

# Sorting III: Heapsort

---

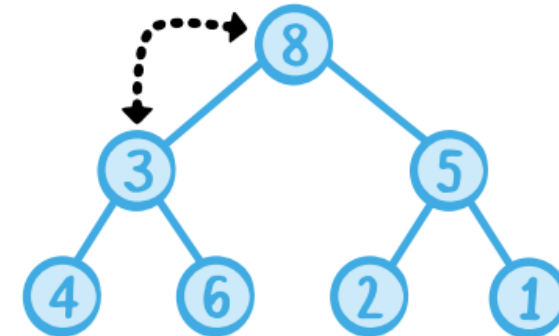
COMPSCI 220: WEEK 8.2

Instructor: Meng-Fen Chiang



# OUTLINE

- Heapsort
  - Algorithms
  - Illustrating Example
- Time Complexity Analysis



# Heapsort

Given an array:

1. Build the heap
2. Delete the maximum repeatedly (arranging the elements in the output list in reverse order of deletion) until the heap is empty.

# Step1: Building a Heap

- A straightforward method: sequentially insert nodes to the heap and maintain the heap property
- Inserting a node to a heap of height  $h$ , we need  $h = \lfloor \log_2 n \rfloor$  comparisons at most. Then the worst-case running time complexity of building a heap of size  $n$  is:

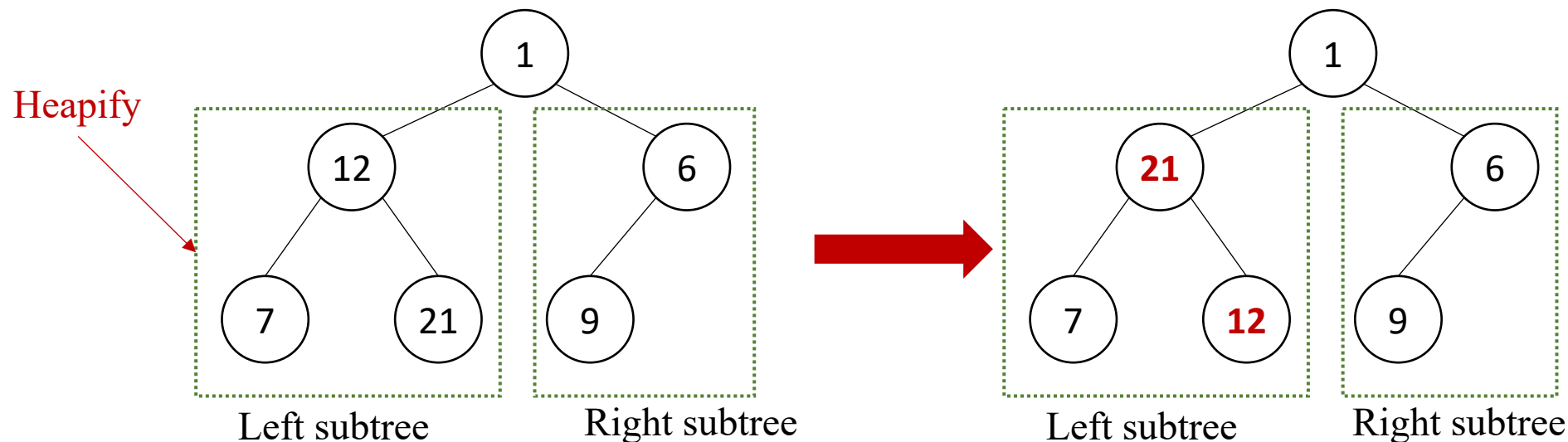
$$\begin{aligned} \bullet T(n) &= \lfloor \log_2 1 \rfloor + \lfloor \log_2 2 \rfloor + \lfloor \log_2 3 \rfloor + \cdots + \lfloor \log_2 n \rfloor \\ &\leq \log_2 n + \log_2 n + \log_2 n + \cdots + \log_2 n = n \log_2 n \end{aligned}$$

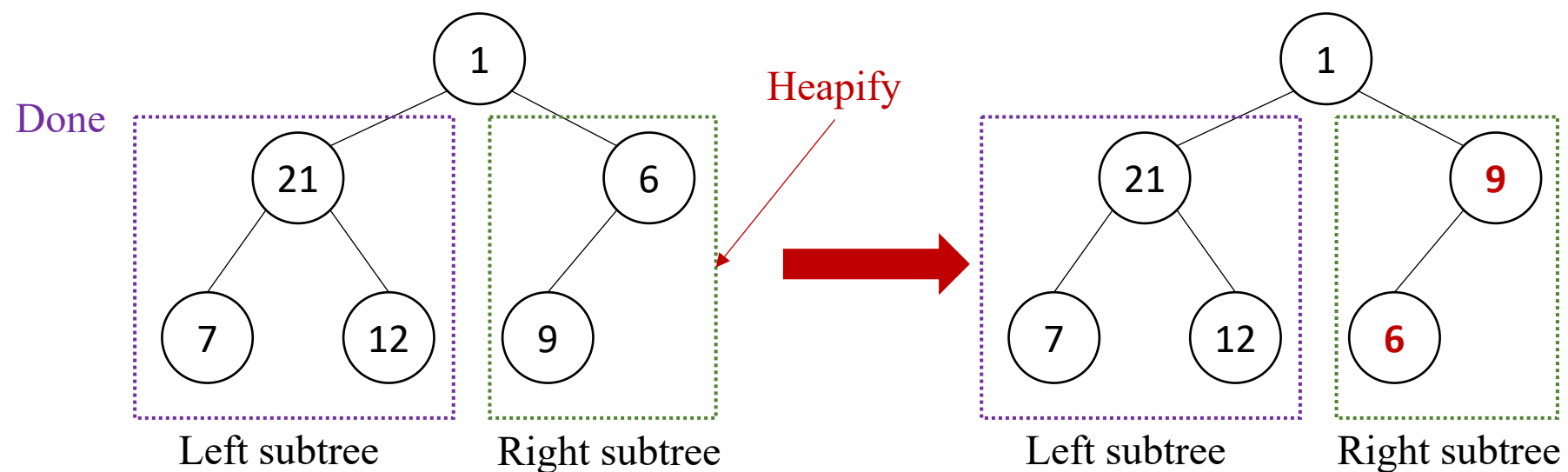
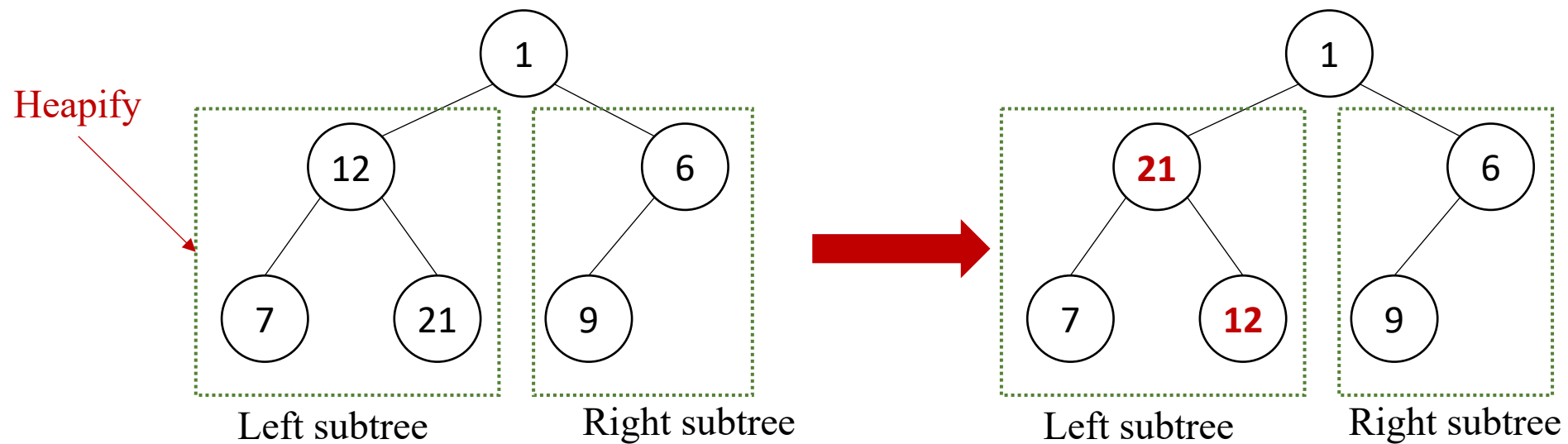
$$\begin{aligned} \bullet T(n) &= \lfloor \log_2 1 \rfloor + \lfloor \log_2 2 \rfloor + \cdots + \left\lfloor \log_2 \left\lfloor \frac{n}{2} \right\rfloor \right\rfloor + \left\lfloor \log_2 \left( \left\lfloor \frac{n}{2} \right\rfloor + 1 \right) \right\rfloor + \cdots + \lfloor \log_2 n \rfloor \\ &\geq \left\lfloor \log_2 \left( \left\lfloor \frac{n}{2} \right\rfloor \right) \right\rfloor + \left\lfloor \log_2 \left( \left\lfloor \frac{n}{2} \right\rfloor \right) \right\rfloor + \cdots + \left\lfloor \log_2 \left( \left\lfloor \frac{n}{2} \right\rfloor \right) \right\rfloor \\ &= \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \log_2 \left( \left\lfloor \frac{n}{2} \right\rfloor \right) \right\rfloor = \left\lfloor \frac{n}{2} \right\rfloor \left( \left\lfloor \log_2 \left( \left\lfloor \frac{n}{2} \right\rfloor + 1 \right) \right\rfloor - 1 \right) \geq \frac{n}{2} \log_2 \left( \left\lfloor \frac{n+2}{2} \right\rfloor \right) \geq \frac{n}{2} \log_2 \left( \frac{n}{2} \right) \end{aligned}$$

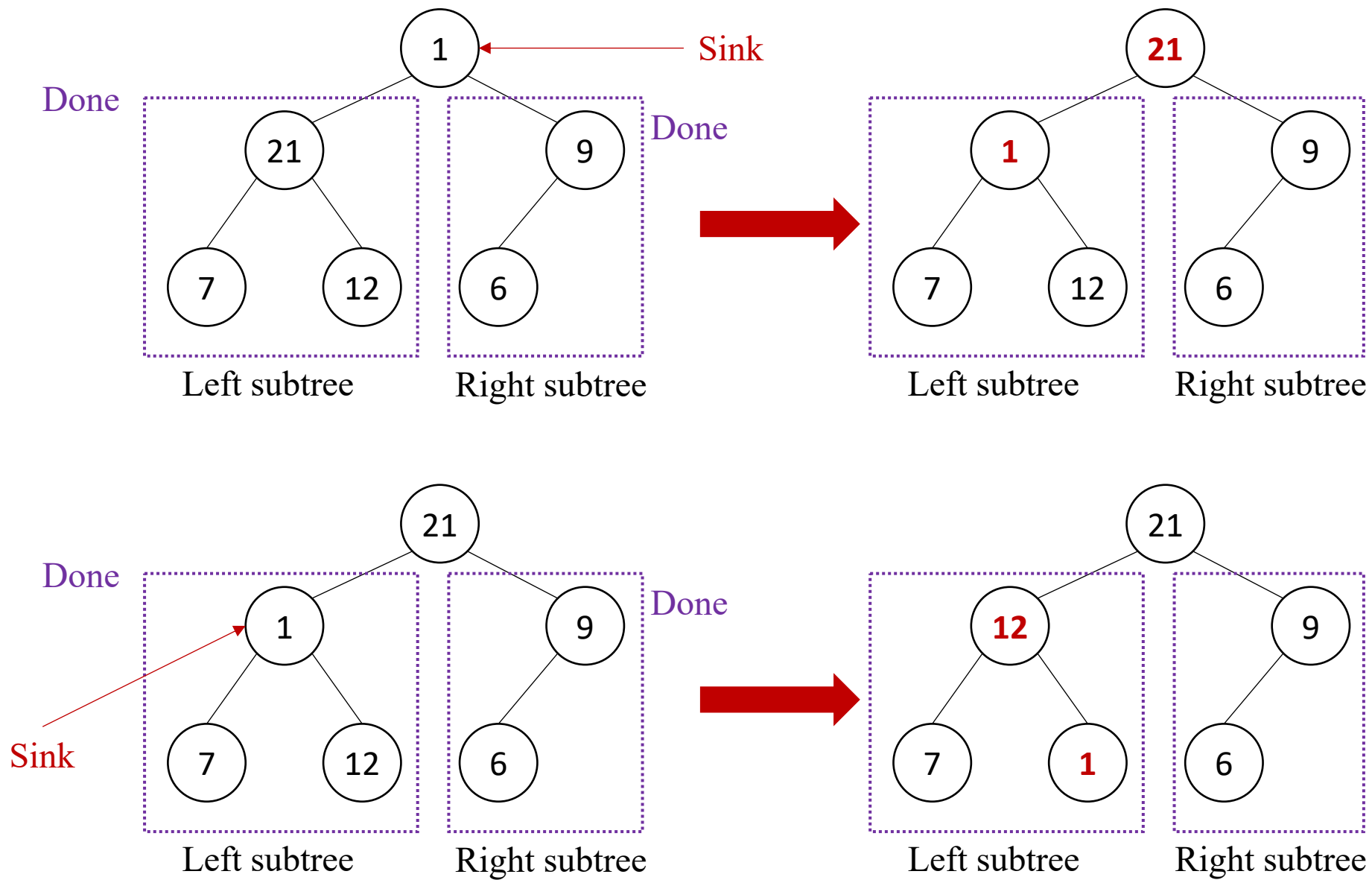
The above gives:  $T(n) \in \Theta(n \log n)$

# Can we do better?

- **Motivation of an alternative method.** We only need the heap property at the end, no need to keep it whenever an element is inserted.
  1. Build a complete binary tree without keeping the heap property – this is especially easy for array implementation which only needs to insert all elements to a list.
  2. After all elements are inserted, reorganize the positions of the elements to follow heap property, this operation is called **heapify** – recursively heapify the left and right subtrees, then put the root down to the right position.

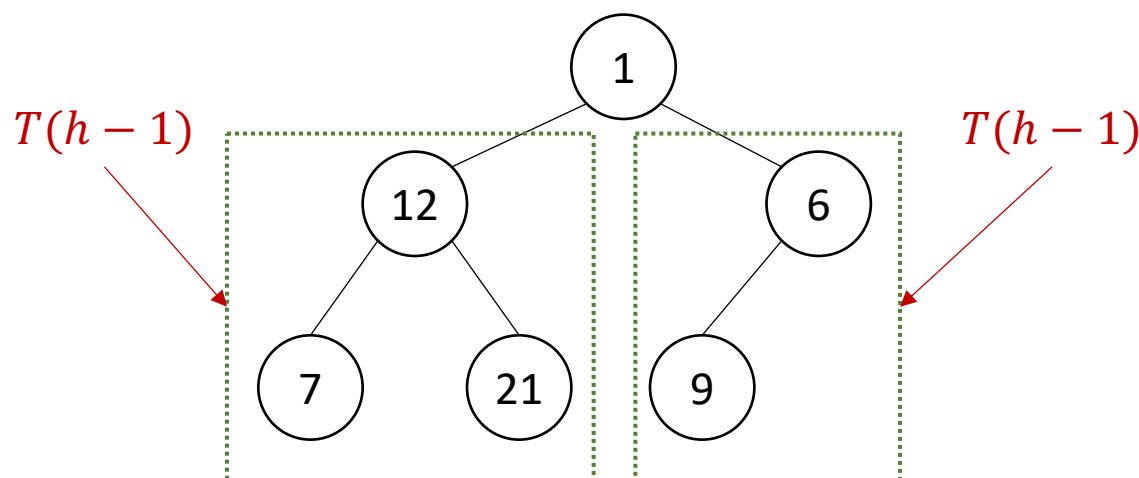






# Time Complexity Analysis: Heapifying

- Let  $T(h)$  is the running time complexity for building a heap with height at most  $h$ . It comprises of
  - Running time of heapifying the two subtrees recursively  $T(h-1)$
  - Running time of sinking the root to the right position needs at most  $h$  comparisons/swaps
  - Thus,  $T(h) = 2T(h-1) + h$ ,  $T(0) = 0$



We can show that  $T(h)$  is  $\Theta(2^h)$

Because  $h = \lfloor \log_2 n \rfloor$ , then  $T(n)$  is  $\Theta(n)$

$$T(h) - 2T(h-1) = h$$

$$2T(h-1) - 4T(h-2) = 2(h-1)$$

$$4T(h-2) - 8T(h-3) = 4(h-2)$$

... ..

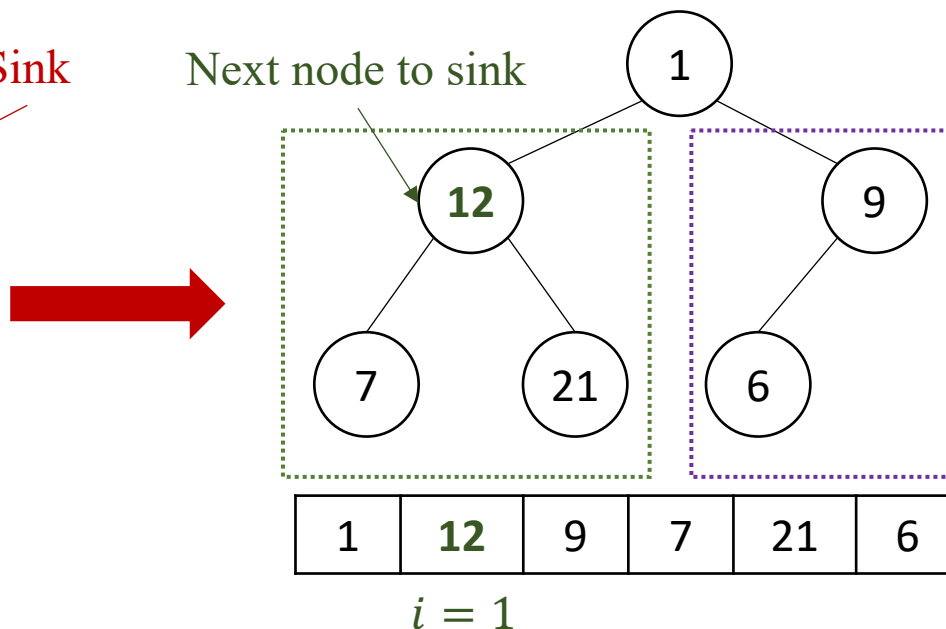
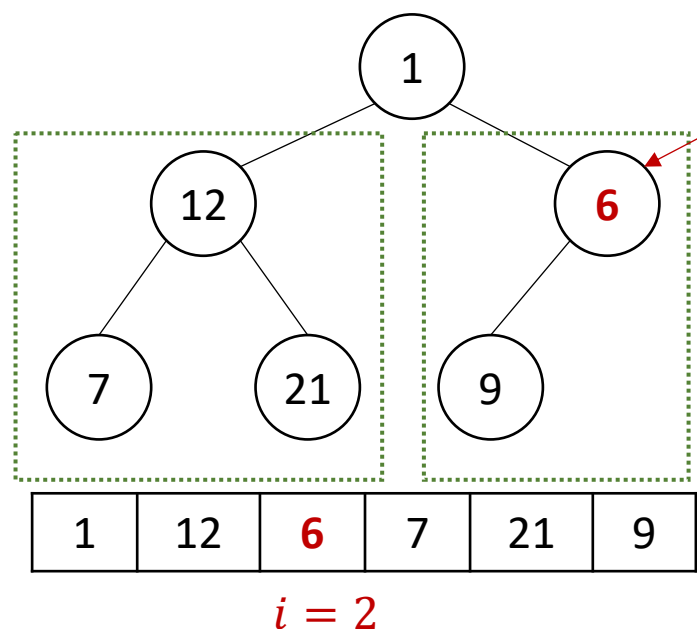
$$2^{h-1}T(1) - 2^hT(0) = 2^{h-1} \cdot 1$$

$$T(h) = (2^{h-1} \cdot 1 + 2^{h-2} \cdot 2 + \dots + 2^0 \cdot h)$$



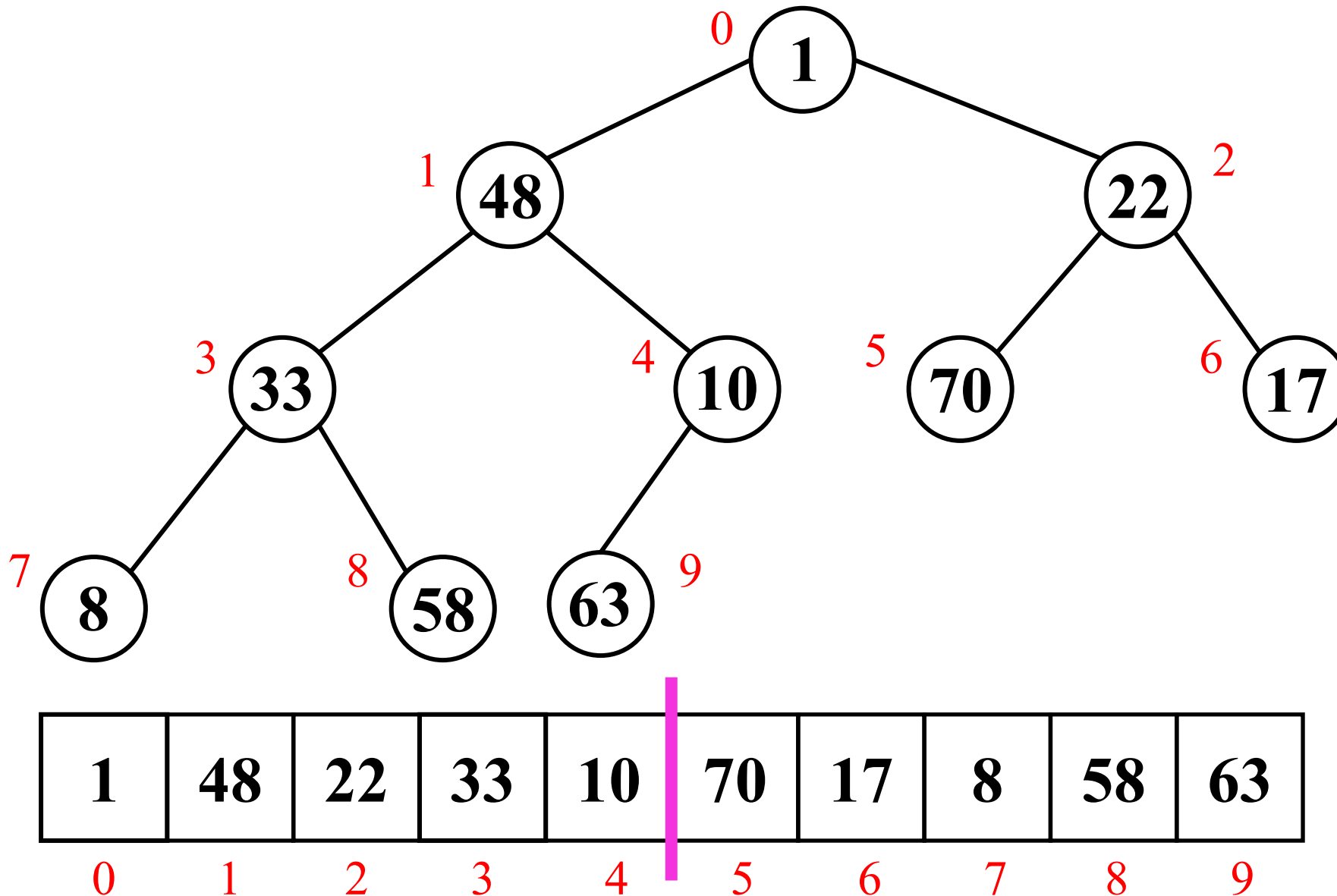
# Heapifying without Recursions

- Sink the nodes in a bottom-up manner. Consider the array implementation of heap.
  - Find the lowest non-leaf node, which is at position  $\lfloor (n-1)/2 \rfloor$ . Set the current node index  $i = \lfloor (n-1)/2 \rfloor$
  - Sink the current node at position  $i$  to the correct position.
  - Find the next non-leaf node, that is the node at position  $i-1$ , go to step 2.

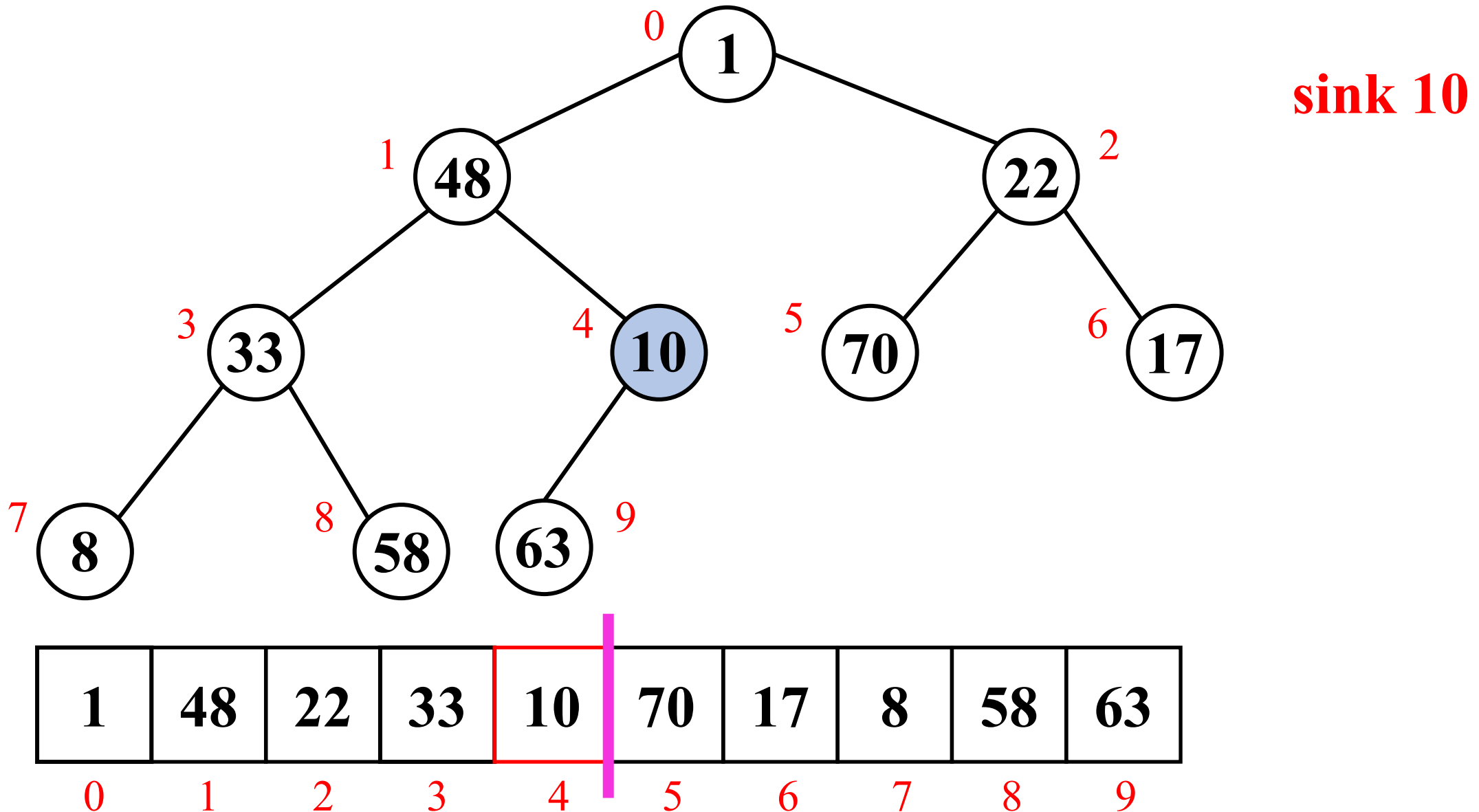


Repeat until we  
complete  
sinking the root

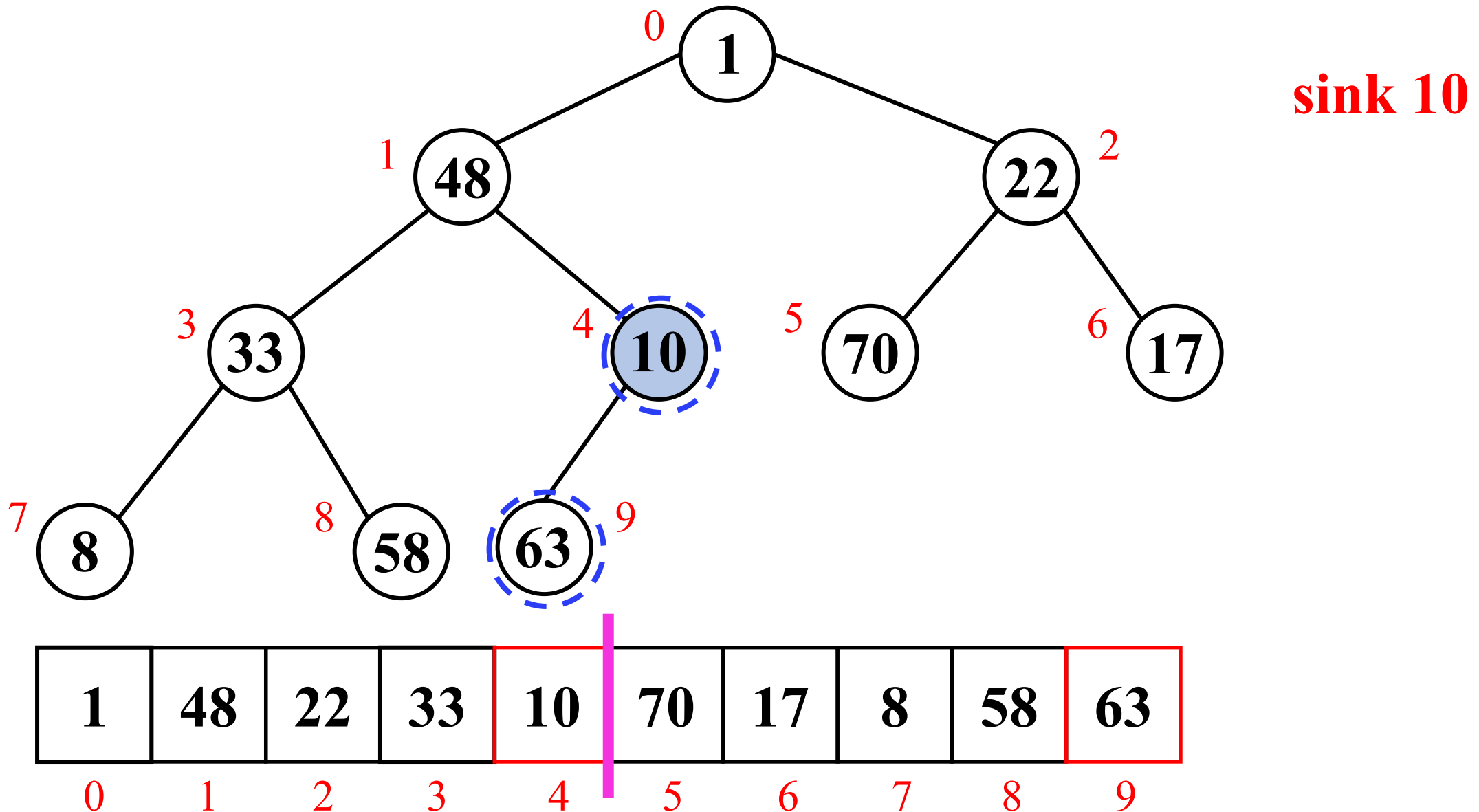
# Heapsort algorithms: array implementation



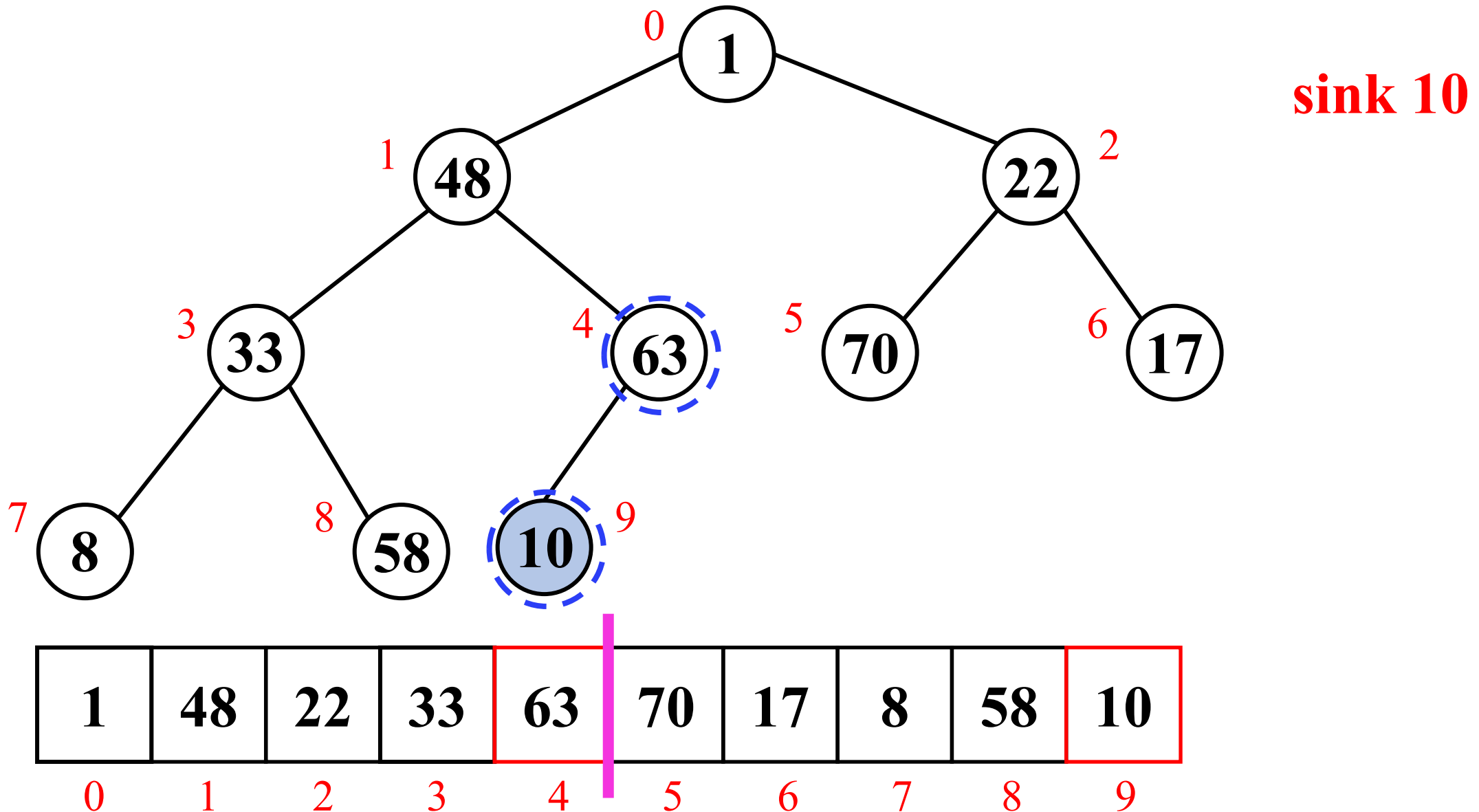
# Heapsort algorithms: array implementation



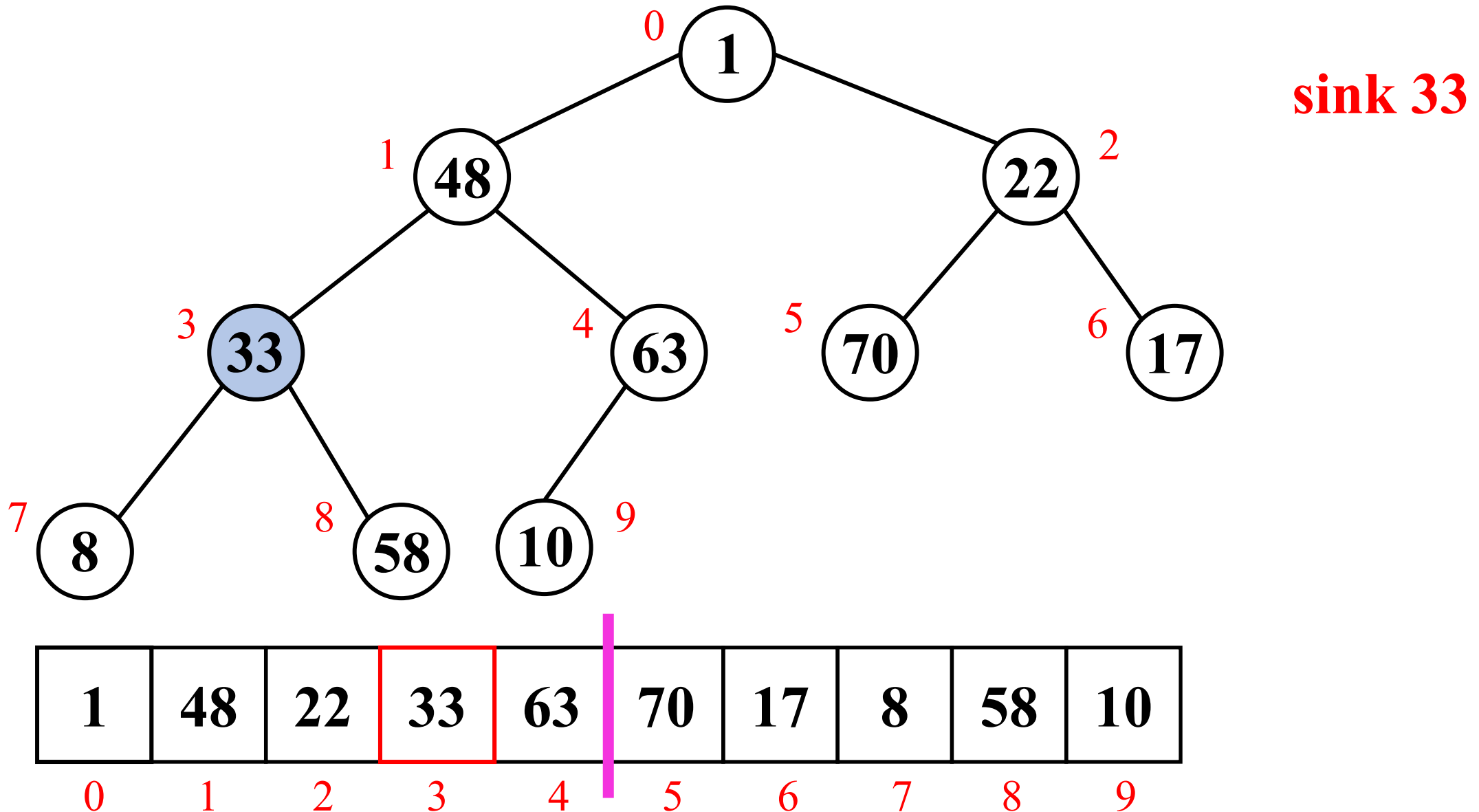
# Heapsort algorithms: array implementation



