

EECS2040 Data Structure Hw #5 (Chapter 6 Graph)

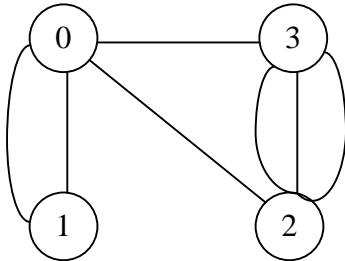
due date 5/30/2022

by 108011235 陳昭維

(Part 1: 2% of final Grade)

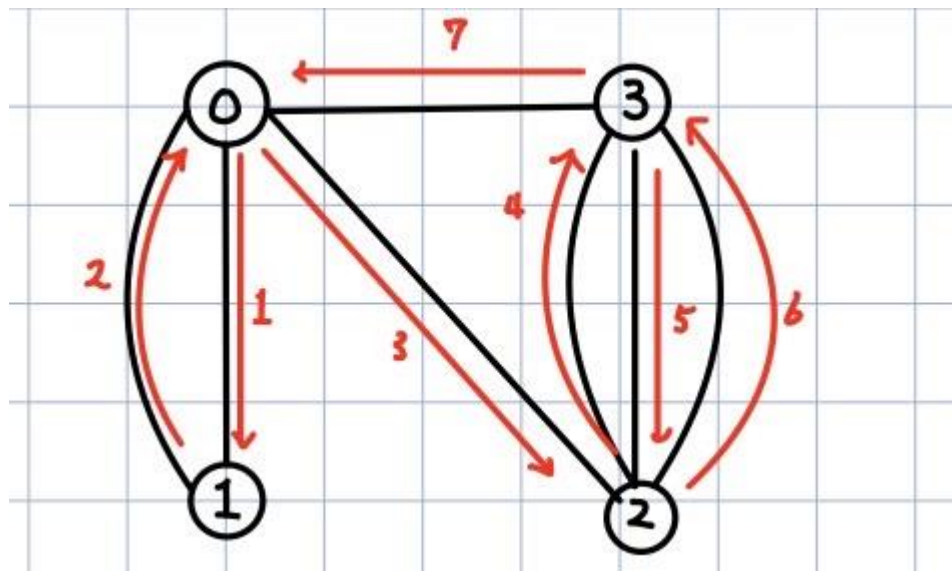
Part 1

1. (10%) Does the multigraph below have an Eulerian walk? If so, find one.

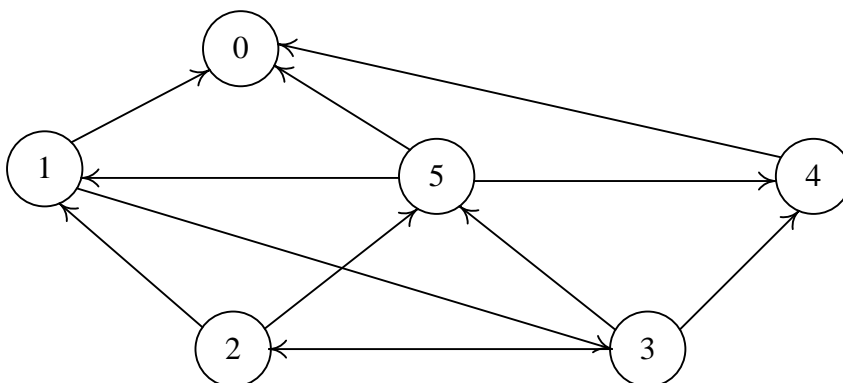


<answer>

Yes there is an Eulerian walk



2. (10%) For the digraph below obtain
- The in-degree and out-degree of each vertex
 - Its adjacency-matrix
 - Its adjacency-list representation
 - Its strongly connected components



<answer>

(a)

