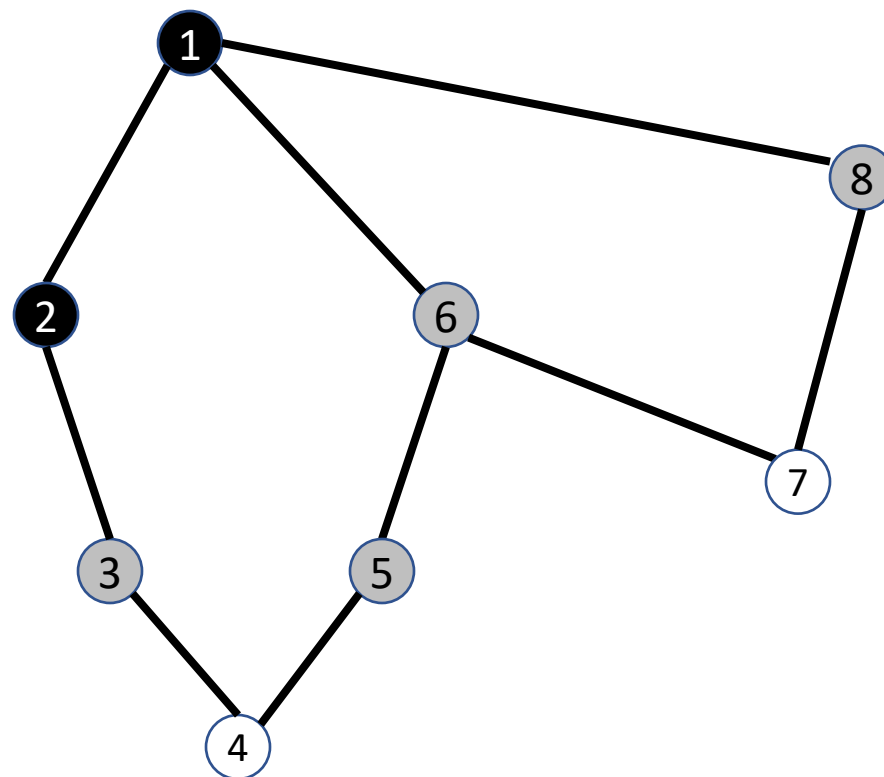
Queue: 2 6 8 3

# Exercise: Finding the Smallest Cycle involving Vertex 1



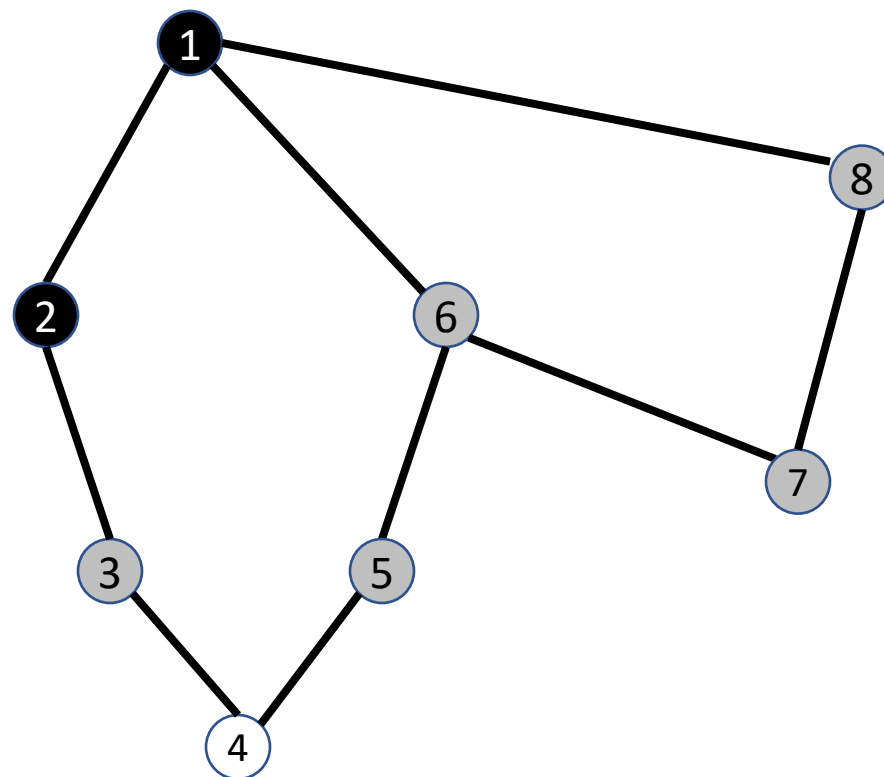
Queue: 6 8 3

# Exercise: Finding the Smallest Cycle involving Vertex 1



Queue: 6 8 3 5

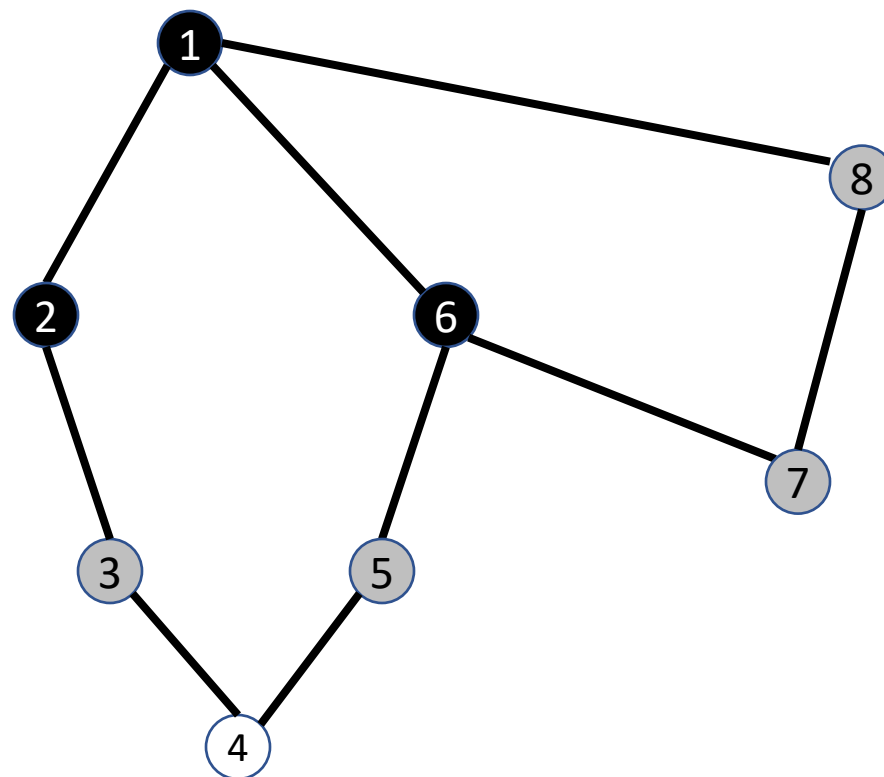
# Exercise: Finding the Smallest Cycle involving Vertex 1