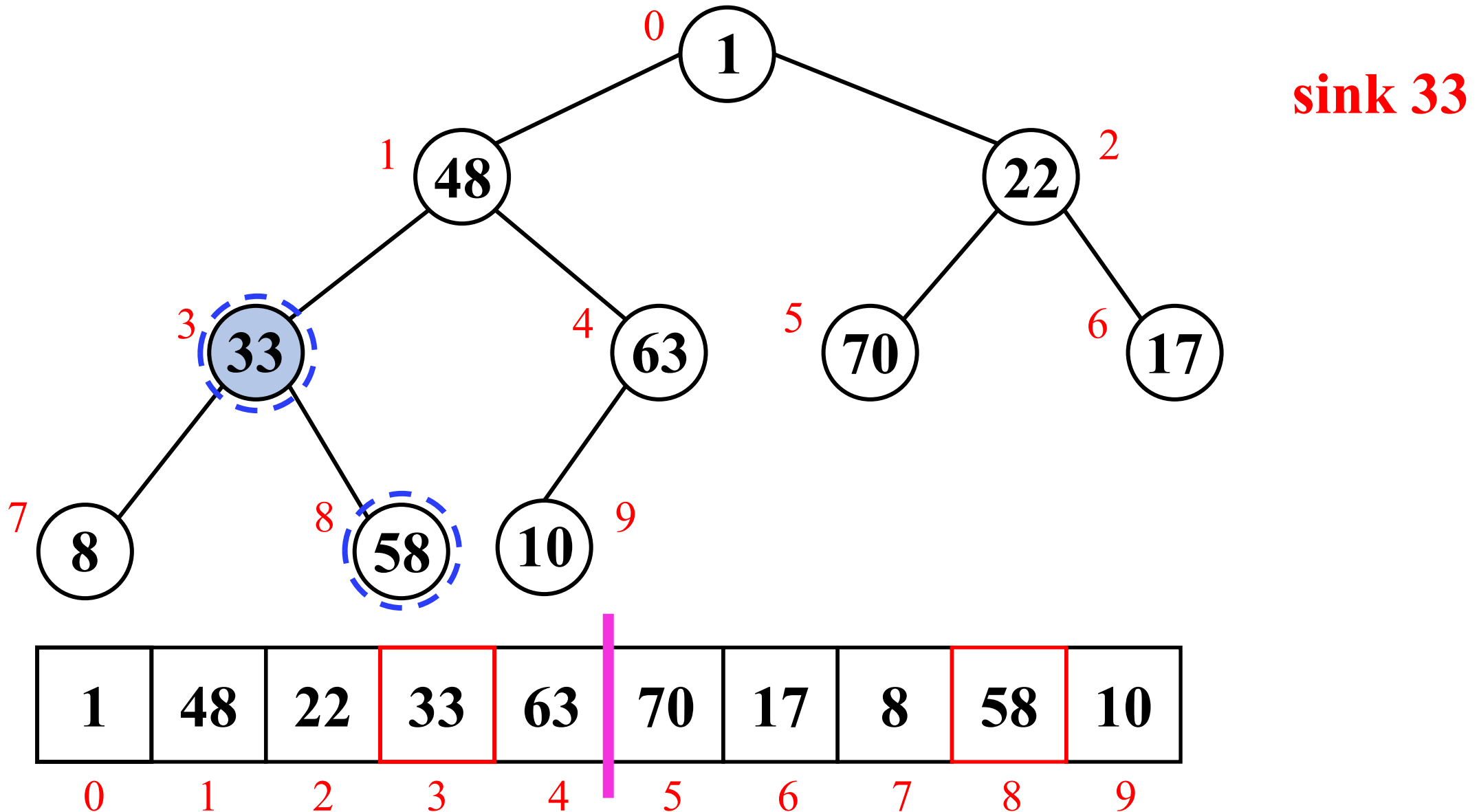
# Heapsort algorithms: array implementation



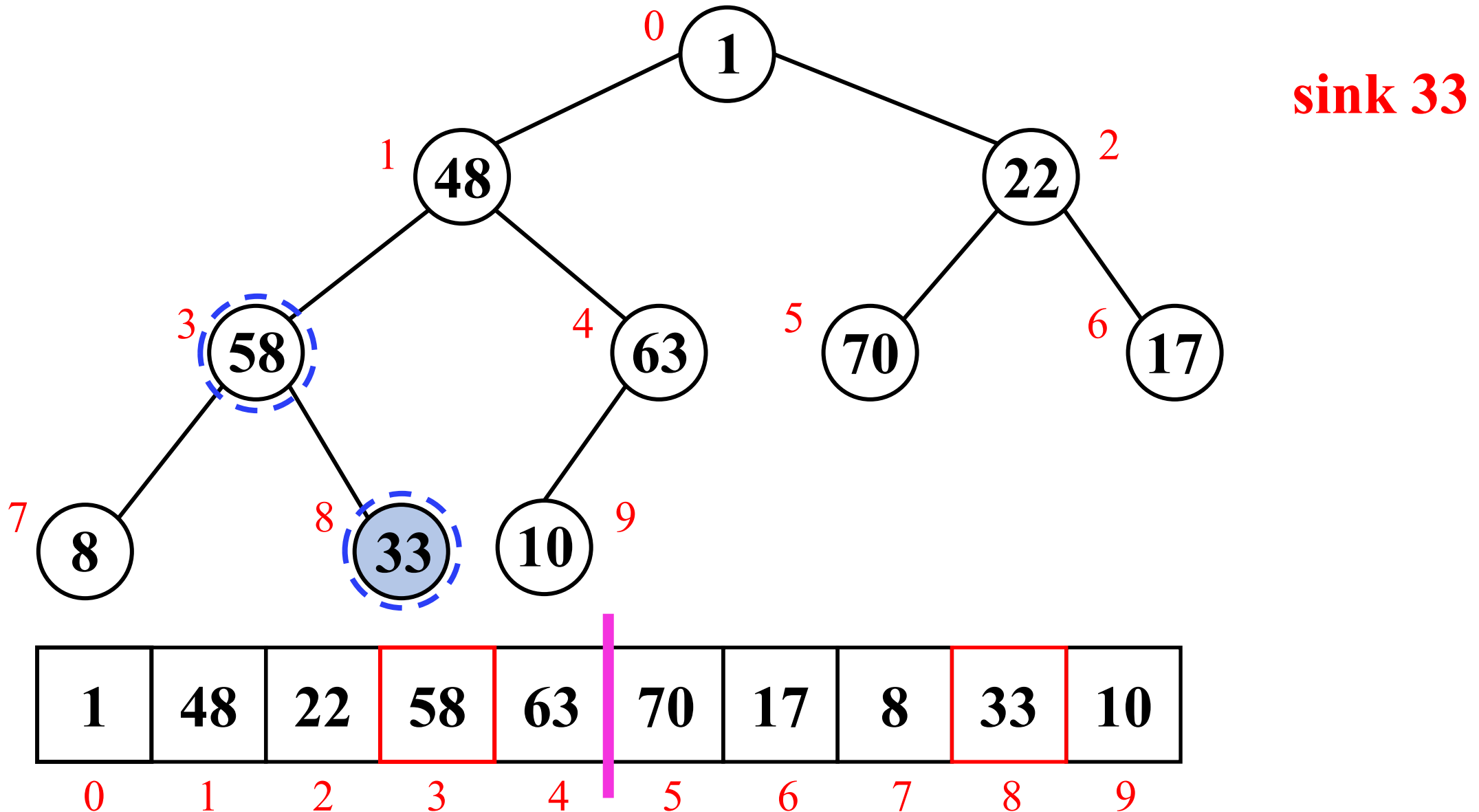
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

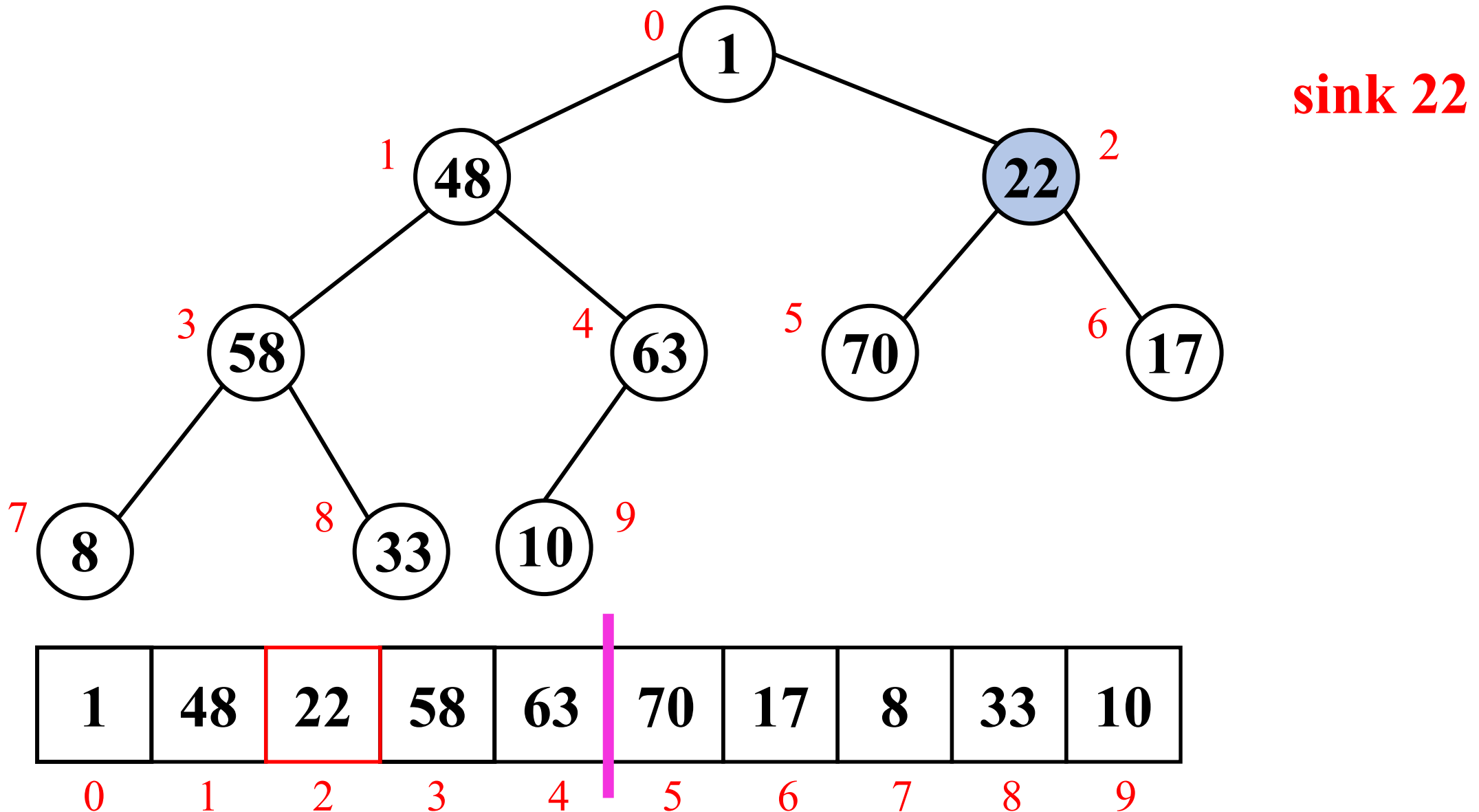


# Heapsort algorithms: array implementation

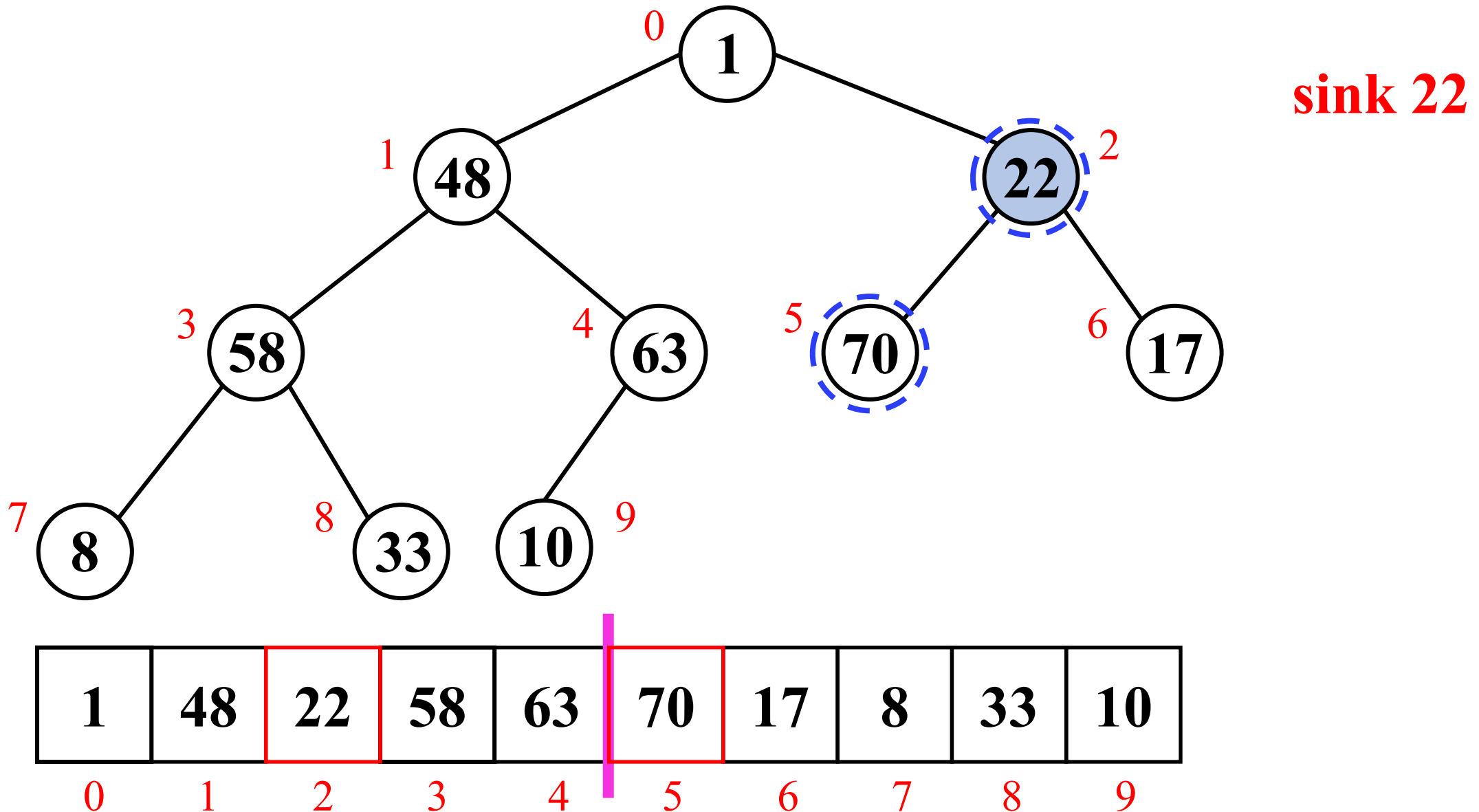




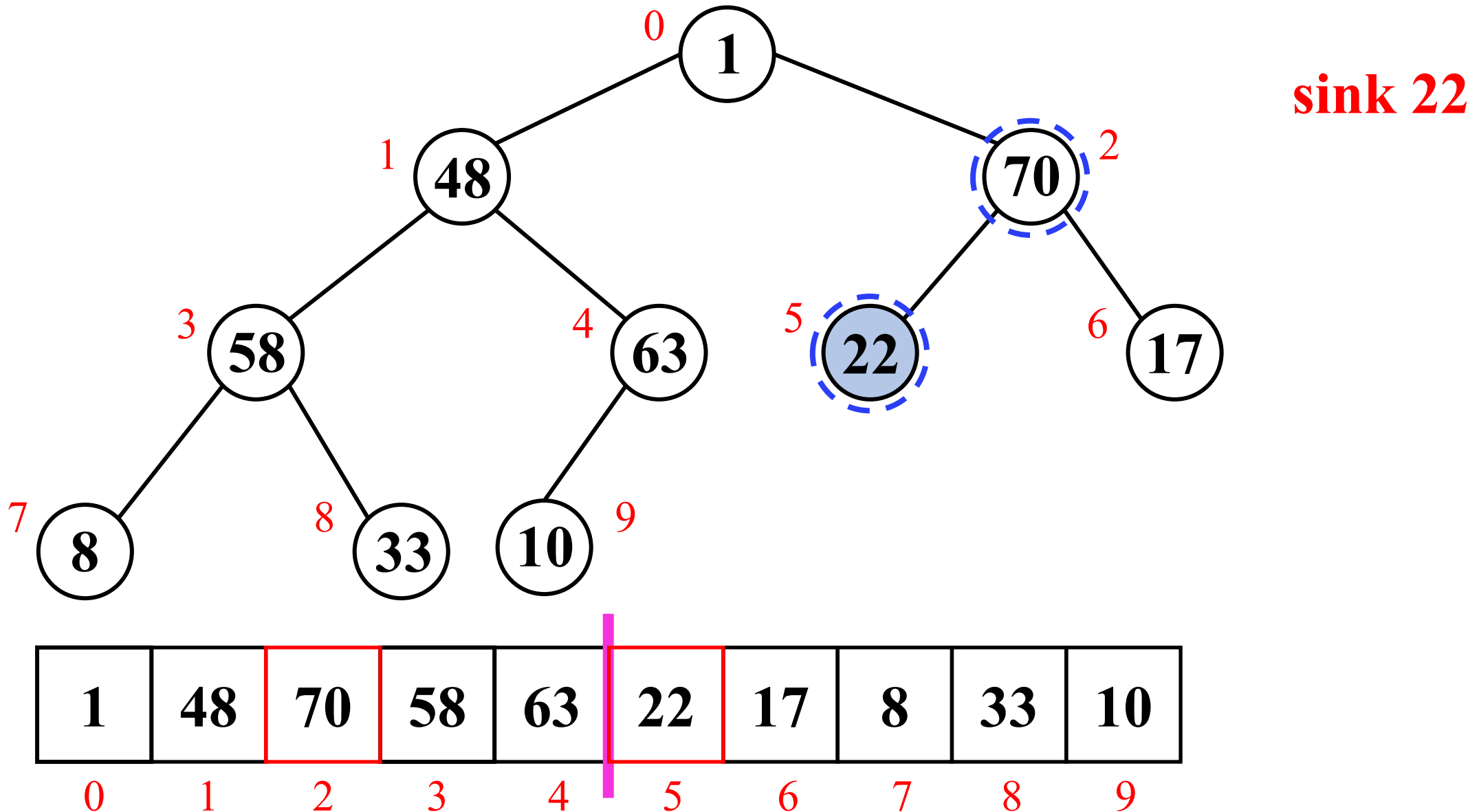
# Heapsort algorithms: array implementation



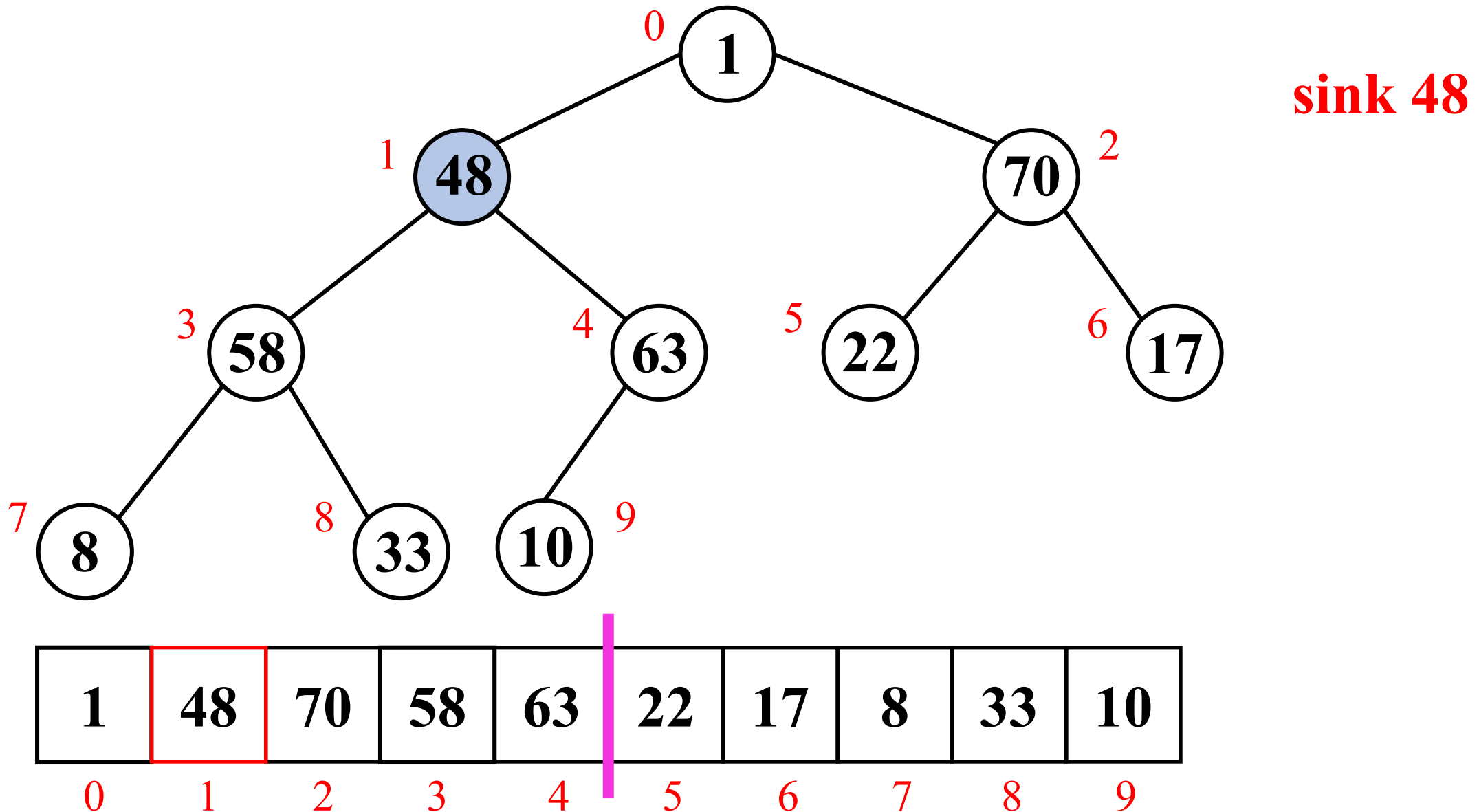
# Heapsort algorithms: array implementation



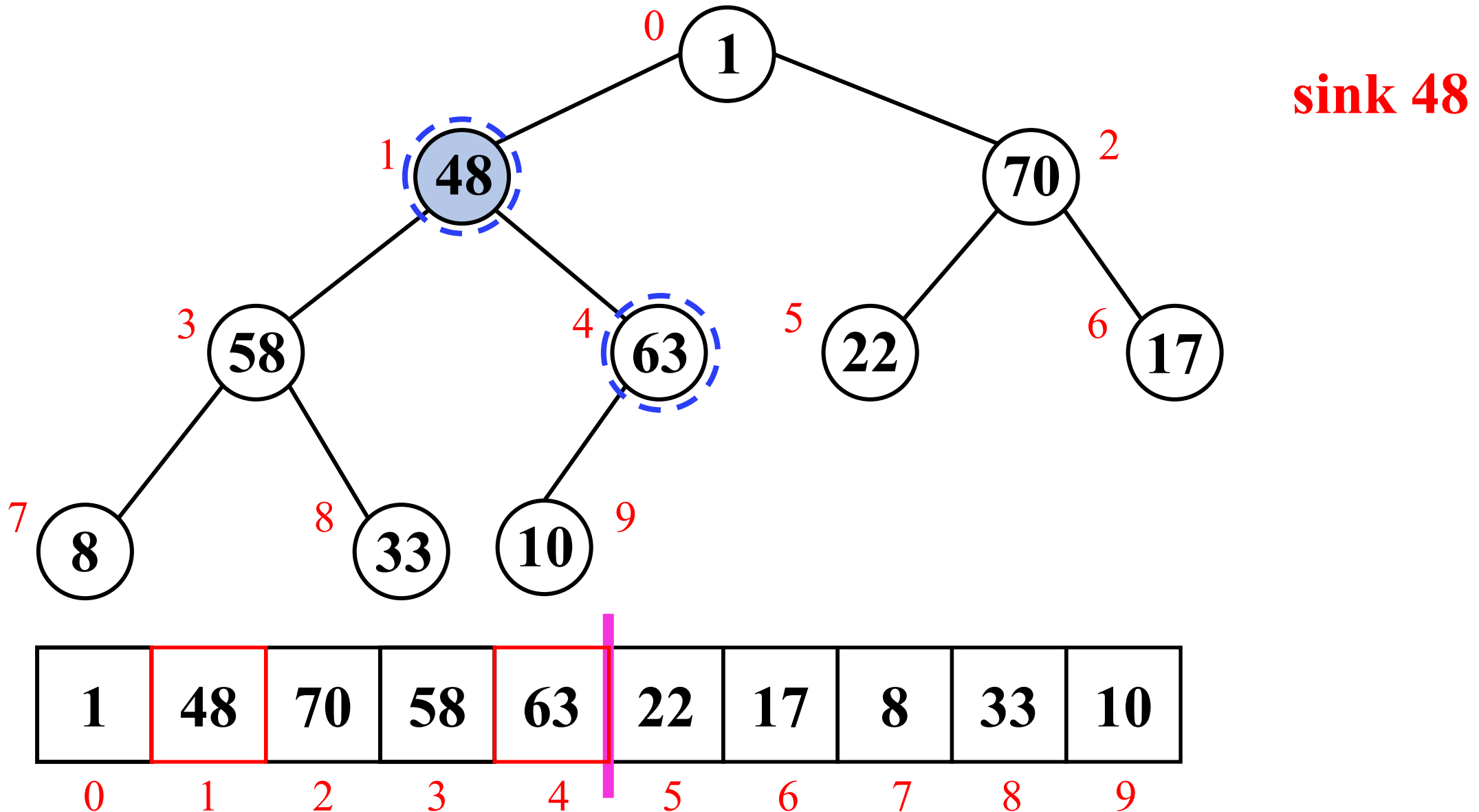
# Heapsort algorithms: array implementation



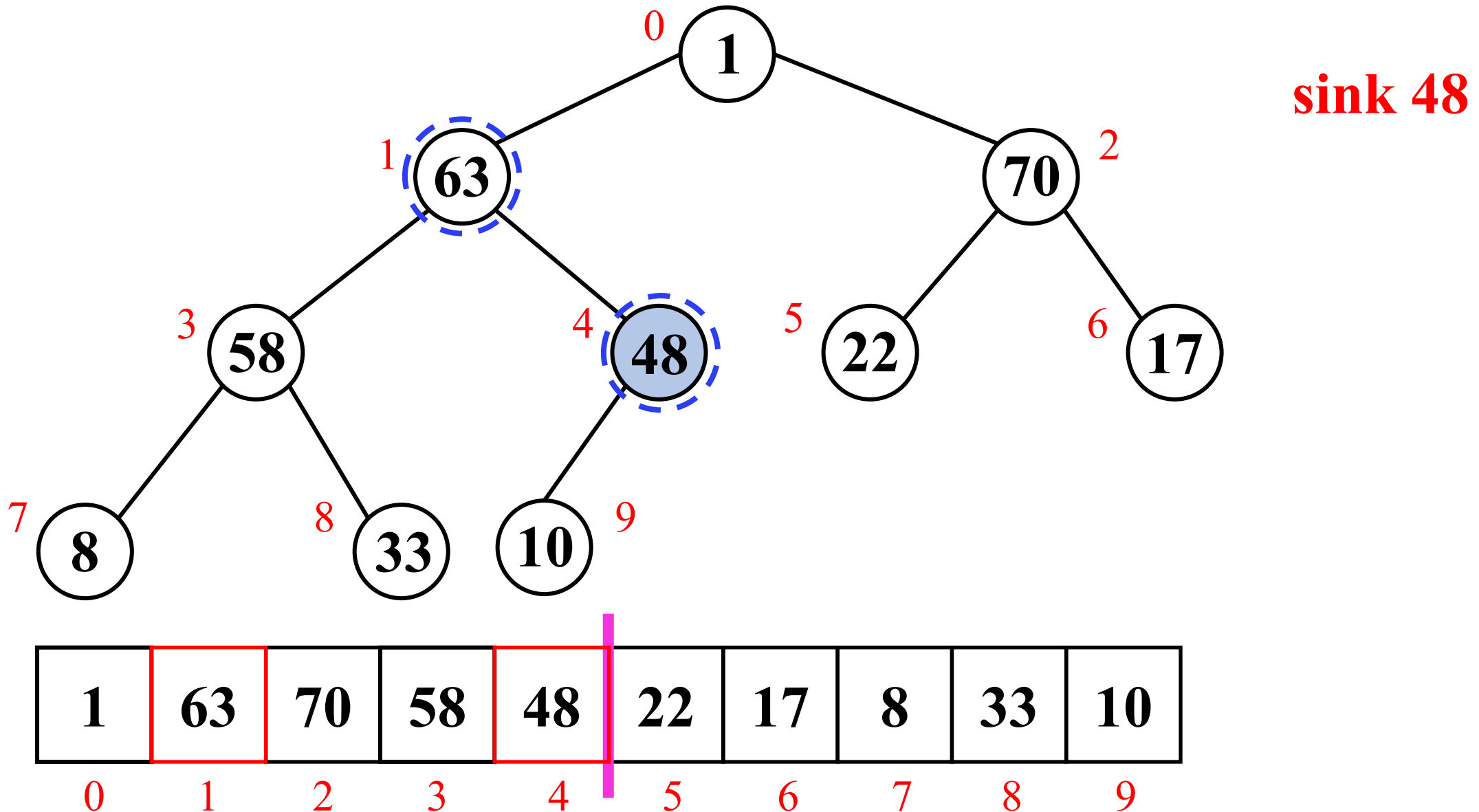
# Heapsort algorithms: array implementation



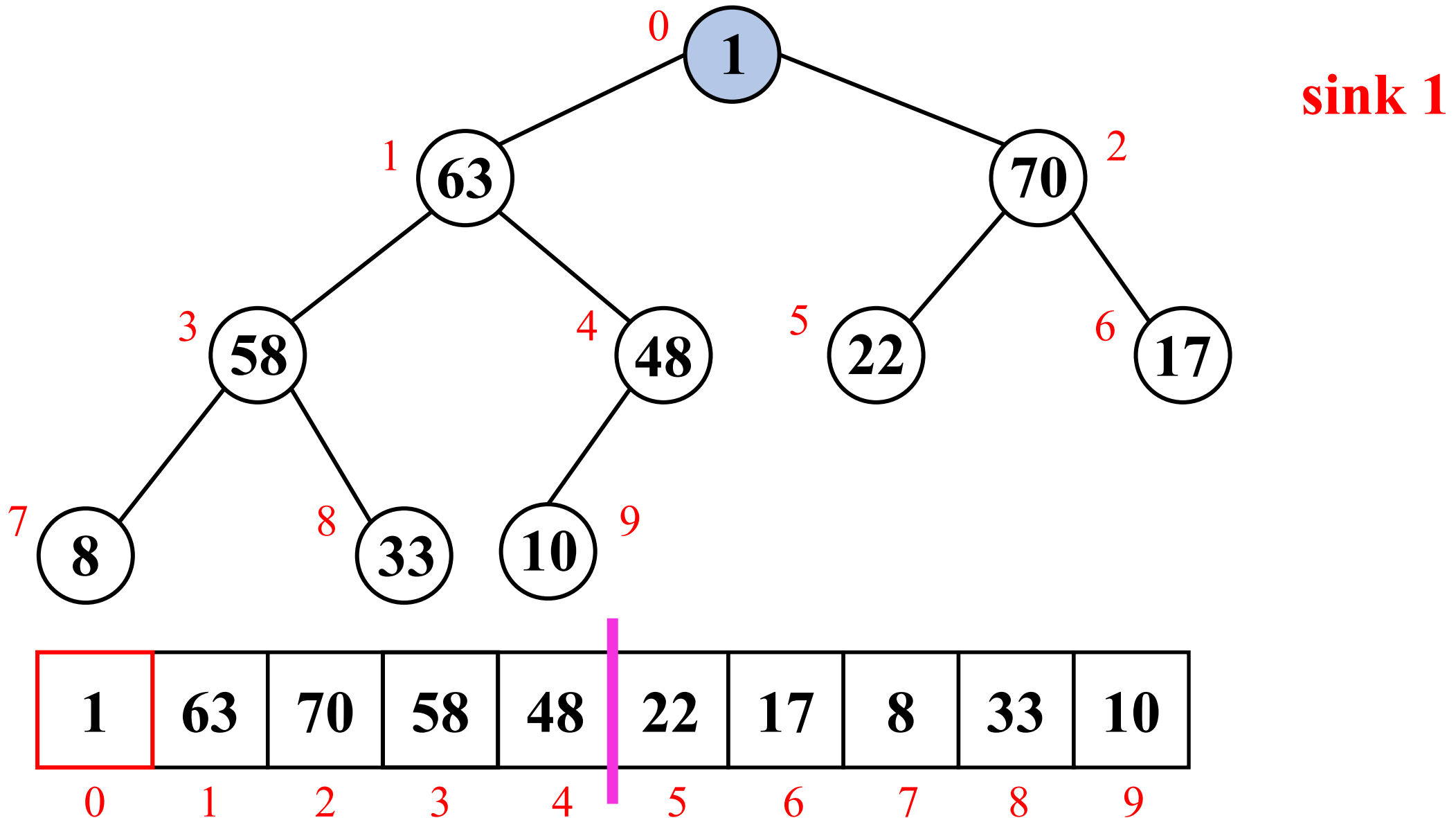
# Heapsort algorithms: array implementation



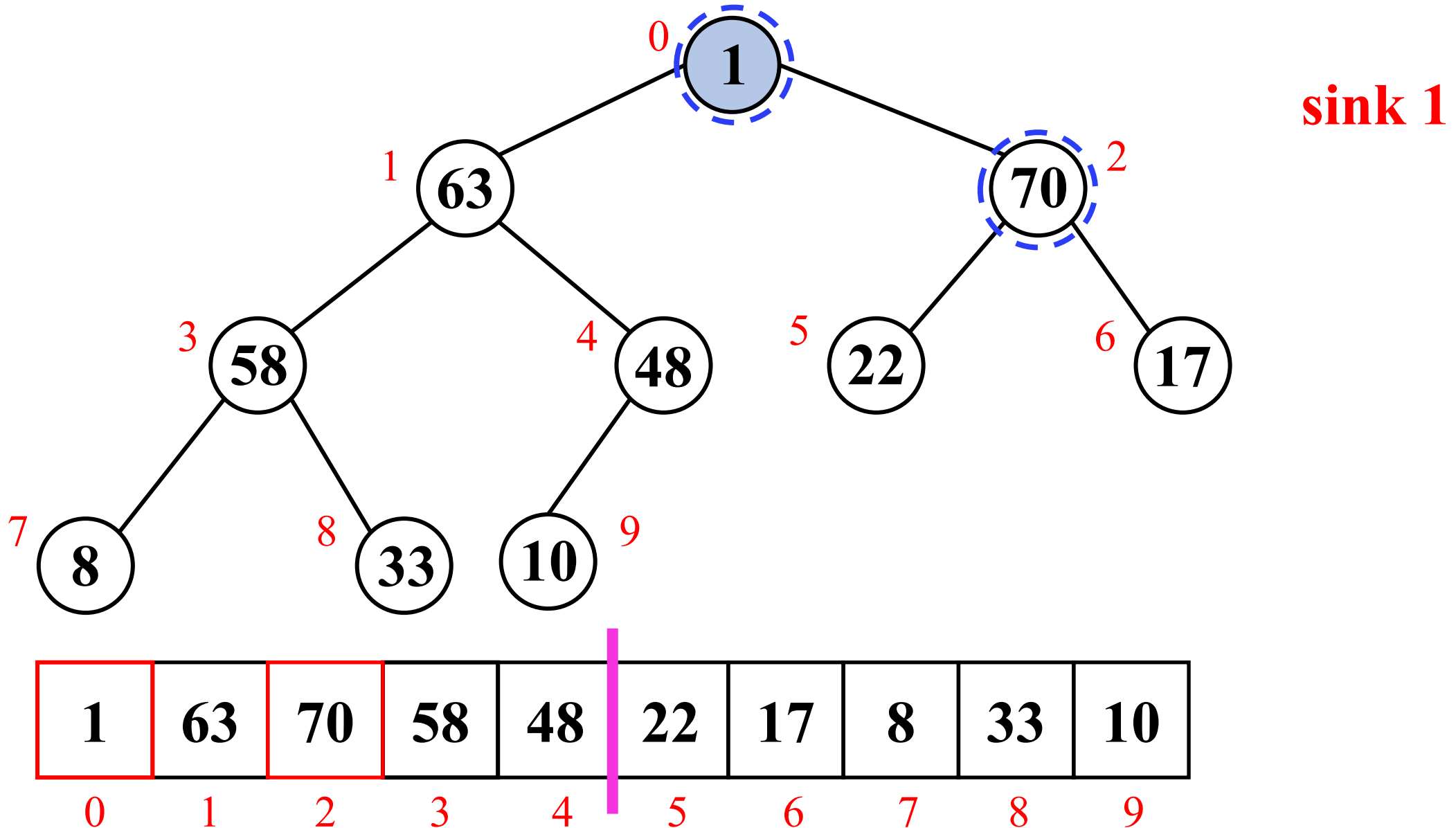
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