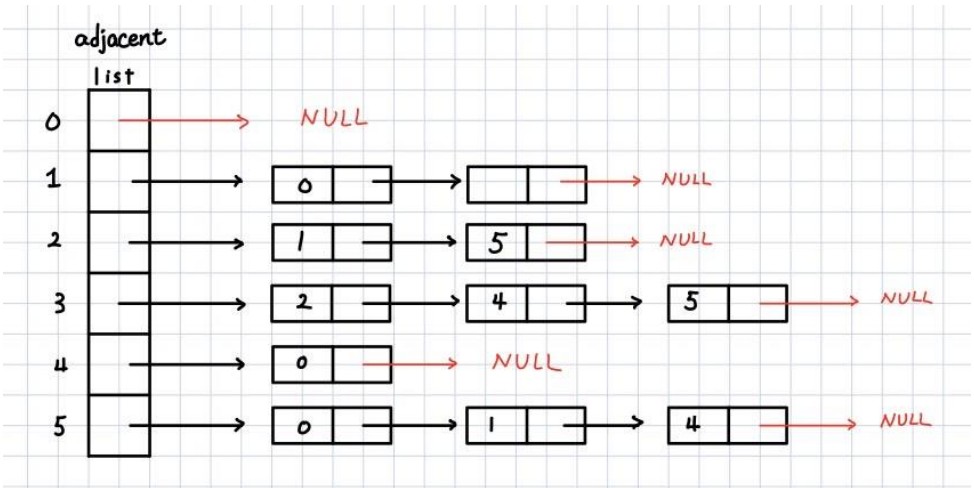
Vertex	In-degree	Out-degree
0	3	0
1	2	2
2	1	2
3	1	3
4	2	1
5	2	3

(b)

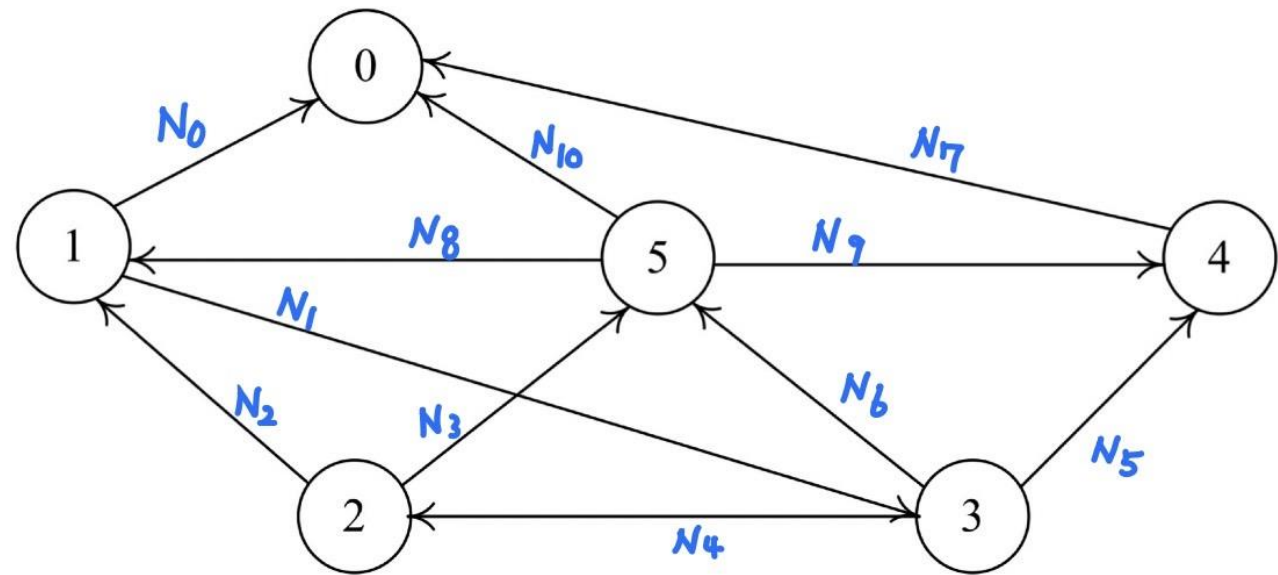
This column and row is indexed by (0~5) and (0~5), the graph is generated by LaTeX