Queue: 6 8 3 5 7

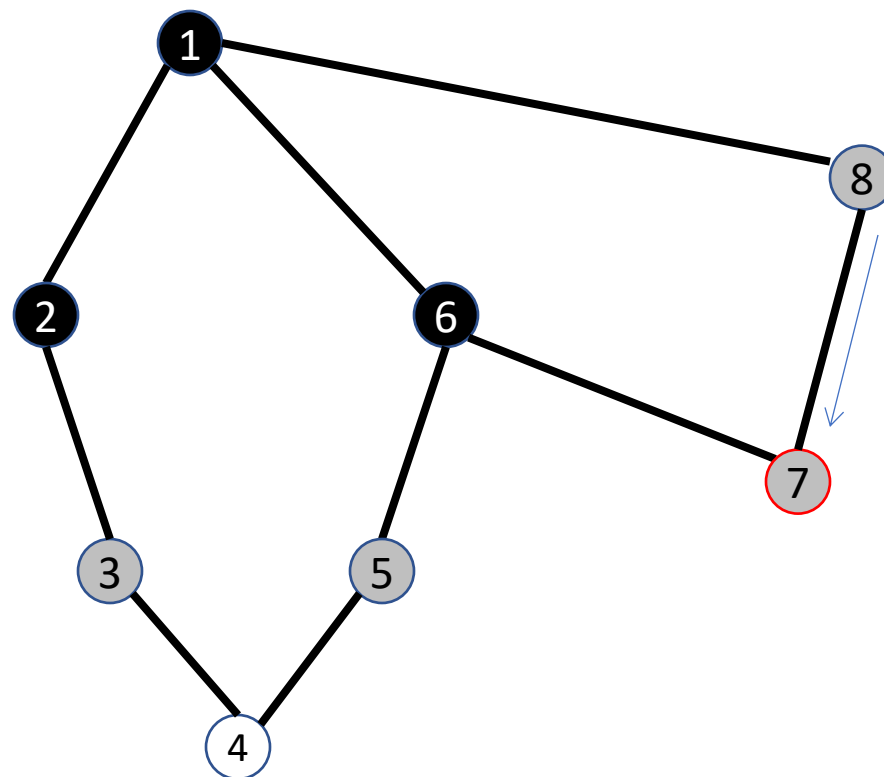


# Exercise: Finding the Smallest Cycle involving Vertex 1



Queue: 8 3 5 7

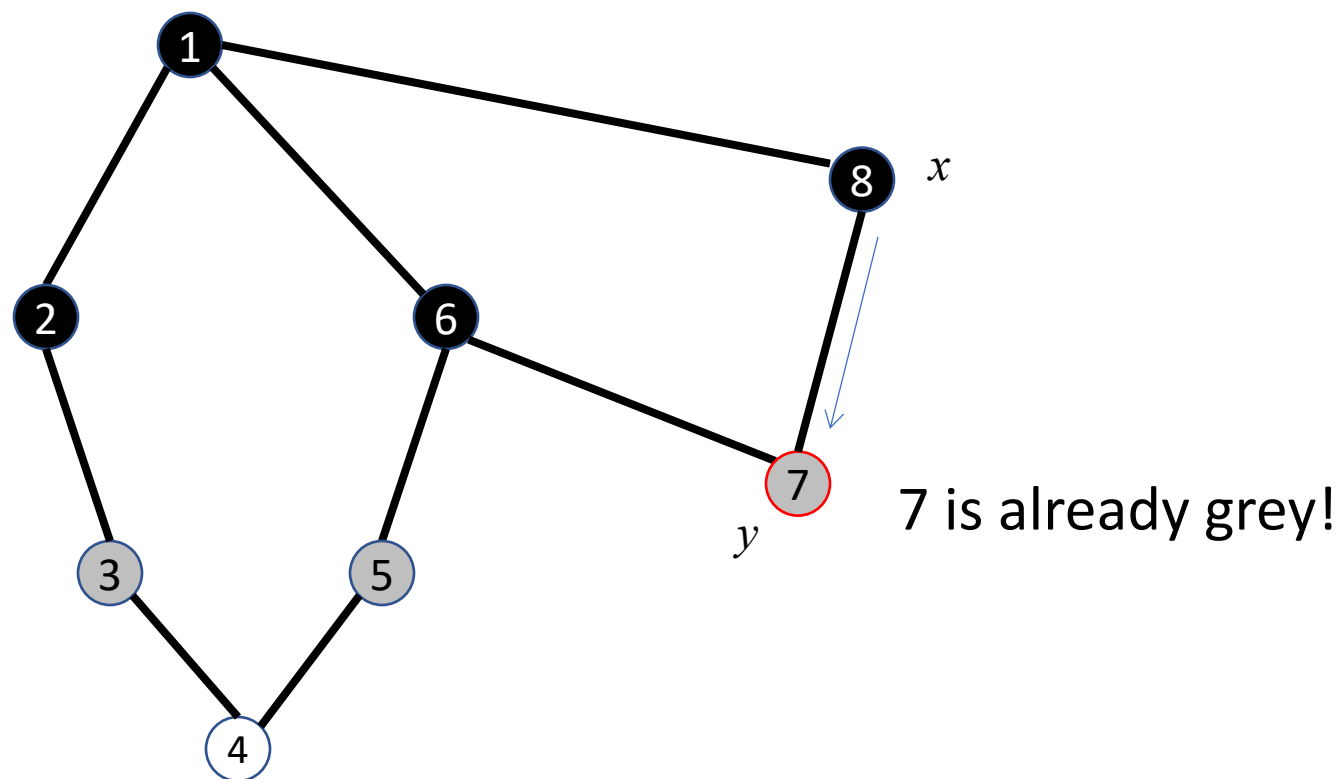
# Exercise: Finding the Smallest Cycle involving Vertex 1



7 is already grey!

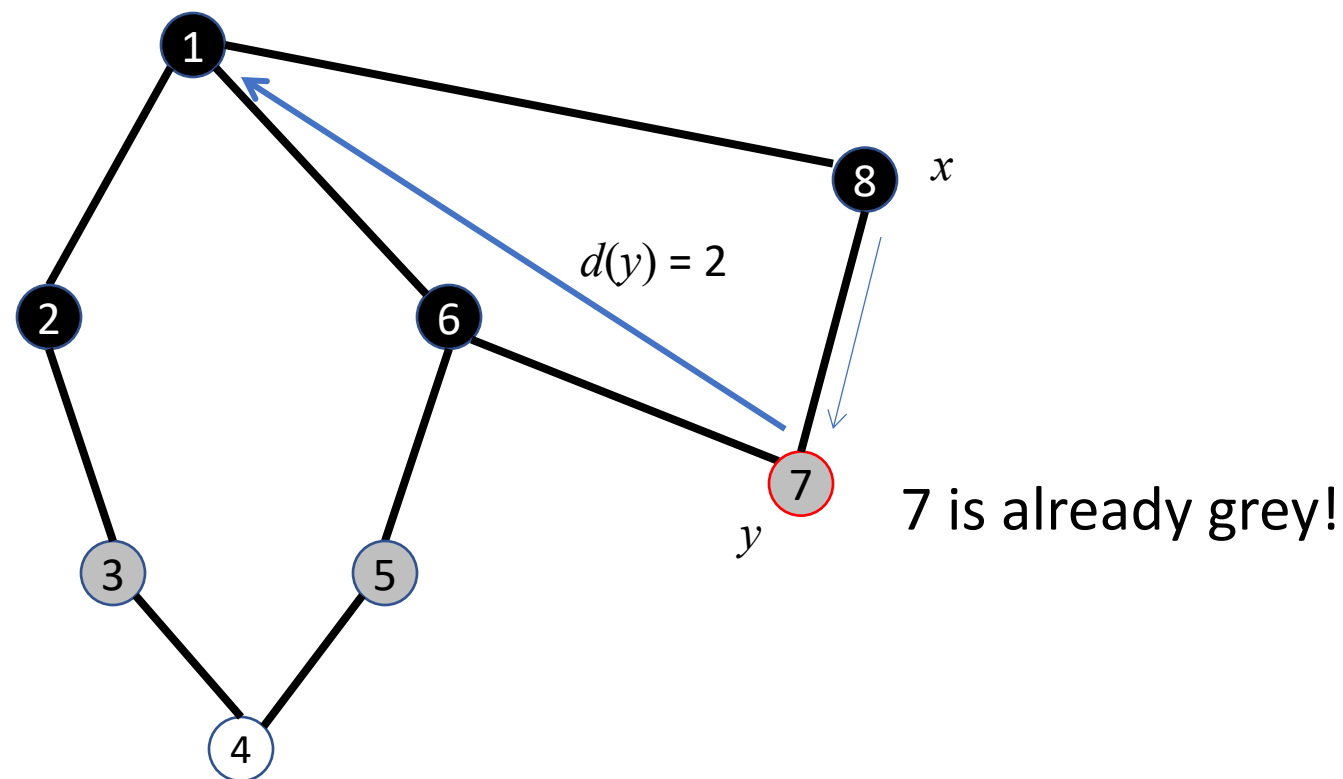
Queue: 8 3 5 7

# Exercise: Finding the Smallest Cycle involving Vertex 1



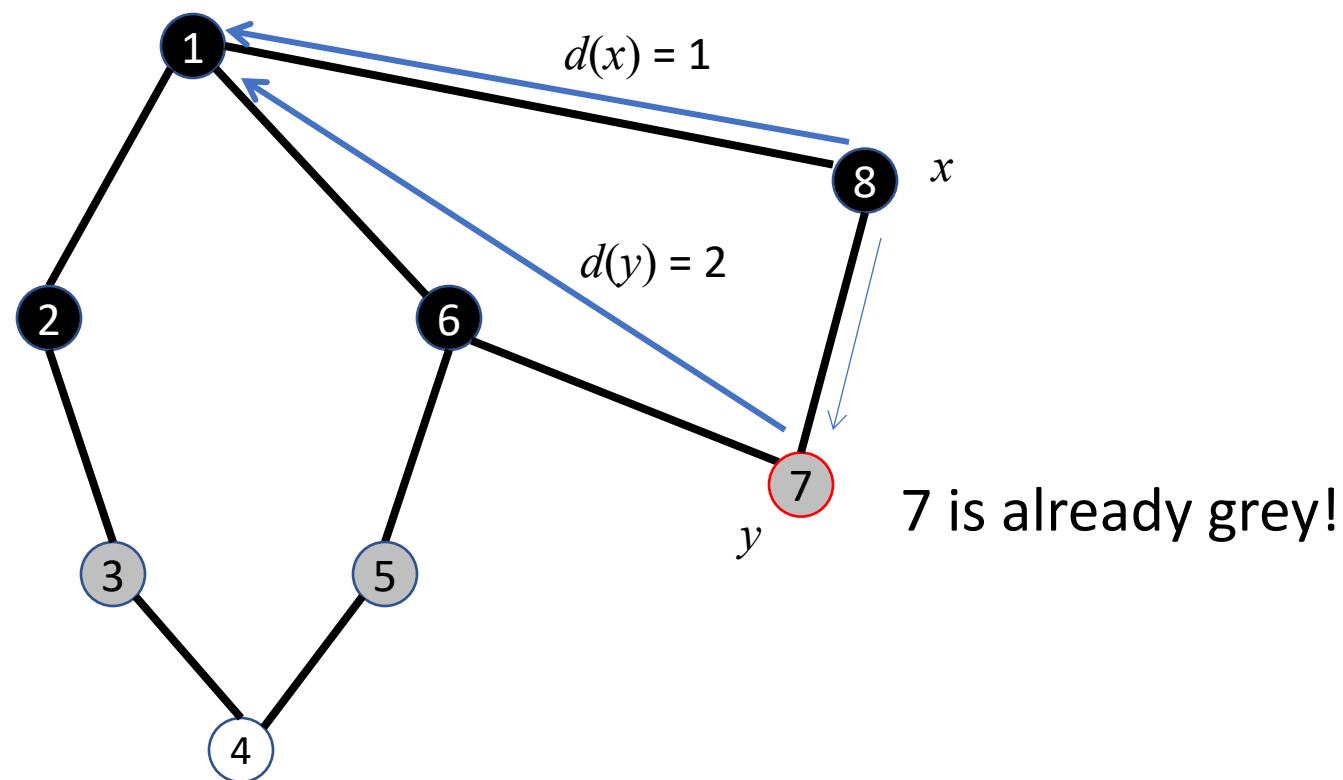
Queue: 3 5 7 (stopped)