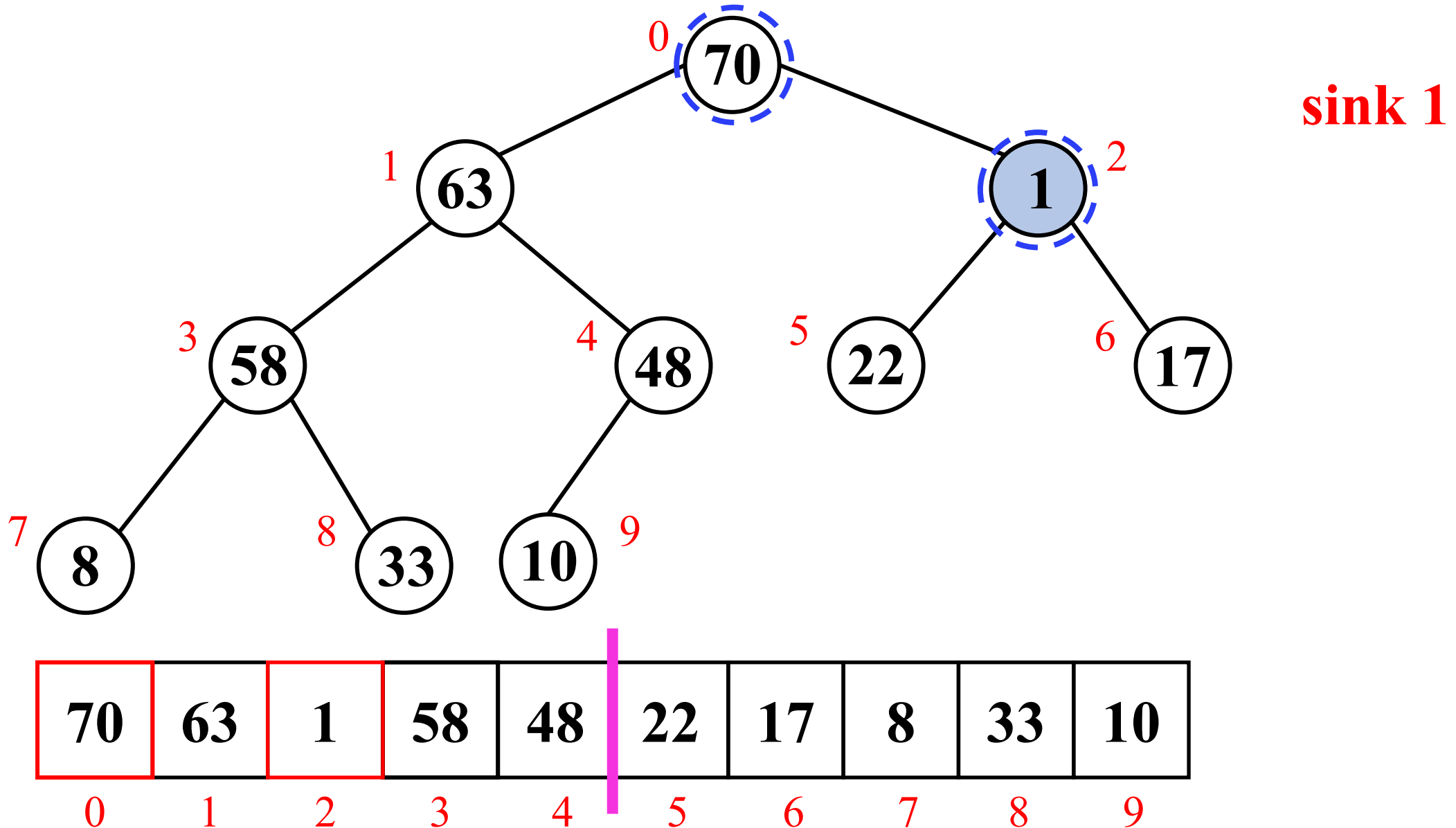


# Heapsort algorithms: array implementation

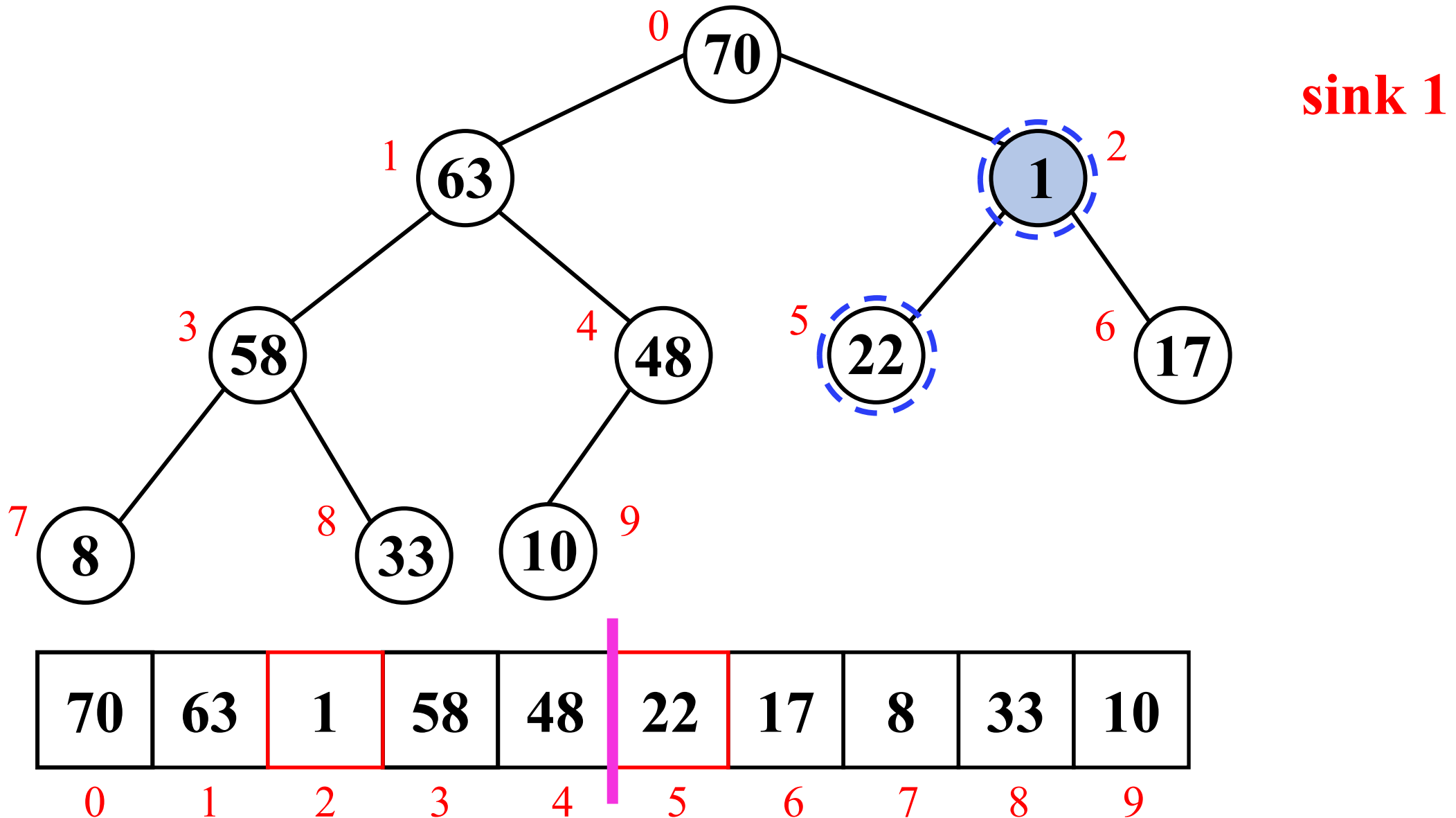




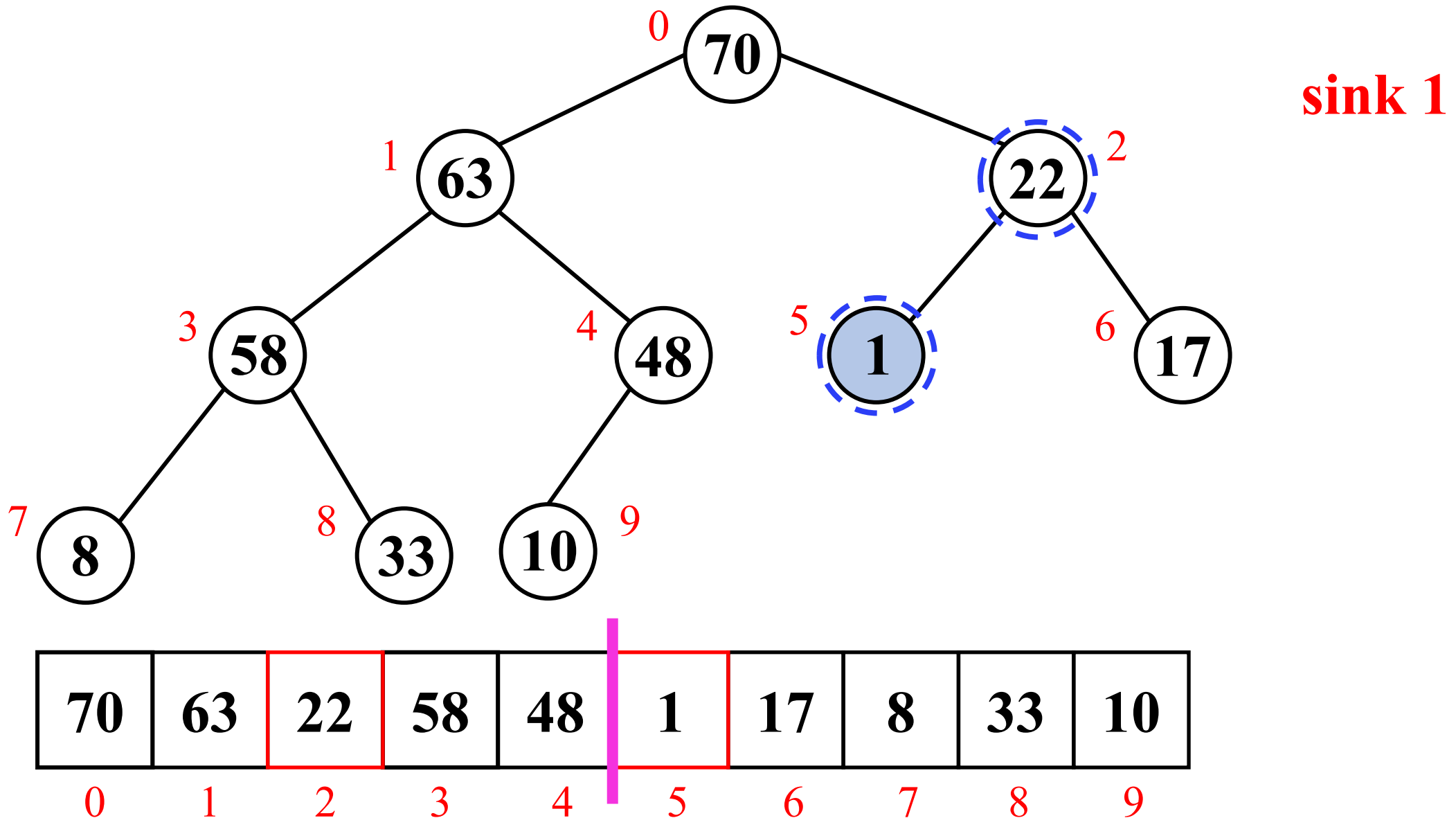
# Heapsort algorithms: array implementation



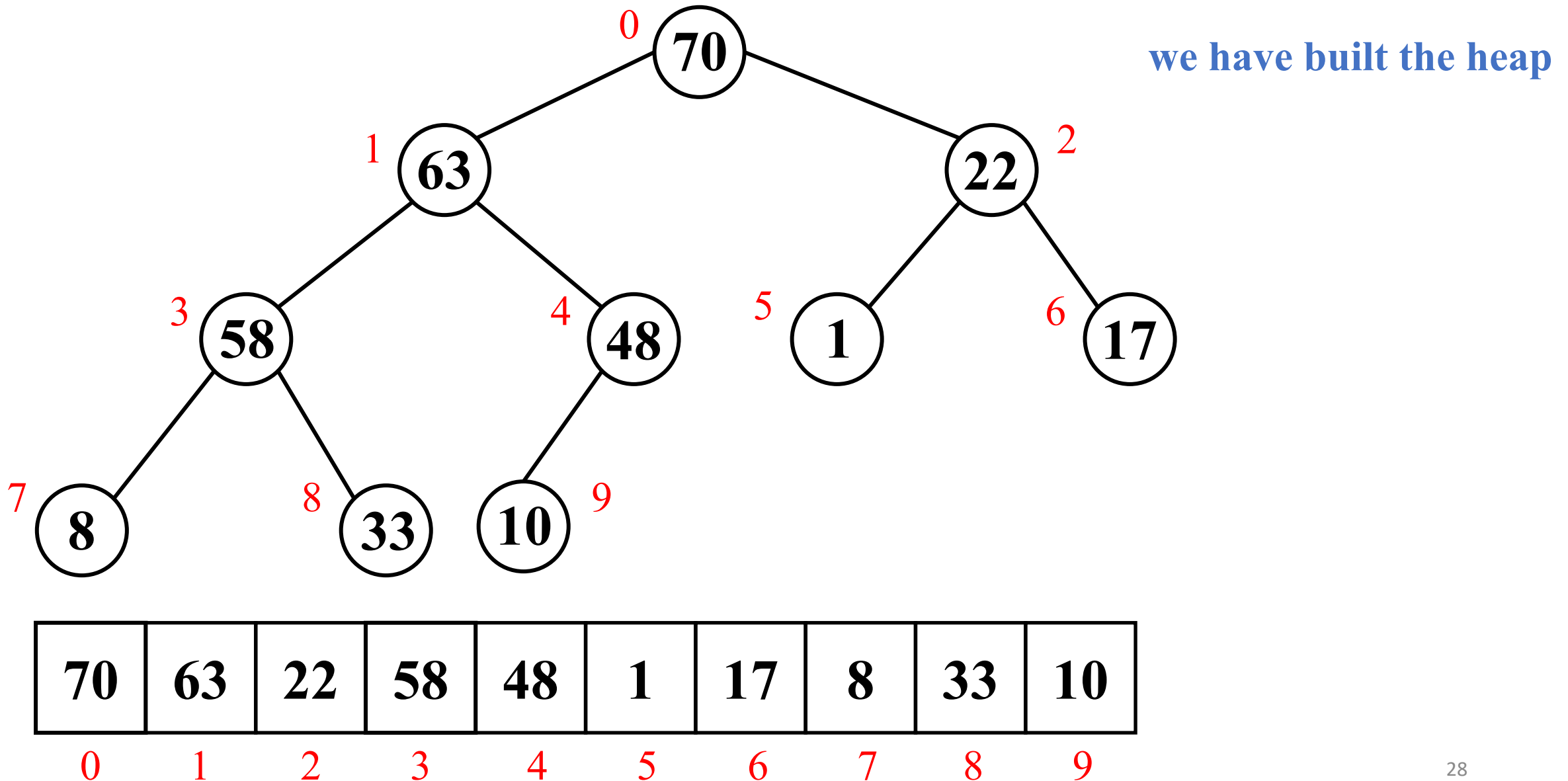
# Heapsort algorithms: array implementation



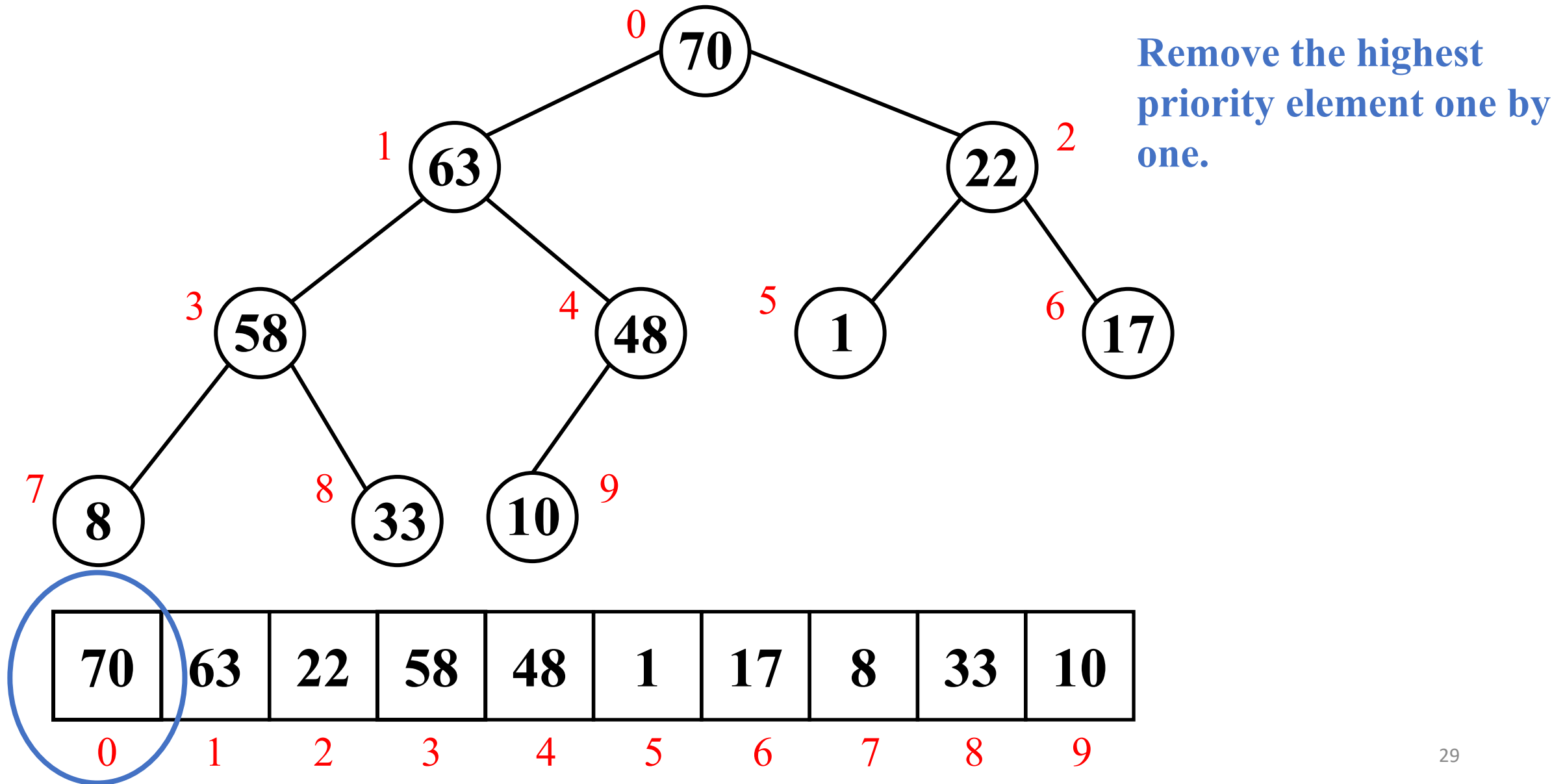
# Heapsort algorithms: array implementation



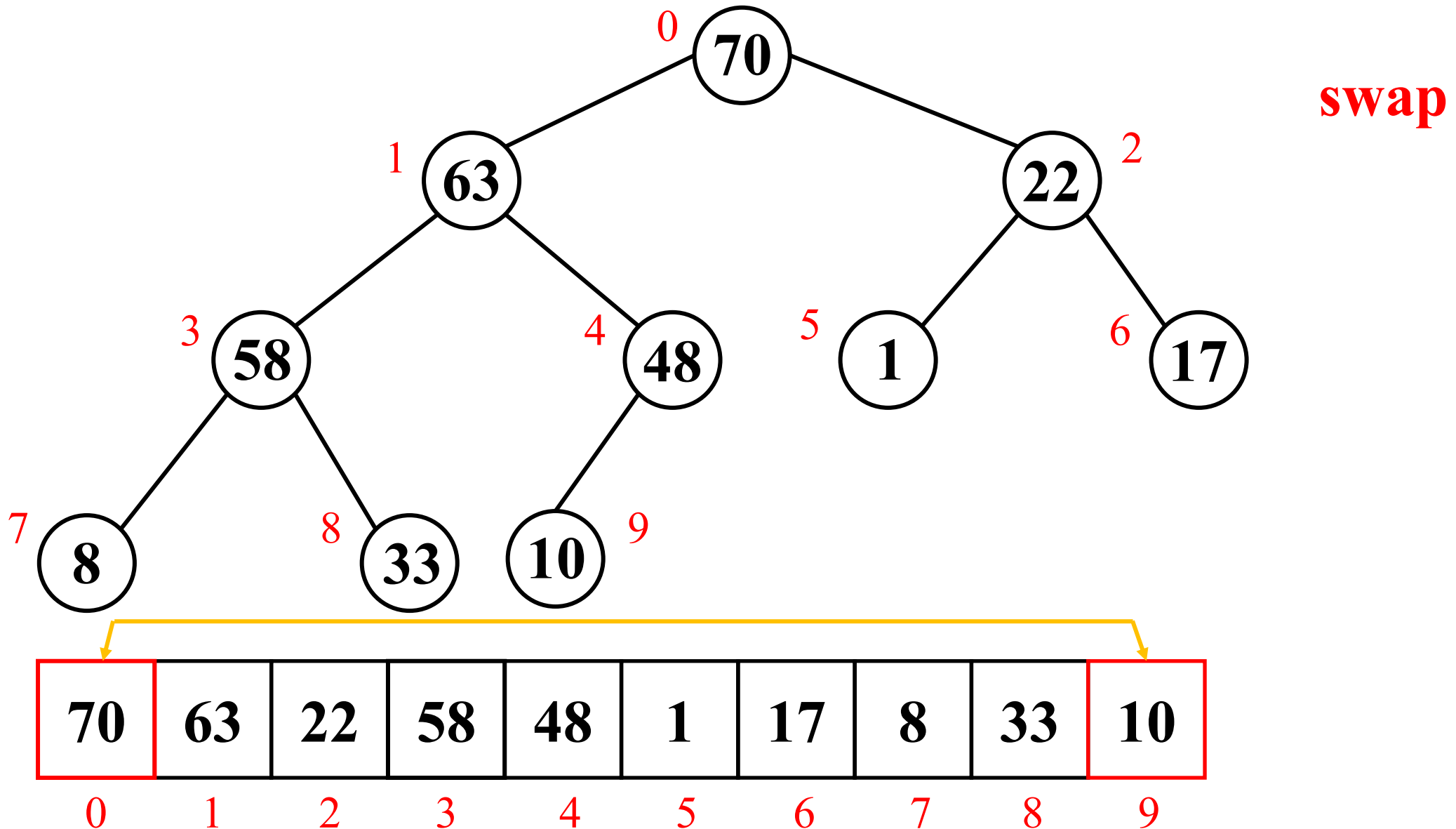
# Heapsort algorithms: array implementation



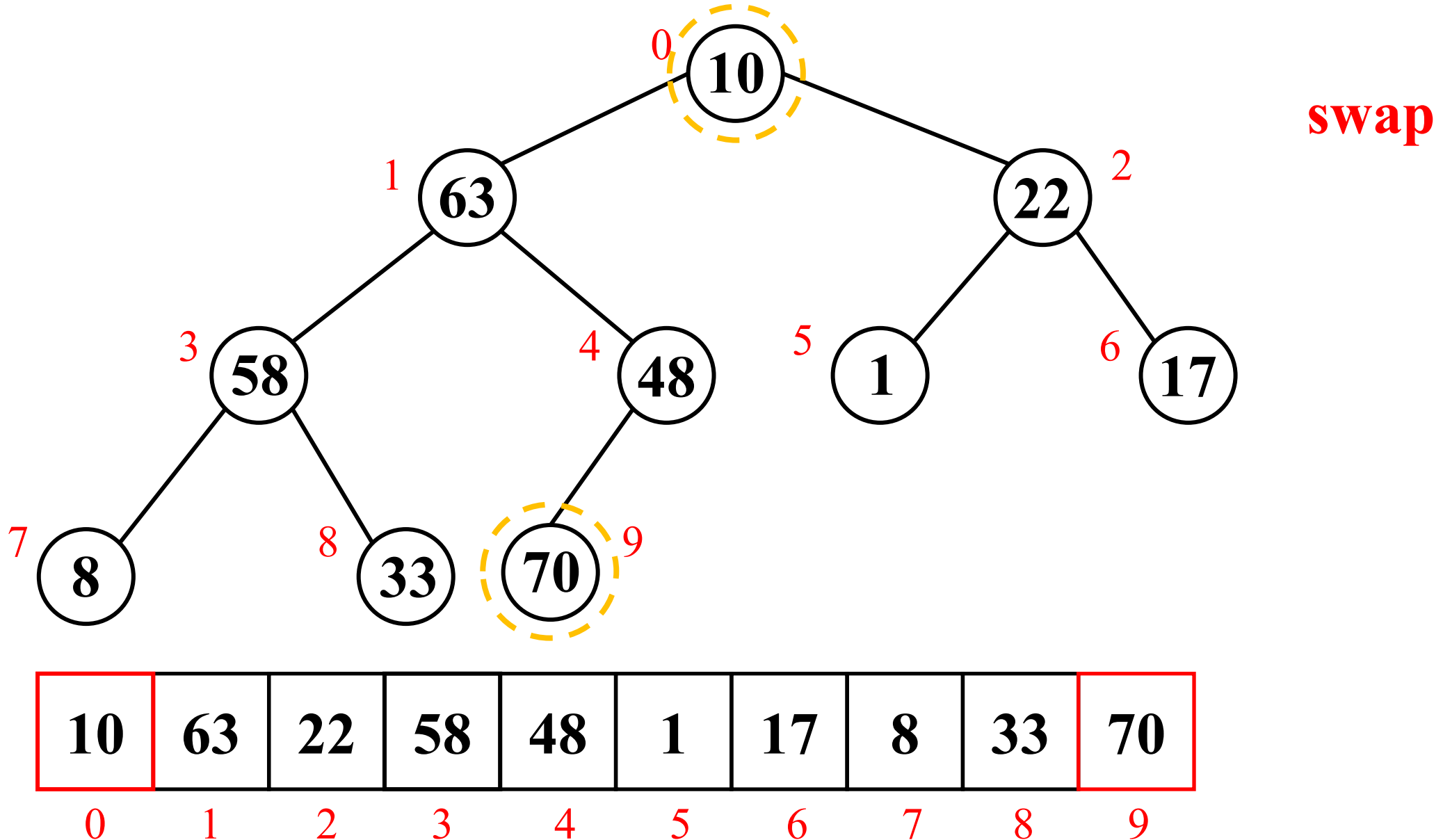
# Heapsort algorithms: array implementation



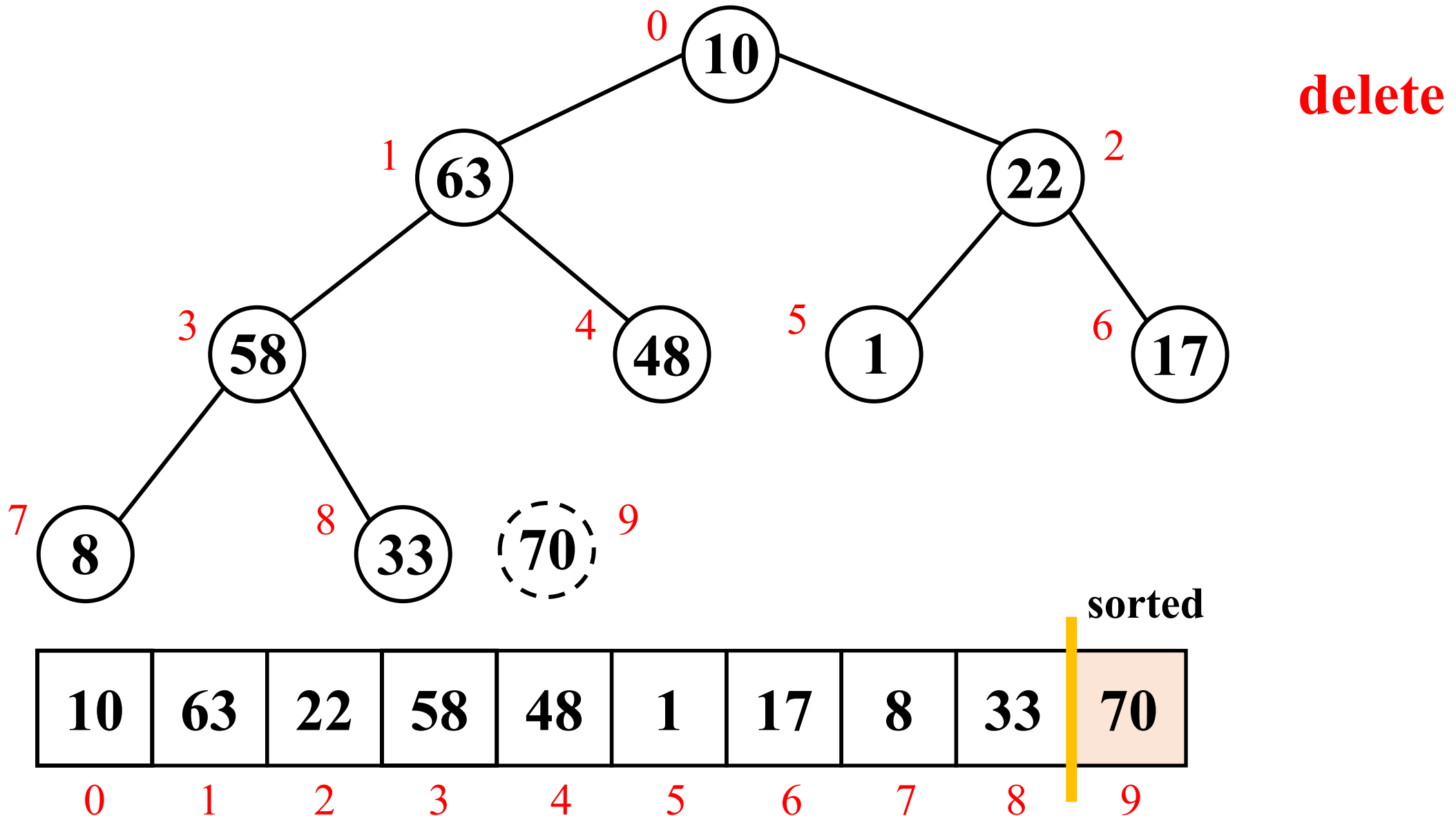
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

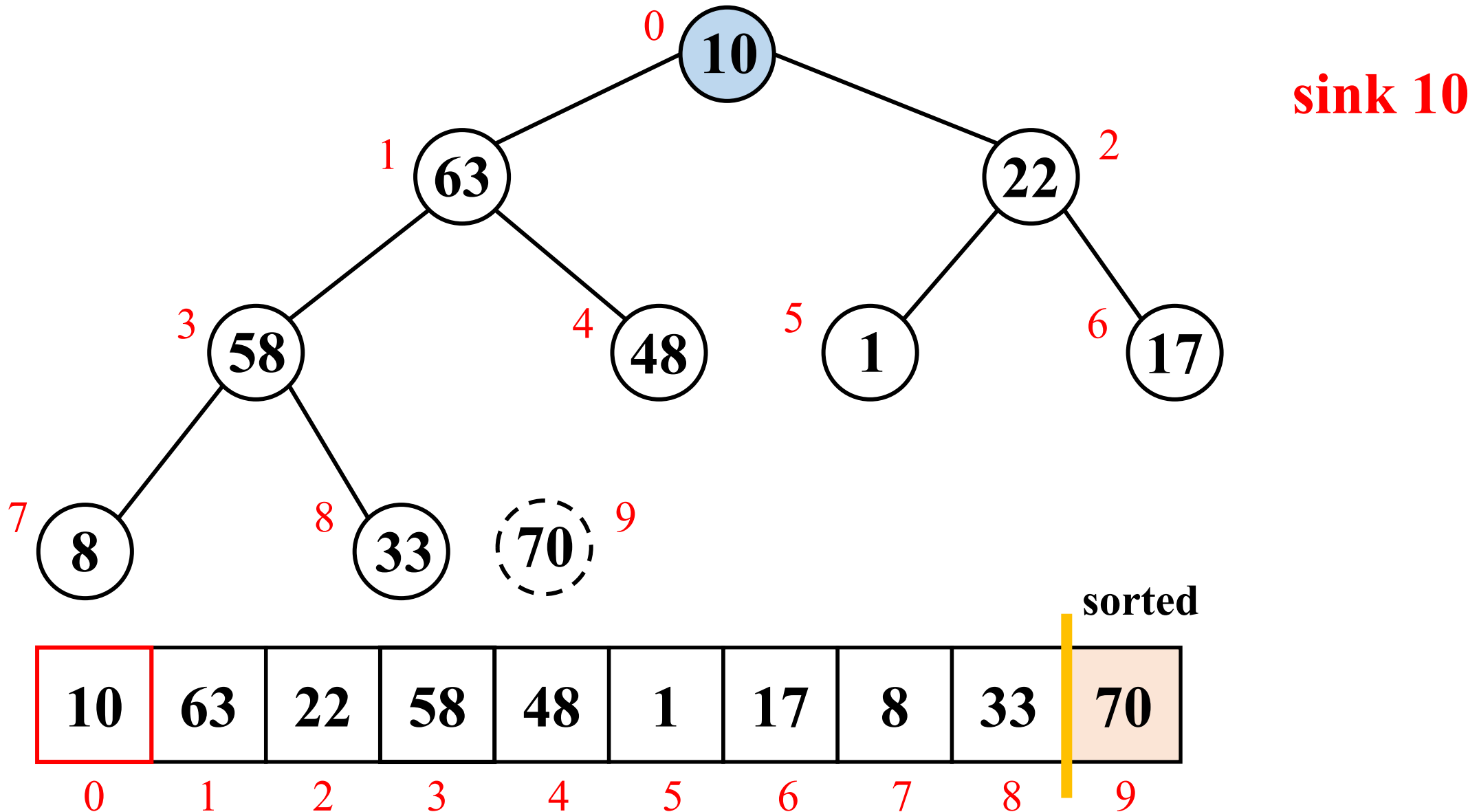


# Heapsort algorithms: array implementation

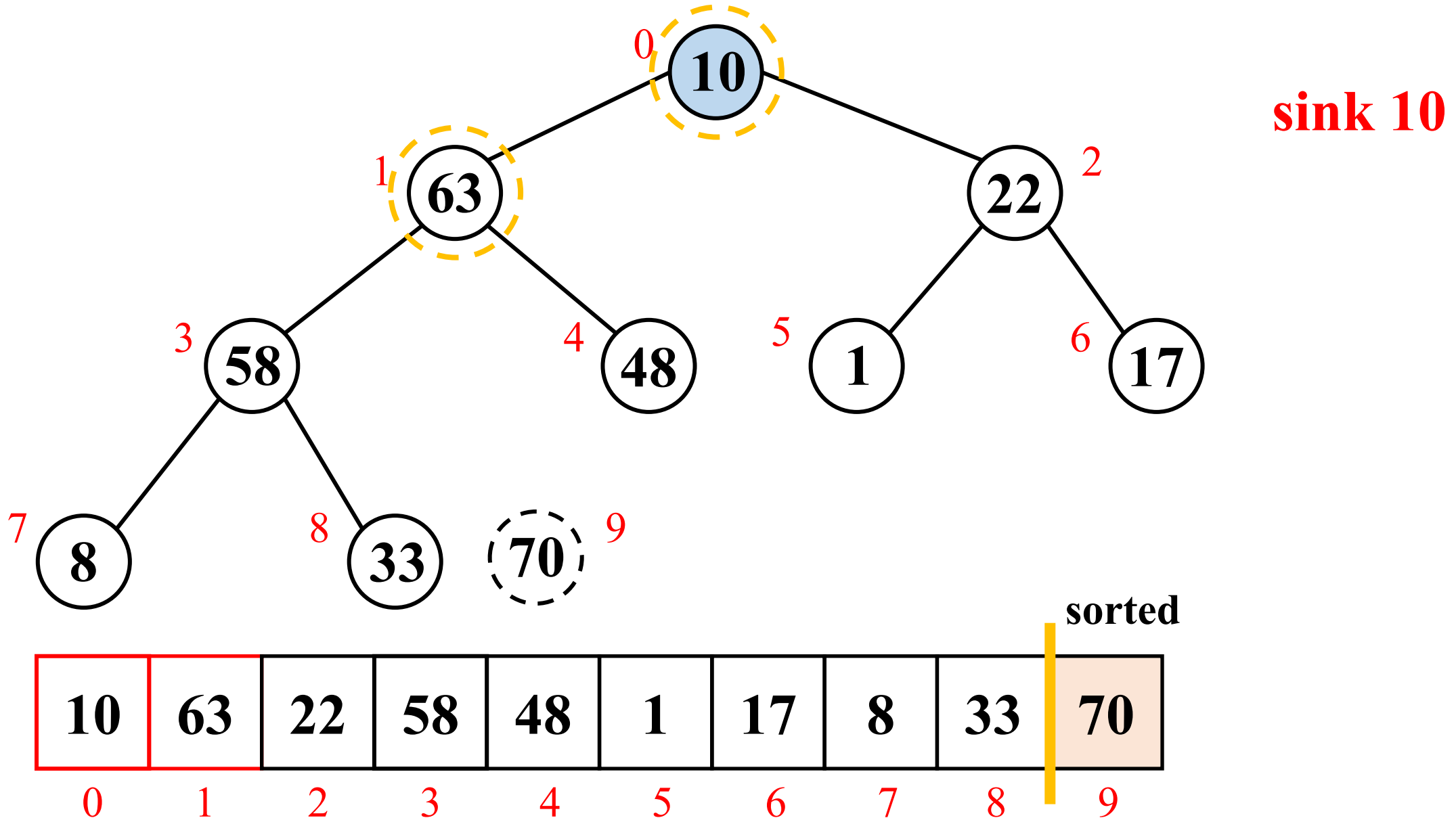




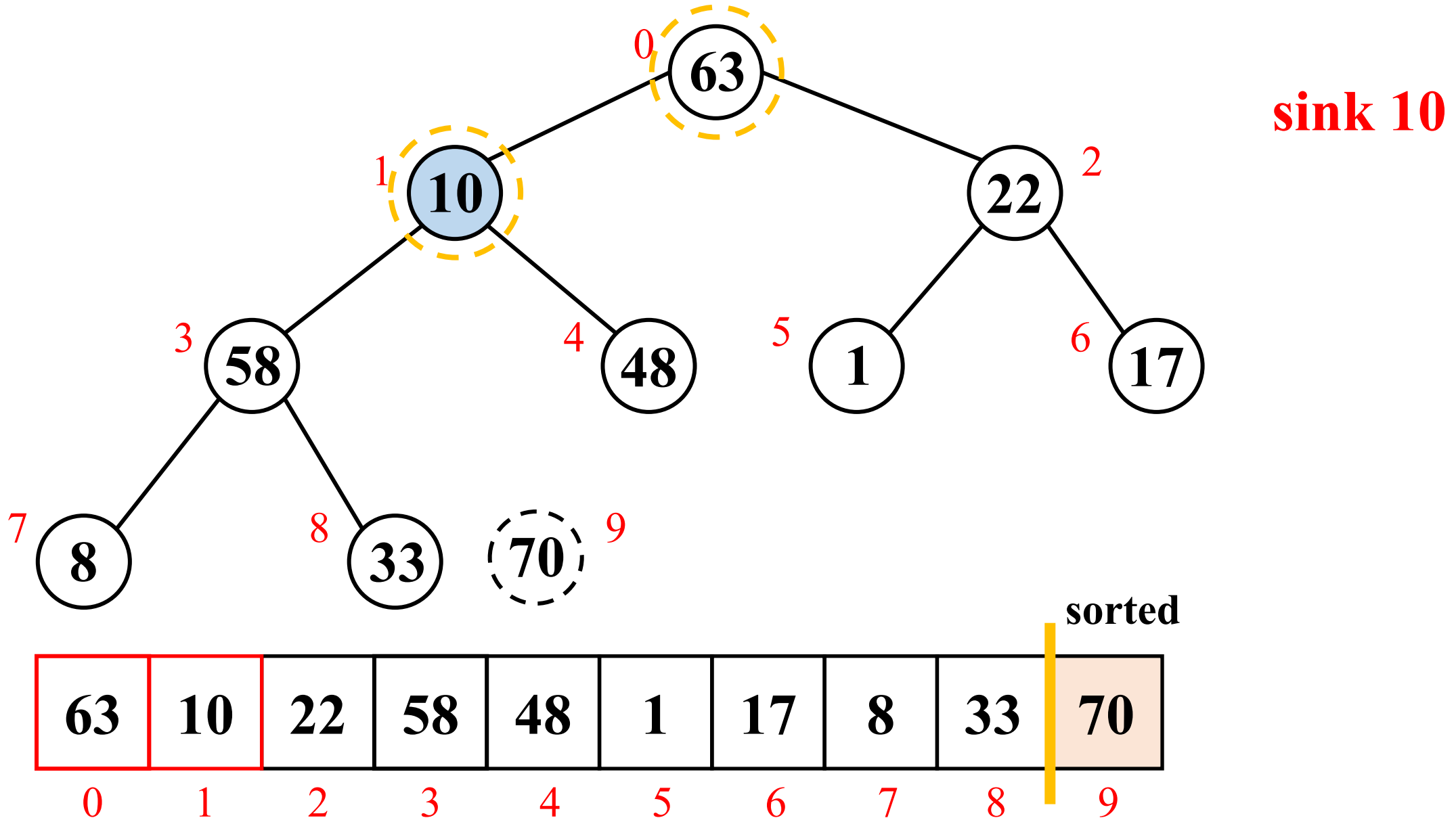
# Heapsort algorithms: array implementation