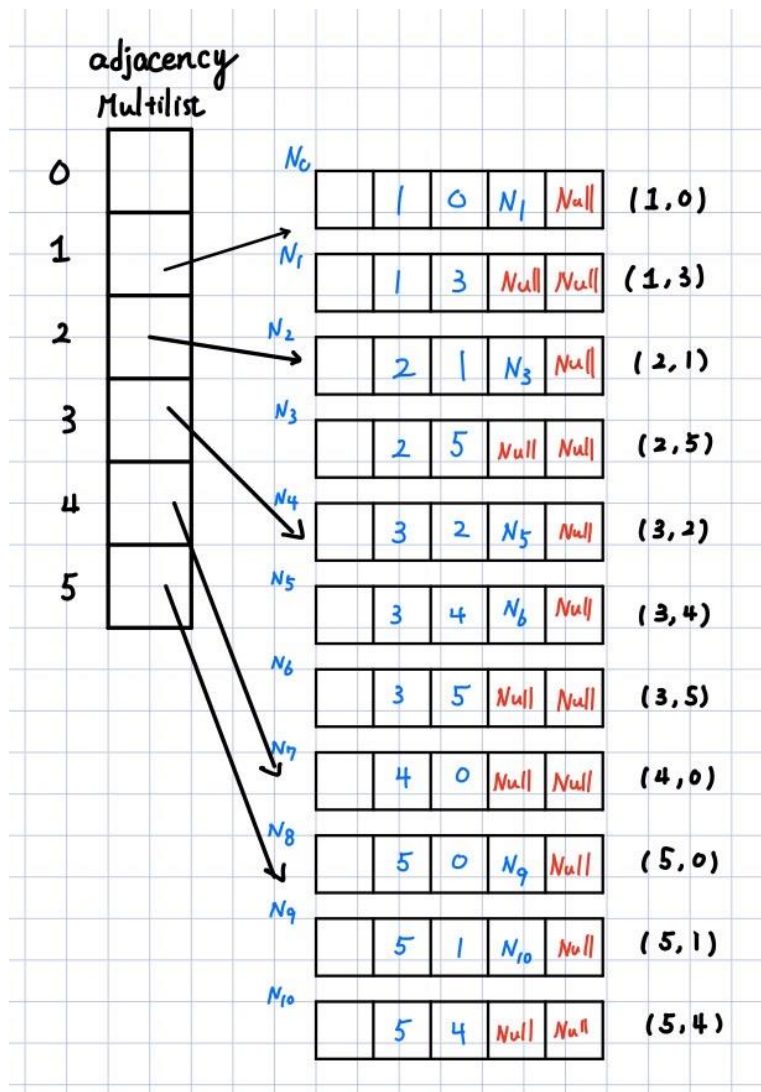
$$\begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

(c)

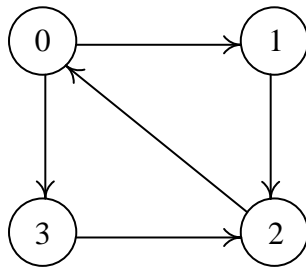


(d)