# Exercise: Finding the Smallest Cycle involving Vertex 1



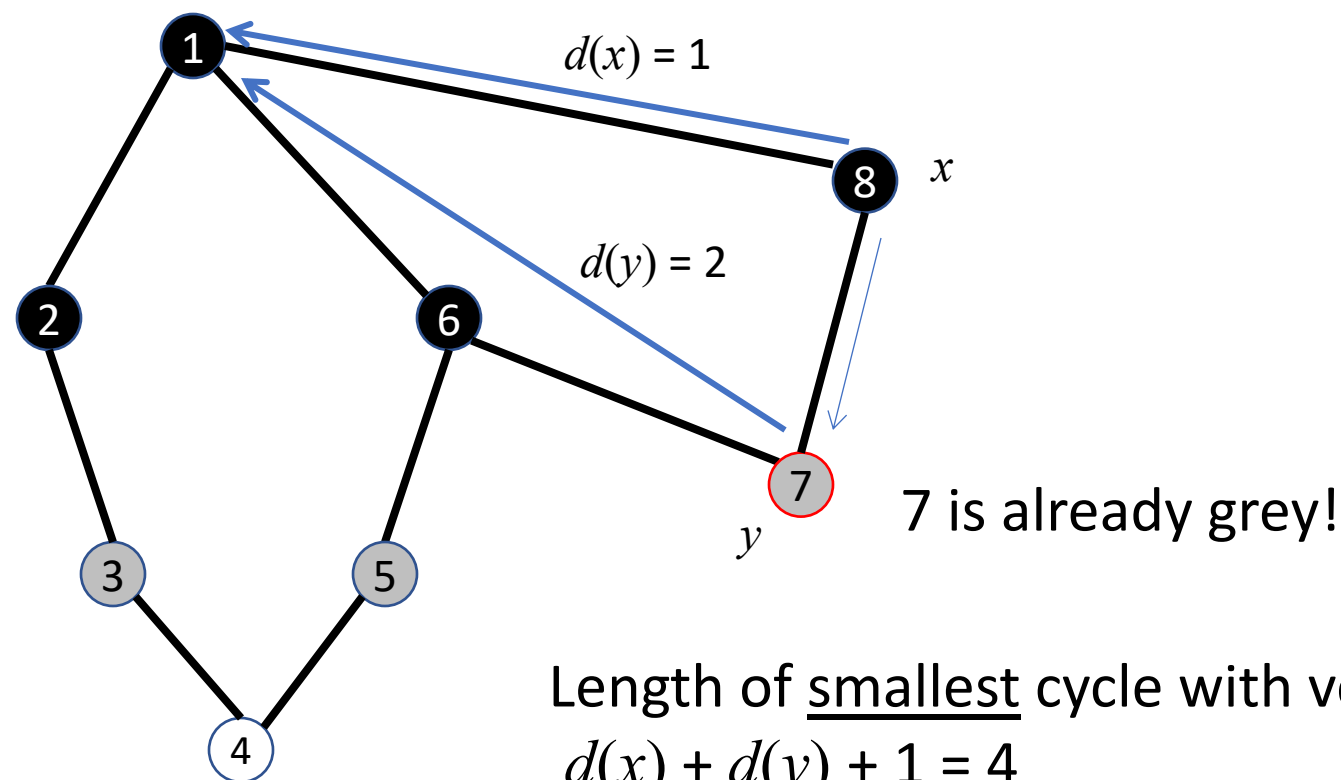
Queue: 3 5 7 (stopped)

# Exercise: Finding the Smallest Cycle involving Vertex 1



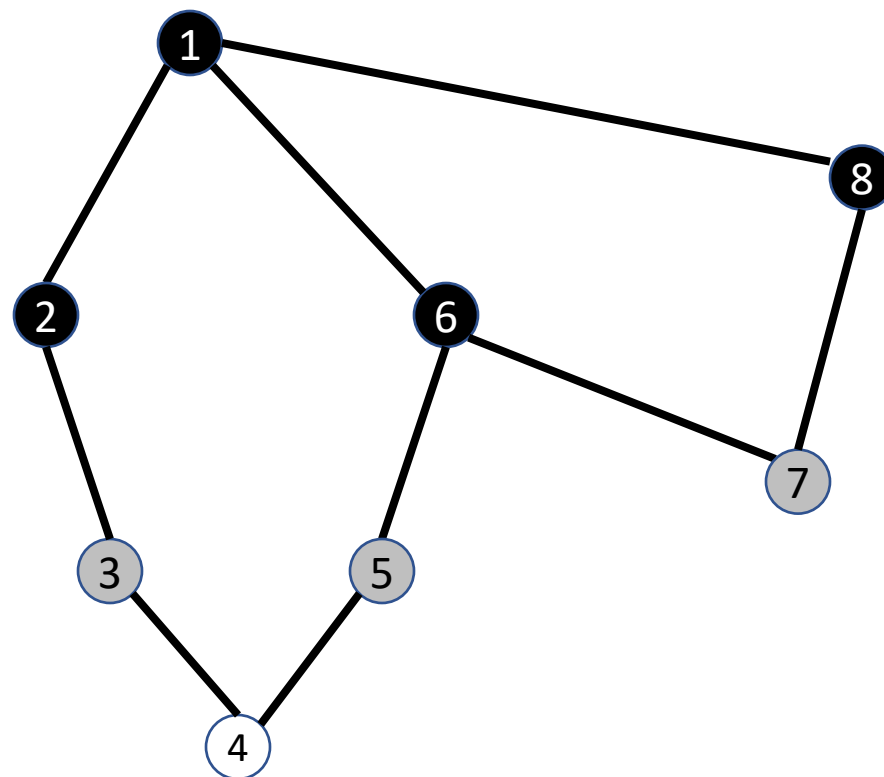
Queue: 3 5 7 (stopped)

# Exercise: Finding the Smallest Cycle involving Vertex 1



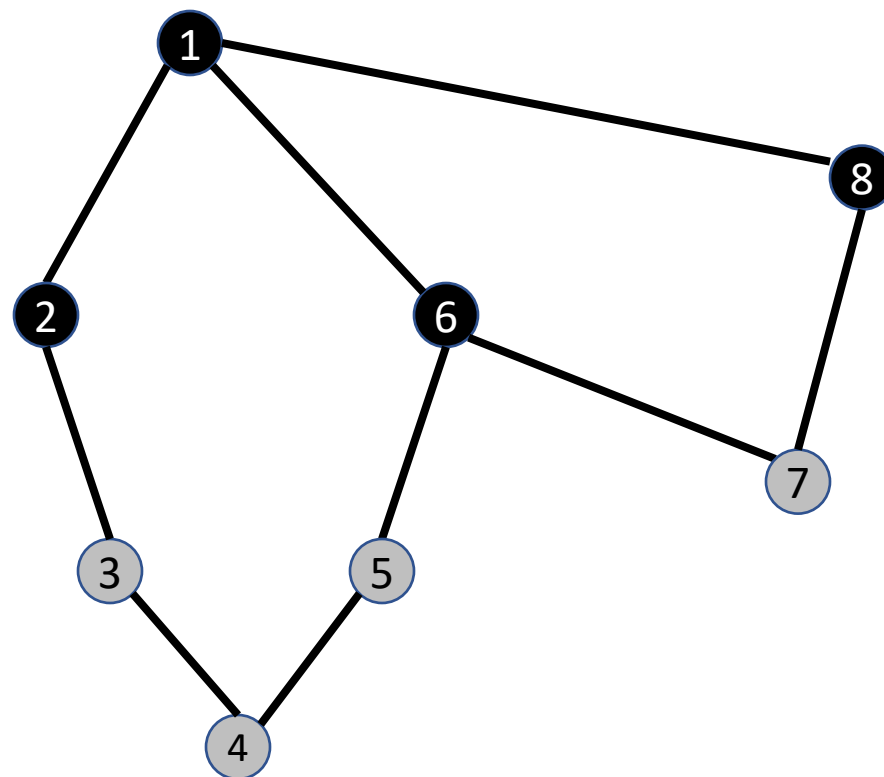
Queue: 3 5 7 (stopped)