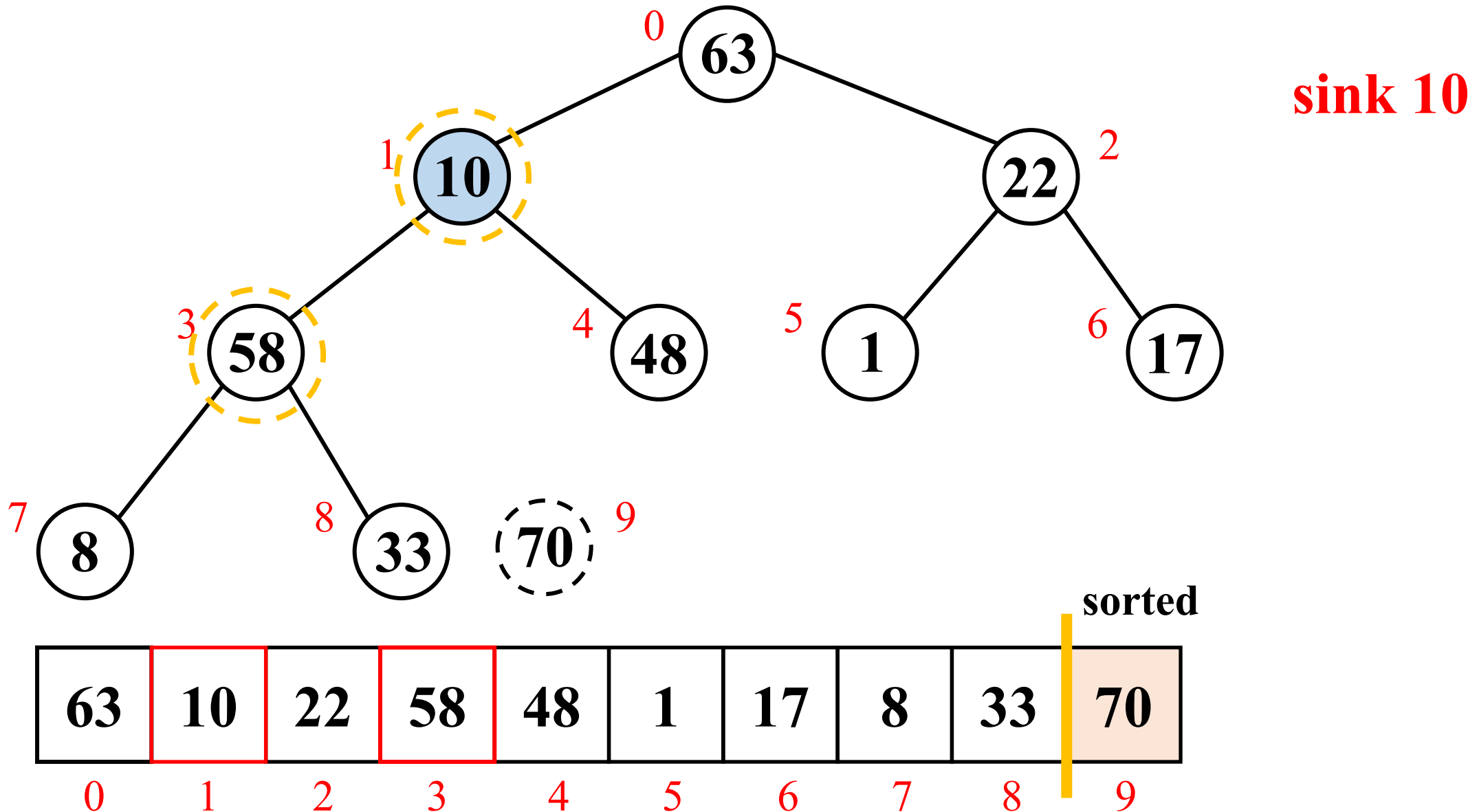
# Heapsort algorithms: array implementation



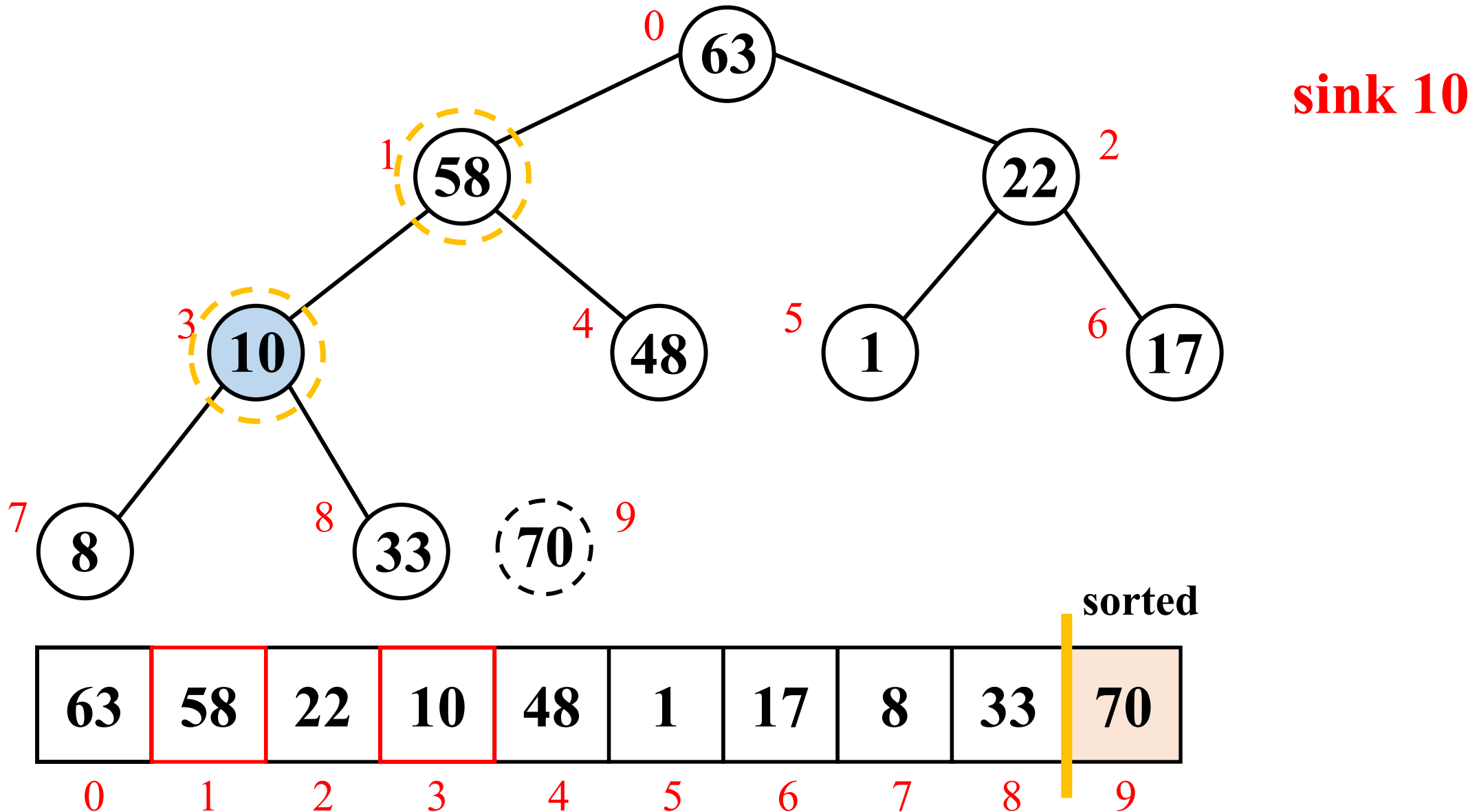
# Heapsort algorithms: array implementation



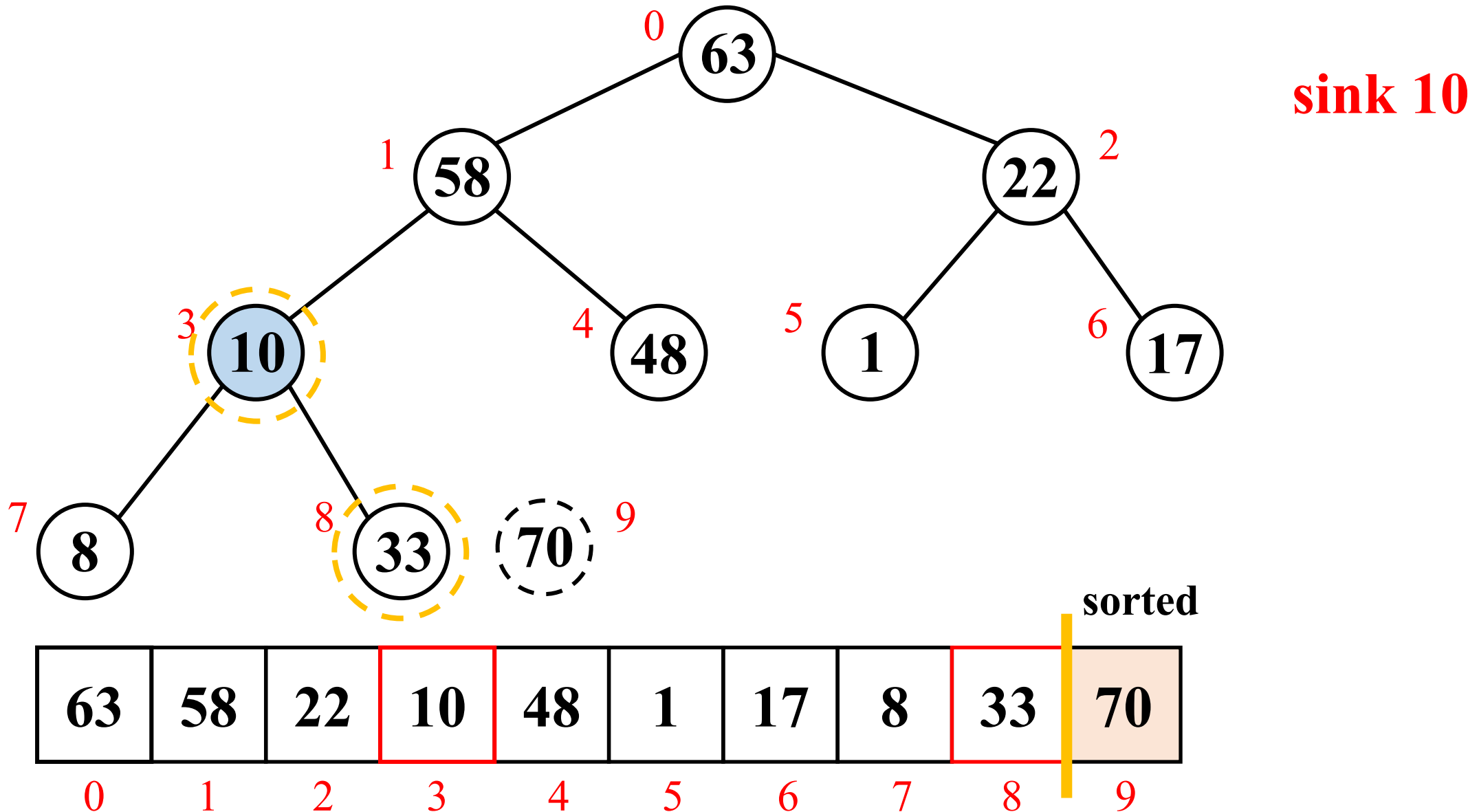
# Heapsort algorithms: array implementation



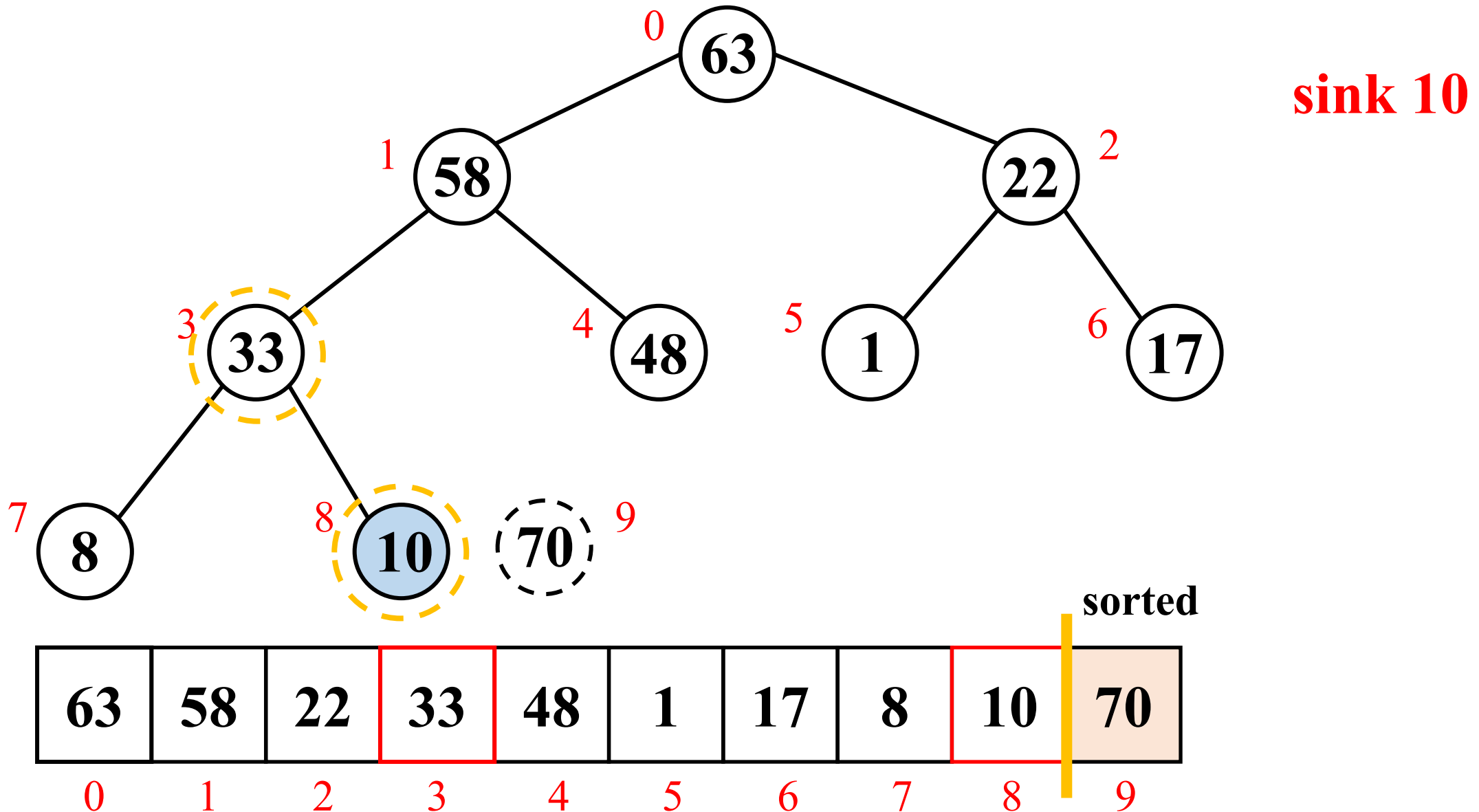
# Heapsort algorithms: array implementation



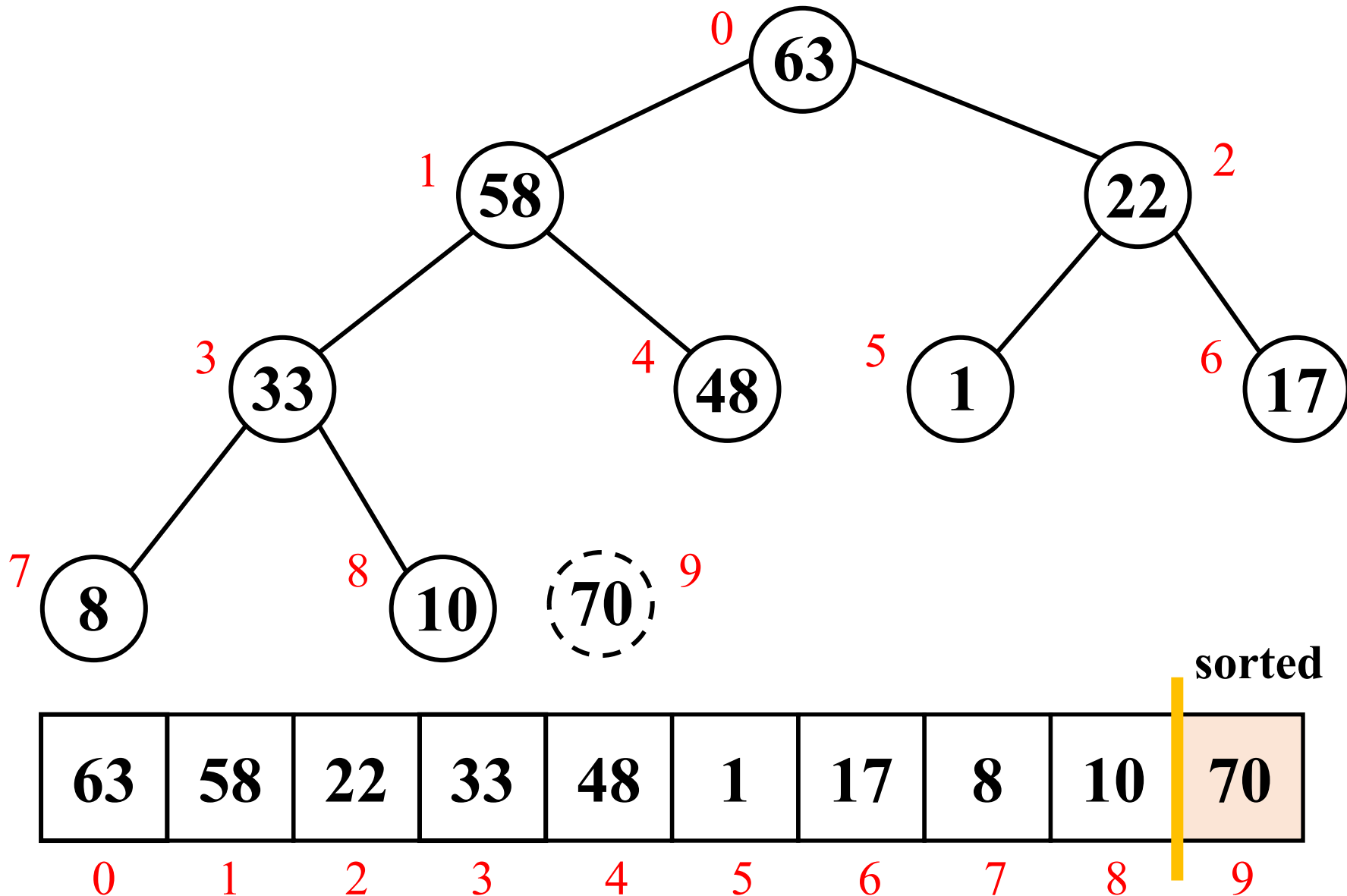
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

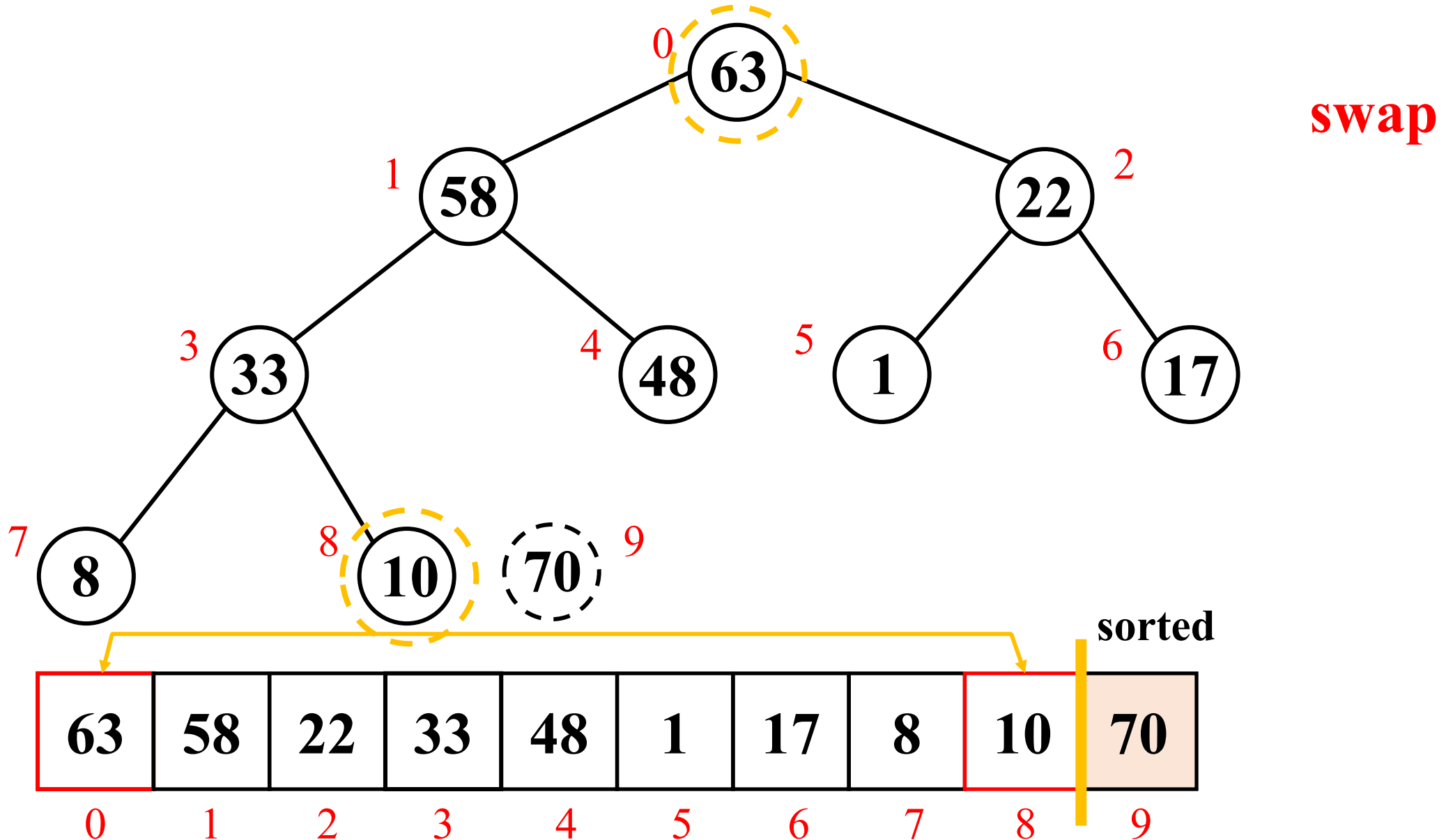


# Heapsort algorithms: array implementation

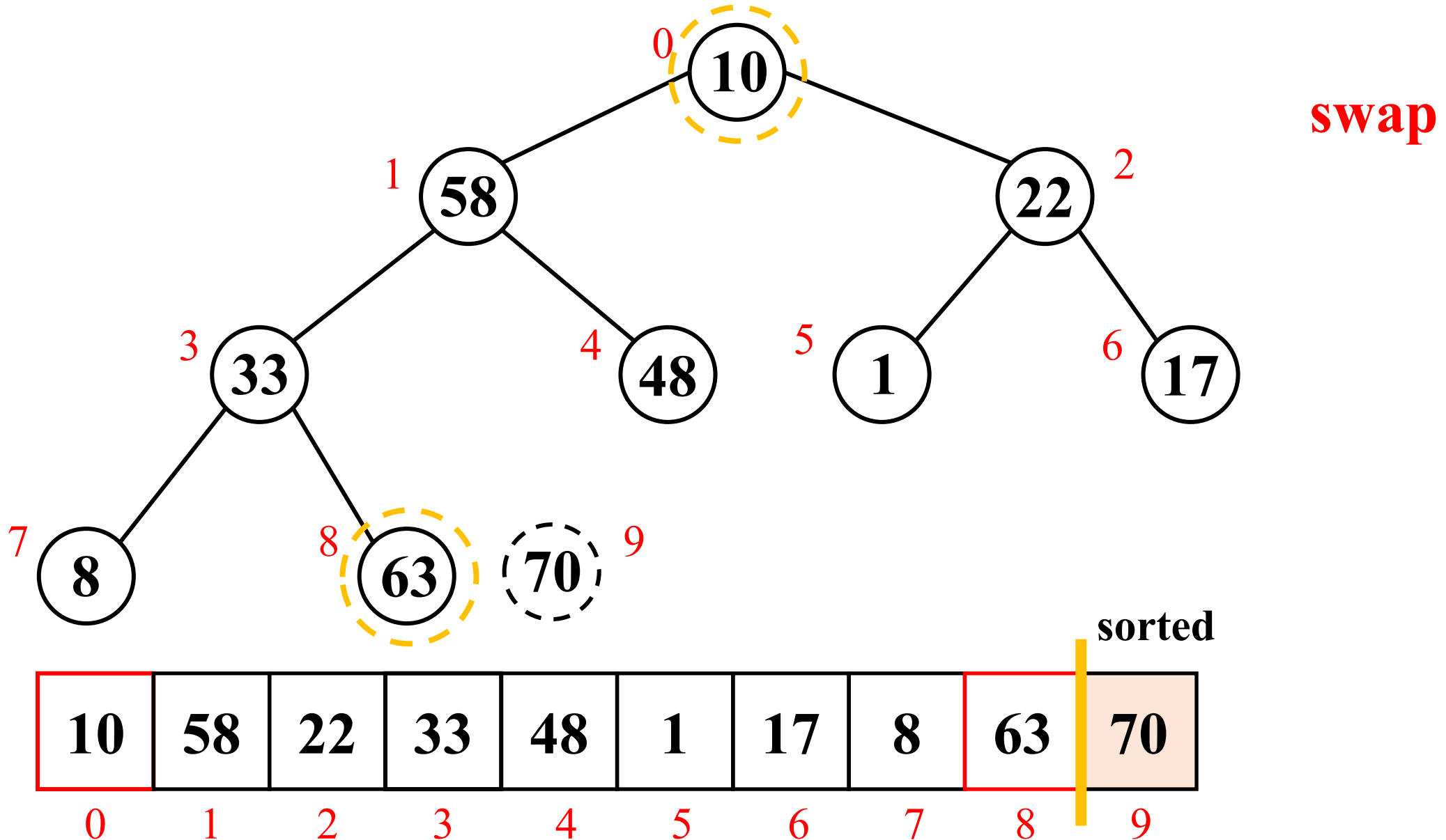




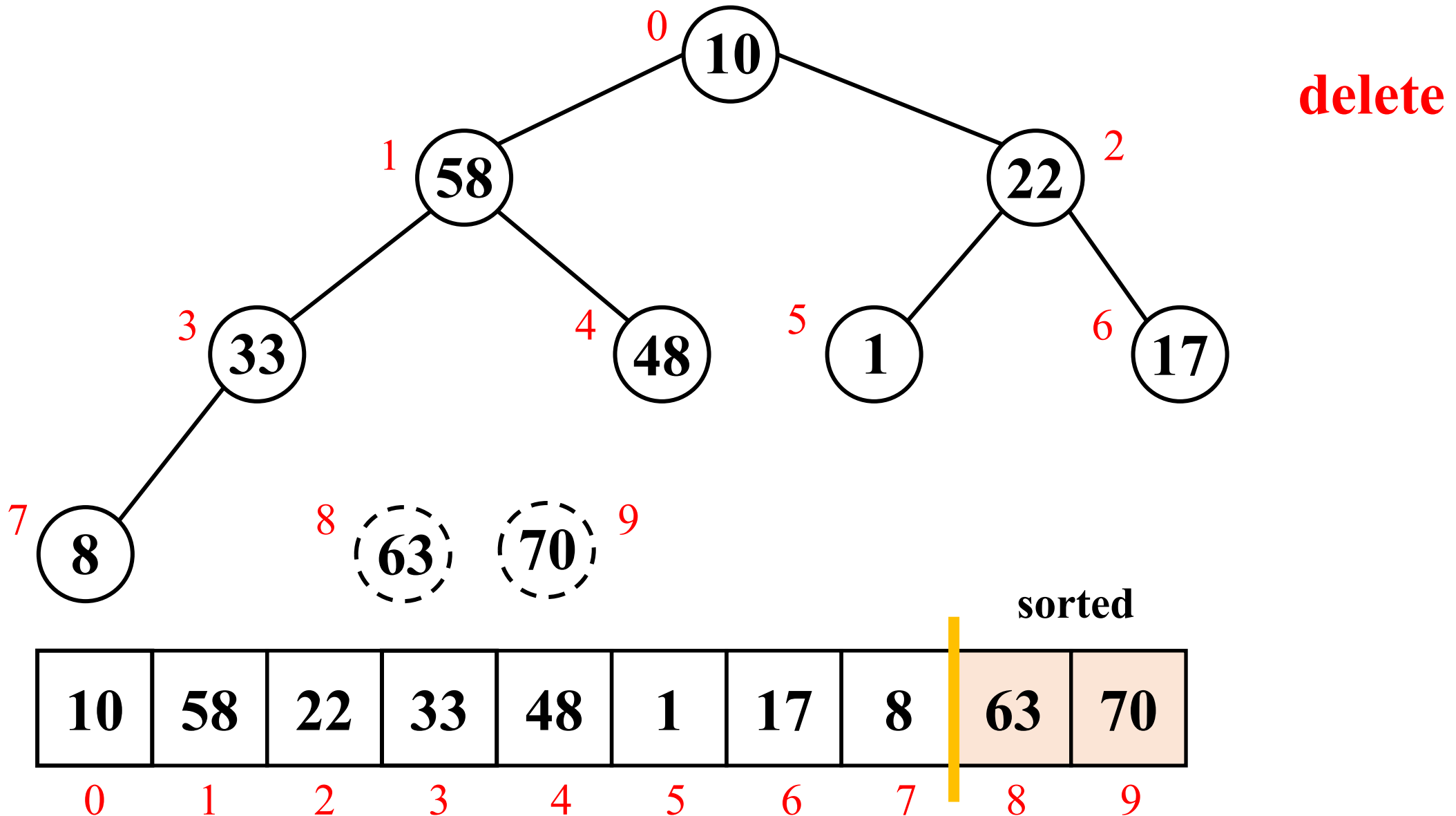
# Heapsort algorithms: array implementation



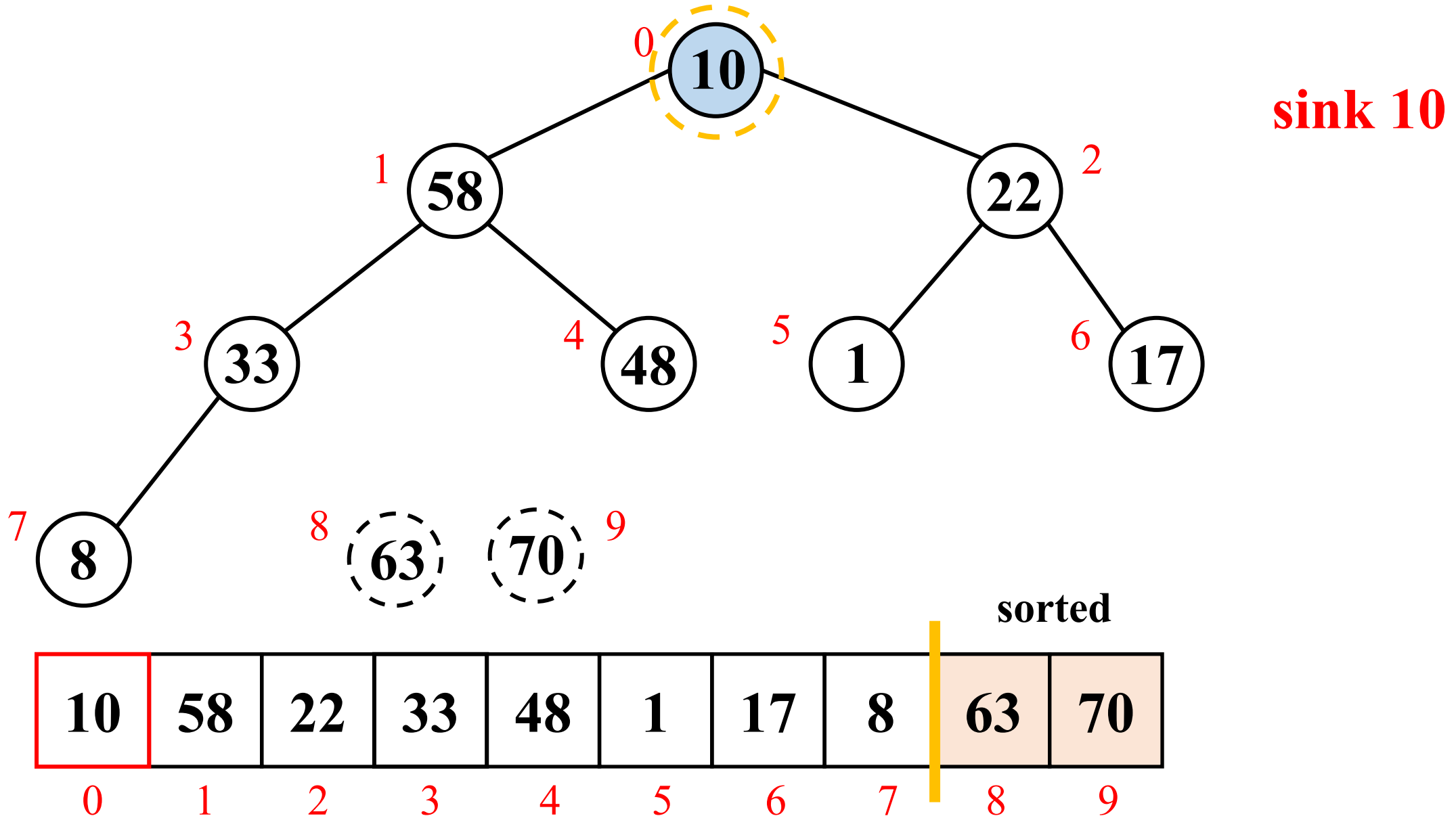
# Heapsort algorithms: array implementation