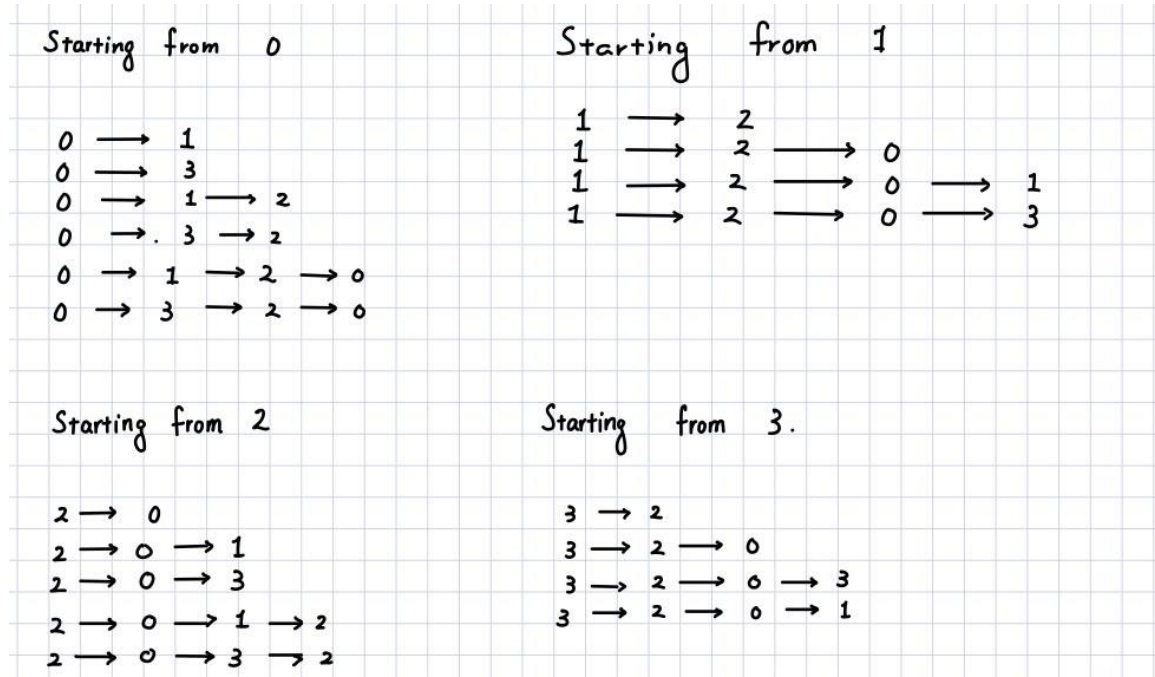
3. (10%) Is the digraph below strongly connected? List all the simple paths.



<answer>

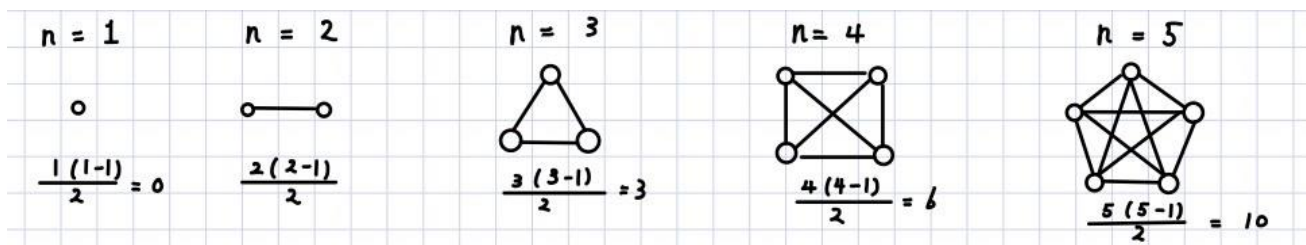
Yes, the graph is strongly connected.

Simple path of the graph:



4. (10%) Draw the complete undirected graphs on one, two, three, four, and five vertices. Prove that the number of edges in an n -vertex complete graph is $n(n-1)/2$.