# Exercise: Finding the Smallest Cycle involving Vertex 1



Queue: 3 5 7

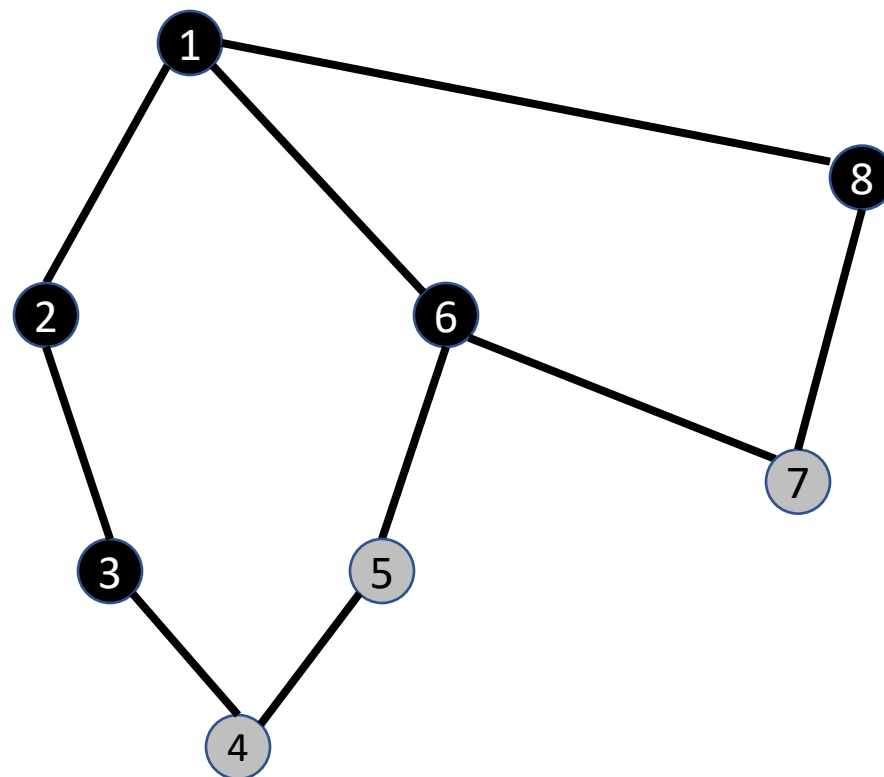
# Exercise: Finding the Smallest Cycle involving Vertex 1