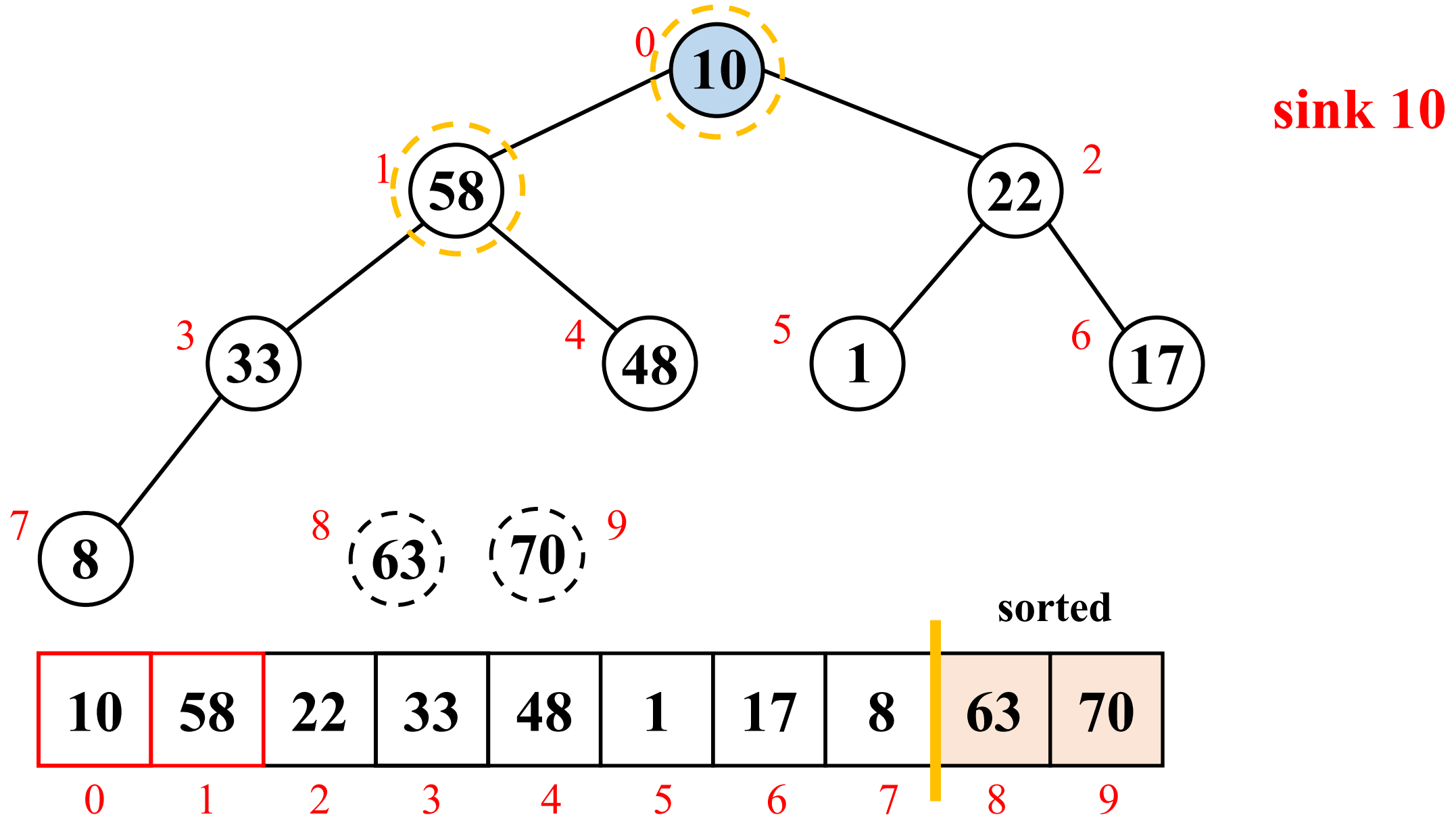
# Heapsort algorithms: array implementation



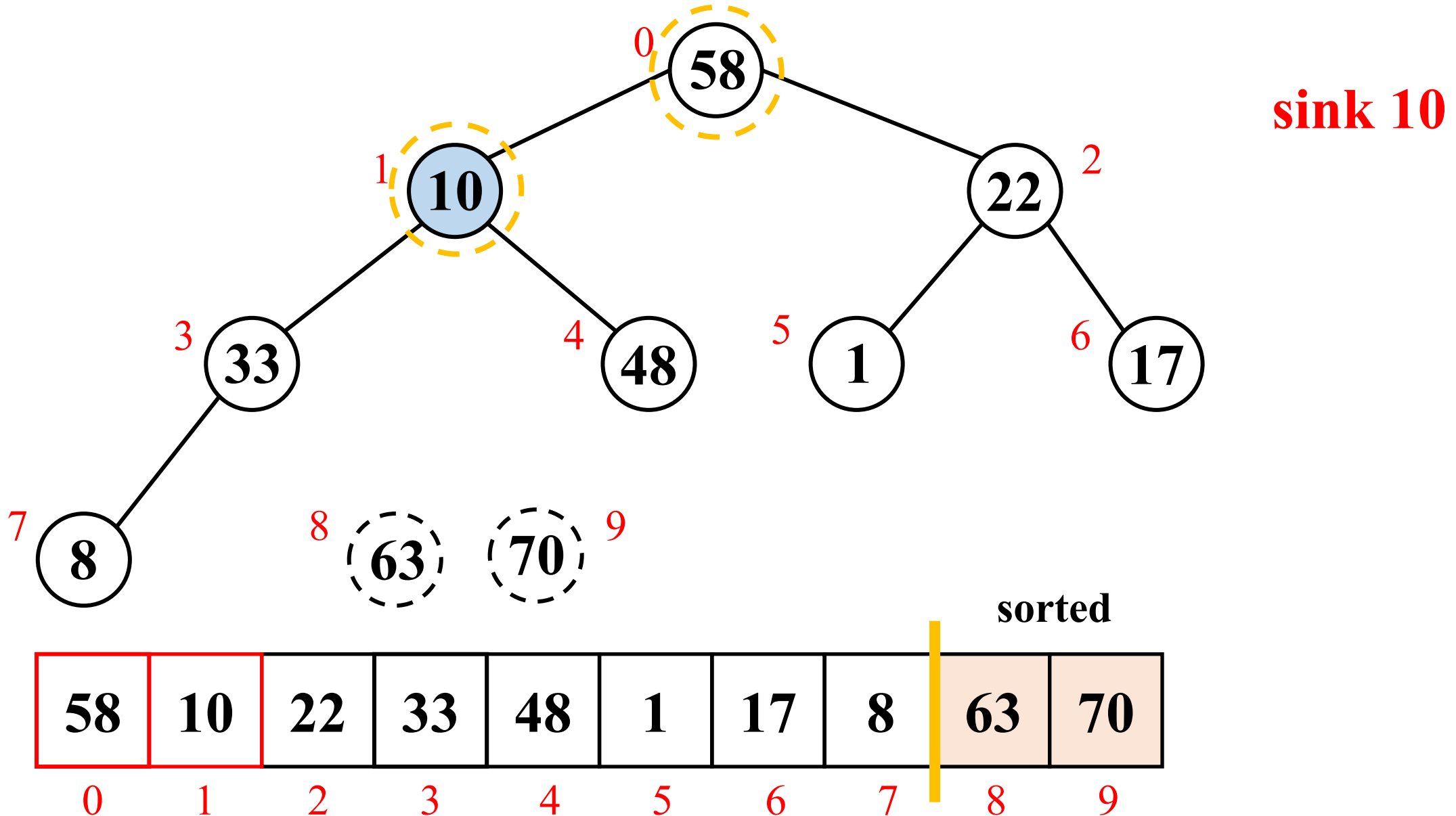
# Heapsort algorithms: array implementation



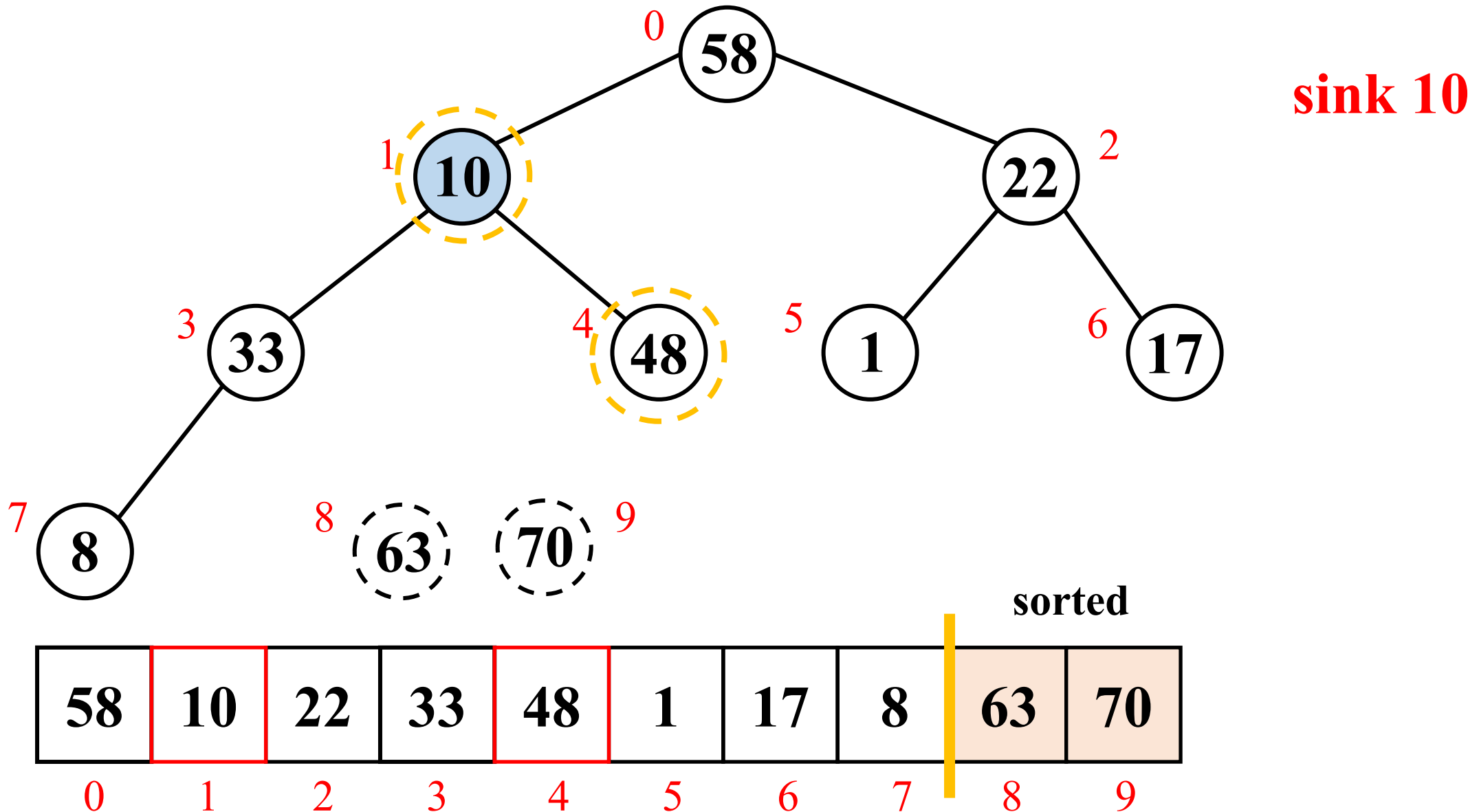
# Heapsort algorithms: array implementation



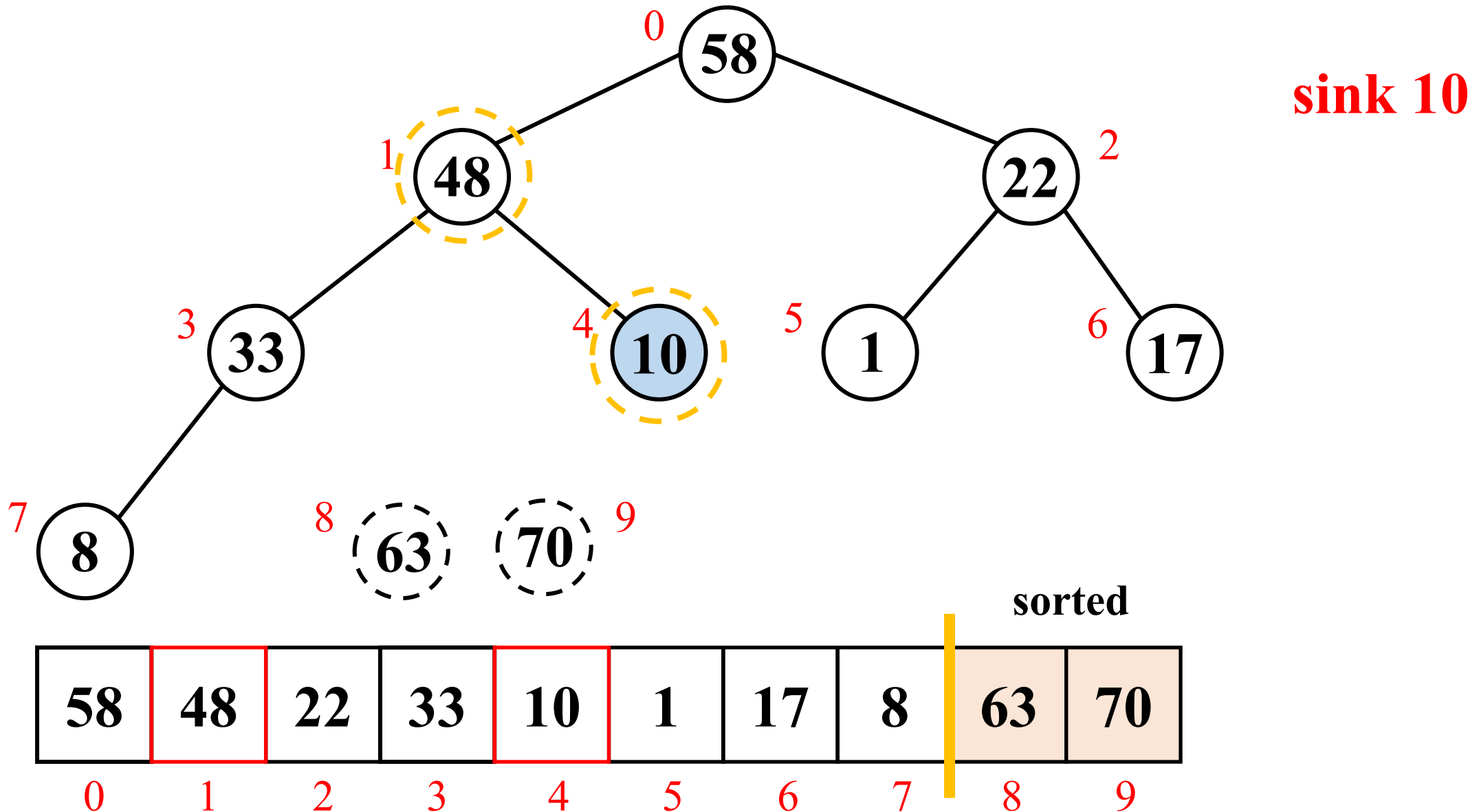
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

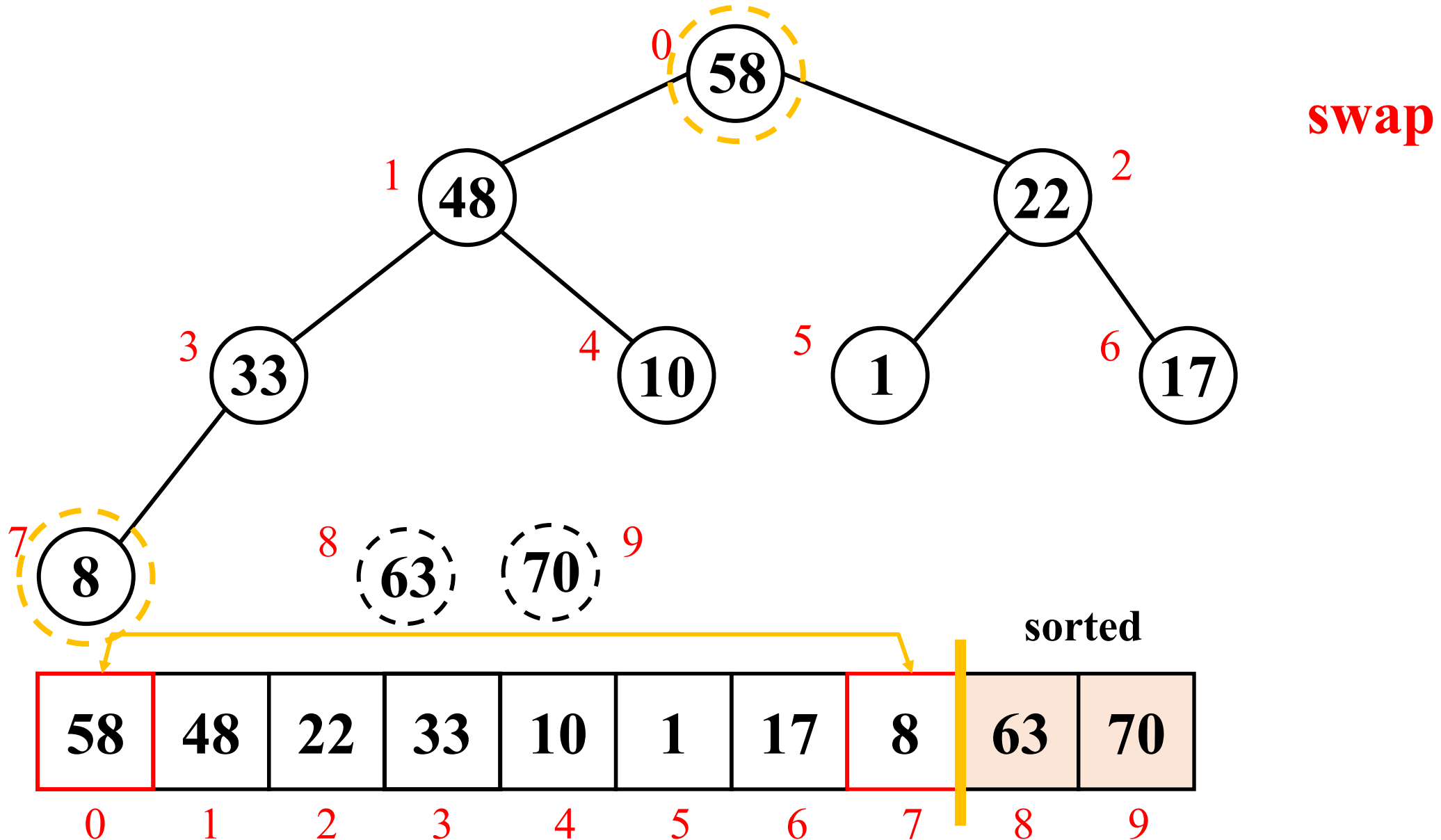


# Heapsort algorithms: array implementation

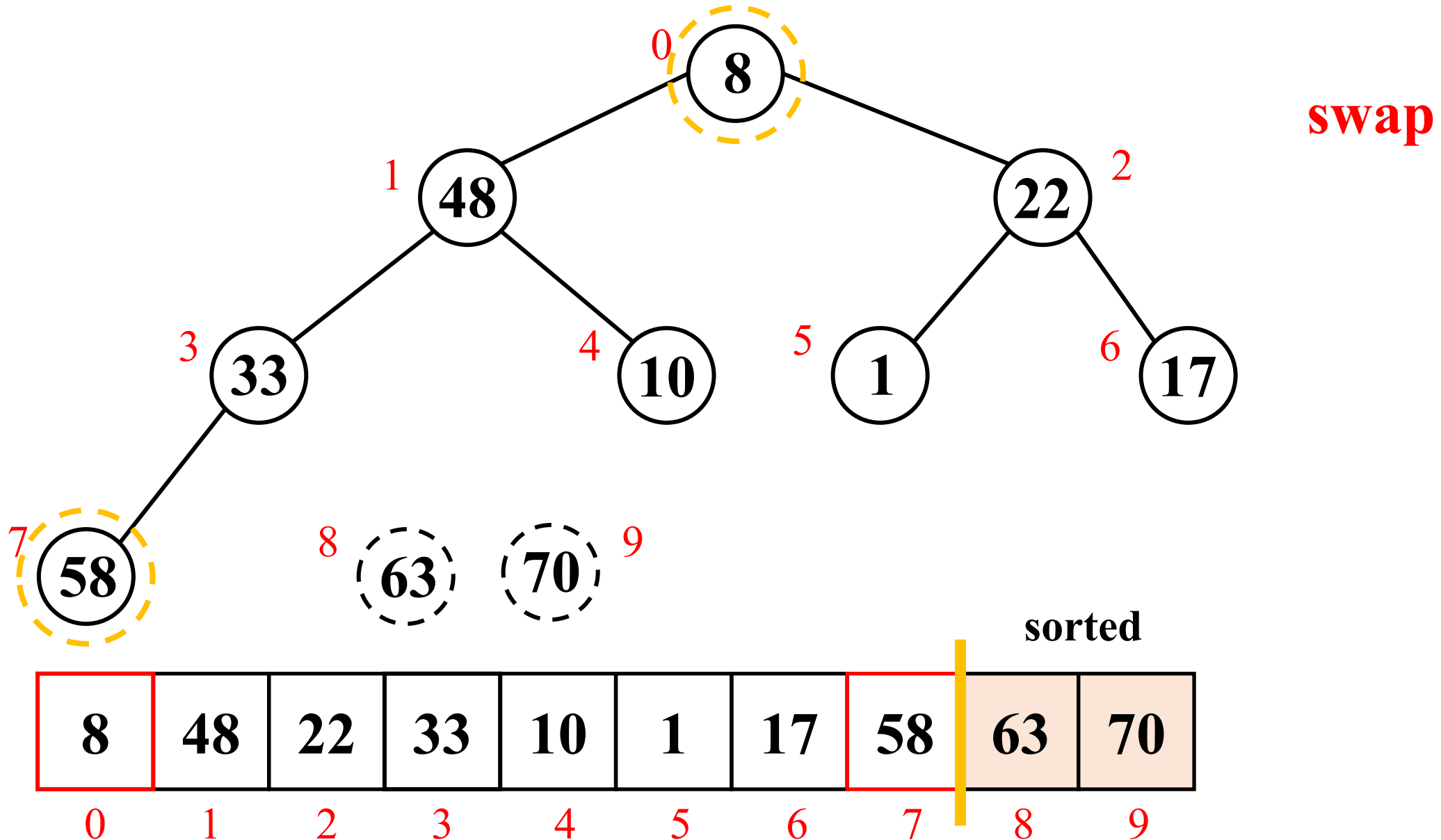




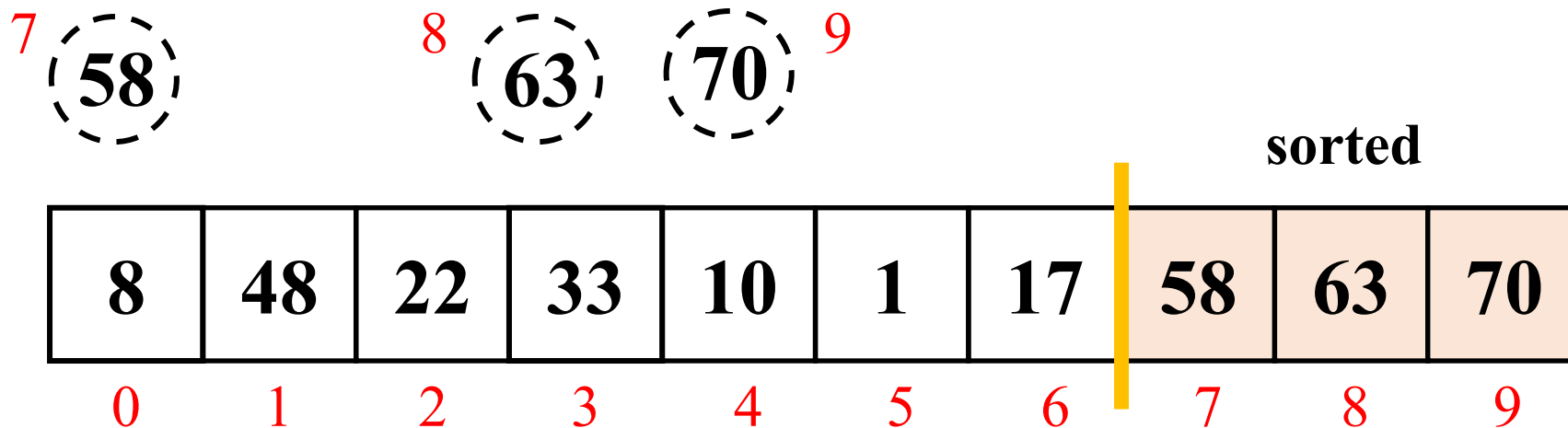
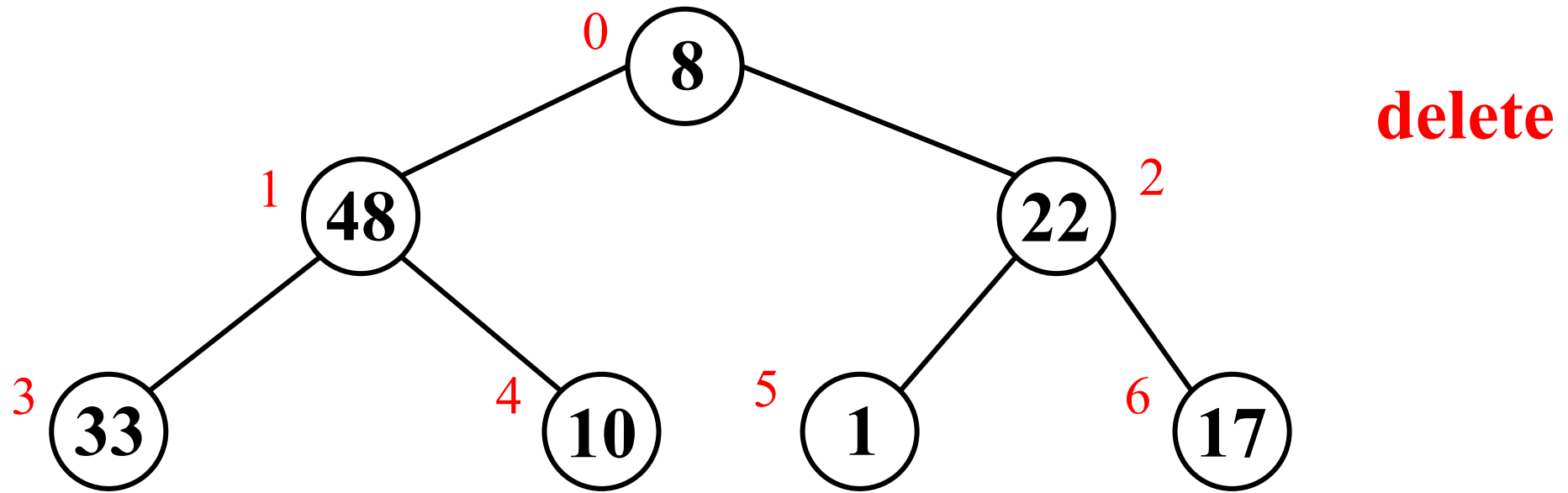
# Heapsort algorithms: array implementation



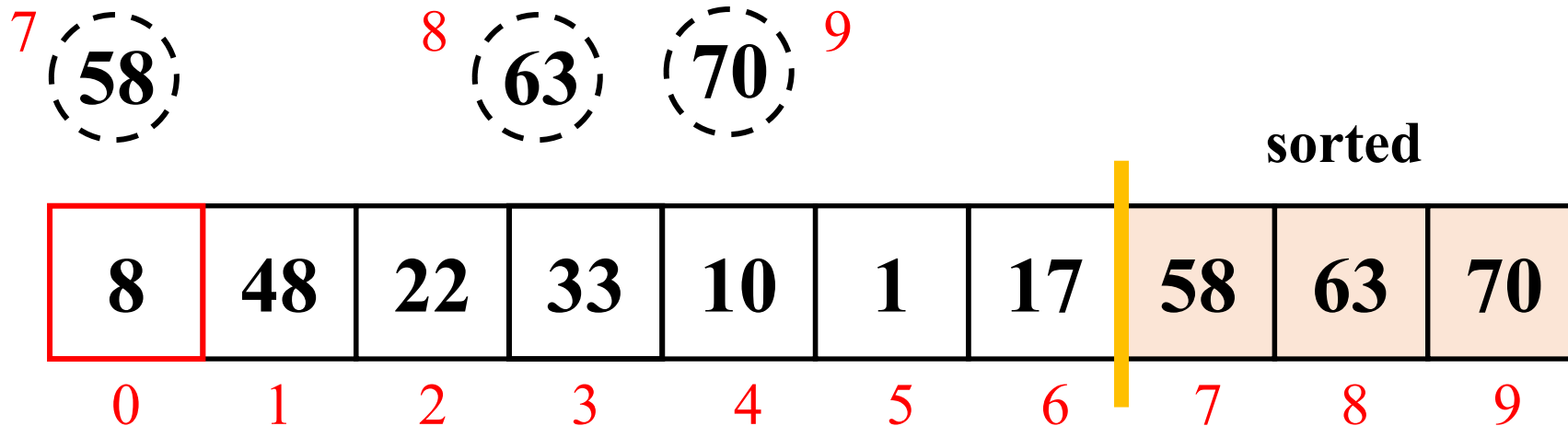
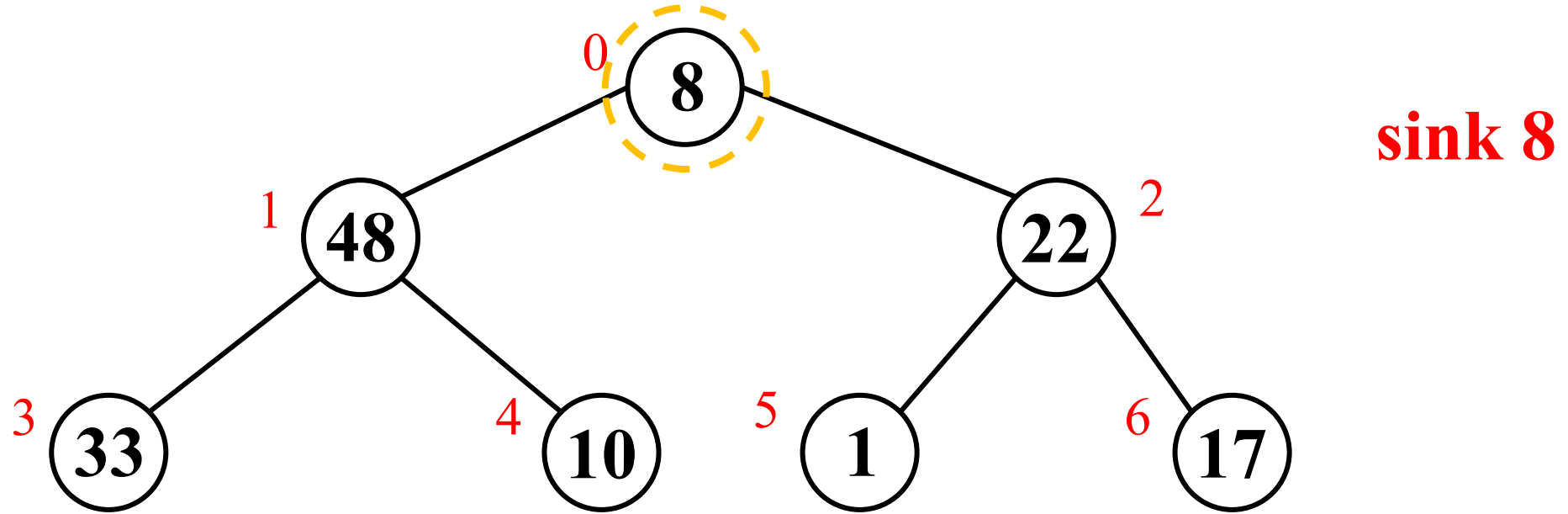
# Heapsort algorithms: array implementation



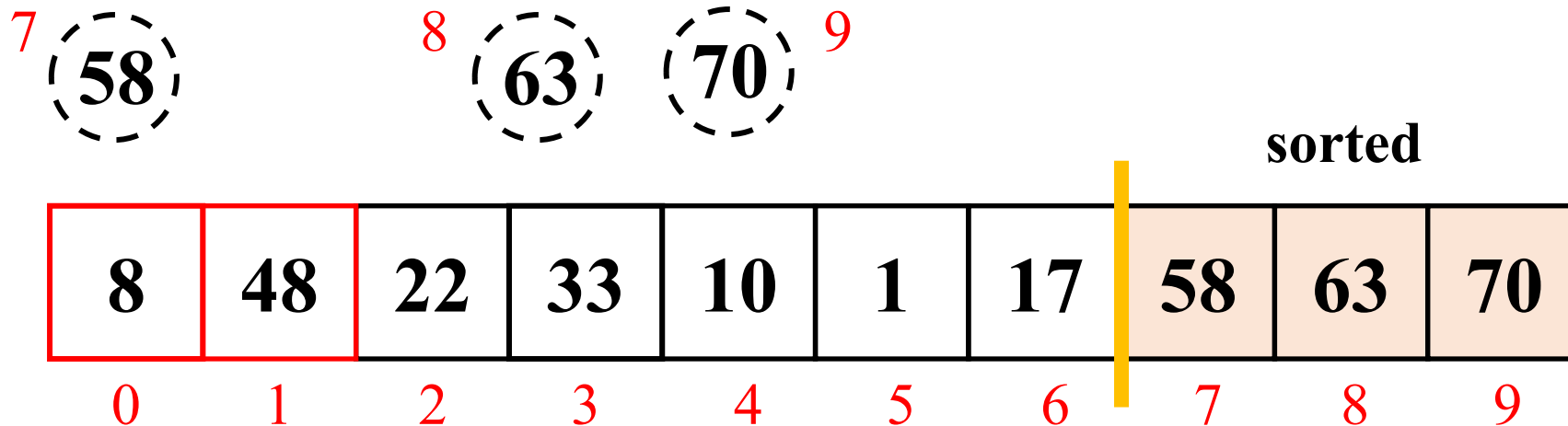
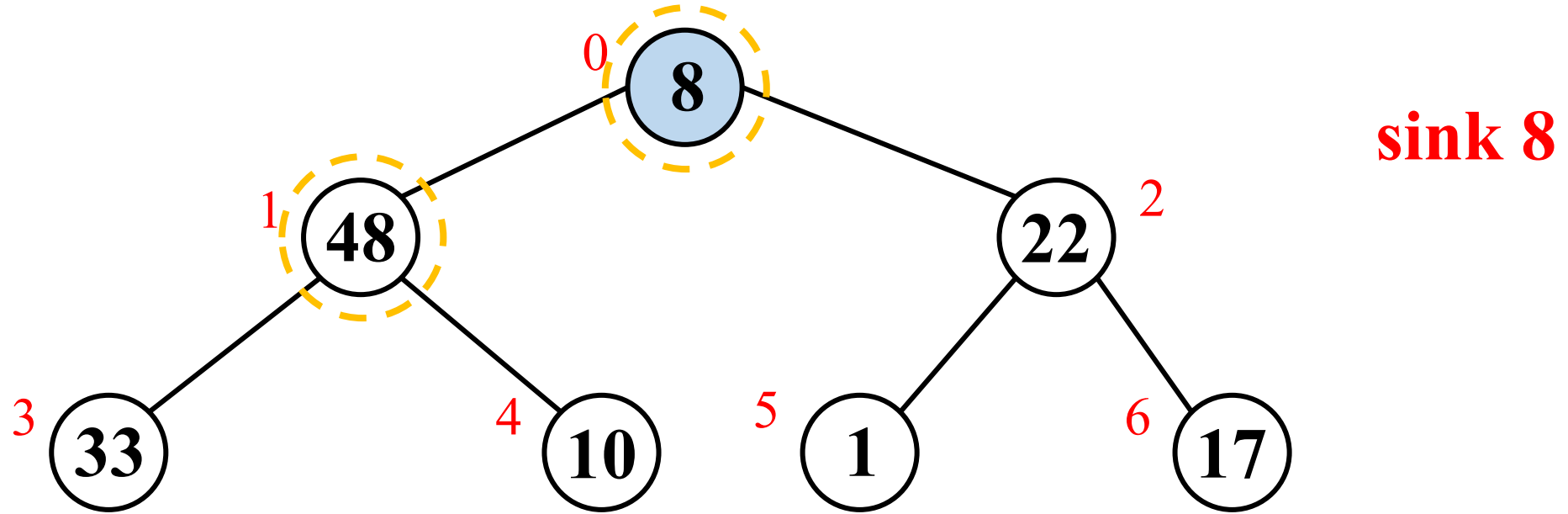
# Heapsort algorithms: array implementation



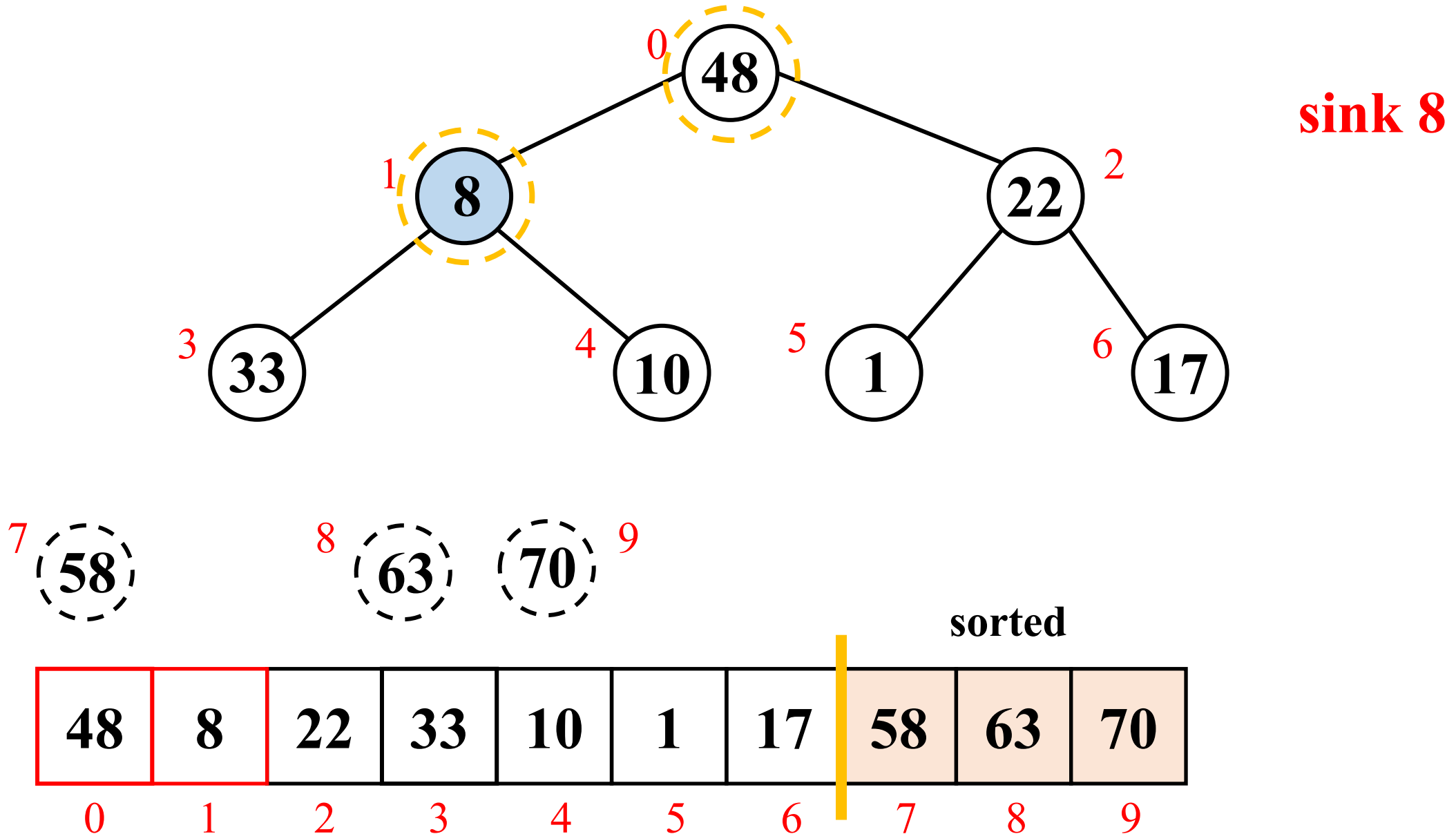
# Heapsort algorithms: array implementation