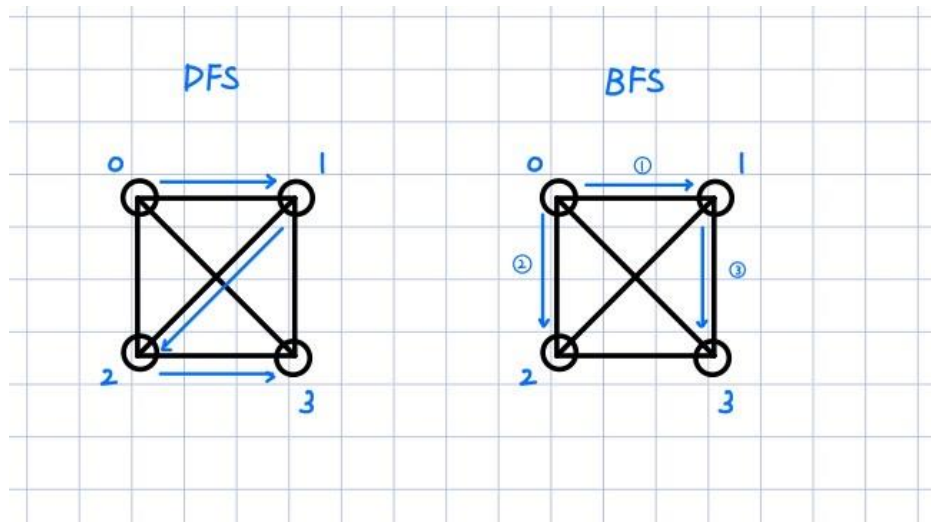
<answer>



For each n vertex to be connected to other vertex there are $n - 1$ edges, there are n vertices but so there are $\frac{n \times (n-1)}{2}$ edges, since 2 of each are replicates.

5. (4%) Apply **depth-first** and **breadth-first** searches to the **complete graph on four vertices**. Assume that vertices are numbered 0 to 3, are stored in increasing order in each list in the adjacency-list representation, and both traversals begin at vertex 0. List the vertices in the order they would be visited.

<answer>



Both the order of traversal for DFS and BFS are $0 \rightarrow 1 \rightarrow 2 \rightarrow 3$

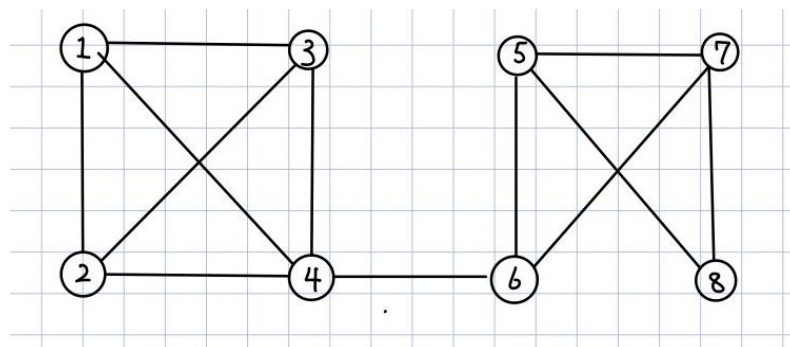
6. (6%) Let G be a graph whose vertices are the integers 1 through 8, and let the adjacent vertices of each vertex be given by the table below:

Vertex	Adjacent Vertices
1	(2, 3, 4)
2	(1, 3, 4)
3	(1, 2, 4)
4	(1, 2, 3, 6)
5	(6, 7, 8)
6	(4, 5, 7)
7	(5, 6, 8)
8	(5, 7)

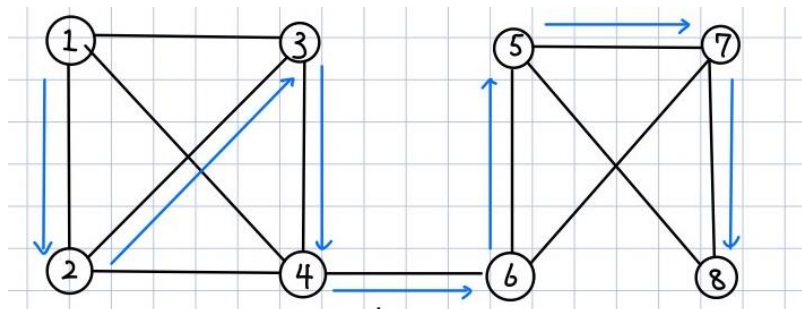
Assume that, in a traversal of G , the adjacent vertices of a given vertex are returned in the same order as they are listed in the table above.