Queue: 3 5 7 4

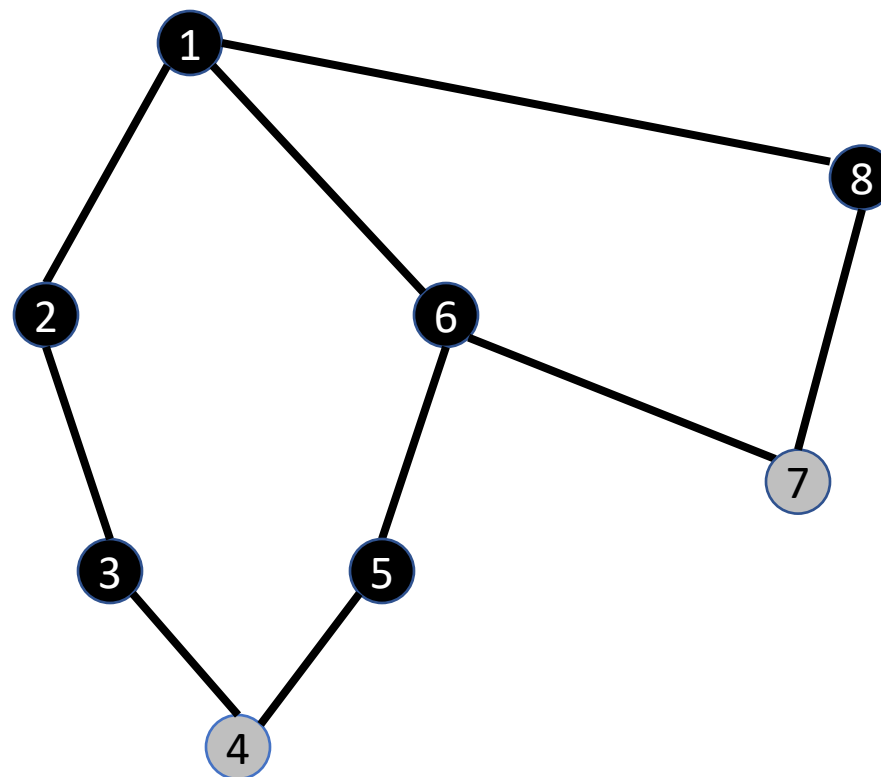


# Exercise: Finding the Smallest Cycle involving Vertex 1



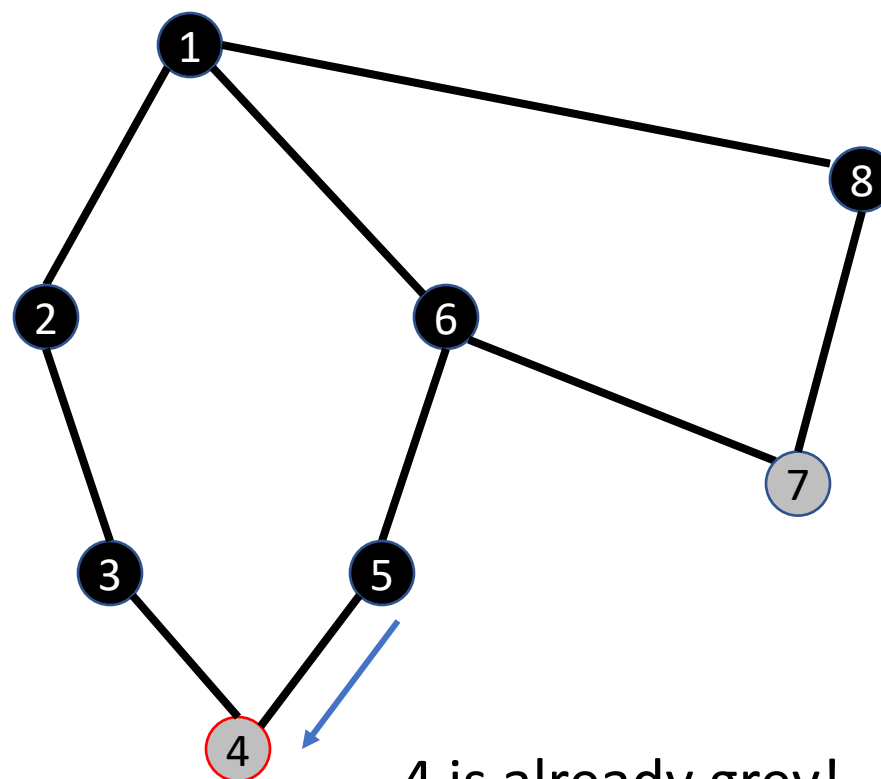
Queue: 5 7 4

# Exercise: Finding the Smallest Cycle involving Vertex 1



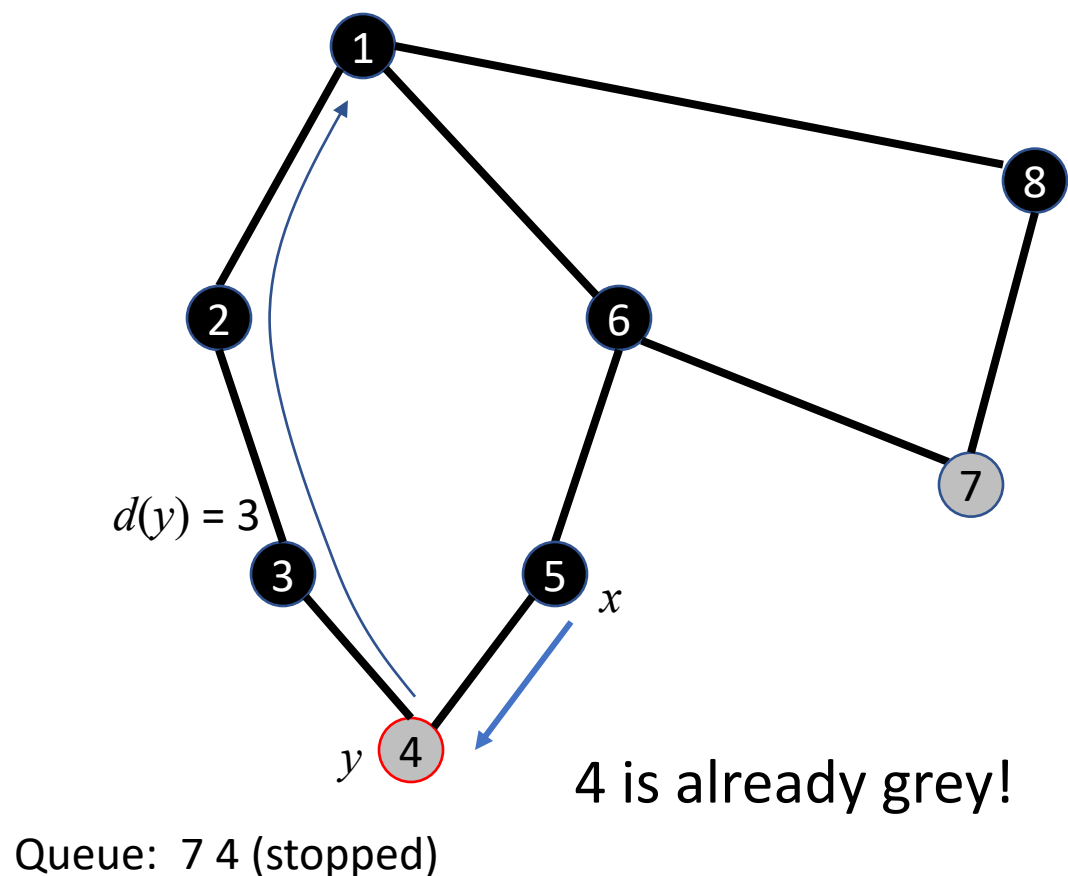
Queue: 7 4

# Exercise: Finding the Smallest Cycle involving Vertex 1

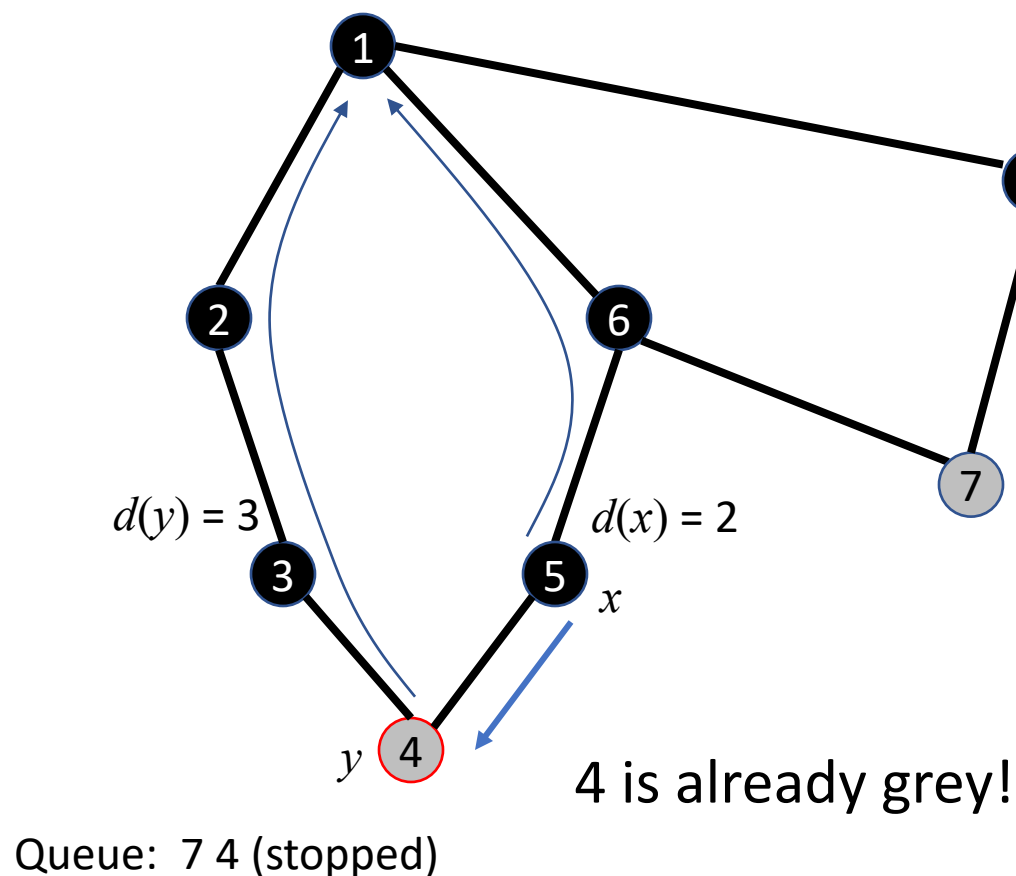


Queue: 7 4 (stopped)

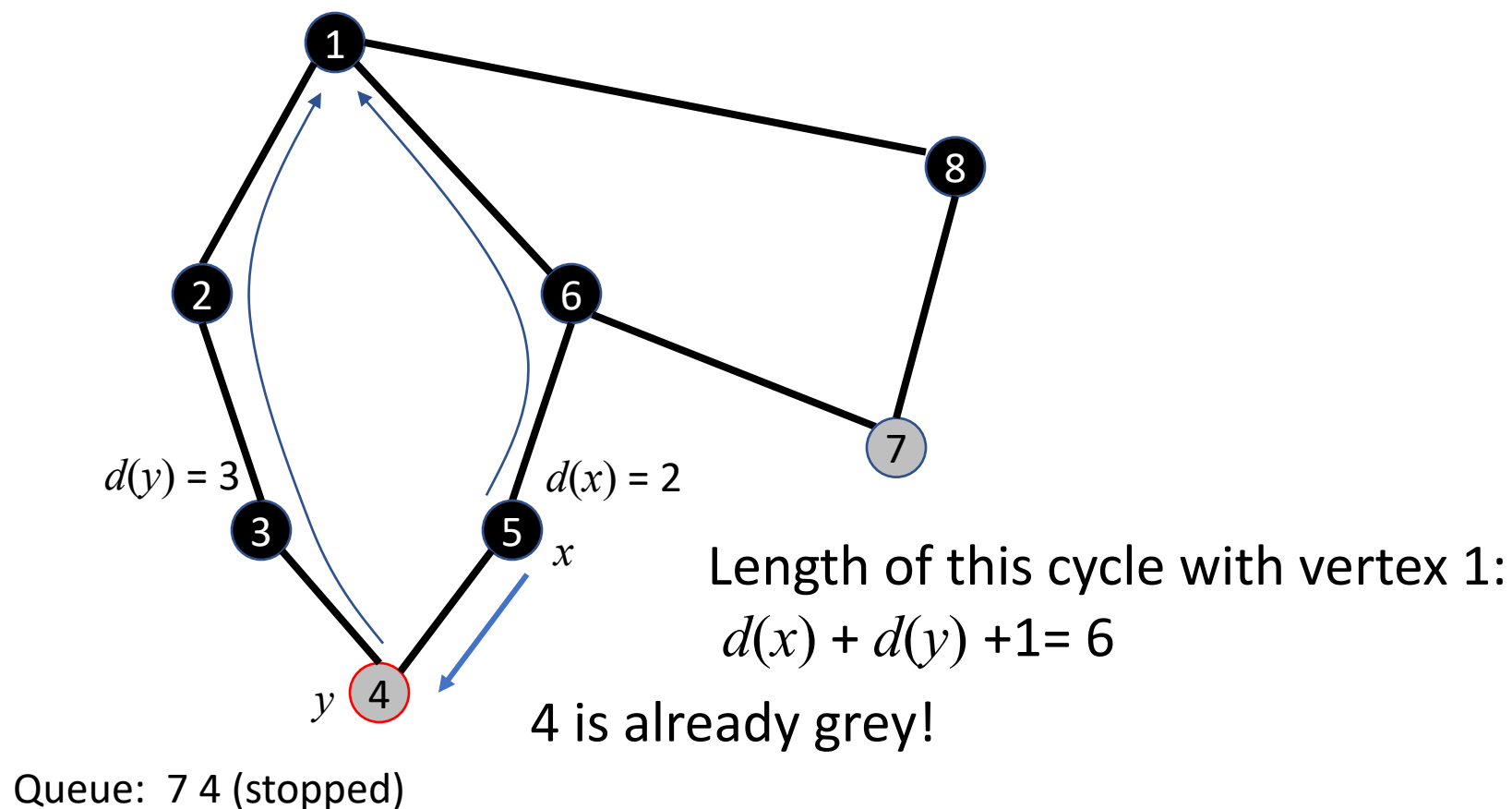
# Exercise: Finding the Smallest Cycle involving Vertex 1



# Exercise: Finding the Smallest Cycle involving Vertex 1



# Exercise: Finding the Smallest Cycle involving Vertex 1

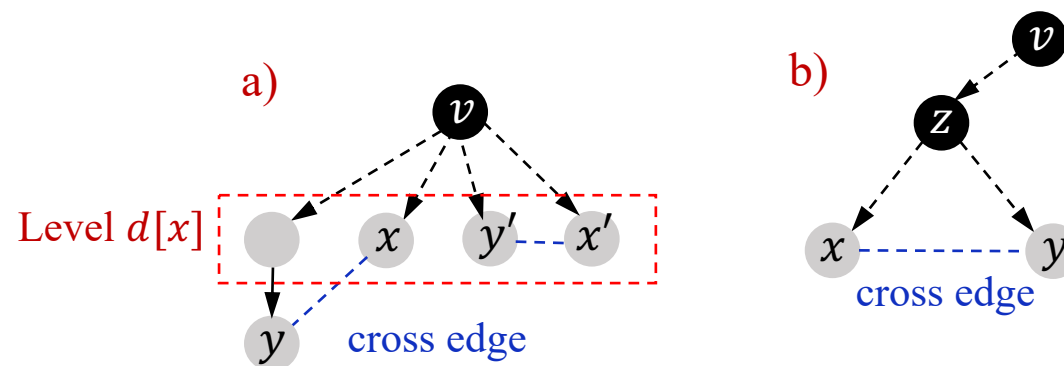


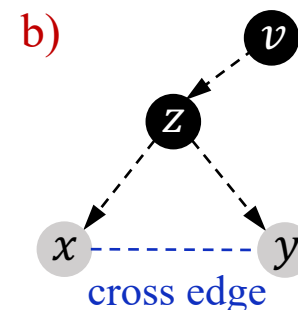
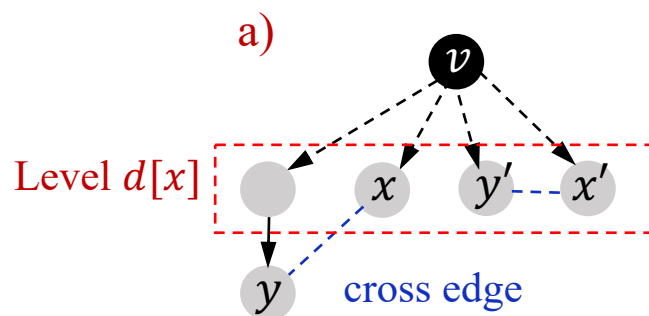
# Correctness of the Algorithm

1. Meeting a grey neighbour means that we have found a cycle;
2. The cycle found in this way will be the shortest.

Running  $\text{BFSVisit}(v)$  above, we have two cases when we first meet a cross edge  $(x, y)$  and let  $d[x] \leq d[y]$ :

- a)  $x$  and  $y$  are in different subtrees of  $v$
- b)  $x$  and  $y$  are in the same subtree of  $v$





Running  $\text{BFSVisit}(v)$  above gives us two cases when we first meet a cross edge  $(x, y)$  and let  $d[x] \leq d[y]$ :

- a)  $x$  and  $y$  are in different subtrees of  $v$ :  $v, x, y$  are in the same cycle with length  $d[x] + d[y] + 1 \leq 2d[x] + 2$ . All cross edges detected at the same level have the same upper bound of cycle length. Then, the algorithm finds the smallest cycle containing  $v$ .
- b)  $x$  and  $y$  are in the same subtree of  $v$ : a smaller cycle (not including  $v$ ) exists. Let  $z$  be the least common ancestor of  $x$  and  $y$ . When  $\text{BFSVisit}$  starts from  $z$ , we should find the smallest cycle containing  $z$  based on a).