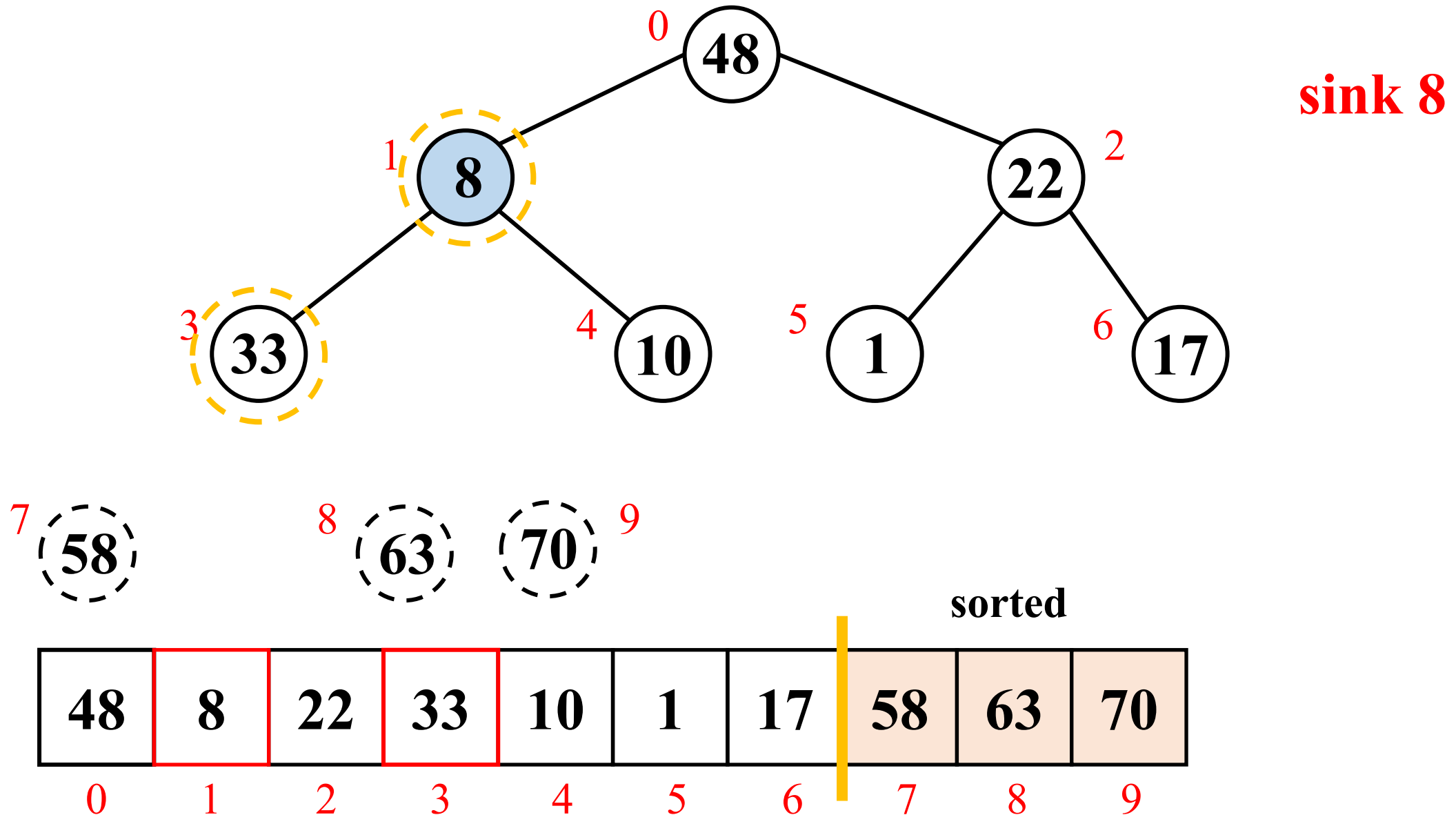
# Heapsort algorithms: array implementation



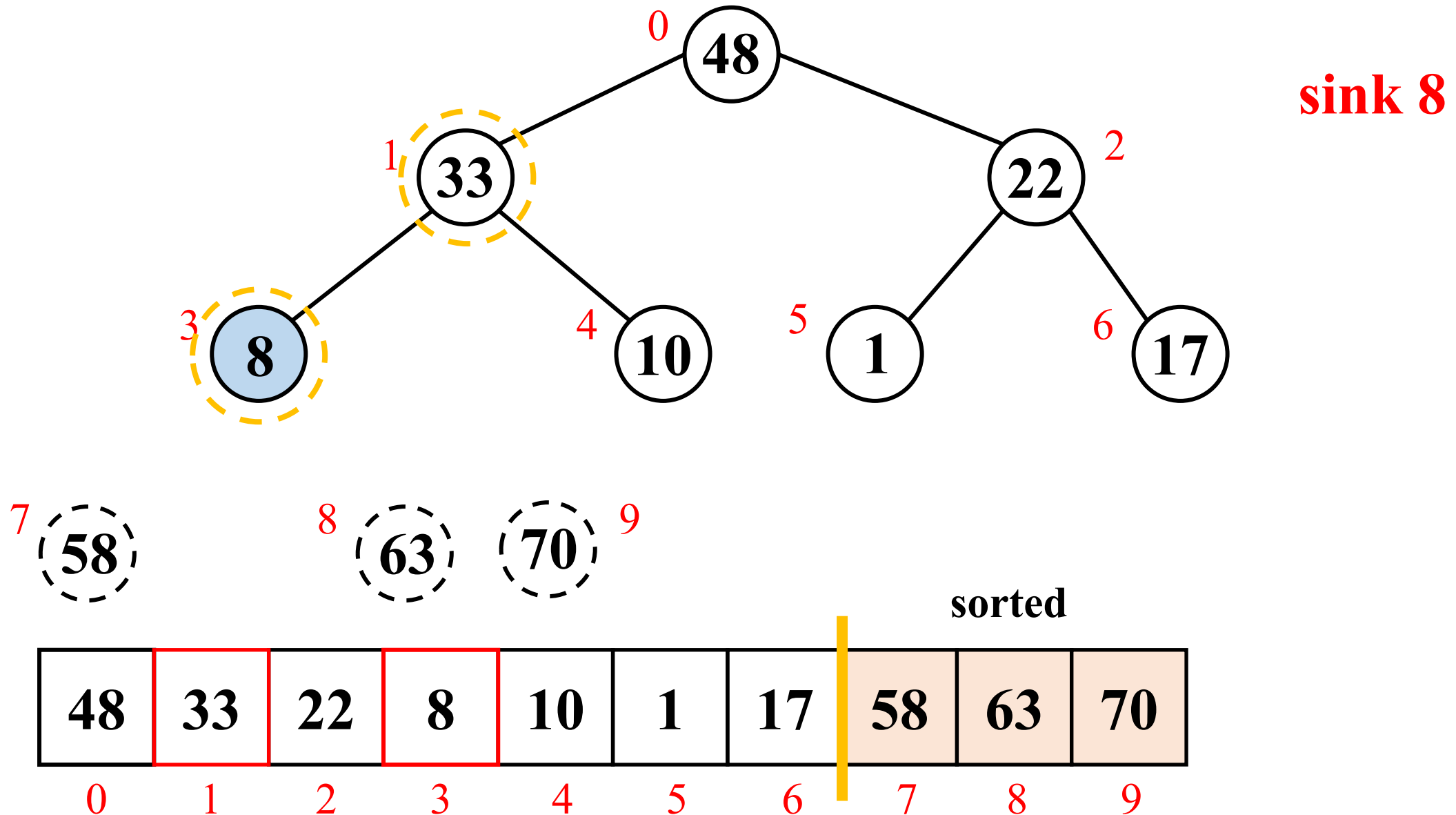
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

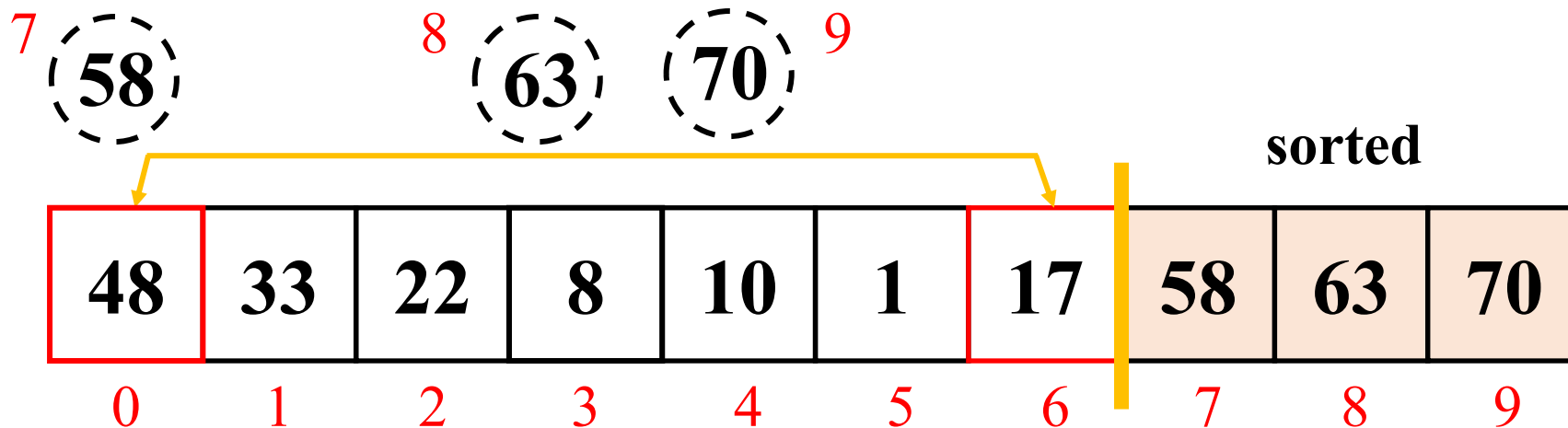
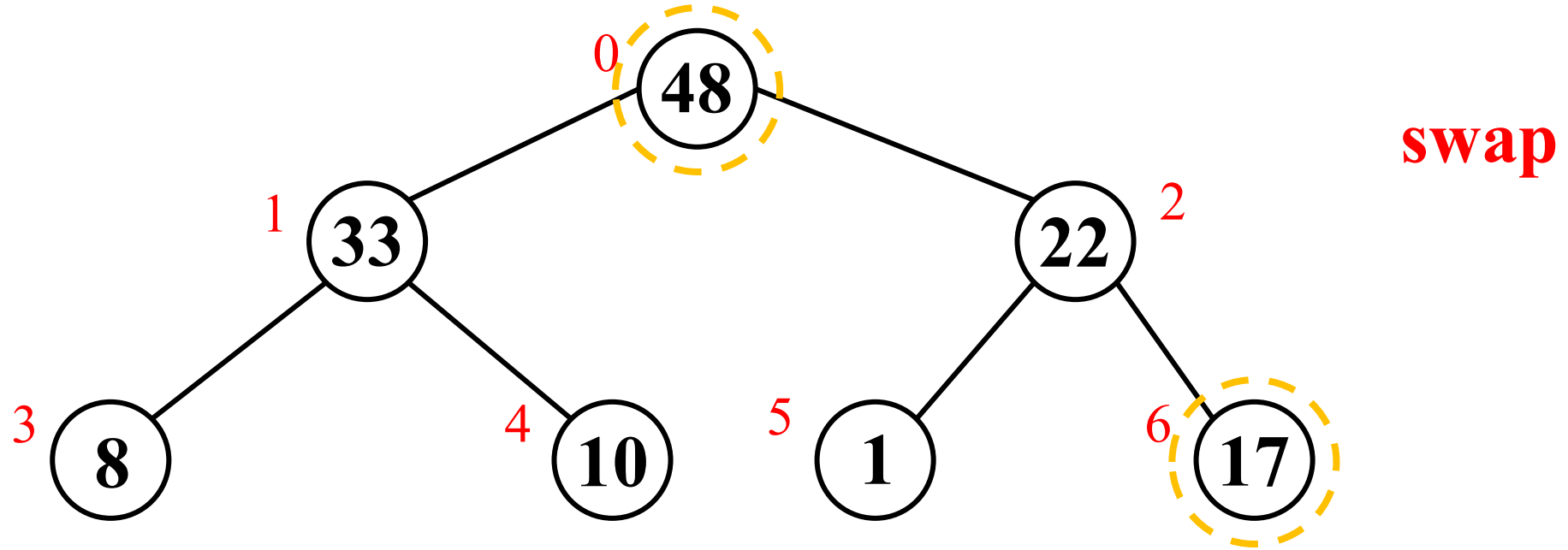


# Heapsort algorithms: array implementation

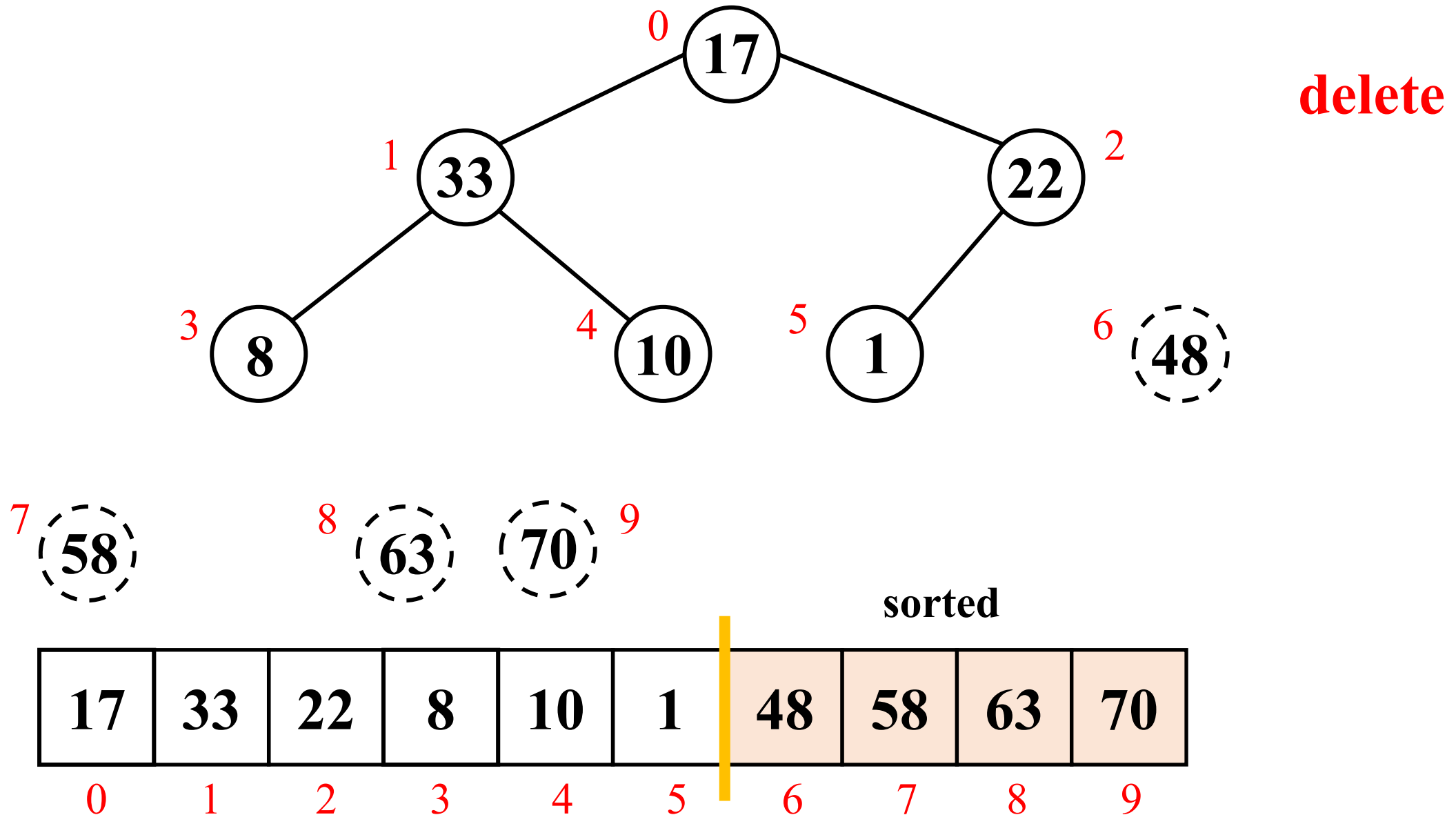




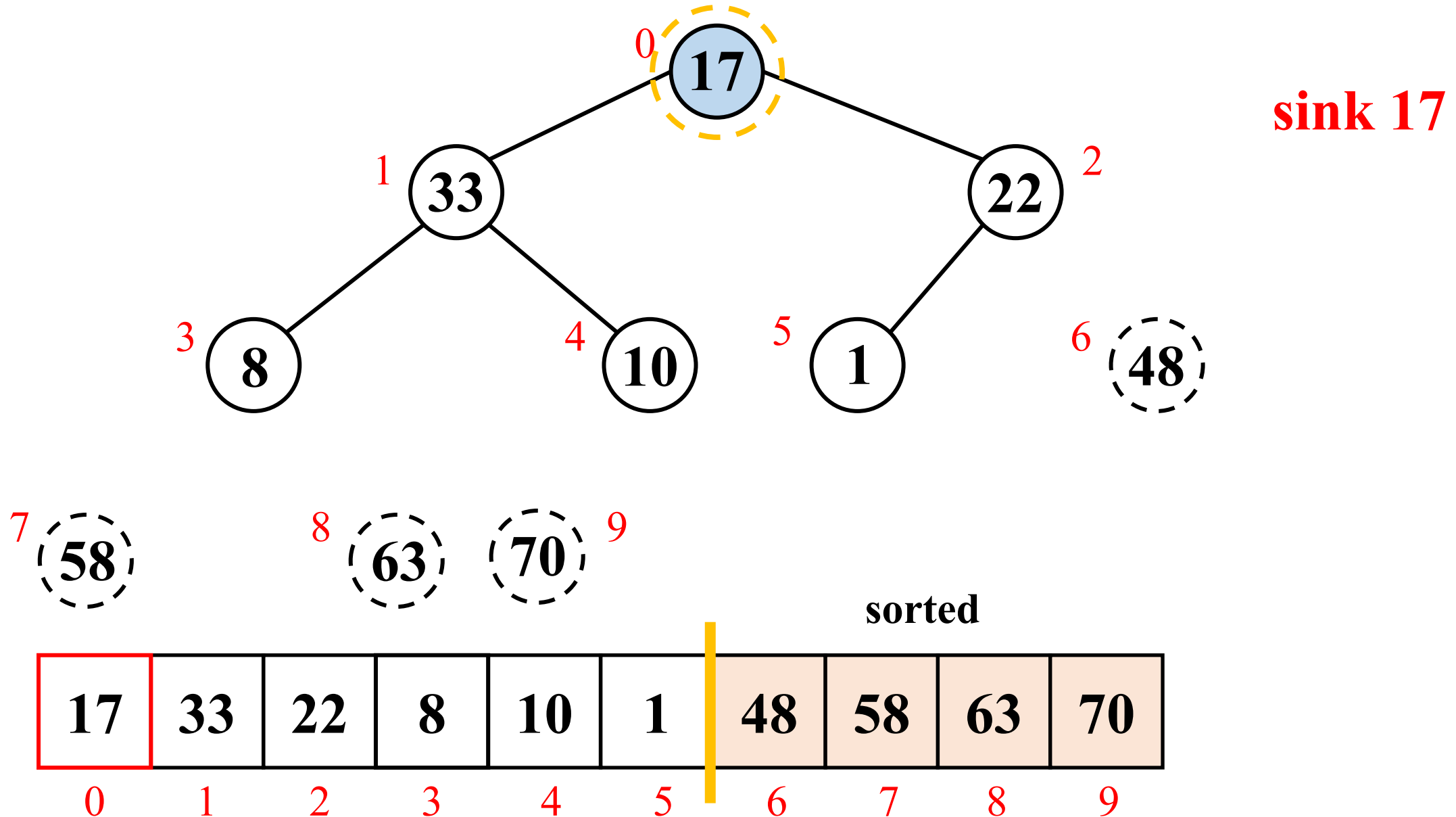
# Heapsort algorithms: array implementation



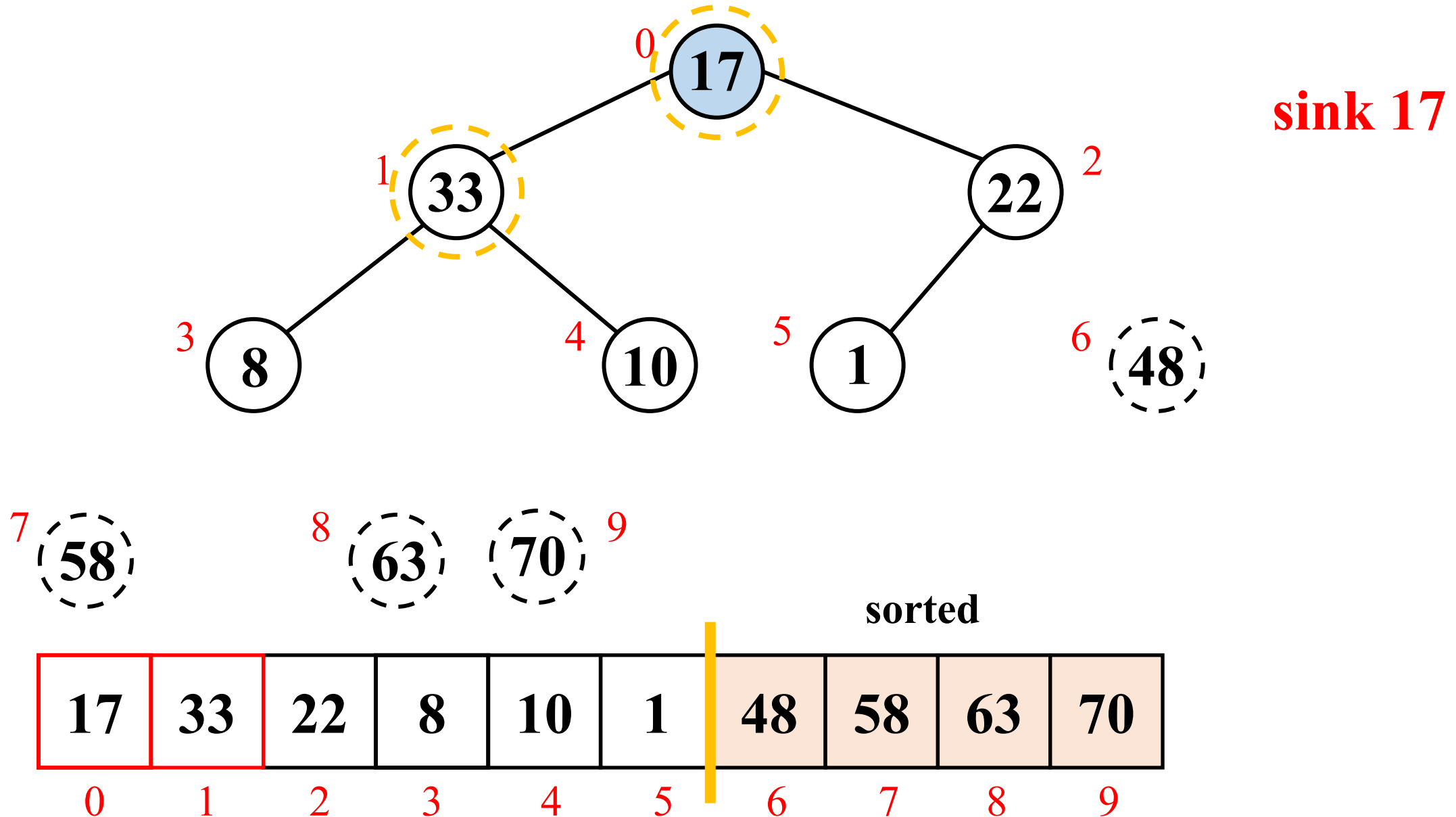
# Heapsort algorithms: array implementation



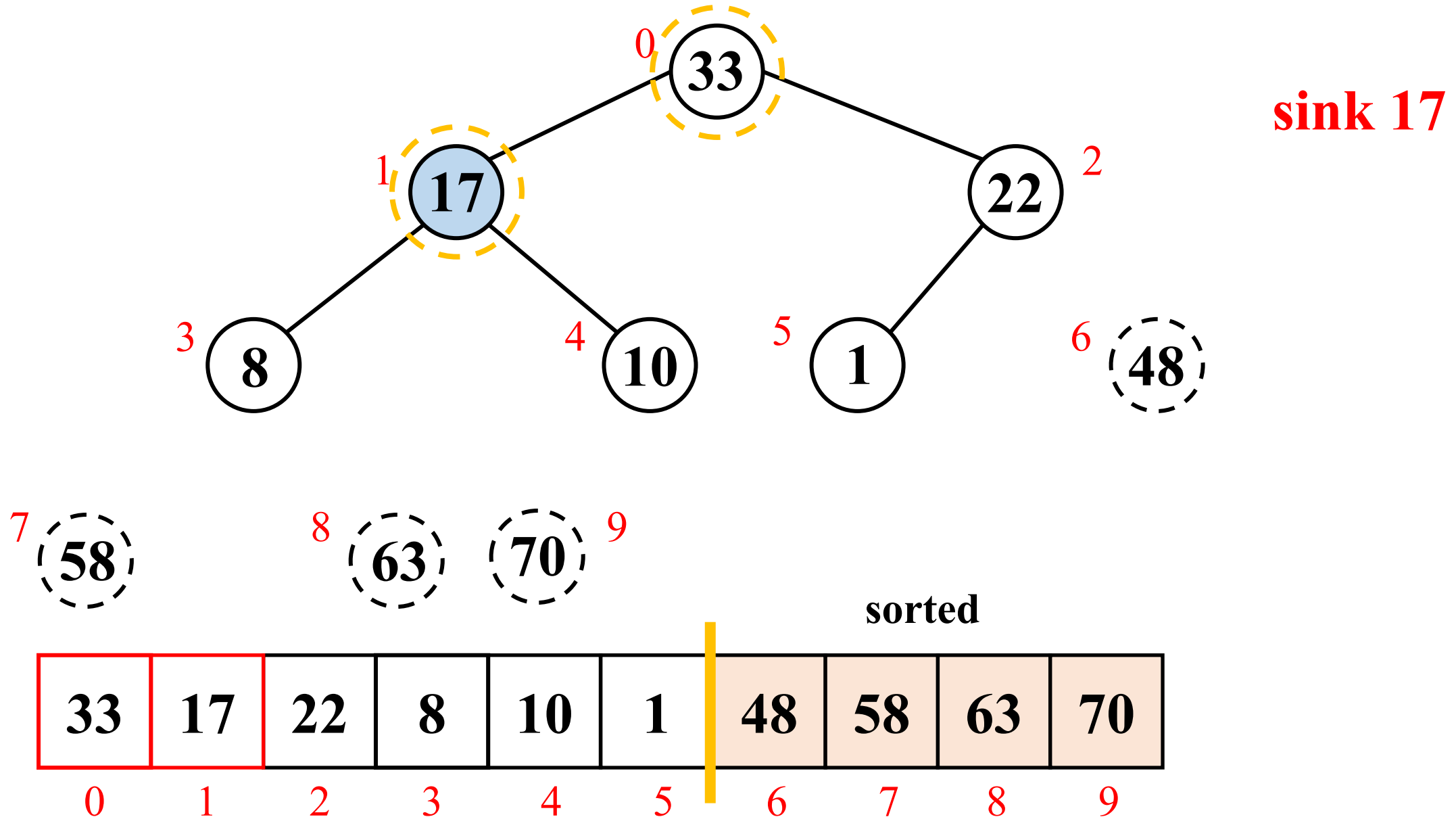
# Heapsort algorithms: array implementation



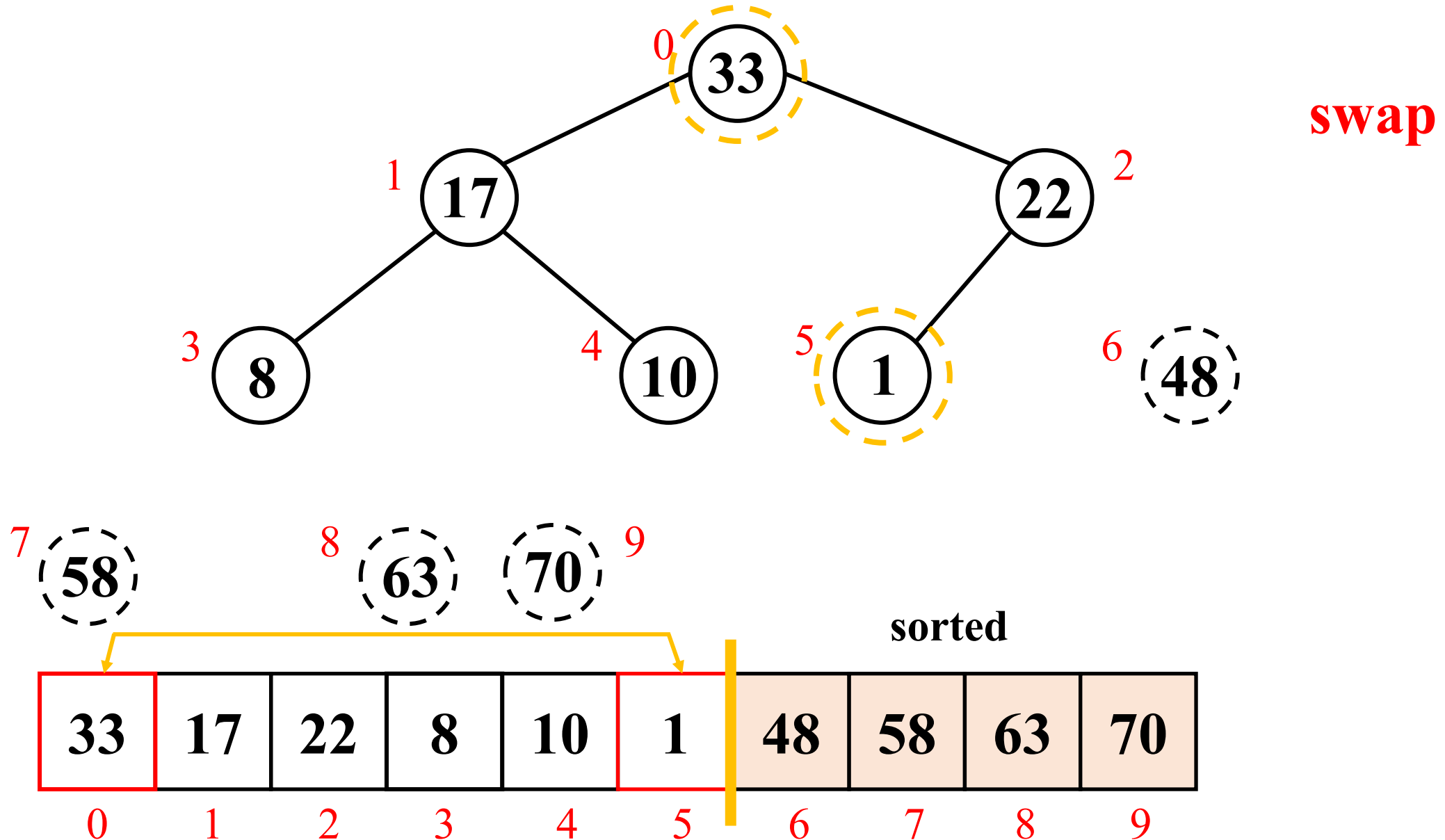
# Heapsort algorithms: array implementation



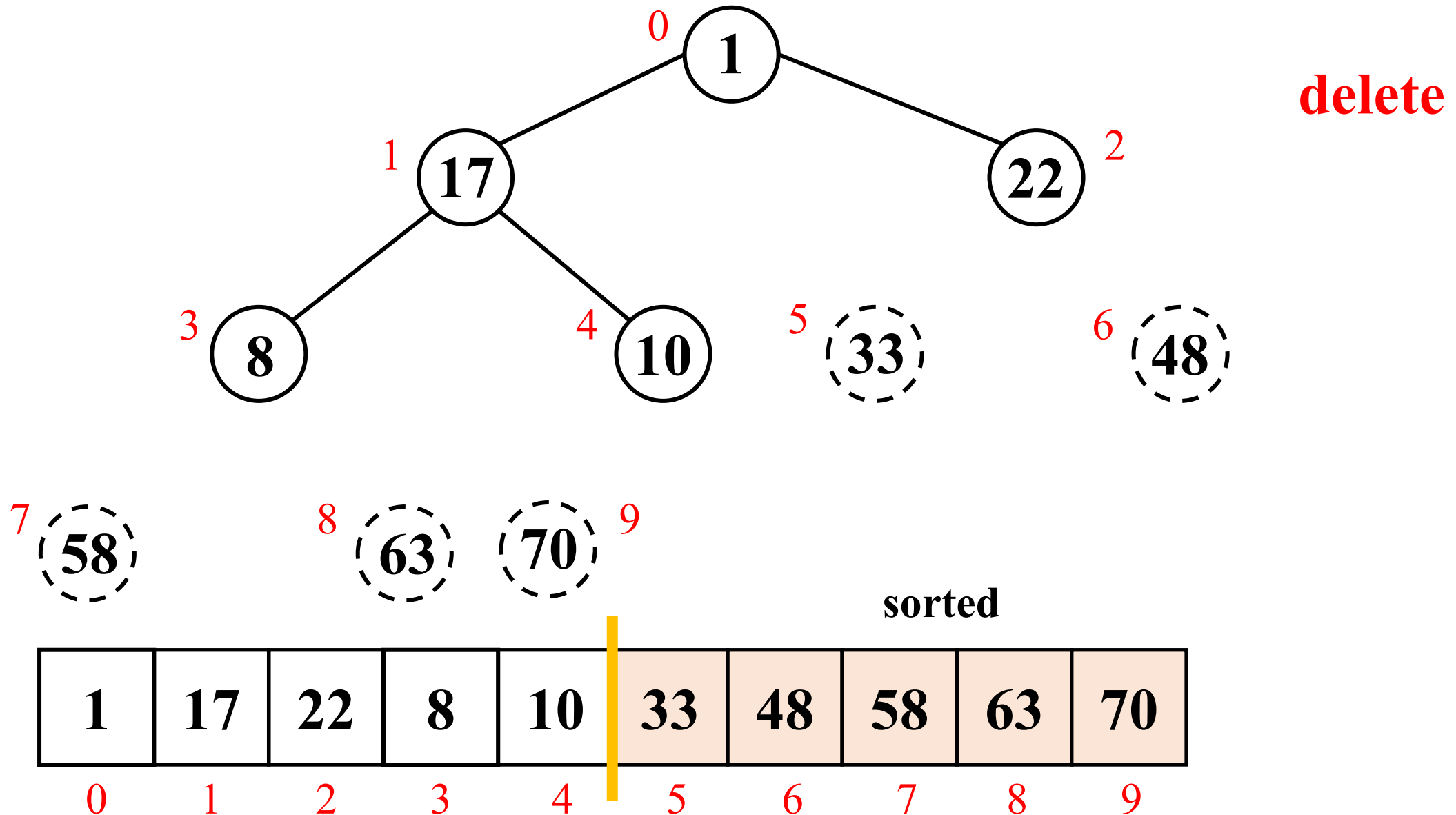
# Heapsort algorithms: array implementation



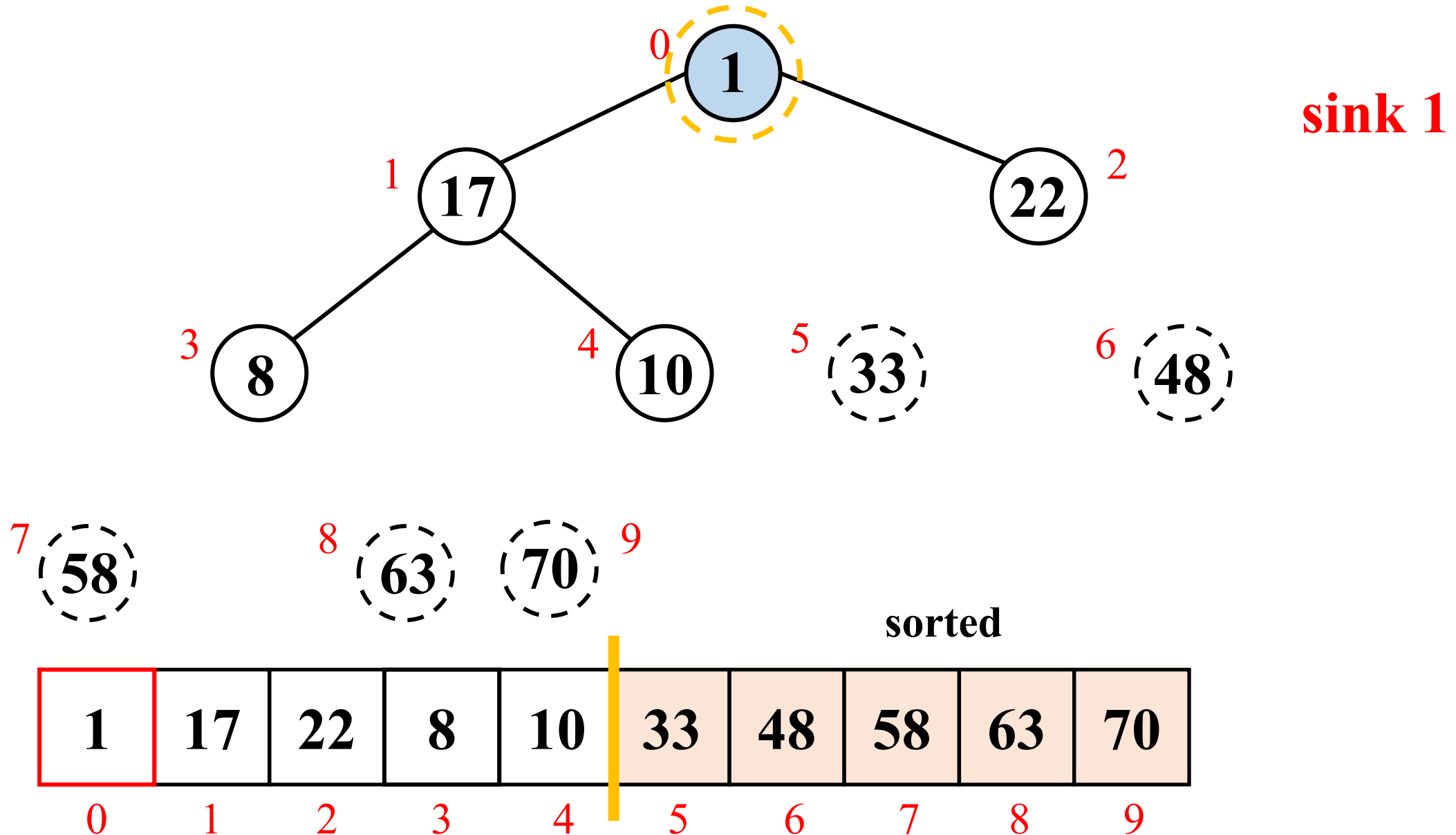
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