<answer>

- (a) Draw G .

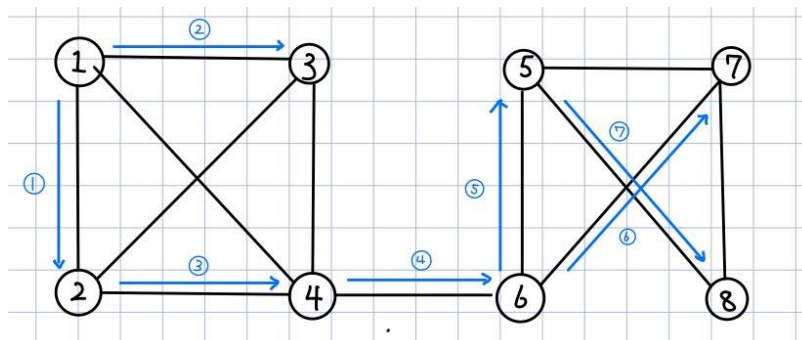


- (b) Give the sequence of vertices of G visited using a DFS traversal starting at vertex 1.



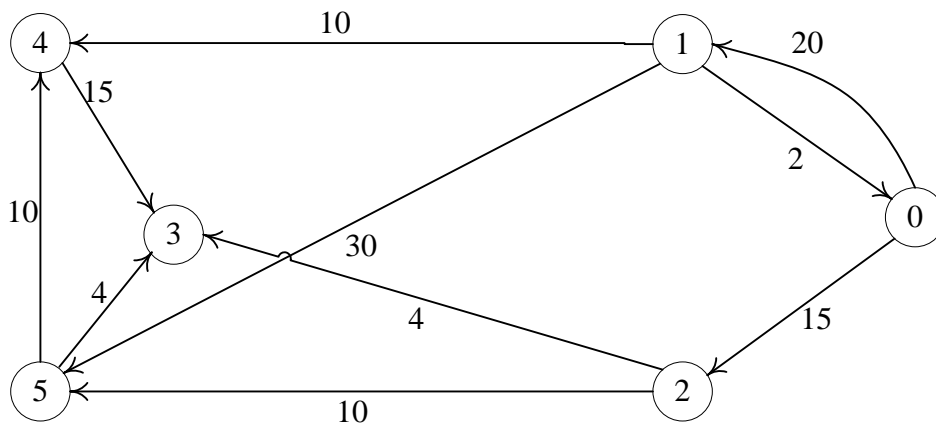
$1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 6 \rightarrow 5 \rightarrow 7 \rightarrow 8$

- (c) Give the sequence of vertices visited using a BFS traversal starting at vertex 1.



$1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 6 \rightarrow 5 \rightarrow 7 \rightarrow 8$

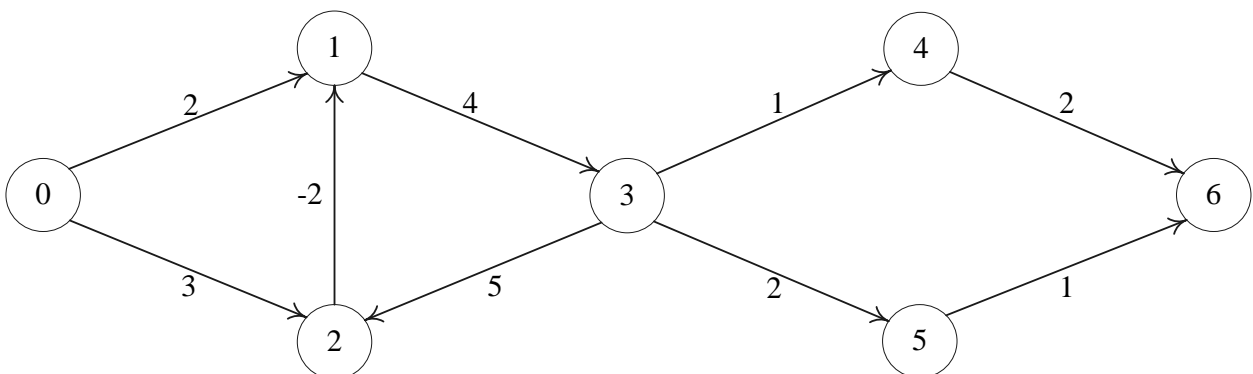
7. (10%) Use ShortestPath (Program 6.8) (Dijkstra's algorithm) to obtain, in nondecreasing order, the **lengths** and the **paths** of the shortest paths from vertex 0 to all remaining vertices in the graph below.



