

CAUSWE Compiler, 2021

# Compiler Project #2

## Documentation

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## How to run

### Requirements

- Python >= 3.8.2
- Ubuntu or MacOS or Windows with Python

### Execution

run.sh file\_name

```
run.sh
#!/bin/bash

python3 analyzer/lexical_analyzer.py $1 $1.lex_out
python3 analyzer/syntax_analyzer.py $1.lex_out
```

“analyzer/lexical\_analyzer.py” (which is made from Assignment #1) reads “file\_name” and saves it’s output to “file\_name.lex\_out”.

Then, “analyzer/syntax\_analyzer.py” reads “file\_name.lex\_out” and checks if it’s syntactically acceptable or not.

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# Context Free Grammar G

## What is Changed?

1. Start Symbol
  - There are three production rules from the starting symbol CODE. So we binded them with CODE0.
2. Removed Ambiguity of EXPR
  - An ambiguity exists between addsub and multdiv.  
Ex) "id addsub id multdiv id" can have 2 parse trees.
  - So we set a disambiguiting rule that multdiv has a higher priority than addsub.
3. Removed Ambiguity of EXPR and COND
  - At EXPR, ambiguity rules exist between addsub and multdiv.  
Ex) "id addsub id addsub id"
  - So we set a disambiguiting rule that right addsub and multdiv has a higher priority than left addsub and multdiv.
  - At COND, ambiguity rule exists between comps.  
Ex) "boolstr comp boolstr comp boolstr"
  - So we set a disambiguiting rule that right comp has a higher priority than left comp.

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## CFG G (Ambiguity Removed)

```
CODE0 -> CODE1
CODE1 -> VDECL CODE1
CODE1 -> FDECL CODE1
CODE1 -> CDECL CODE1
CODE1 -> "
VDECL -> vtype id semi
VDECL -> vtype ASSIGN semi
ASSIGN -> id assign RHS
RHS -> EXPR1
RHS -> literal
RHS -> character
RHS -> boolstr
EXPR1 -> EXPR2 addsub EXPR1
EXPR1 -> EXPR2
EXPR2 -> EXPR3 multdiv EXPR2
EXPR2 -> EXPR3
EXPR3 -> lparen EXPR1 rparen
EXPR3 -> id
EXPR3 -> num
FDECL -> vtype id lparen ARG rparen lbrace BLOCK RETURN rbrace
ARG -> vtype id MOREARGS
ARG -> "
MOREARGS -> comma vtype id MOREARGS
MOREARGS -> "
BLOCK -> STMT BLOCK
BLOCK -> "
STMT -> VDECL
STMT -> ASSIGN semi
STMT -> if lparen COND1 rparen lbrace BLOCK rbrace ELSE
STMT -> while lparen COND1 rparen lbrace BLOCK rbrace
COND1 -> COND2 comp COND1
COND1 -> COND2
COND2 -> boolstr
ELSE -> else lbrace BLOCK rbrace
ELSE -> "
RETURN -> return RHS semi
CDECL -> class id lbrace ODECL rbrace
ODECL -> VDECL ODECL
ODECL -> FDECL ODECL
ODECL -> "
```

# SLR Parsing Table from CFG G

41	#53	#54			F <sub>25</sub>	#51	#52	F <sub>25</sub>	49	50		47	48	
42					F <sub>20</sub>									
43	#55				F <sub>12</sub>									
44		F <sub>12</sub>			F <sub>12</sub>									
45	F <sub>14</sub>				F <sub>14</sub>									
46	F <sub>16</sub>				F <sub>16</sub>									
47								F <sub>57</sub>					56	
48	#53	#54			F <sub>25</sub>	#51	#52	F <sub>25</sub>	49	50		58	48	
49	F <sub>26</sub>	F <sub>26</sub>			F <sub>26</sub>	F <sub>26</sub>	F <sub>26</sub>	F <sub>26</sub>						
50		#59												
51				#60										
52				#61										
53		#62											11	
54		#15												
55		#63												
56														
57	#28	#22	#23	#24	#27	#29	#64				65	21	25	26
58								F <sub>24</sub>						
59	F <sub>27</sub>	F <sub>27</sub>				F <sub>27</sub>	F <sub>27</sub>	F <sub>27</sub>						
60				#68				F <sub>27</sub>					66	67
61				#68									69	67
62		#13	#15											
63					F <sub>23</sub>		#43							70
64	F <sub>19</sub>				F <sub>19</sub>			F <sub>19</sub>						
65		#71												
66														
67		#72												
68					F <sub>31</sub>			F <sub>31</sub>						
69								F <sub>31</sub>						
70														
71														
72														
73				#75									76	67
74														
75	#53	#54			#77									
76					F <sub>25</sub>	#51	#52	F <sub>25</sub>	49	50		78	48	
77	#53	#54												
78														
79														
80	F <sub>34</sub>	F <sub>34</sub>			F <sub>34</sub>	F <sub>34</sub>	F <sub>34</sub>	#83	F <sub>34</sub>					82
81	F <sub>29</sub>	F <sub>29</sub>			F <sub>29</sub>	F <sub>29</sub>	F <sub>29</sub>	F <sub>29</sub>						
82	F <sub>28</sub>	F <sub>28</sub>			F <sub>28</sub>	F <sub>28</sub>	F <sub>28</sub>	F <sub>28</sub>						
83														
84	#53	#54			F <sub>25</sub>	#51	#52	F <sub>25</sub>	49	50		85	48	
85														
86	F <sub>33</sub>	F <sub>33</sub>			F <sub>33</sub>	F <sub>33</sub>	F <sub>33</sub>	F <sub>33</sub>						