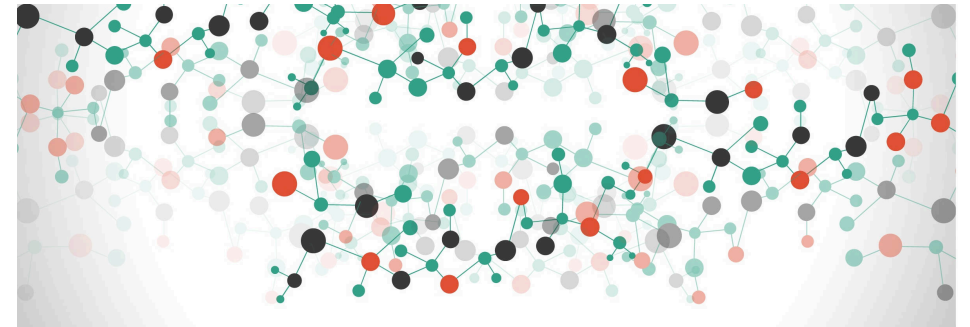
# Finding Directed Girth of Digraph

Following procedure finds the (length of the) shortest directed cycle through node  $v$  in a digraph.

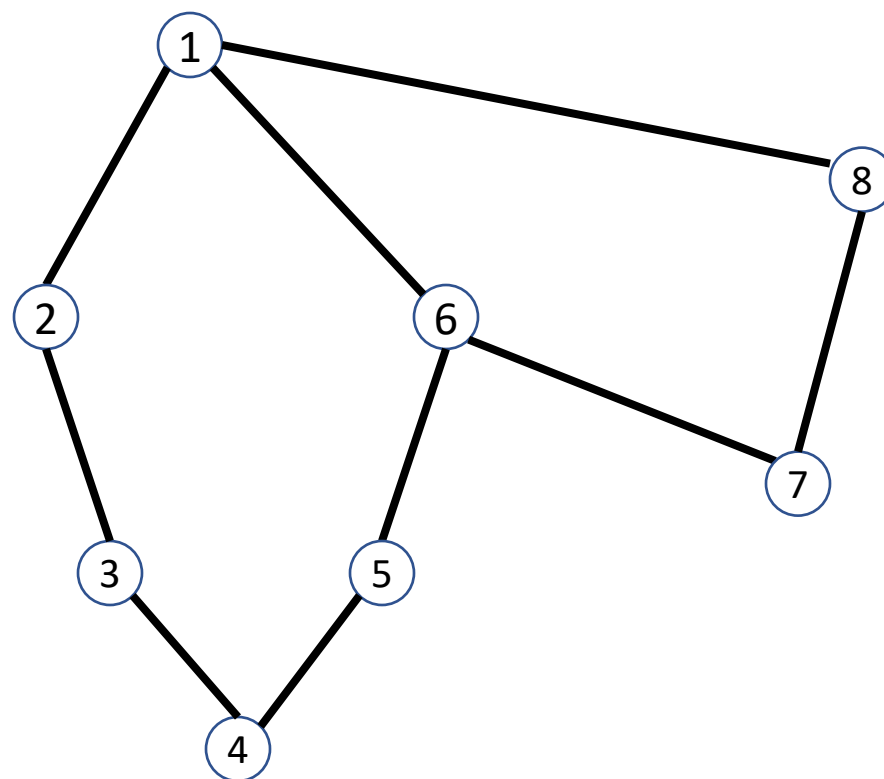
1. Run BFSVisit( $v$ ) starting at node  $v$ .
  2. The first time a **back-arc** of the form  $(x, v)$  is found, we have found a cycle of length equal to  $(1 + \text{the depth of } x \text{ in the search tree})$ .
  3. Stop and Return the length found.
- Run the above procedure on every node and pick the length of the shortest cycle

# SUMMARY

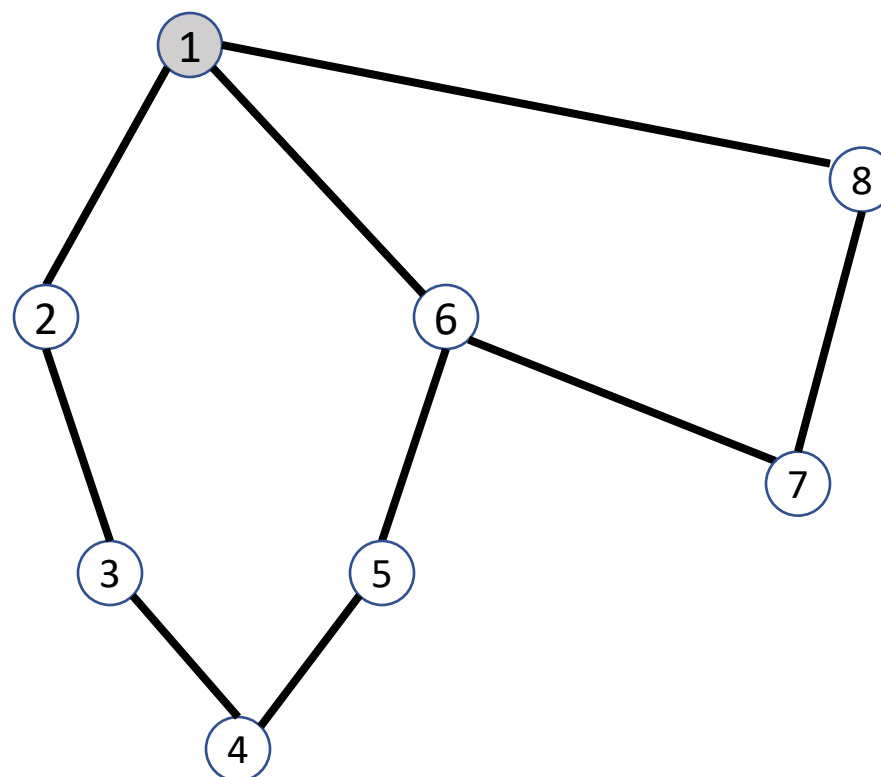
- Terminology
  - Cycle
  - Girth
- DFS Application: Cycle Detection
- BFS Application: Girth Detection
- Correctness of Girth Computation
- Illustrative Examples



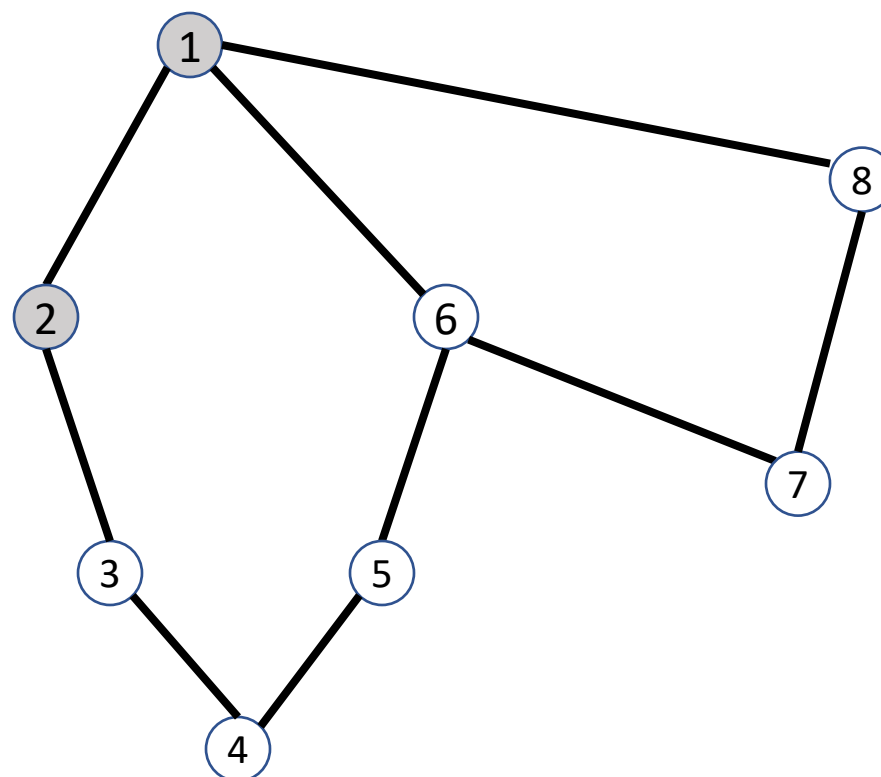
# Exercise: Finding Cycle involving Vertex 1 using DFS