<answer>

	path	lengths
1	0 → 1	20
2	0 → 2	15
3	0 → 2 → 3	19
4	0 → 1 → 4	30
5	0 → 2 → 5	25

8. (10%) Using the directed graph below, explain why ShortestPath (Program 6.8) will not work properly. What is the shortest path between vertices 0 and 6?

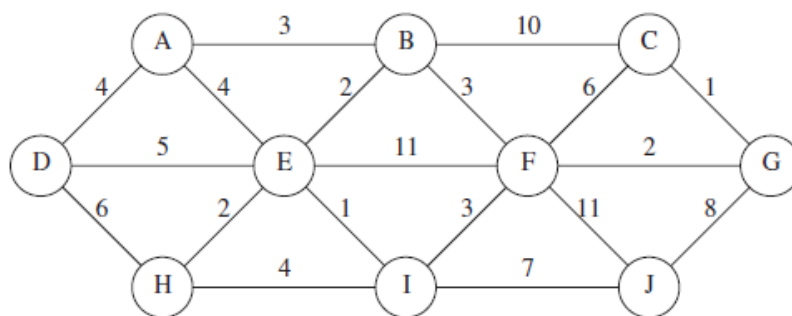


<answer>

The shortest path from 0 to 6 is $0 \rightarrow 2 \rightarrow 1 \rightarrow 3 \rightarrow 4 \rightarrow 6$ with length 8.

The Dijkstra's algorithm wouldn't work because, negative weight will not guaranteed a closed vertex is indeed minimal.

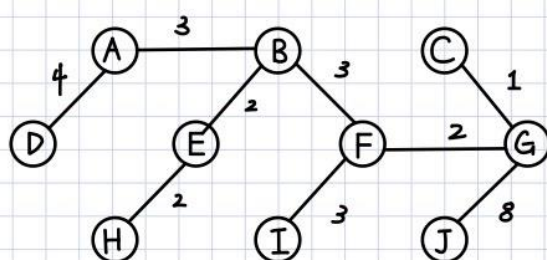
9. (10%) For the weighted graph G shown below,



(a) Find a minimum spanning tree for the graph using both Prim's and Kruskal's algorithms.

<answer>

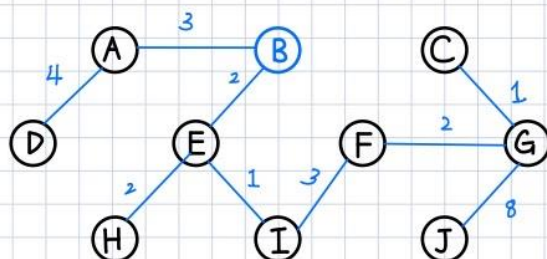
Kruskal's



total cost : 28

Prim's

Suppose we pick (B) as first vertex in tree



total cost : 28

(b) Is this minimum spanning tree unique? Why?

<answer>

No the minimum spanning tree is not unique, it varies with tie-break mechanism, but the total cost shouldn't vary for distinct minimal spanning tree.

10. (10%) Does the following set of precedence relations ($<$) define a **partial order** on the elements 0 through 4? Why?