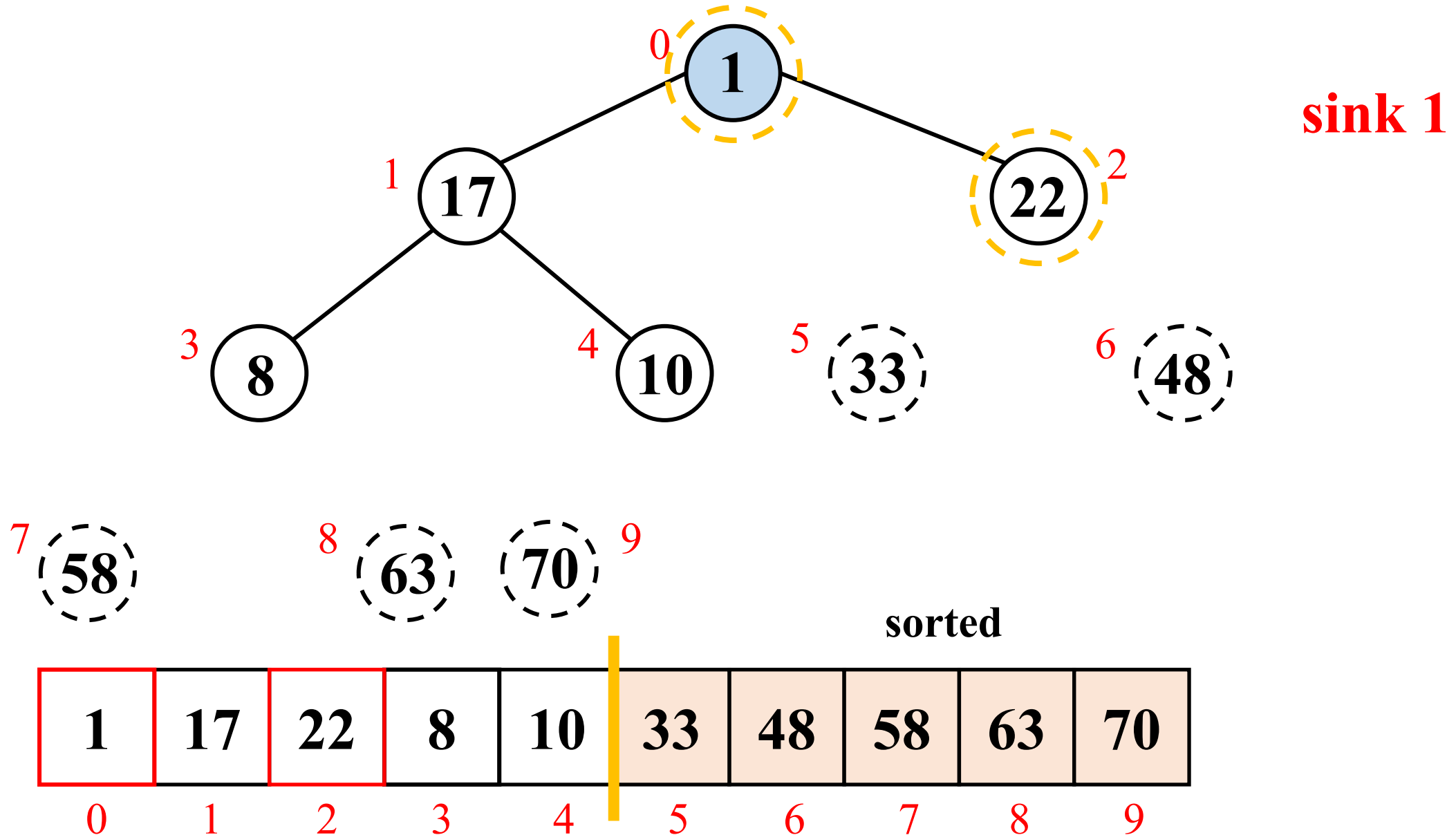


# Heapsort algorithms: array implementation

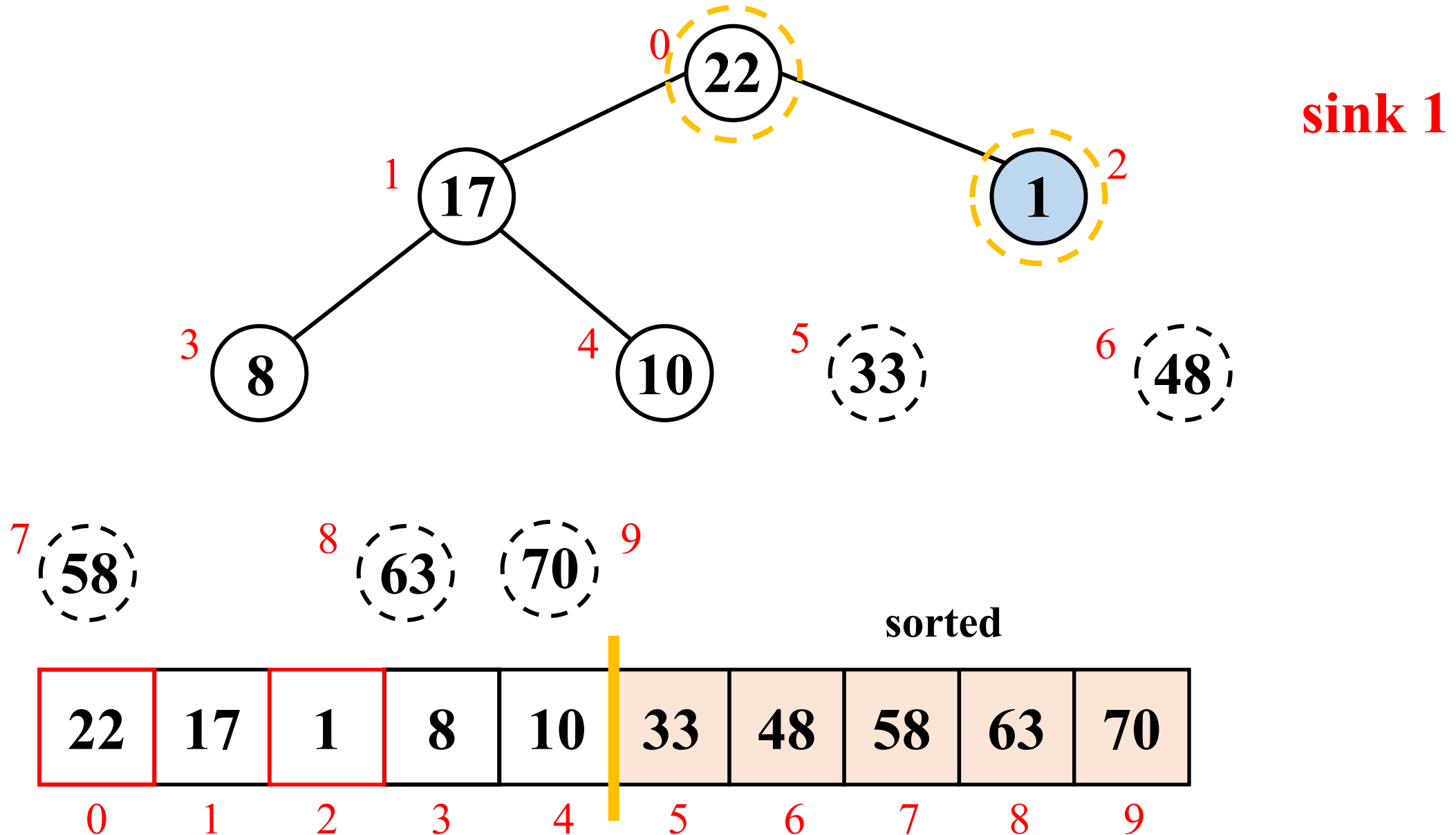




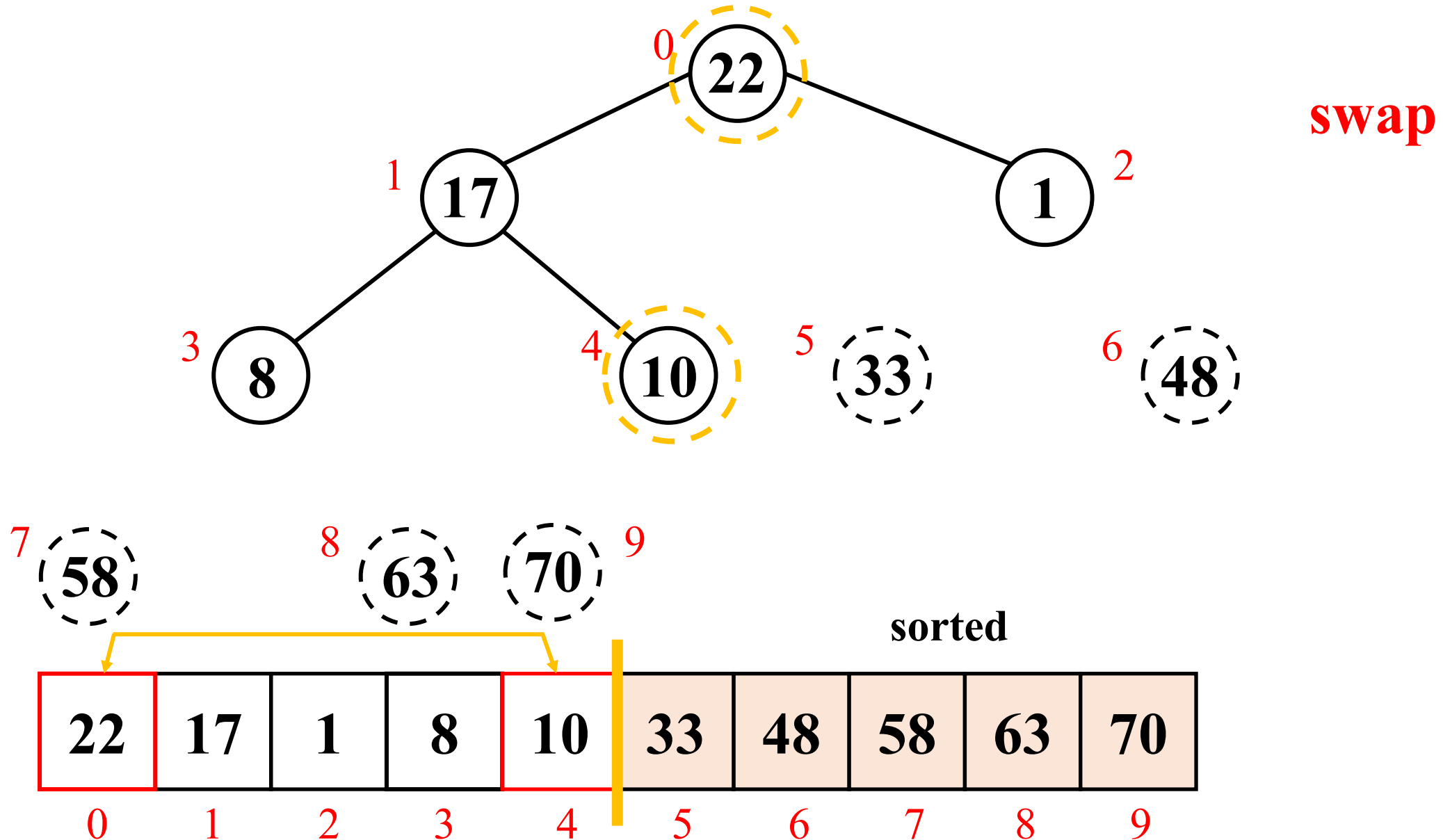
# Heapsort algorithms: array implementation



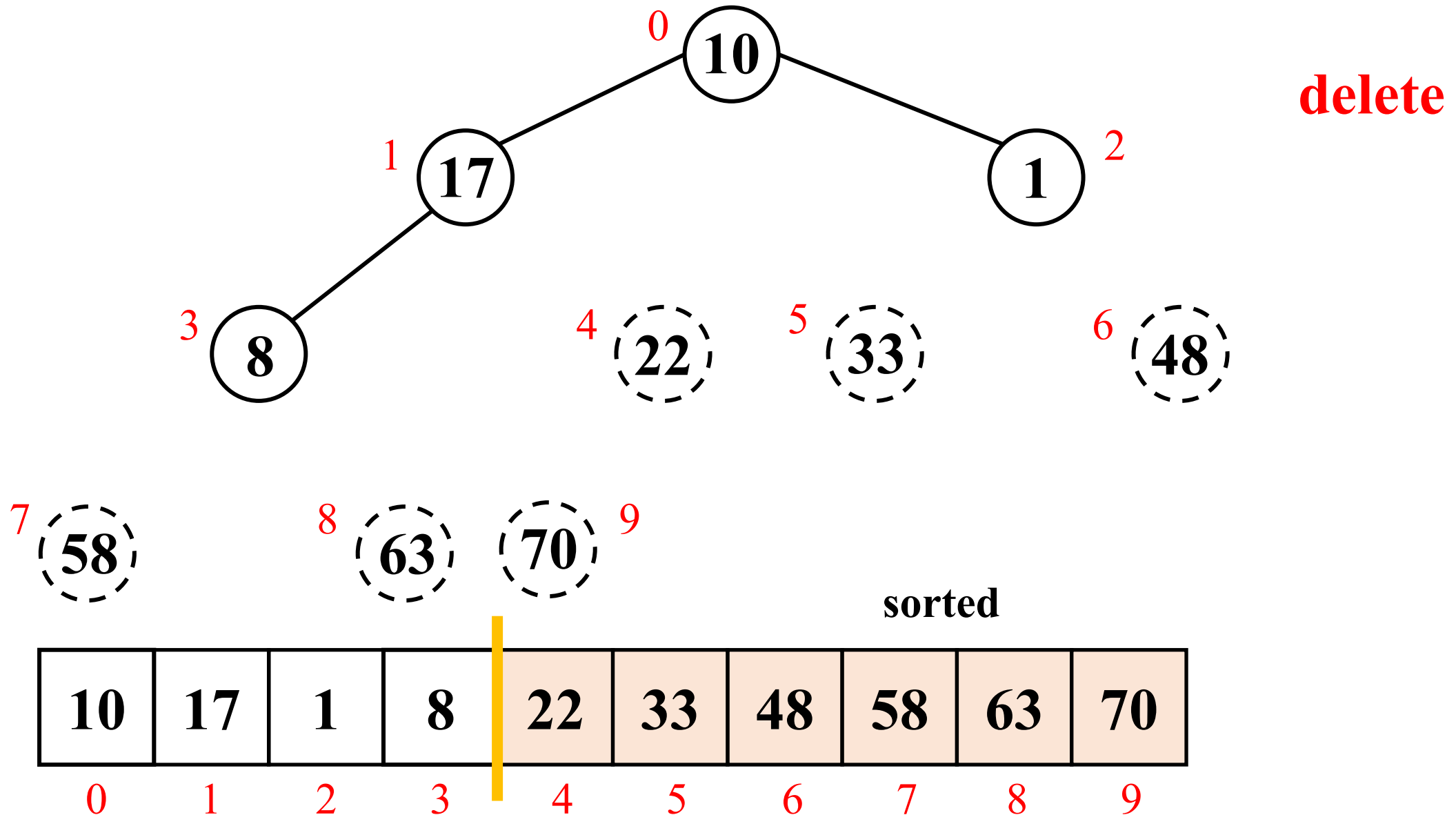
# Heapsort algorithms: array implementation



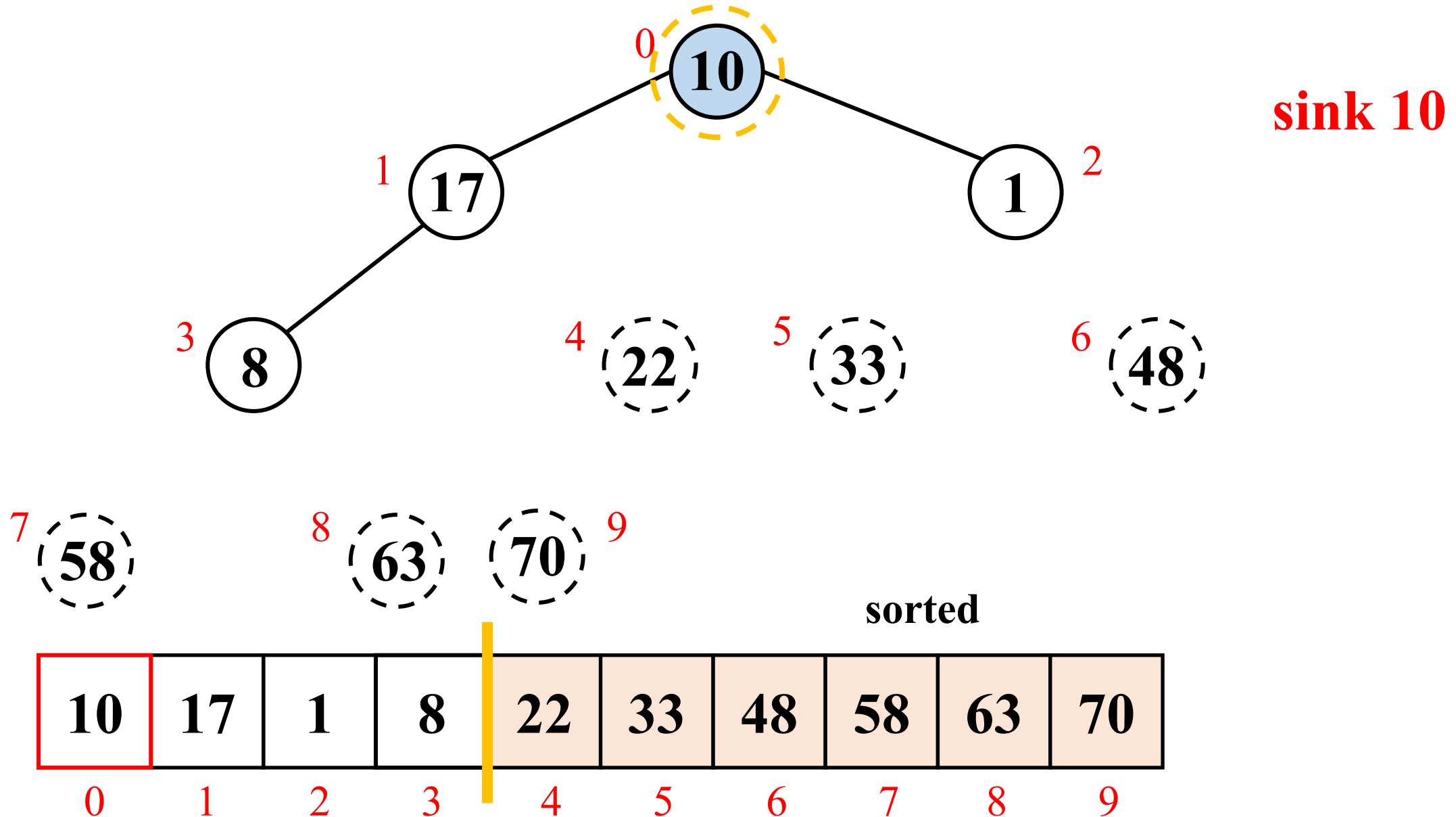
# Heapsort algorithms: array implementation



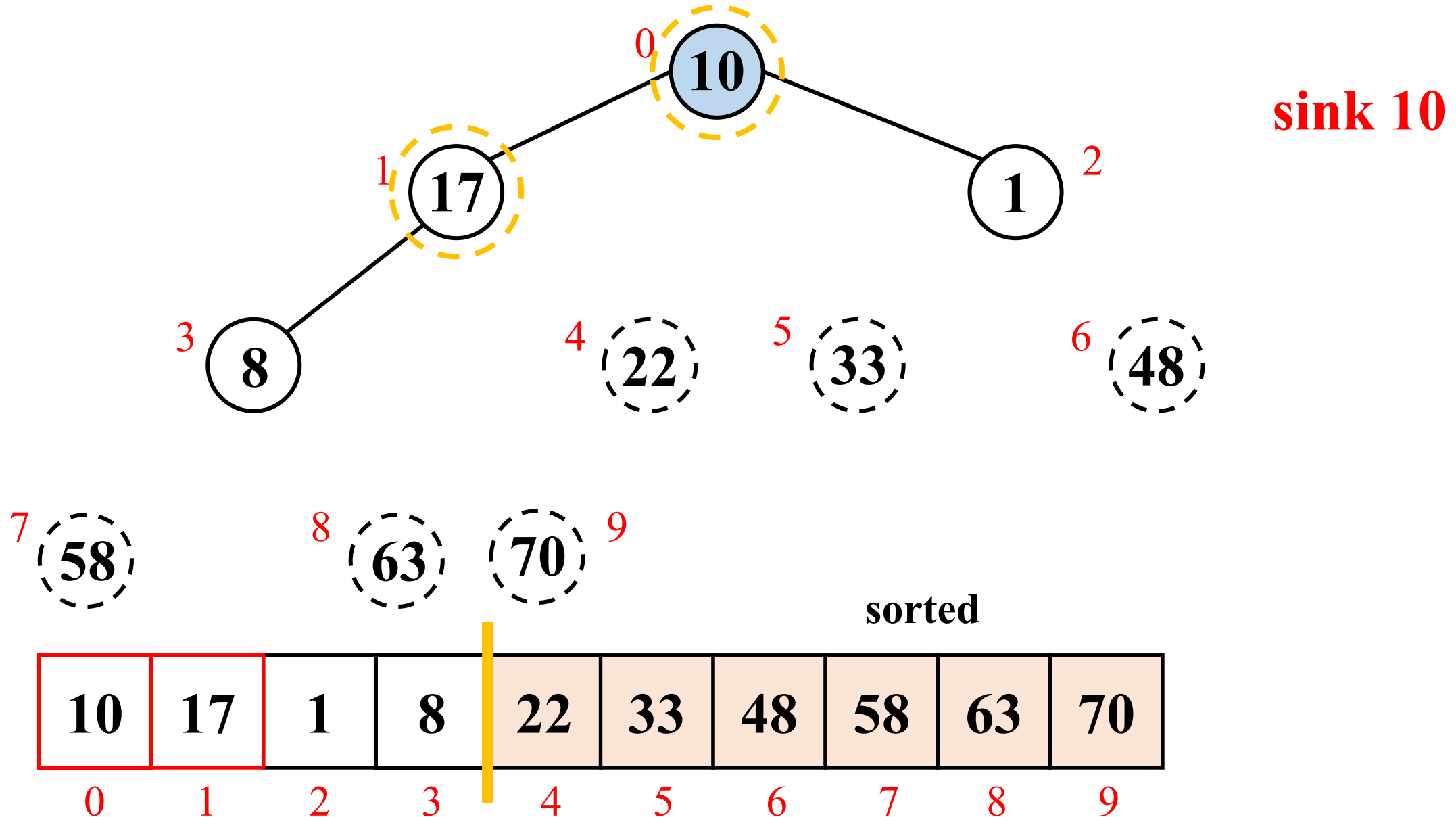
# Heapsort algorithms: array implementation



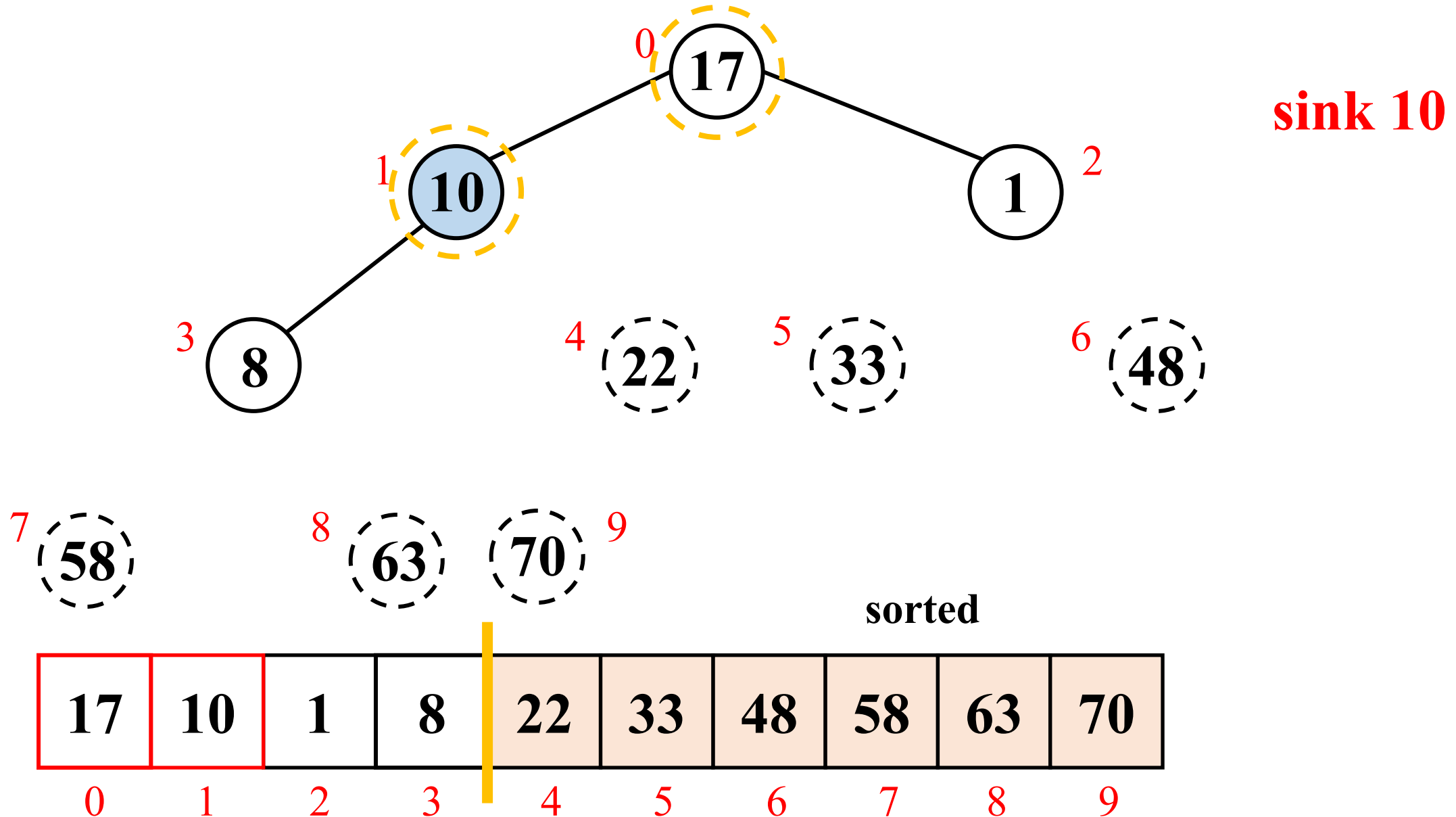
# Heapsort algorithms: array implementation



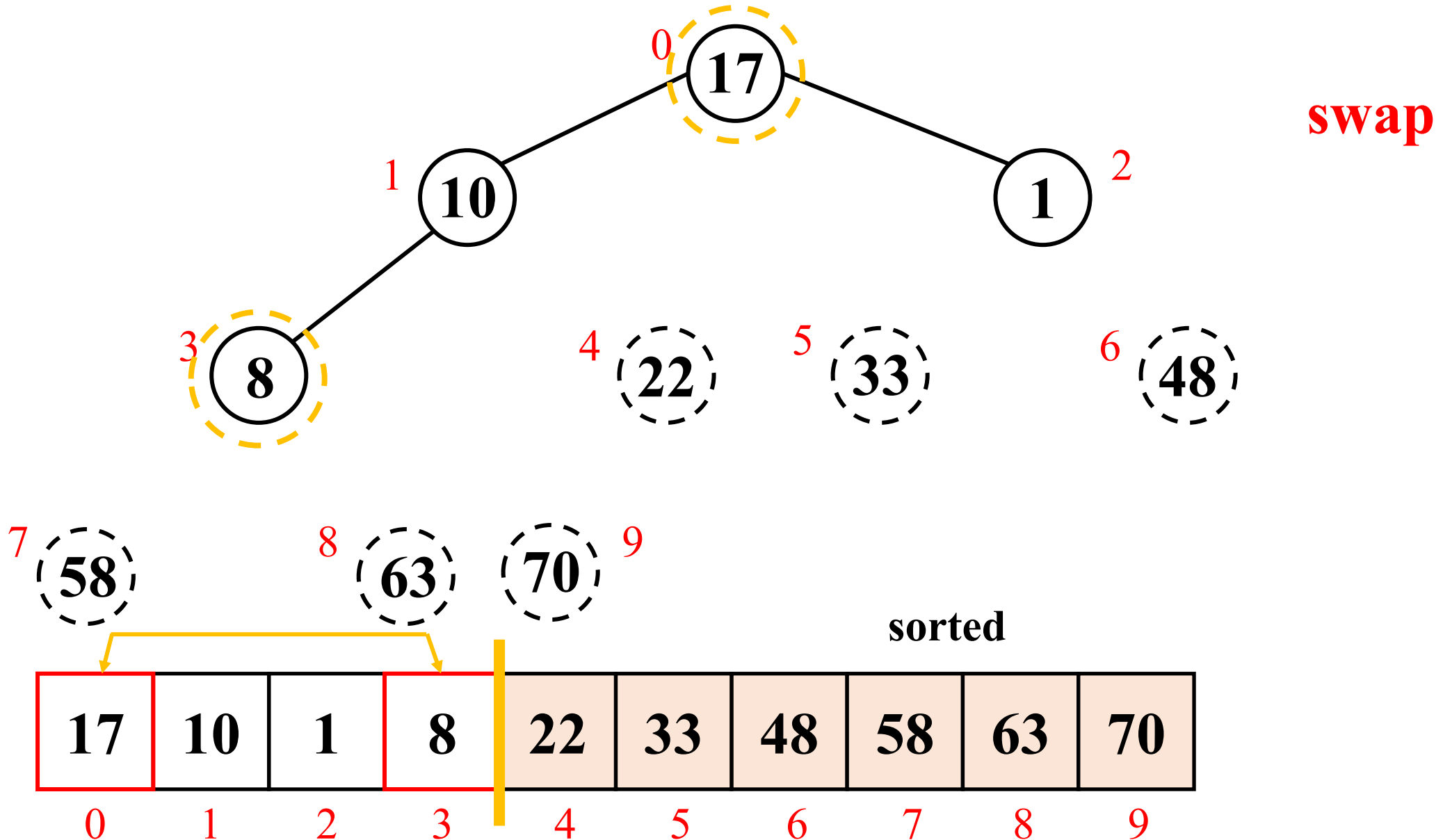
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

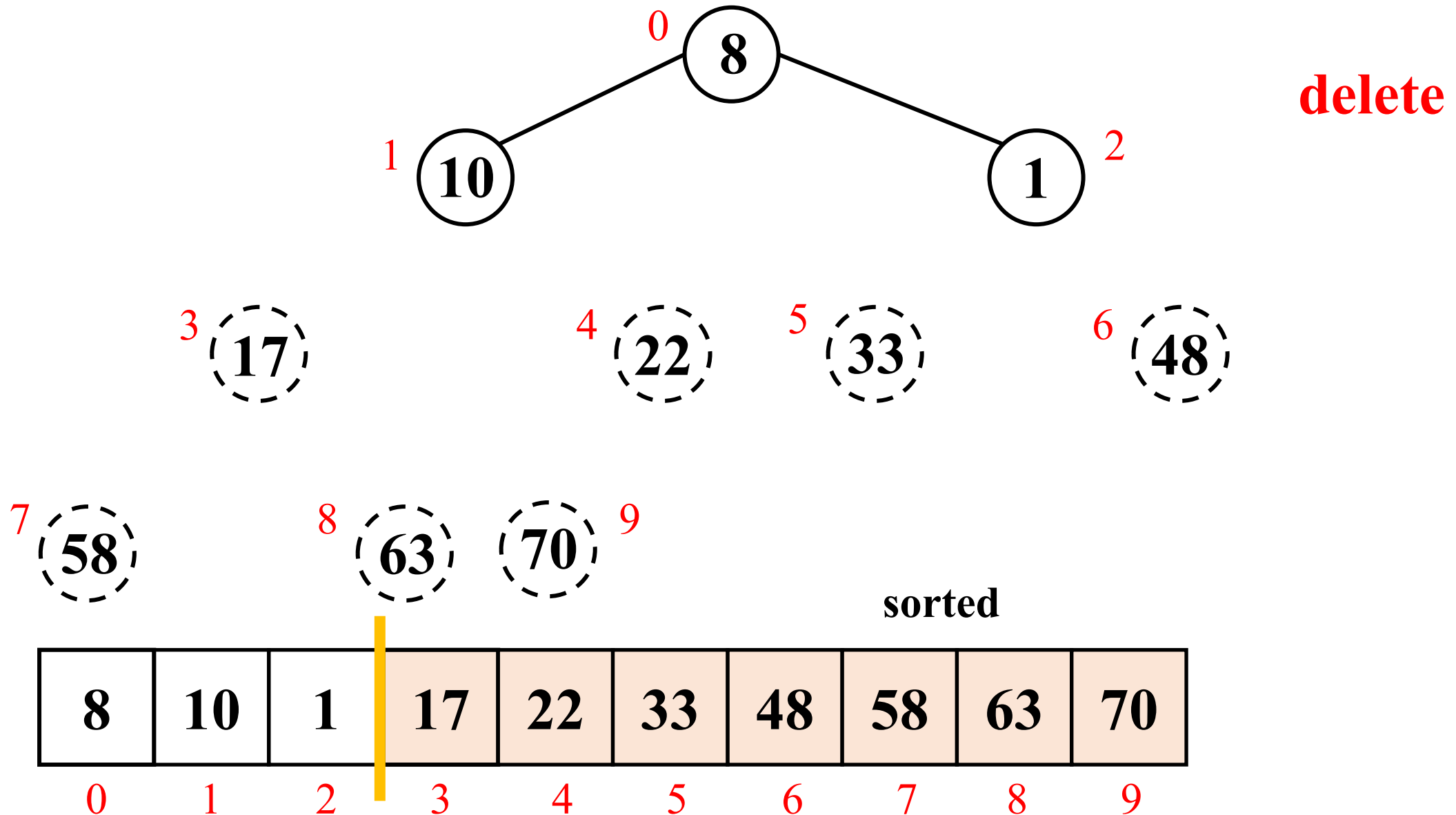


# Heapsort algorithms: array implementation

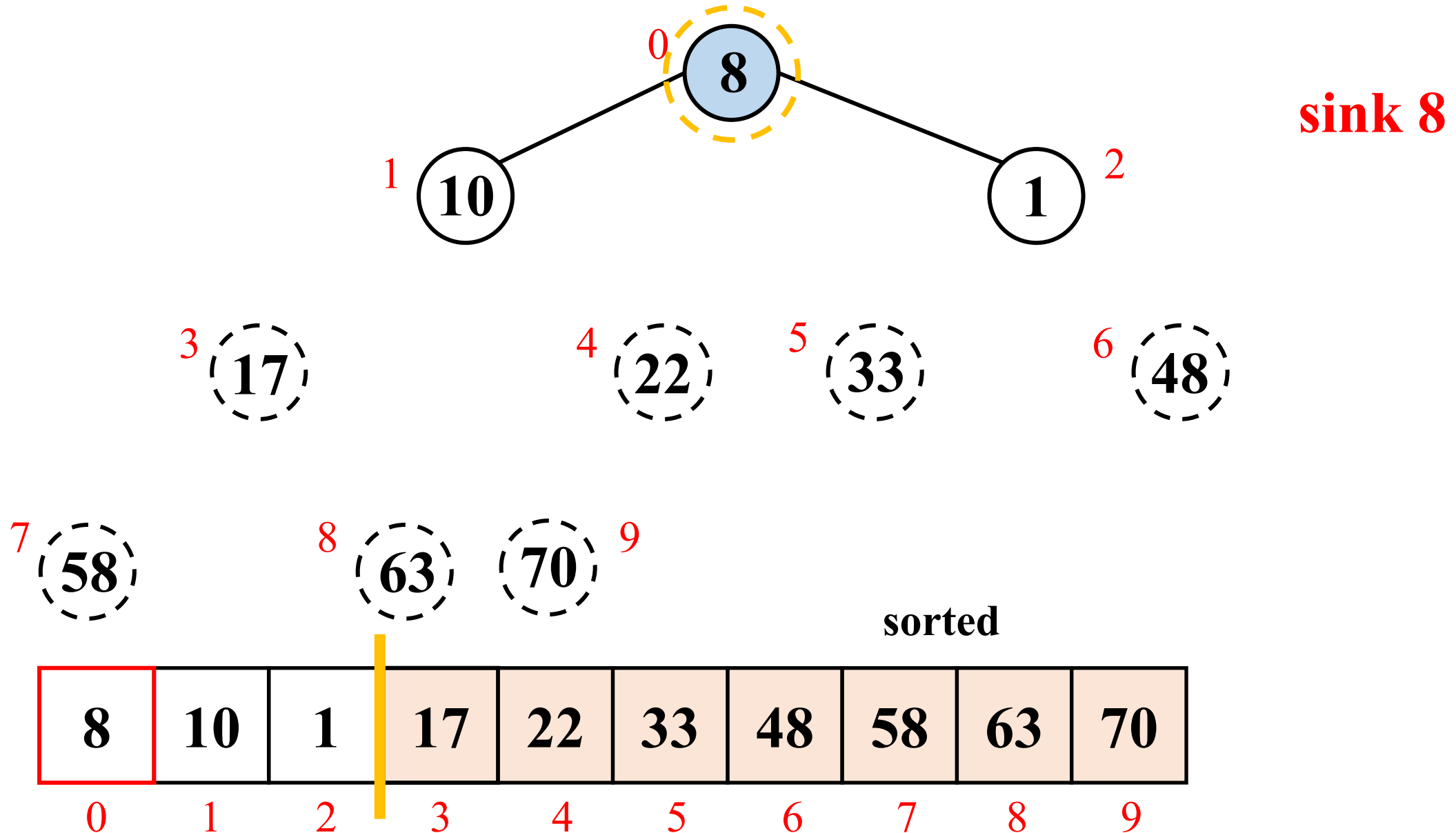




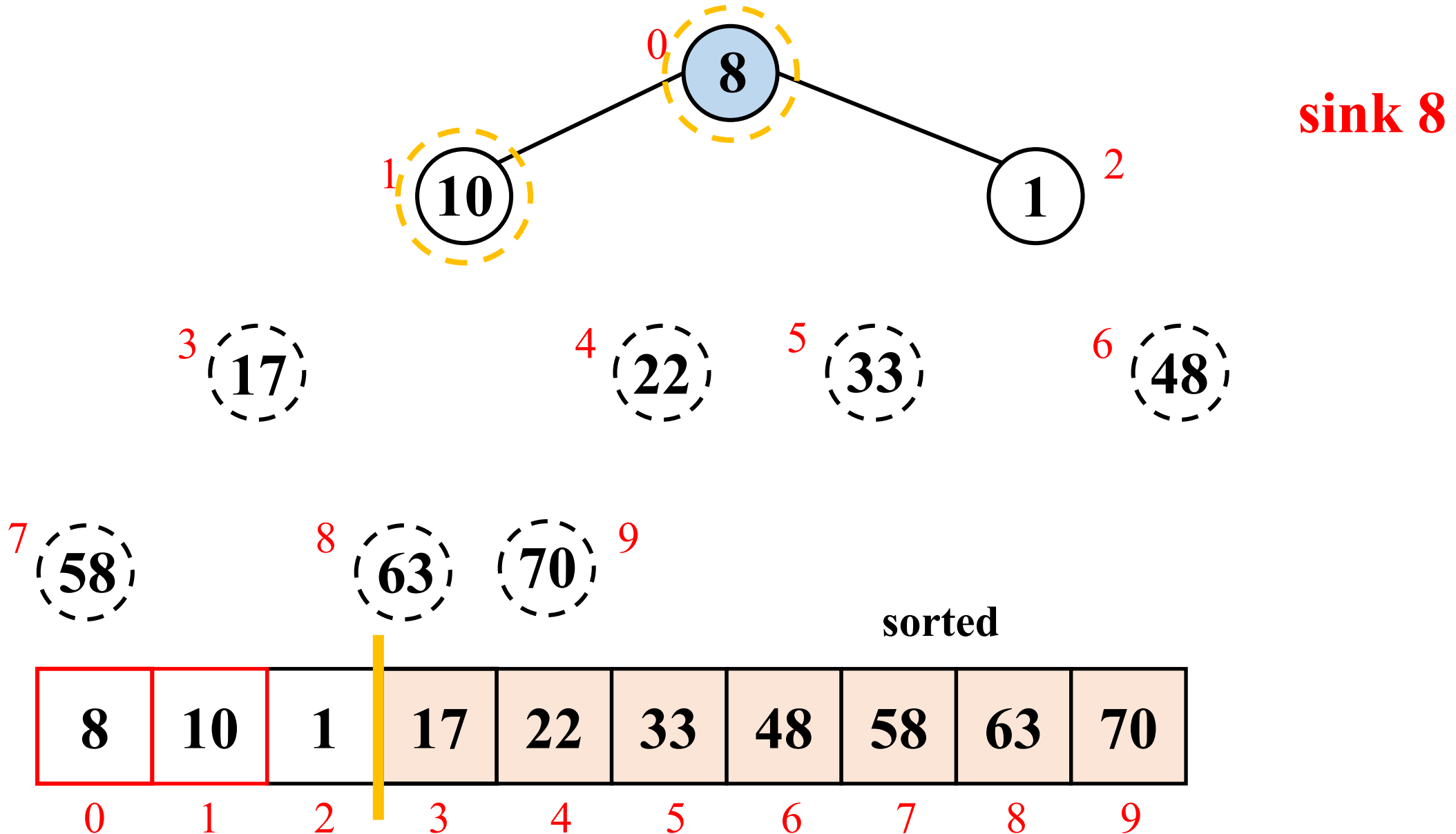
# Heapsort algorithms: array implementation