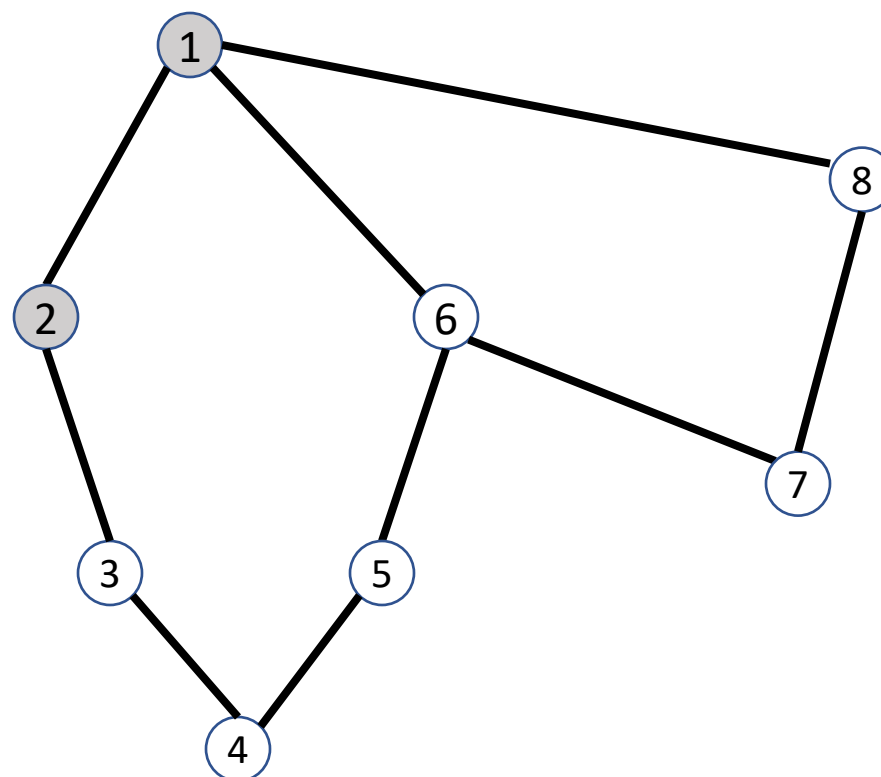
# Exercise: Finding Cycle involving Vertex 1 using DFS



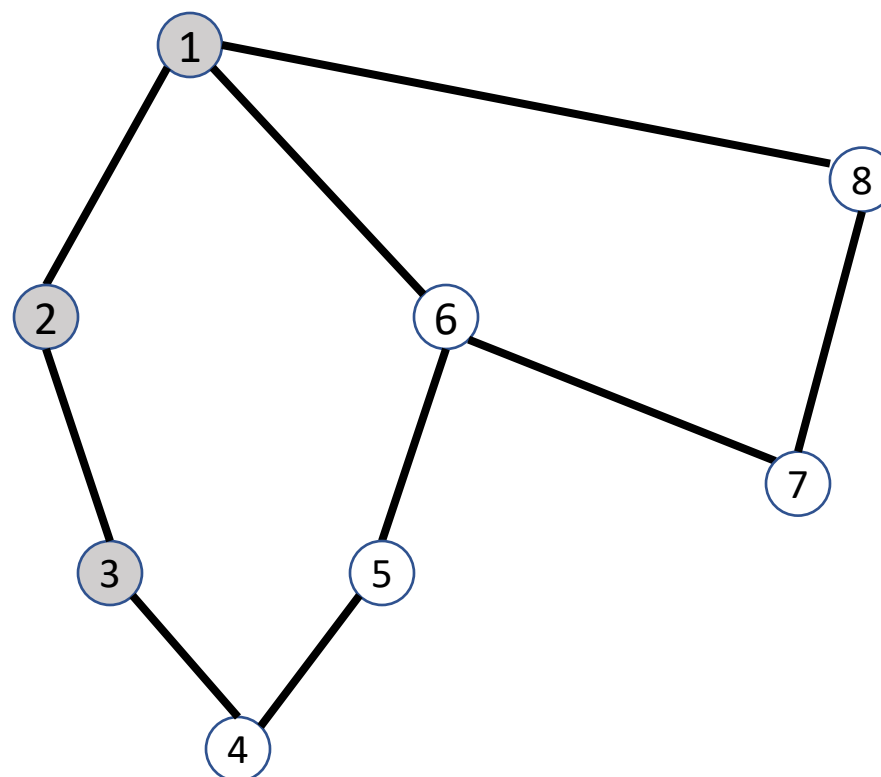
# Exercise: Finding Cycle involving Vertex 1 using DFS



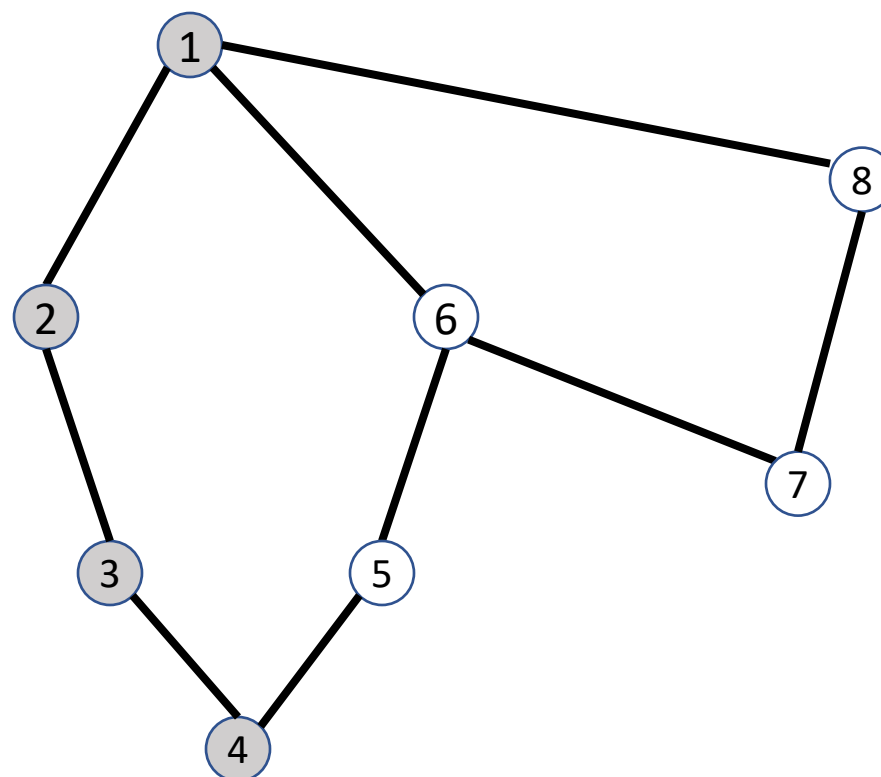
# Exercise: Finding Cycle involving Vertex 1 using DFS