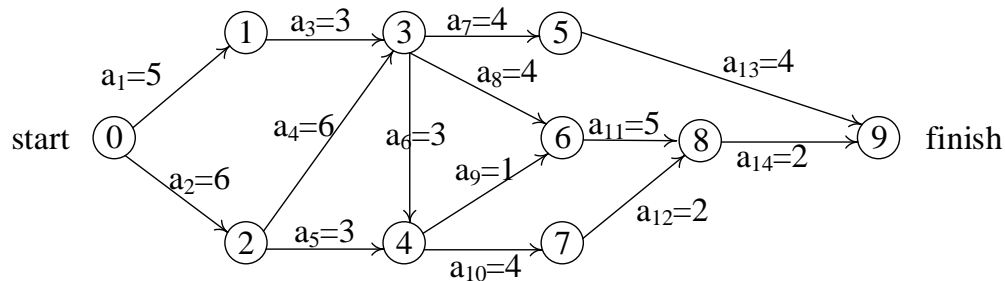
$0 < 1; 1 < 3; 1 < 2; 2 < 3; 2 < 4; 4 < 0$

<answer>

Assume a partial order is defined on the relations, since partial order guarantees transitive, thus from above relations, we get $0 < 0$ which is a mapping from 0 to 0, which violates the irreflexive properties of partial order. Thus, it doesn't define a partial order.

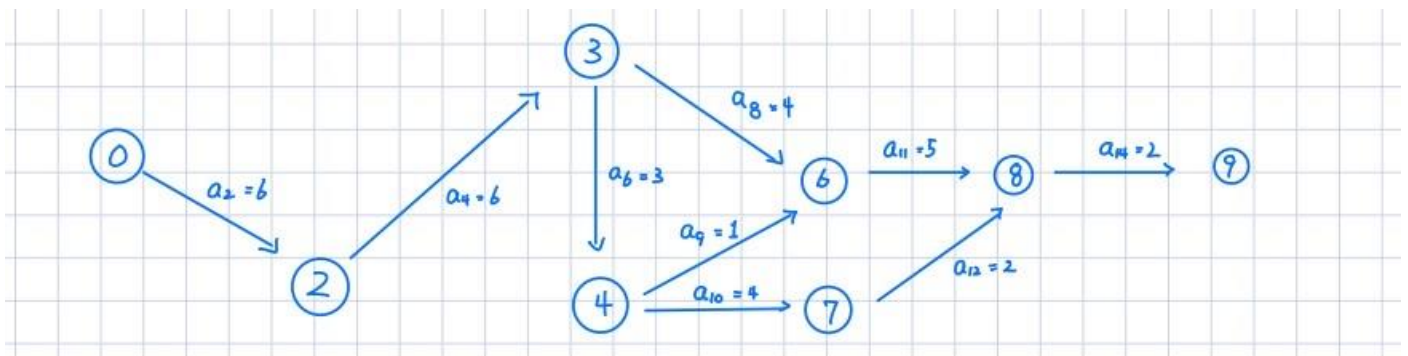
11. (10%) For the AOE network shown below,

- (a) Obtain the early, $e(a_i)$, and late, $l(a_i)$, start times for each activity. Use the forward-backward approach.
- (b) What is the earliest time the project can finish?
- (c) Which activities are critical? Fill the table below for answers to (a), (b), and (c).



activity	Early time	Late time	slack	critical
	$e(a_i)$	$l(a_i)$	$l(a_i) - e(a_i)$	
a1	0	4	4	No
a2	0	0	0	Yes
a3	5	9	4	No
a4	6	6	0	Yes
a5	6	12	6	No
a6	12	12	0	Yes
a7	12	15	3	No
a8	12	12	0	Yes
a9	15	15	0	Yes
a10	15	15	0	Yes
a11	16	16	0	Yes
a12	19	19	0	Yes
a13	16	19	3	No
a14	21	21	0	Yes

- (d) Is there any single activity whose speed-up would result in a reduction of the project finish time?



The single activity that can speed up on the critical path is a2, a4, a14