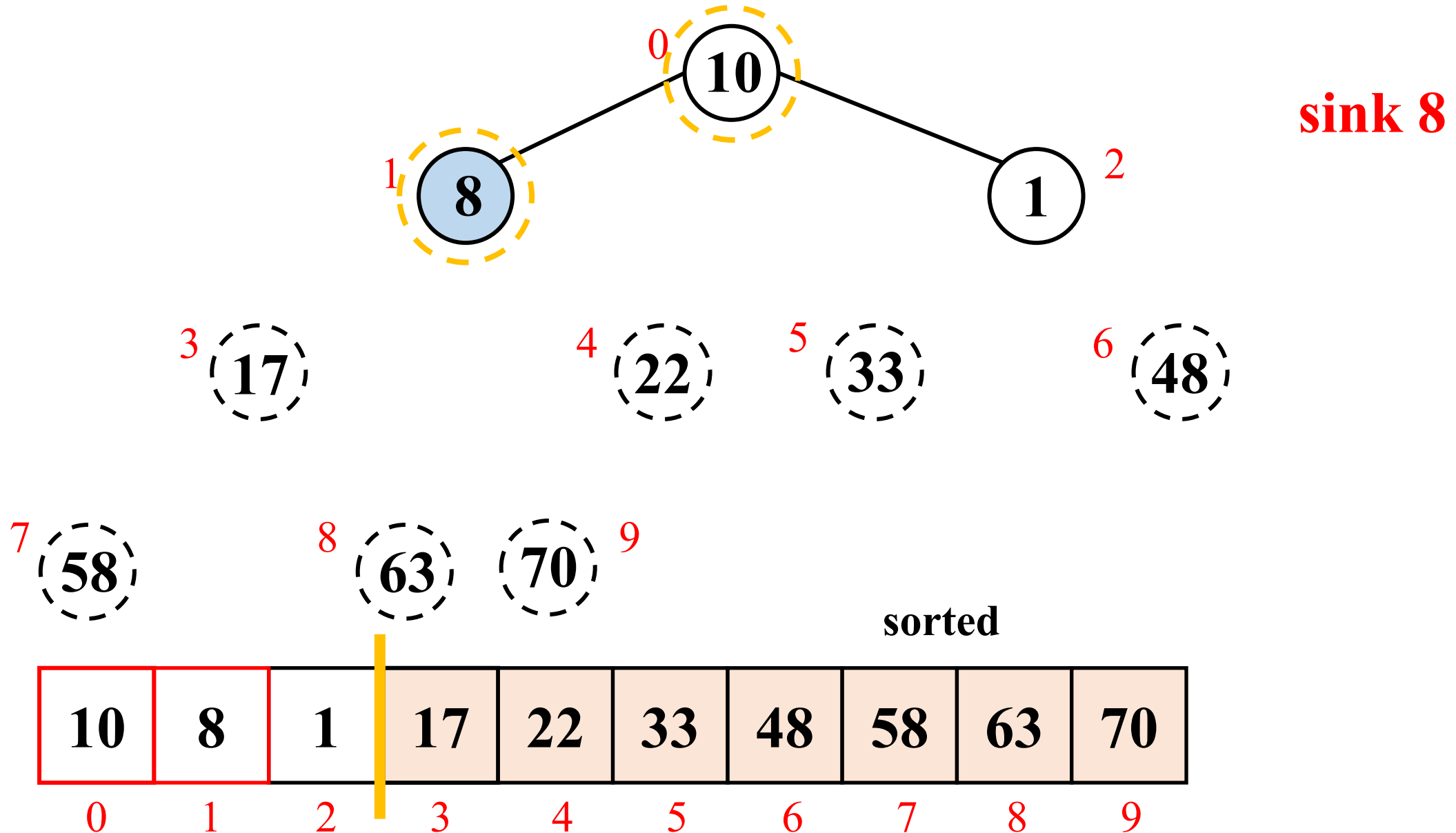
# Heapsort algorithms: array implementation



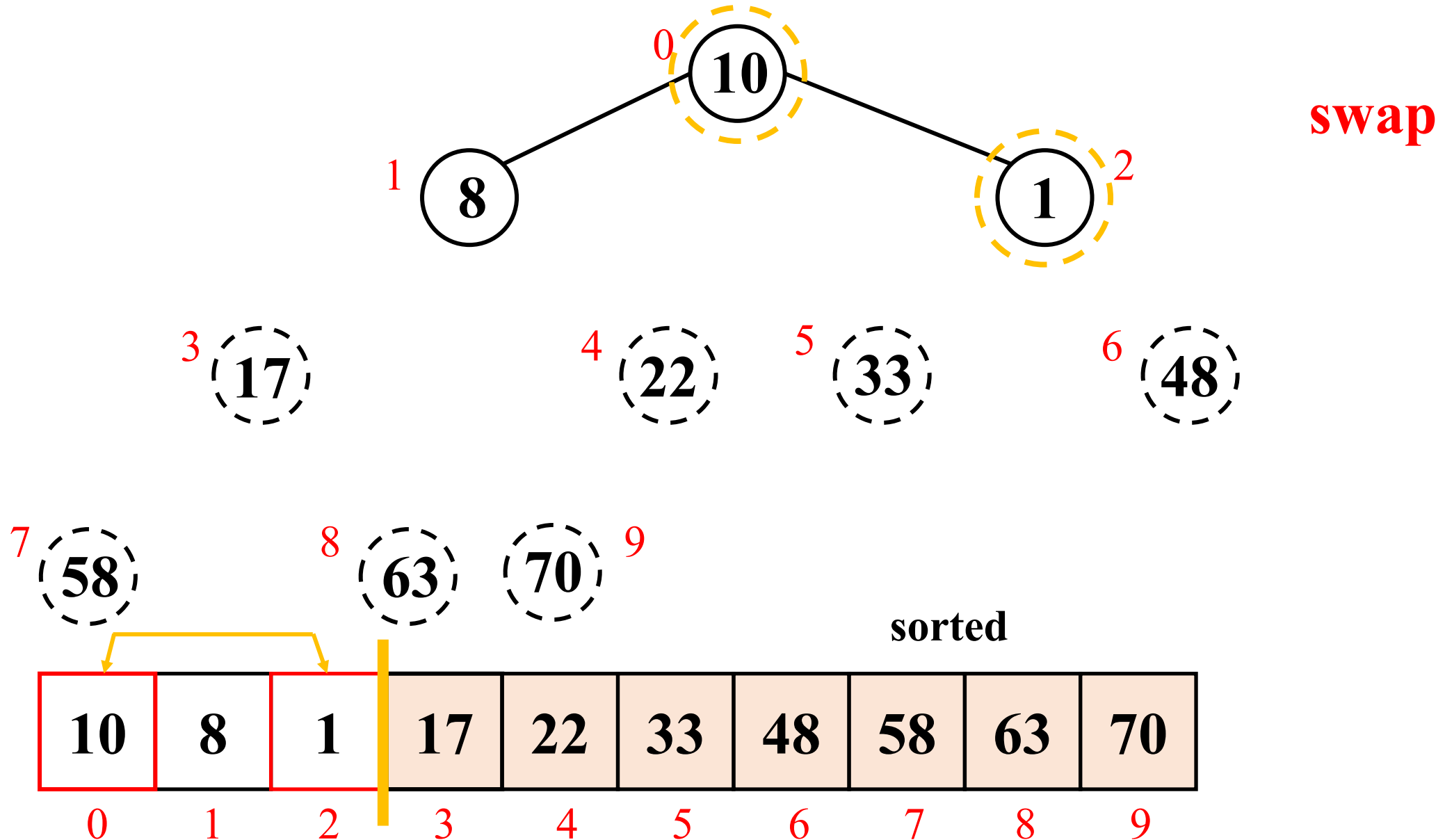
# Heapsort algorithms: array implementation



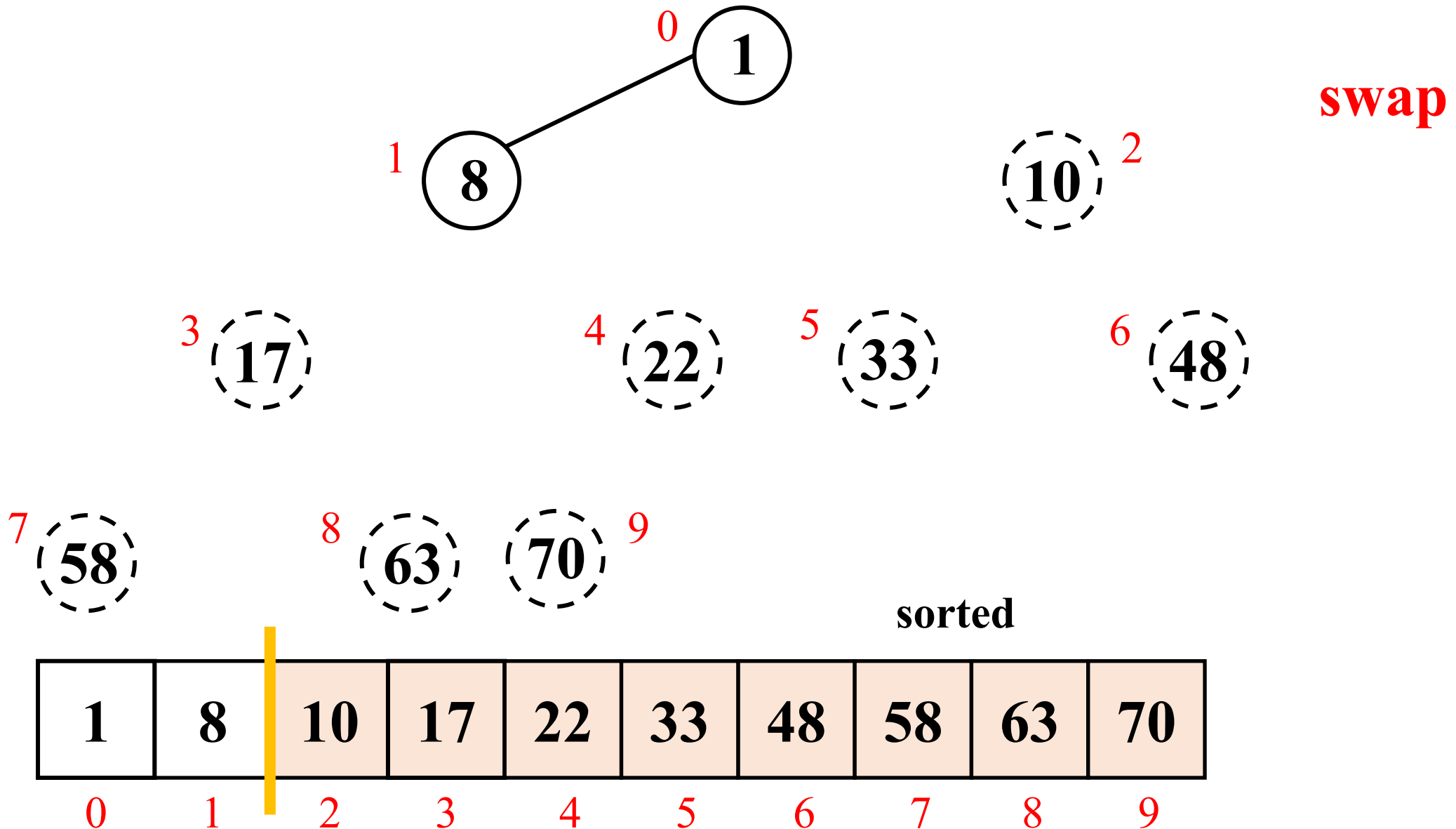
# Heapsort algorithms: array implementation



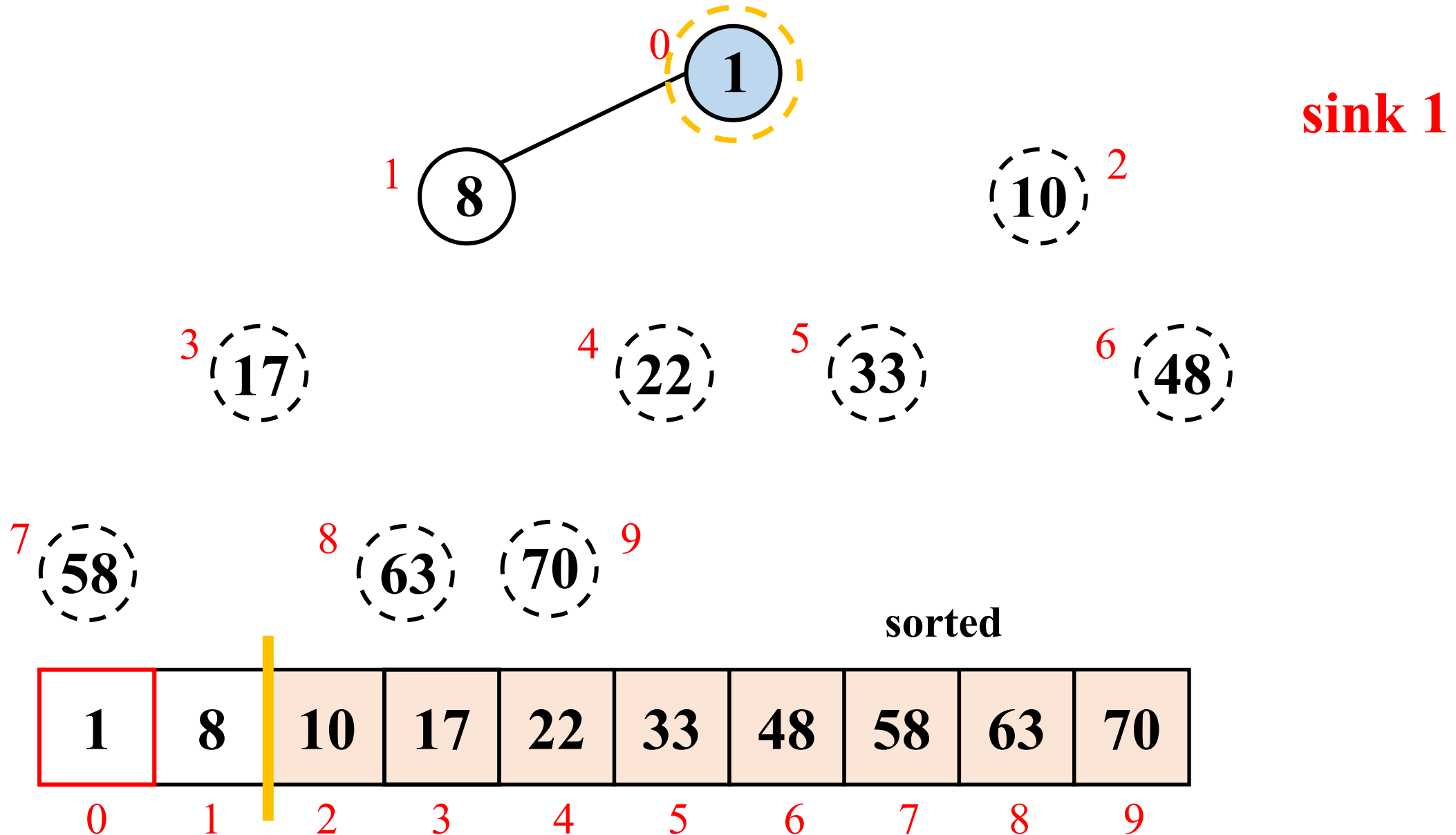
# Heapsort algorithms: array implementation



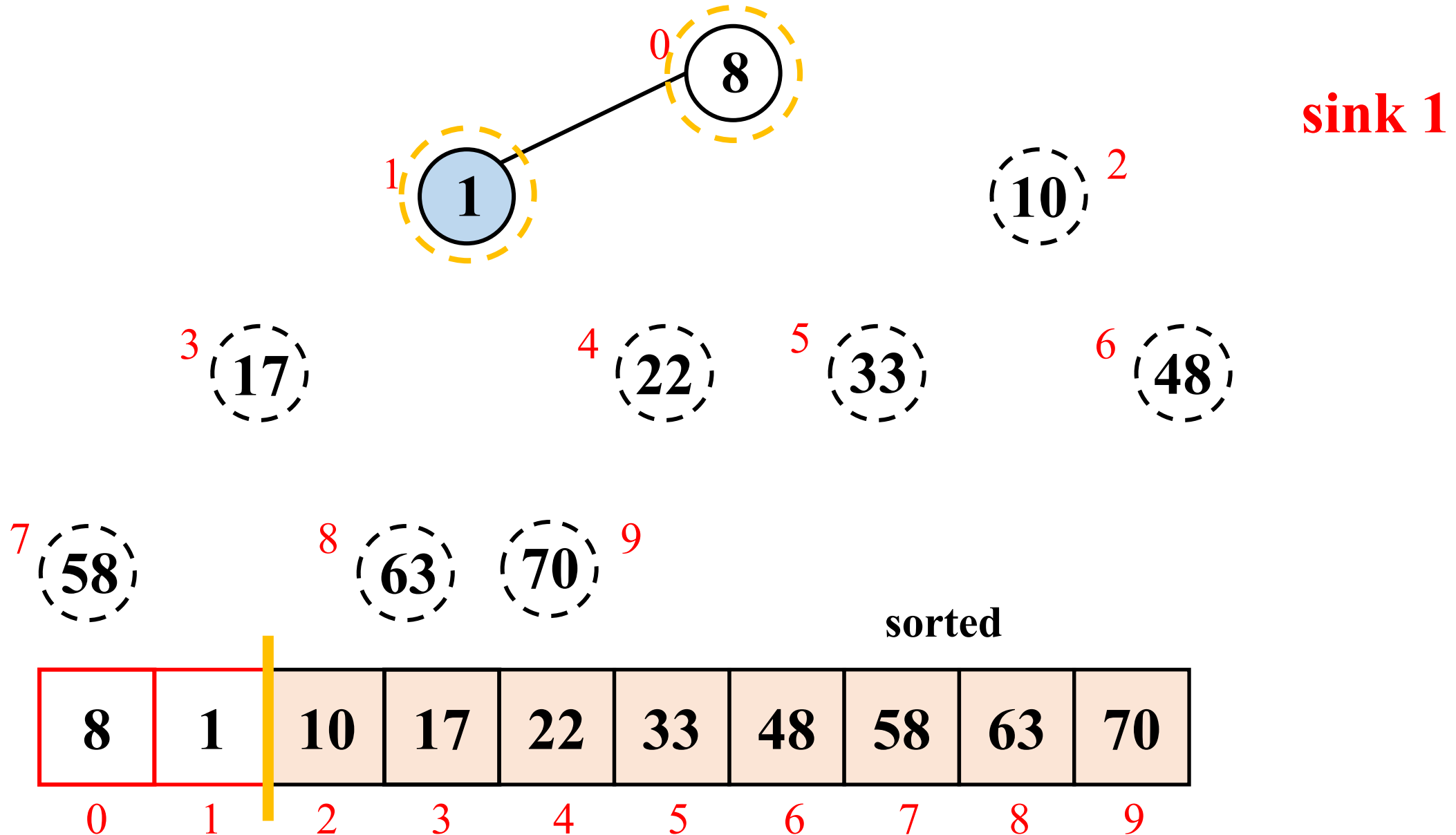
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation

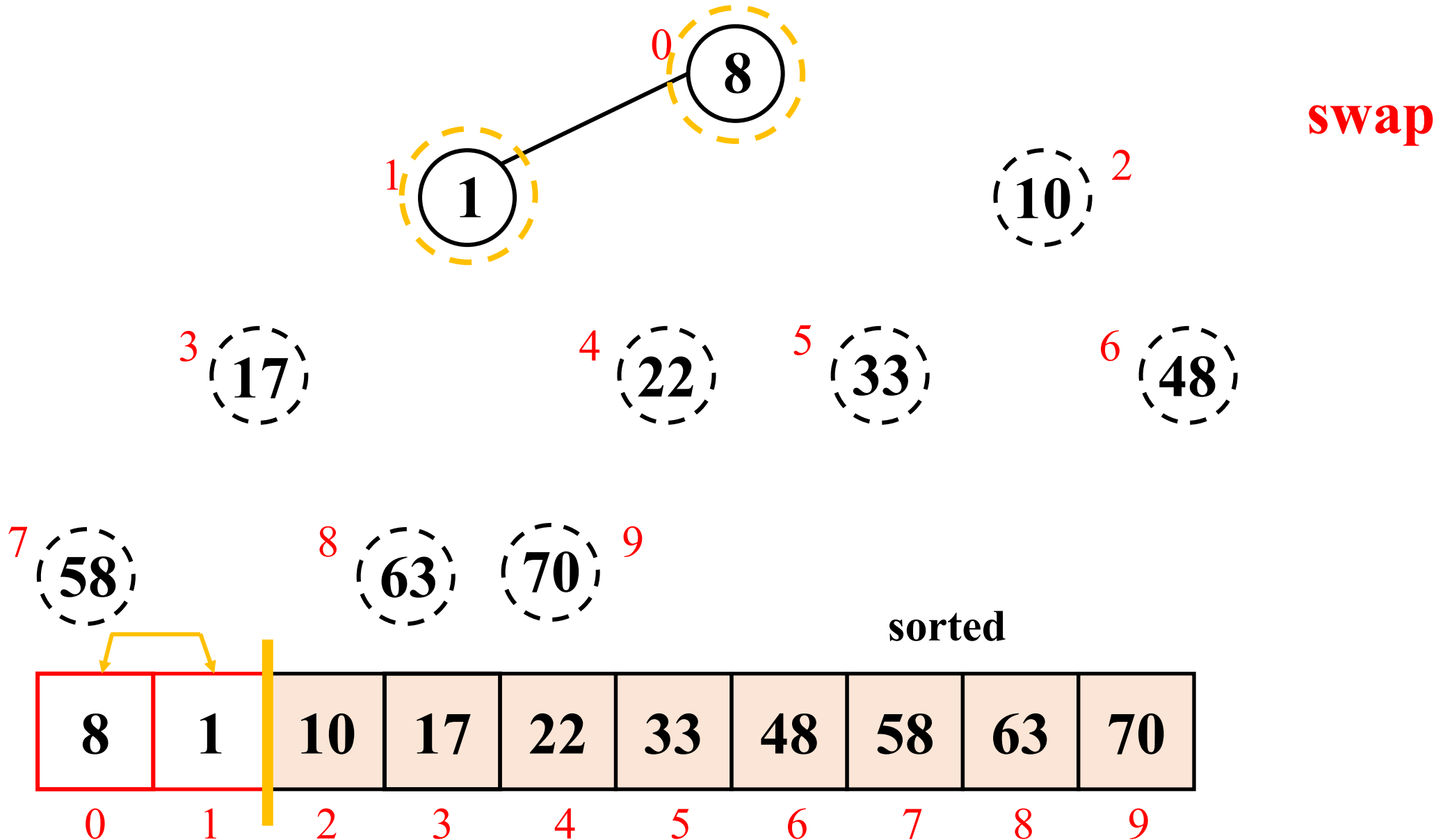


# Heapsort algorithms: array implementation

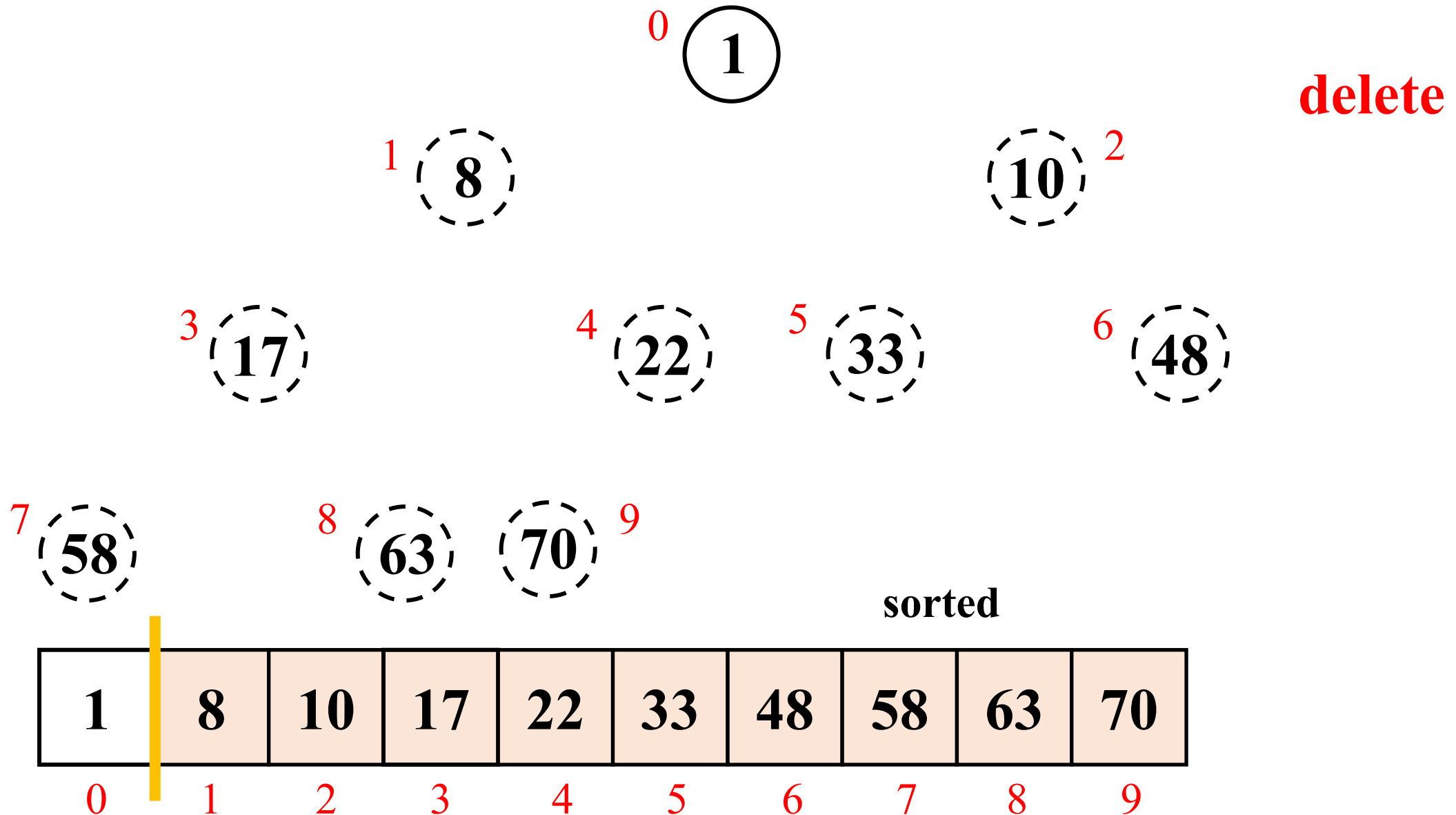




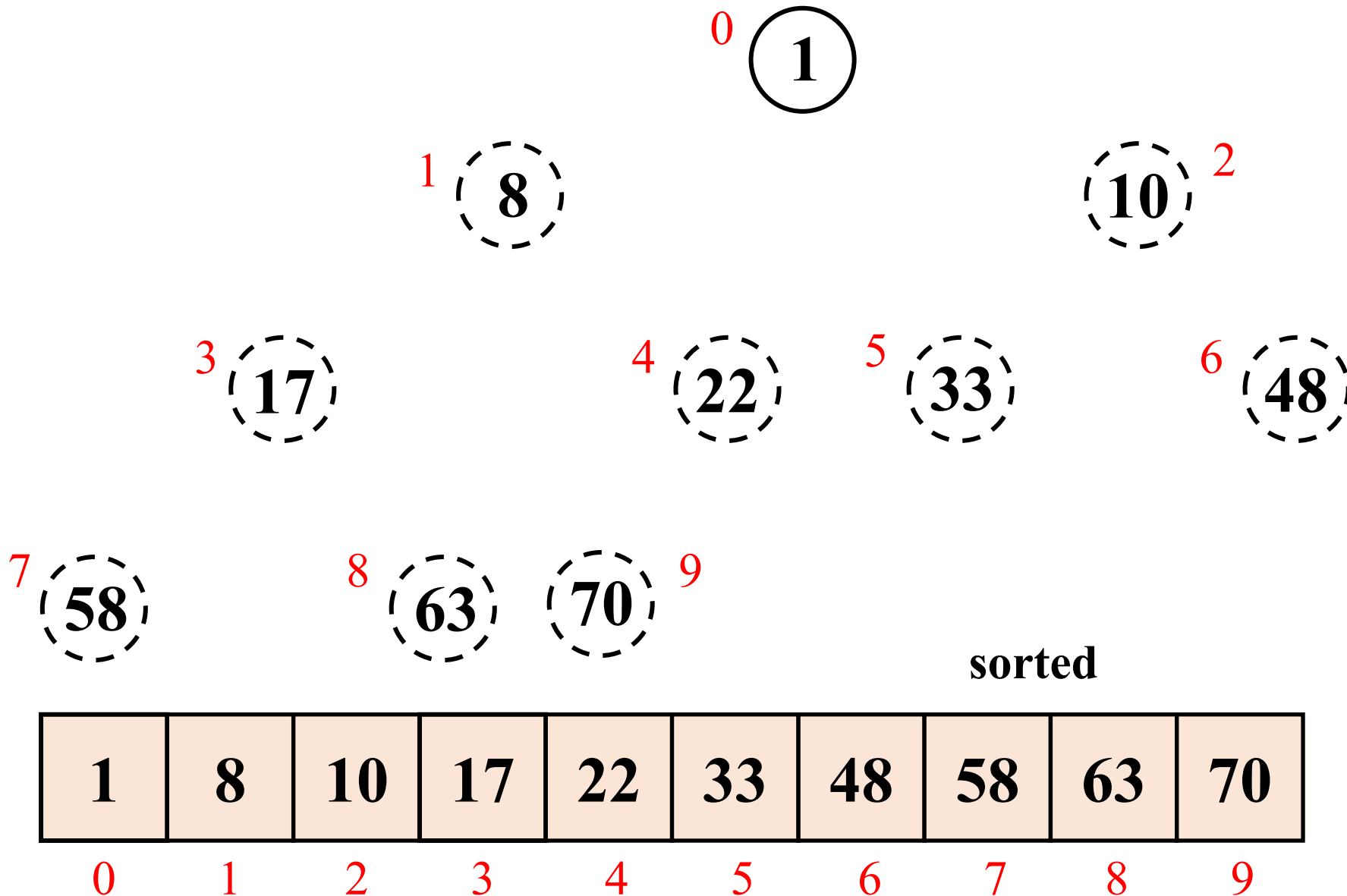
# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation



# Heapsort algorithms: array implementation



# Heapsort Algorithms: Array Implementation

---

**Algorithm 1** Heapsort.

---

```
1: function HEAPSORT(array  $a[0..n-1]$ )
2:      $k \leftarrow n - 1$                                 ▷ index of last leaf in heap
3:      $i \leftarrow \left\lfloor \frac{n-1}{2} \right\rfloor$ 
4:     while  $i \geq 0$  do
5:          $\text{sink}(a, i, k)$                                 ▷ let element sink to its correct place
6:          $i \leftarrow i - 1$                                 ▷ heap now build
7:     while  $k > 0$  do
8:          $\text{swap}(a[k], a[0])$ 
9:          $k \leftarrow k - 1$ 
10:         $\text{sink}(a, 0, k)$                                 ▷ let element at root sink to its correct place
11:    return  $a$                                             ▷ Array is now sorted
```

**(Build the heap)**

---

# Heapsort Algorithms: Array Implementation (Contd.)

---

**Algorithm 1** Heapsort.

---

```
1: function HEAPSORT(array  $a[0..n-1]$ )
2:      $k \leftarrow n - 1$                                 ▷ index of last leaf in heap
3:      $i \leftarrow \left\lfloor \frac{n-1}{2} \right\rfloor$ 
4:     while  $i \geq 0$  do
5:          $\text{sink}(a, i, k)$                                 ▷ let element sink to its correct place
6:          $i \leftarrow i - 1$                                 ▷ heap now build
7:     while  $k > 0$  do
8:          $\text{swap}(a[k], a[0])$ 
9:          $k \leftarrow k - 1$                                 (Delete the maximum repeatedly)
10:         $\text{sink}(a, 0, k)$                                 ▷ let element at root sink to its correct place
11:    return  $a$                                             ▷ Array is now sorted
```

---

# Time Complexity Analysis

- Heapsort runs in time in  $\Theta(n \log n)$  in the worst and average case.
- **Proof:**
  1. Building the heap:
    - The heap can be constructed in time  $\Theta(n \log n)$  with straightforward method. We can do better using the alternative bottom-up method with  $\Theta(n)$  time.
  2. Removing the maximum key:
    - Then heapsort repeats  $n$  times the deletion of the maximum key and restoration of the heap property (each restoration is logarithmic in the worst and average case).
    - The running time is  $\log(n) + \log(n - 1) + \log(n - 2) + \dots + 1 = \log(n!) \in \Theta(n \log n)$ .
- Thus, the total running time is  $\Theta(n \log n)$

# Notes on Heapsort

- Heapsort is not stable. Find an example that shows Heapsort is not stable.
- Heapsort is not in-place. we need extra space for the heap. Can we make it in-place?

# SUMMARY

- Heapsort
  - Algorithms
  - Illustrating Example
- Time Complexity Analysis