# Exercise: Finding Cycle involving Vertex 1 using DFS

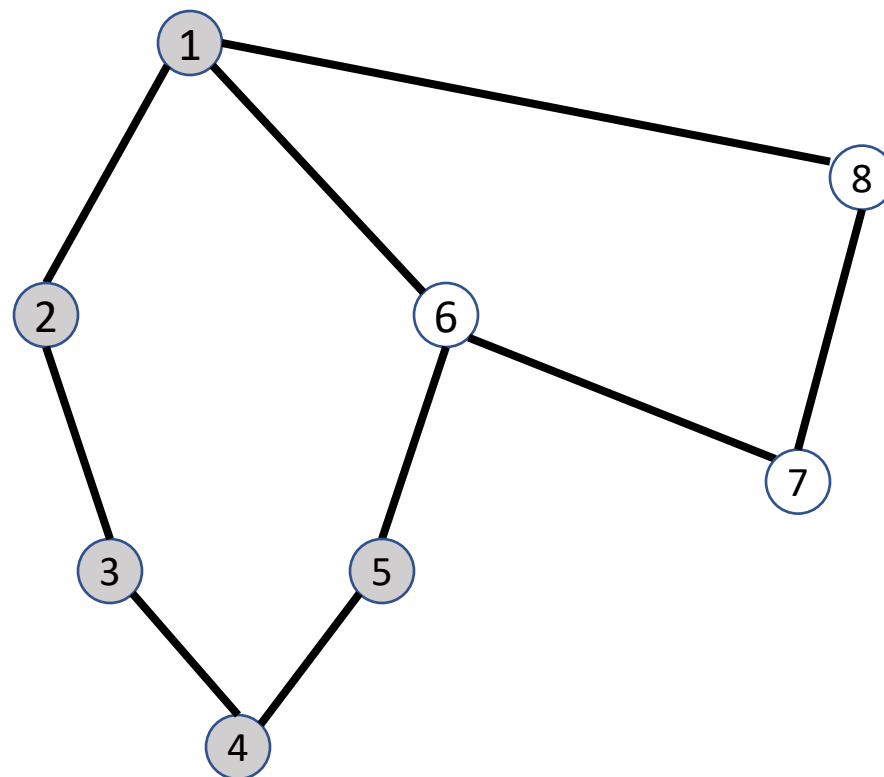


# Exercise: Finding Cycle involving Vertex 1 using DFS

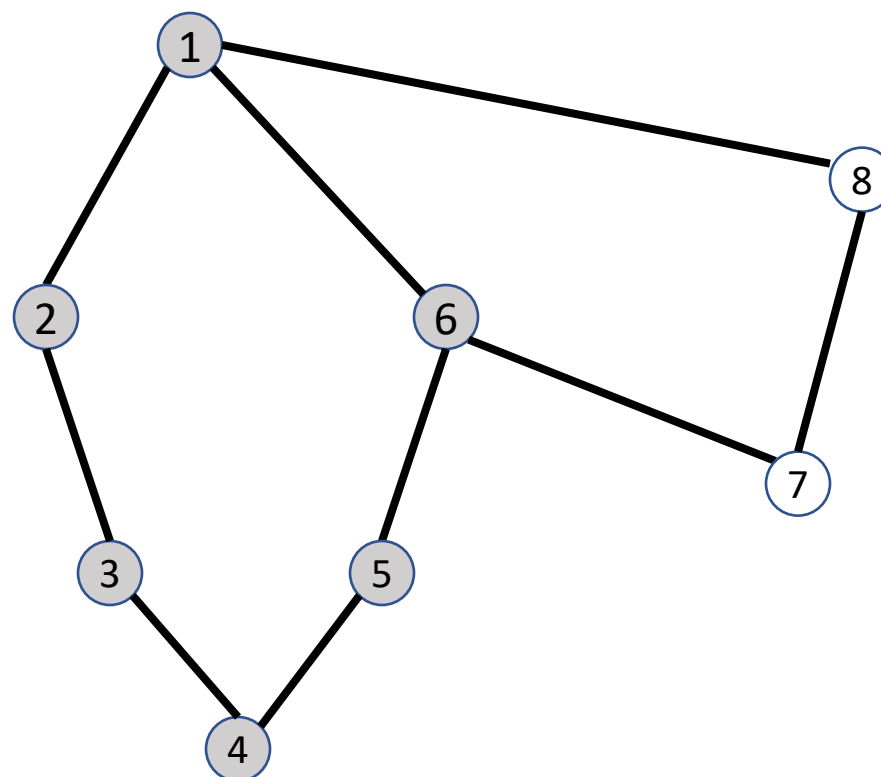




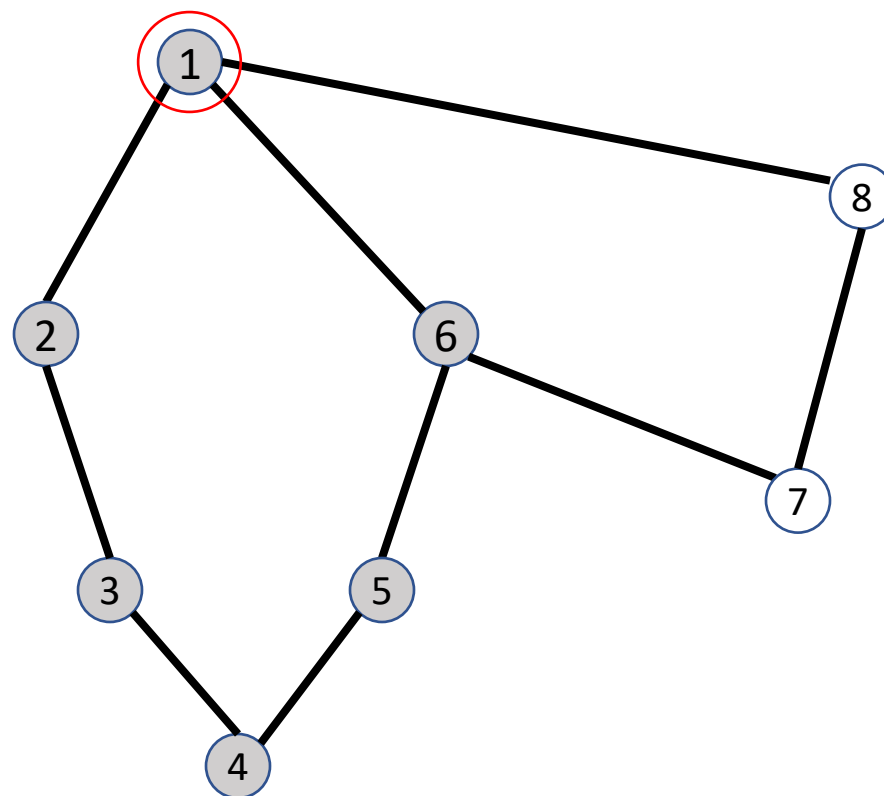
# Exercise: Finding Cycle involving Vertex 1 using DFS



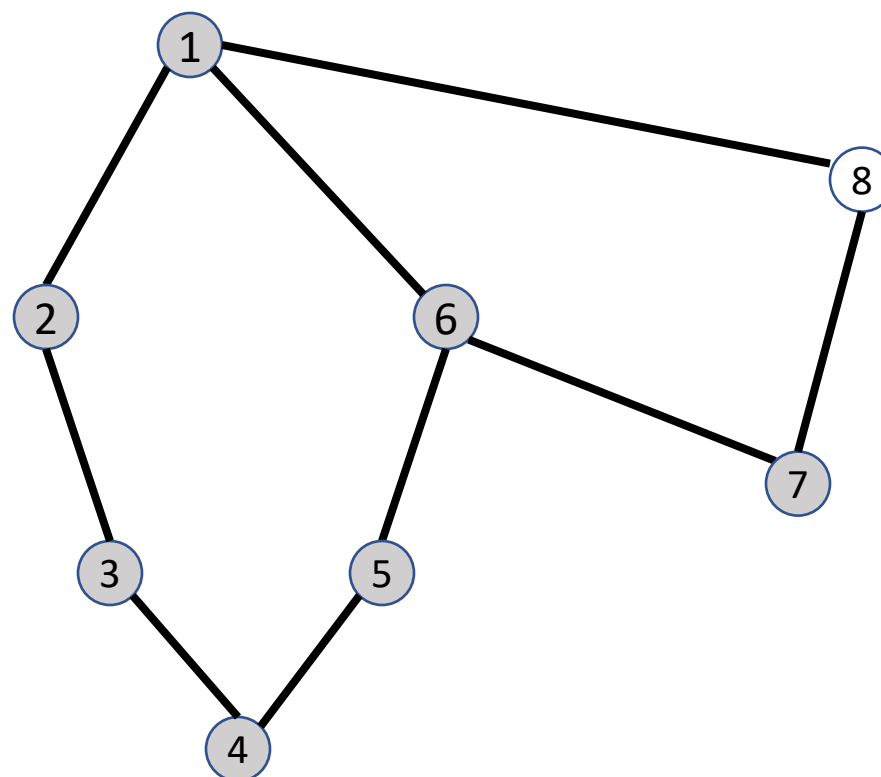
# Exercise: Finding Cycle involving Vertex 1 using DFS



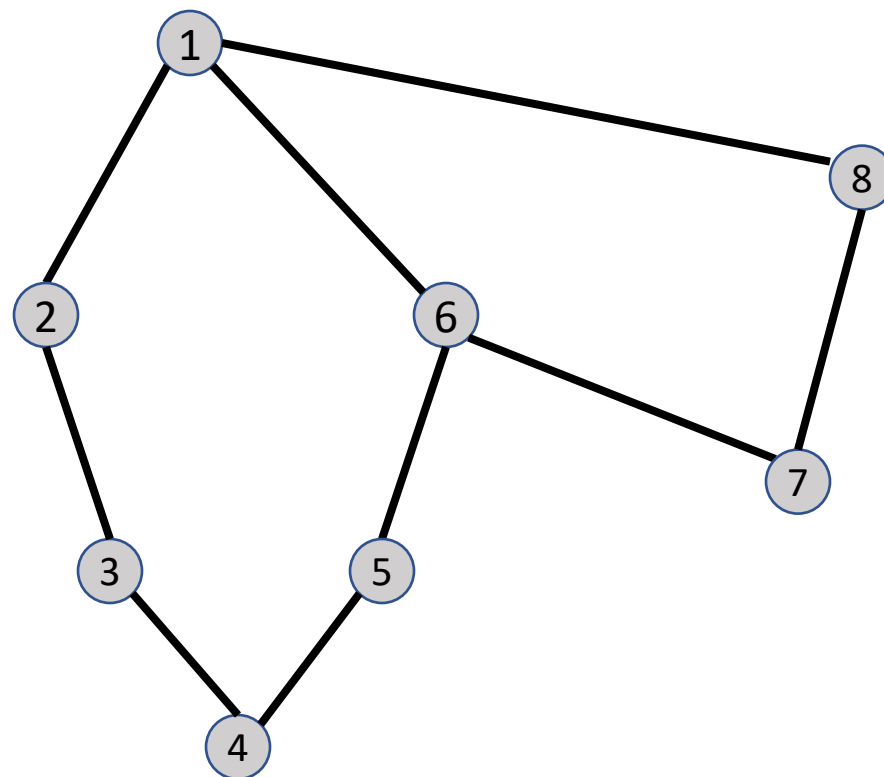
# Exercise: Finding Cycle involving Vertex 1 using DFS



# Exercise: Finding Cycle involving Vertex 1 using DFS



# Exercise: Finding Cycle involving Vertex 1 using DFS



# Exercise: Finding Cycle involving Vertex 1 using DFS