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## + For Addition

We made a javascript script that parses <http://jsmachines.sourceforge.net/machines/sl.r.html>'s LR Table and converts it into python dict.

```
const table = document.getElementById("lrTableView")
const thead = Array.from(table.getElementsByTagName("thead")[0].children[2].children)
const rows = Array.from(table.getElementsByTagName("tbody")[0].children)
var dictionary = rows.map((rowHTML)=>{
    var result = {}
    const row = Array.from(rowHTML.children)
    row.forEach((colHTML, colIndex)=>{
        if(colIndex==0){
            return
        } else if (colHTML.innerHTML == "&nbsp;") {
            result[thead[colIndex-1].innerText] = null
        } else {
            result[thead[colIndex-1].innerText] = colHTML.innerText
        }
    })
    return result
})
dictionary = JSON.stringify(dictionary)
dictionary = dictionary.replaceAll("null", "None")
console.log(dictionary)
```

This script runs on Chrome/Developer Tools/Sources/Snippets

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# How this Syntax Analyzer Work

## Architecture

There's a Stack class. This class is a normal stack that has a push and pop method. It will store state and left substring.

There's a SLRtable class. This class contains a SLR-parsing table to parse given CFG. If it gets a symbol, The corresponding symbol of the state that the first number of stacks points to returns Decision. And follow Decision like shift and goto or reduce. Of course, if the table's return is None which means it can't choose any decision, it prints a fail and reports an error.

There's a CFG class. This class contains list when reduced by following rules. List return grammars to change and how much pop the stack.

For the main function,

1. It only reads Token in its output from the lexical analyzer.
2. Then push the sub string and state in each stack by moving pipe.
3. Repeat it until the table returns 'acc' or 'None'.
4. If it is 'acc', it prints 'ACCEPT', if it is 'None', it prints 'Fail', and it prints an error report.

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## Test Inputs/Outputs

### Successfully Accepted Tokens

input file(test.txt)	.lex_out file (test.txt.lex_out)	console
int a; int c() { int a; return a; }	vtype int id a semi ; vtype int id c lparen ( rparen ) lbrace { vtype int id a semi ; return return id a semi ; rbrace }	Lexical Analyzer : Accepted Lexical Analyzer : Token file saved at test.txt.lex_out successfully Syntax Analyzer : Accepted

input file(test.txt)	.lex_out file (test.txt.lex_out)	console
int divide(int a, int b) { int result; if(true == true) { result = 0; } else { result = a/b; } return c; }	vtype int id divide lparen ( vtype int id a comma , vtype int id b rparen ) lbrace { vtype int id result semi ; if if lparen ( boolstr true comp == boolstr true rparen ) lbrace { id result assign = num 0 semi ; rbrace } else else rbrace {	Lexical Analyzer : Accepted Lexical Analyzer : Token file saved at test.txt.lex_out successfully Syntax Analyzer : Accepted

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	<pre> id      result assign  = id      a multdiv / id      b semi    ; rbrace  } return return id      c semi    ; rbrace  } </pre>	
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input file(test.txt)	.lex_out file (test.txt.lex_out)	console
<pre> int divide(int a, int b) {     int result;     return result; }  class Main {     int state = 0;      int main(int arg1, int arg2) {         int state = 1;         int result = 0;         int a = 20;         int b = 2;         char c = 't';         char str = "some_string";          if(true == false) {             result = a/(a+b/a);             result = a/b;         } else {             result = -1;         }          return result;     } } </pre>	<pre> vtype  int id     divide lparen ( vtype  int id     a comma , vtype  int id     b rparen ) lbrace { vtype  int id     result semi   ; return return id     result semi   ; rbrace } class  class id     Main lbrace { vtype  int id     state assign = num   0 semi   ; vtype  int id     main lparen ( vtype  int id     arg1 comma , vtype  int id     arg2 rparen ) lbrace { vtype  int id     state assign = num   1 semi   ; vtype  int </pre>	<p>Lexical Analyzer : Accepted      Lexical Analyzer : Token file saved at      test.txt.lex_out successfully      Syntax Analyzer : Accepted</p>

```
id      result
assign =
num    0
semi   ;
vtype  int
id     a
assign =
num    20
semi   ;
vtype  int
id     b
assign =
num    2
semi   ;
vtype  char
id     c
assign =
character 't'
semi   ;
vtype  char
id     str
assign =
literal "some_string"
semi   ;
if     if
lparen (
boolstr true
comp   ==
boolstr false
rparen )
lbrace {
id     result
assign =
id     a
multdiv /
lparen (
id     a
addsub +
id     b
multdiv /
id     a
rparen )
semi   ;
id     result
assign =
id     a
multdiv /
id     b
semi   ;
rbrace }
else   else
lbrace {
id     result
assign =
num   -1
semi   ;
rbrace }
return return
id     result
semi   ;
```

	rbrace }	
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## Failed Tokens

input file(test.txt)	.lex_out file (test.txt.lex_out)	console
int divide(int a, int b) { int result; if(b == 0) { result = 0; } else { result = a/b; } return c; }	vtype int id divide lparen ( vtype int id a comma , vtype int id b rparen ) lbrace { vtype int id result semi ; if if lparen ( id b comp == num 0 rparen ) lbrace { id result assign = num 0 semi ; rbrace } else else lbrace { id result assign = id a multdiv / id b semi ; rbrace } return return id c semi ; rbrace }	Lexical Analyzer : Accepted Lexical Analyzer : Token file saved at test.txt.lex_out successfully Syntax Analyzer : Failed Syntax Analyzer : 15's token lparen can't choice id