

Theories of Computation: Summative Assignment 2

To be handed in on Canvas before **Thursday 31st March, 5pm GMT**

Exercise 1 Let $\Sigma = \{a, b\}$ and $u, v \in \Sigma^+$, where Σ^+ is the set of nonempty words over Σ . We say that u is present in v if u can be obtained by deleting letters from v . For example, $abbba$ is present in $aabbababa$. We write $|u|$ to denote the length of word u , i.e., the number of characters. For example, $|abbba| = 5$.

The goal of this exercise is to show that the following decision problem is in NP.

Input: $w_0\#w_1\#\dots\#w_k$ such that $w_i \in \Sigma^+$ for $0 \leq i \leq k$.

Problem: is there a word $x \in \Sigma^+$ such that $|x| = |w_0|$ and x is present in each w_i for $0 < i \leq k$?

We will decompose the task into several steps.

1. Let us consider a two-tape deterministic Turing machine \mathcal{M}_1 on the input alphabet $\Sigma = \{a, b, \#\}$ with initial state 0, tape alphabet $T = \{a, b, \#, \sqcup\}$, return values $V = \{\text{True}, \text{False}\}$, and whose transition function is represented as the diagram in Figure 1 below.

Initially,

- the **Main** tape contains a non-empty block of a s and b s (representing a word $w \in \Sigma^+$) in between a $\#$ on the left, on which the head is positioned, and a $\#$ or a \sqcup on the right. (Outside of these $\#$ s and/or \sqcup s there could be any symbol in T .)
- the **Aux** tape contains a non-empty block of a s and b s (representing a word $x \in \Sigma^+$) and is otherwise blank and the head is located on the \sqcup immediately to the left.

For example,

Main	•	#	a	a	b	b	a	b	a	b	a	#
Aux	•	⊔	a	b	b	b	a	⊔	⊔	⊔	⊔	⊔

In the case where the input block on the **Main** tape forms the word $w = a^n$ and the input block on the **Aux** tape forms the word $x = a^m$ for $m, n > 0$, how many steps does the machine \mathcal{M}_1 go through until it returns a value in V ? (Returning counts as a step.) **[3 marks]**

This is in fact the worst case complexity as a function of $n = |w|$ and $m = |x|$ and you can use this fact in the remainder of the exercise.

2. Design a two-tape nondeterministic Turing machine \mathcal{M}_2 that takes as an input a word $w \in \Sigma^+$ and can generate any word $x \in \Sigma^+$ that has the **same length** as w .

Formally, the start configuration is:

- the **Main** tape contains a nonempty block of a s and b s (representing a word $w \in \Sigma^+$), with a \sqcup to the left on which the head is placed and a $\#$ to the right, the rest of the tape is blank to the left and can contain any symbol in T beyond $\#$ on the right;
- the **Aux** tape is blank.

For example,

Main	•	⊔	a	a	b	b	a	b	a	b	a	#
Aux	•	⊔	⊔	⊔	⊔	⊔	⊔	⊔	⊔	⊔	⊔	⊔

The machine \mathcal{M}_2 should stop when reaching a configuration where:

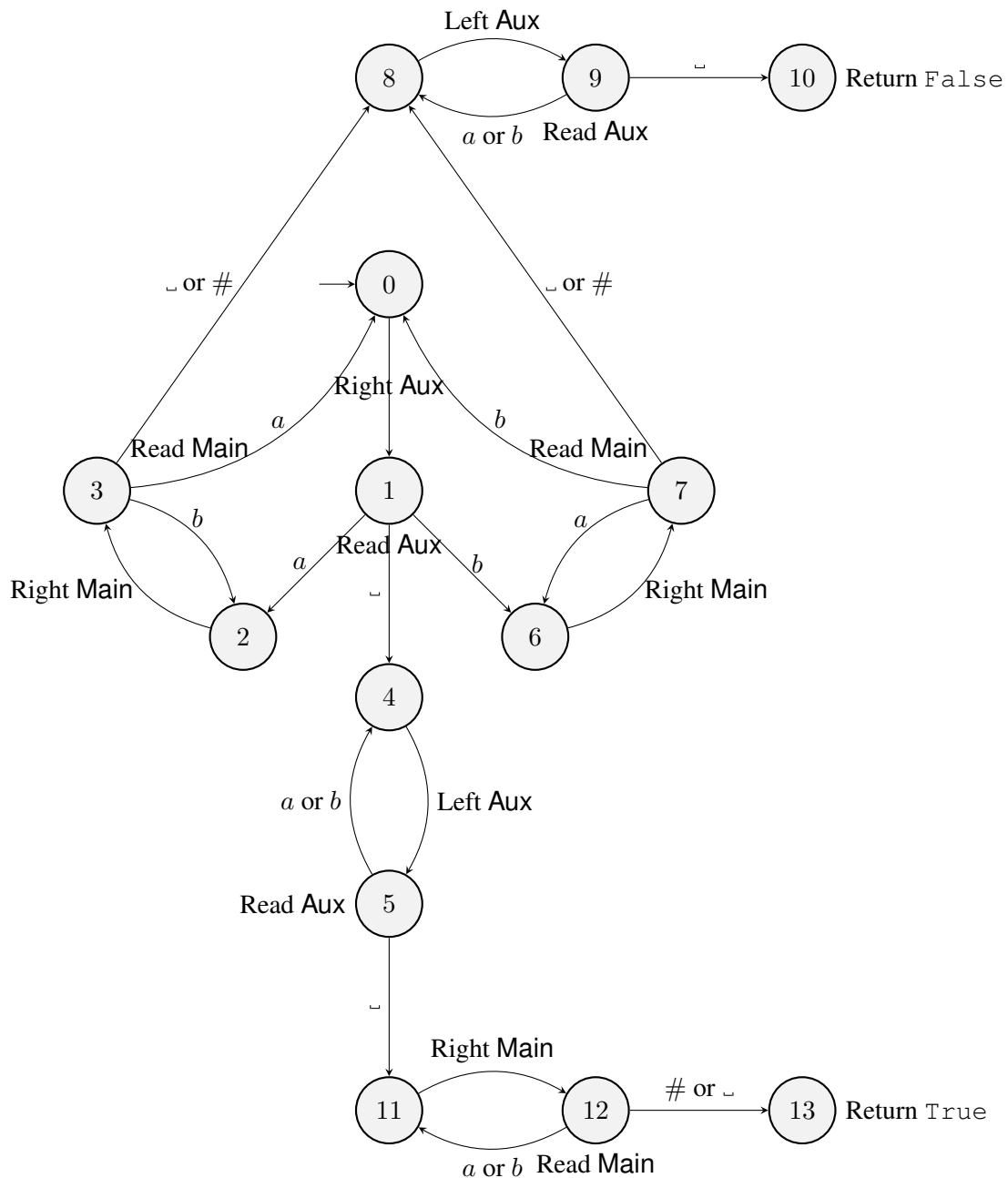


Figure 1: Transition diagram for machine \mathcal{M}_1

- the **Main** tape is unchanged except for the head which should be placed on the $\#$ on the right
- the **Aux** tape contains an arbitrary block of a s and b s (representing a word $x \in \Sigma^+$) of the same length as the input block on the **Main** tape and the head is on the first \sqcup to the left.

For example,

Main	\sqcup	a	a	b	b	a	b	a	b	a	\bullet	$\#$
Aux	\bullet	\sqcup	b	b	a	b	a	a	a	a	b	\sqcup

Give the machine \mathcal{M}_2 and briefly explain your solution.

[3 marks]

(Do not use more than 8 states.)

- Using machines \mathcal{M}_1 and \mathcal{M}_2 as macros, design a two-tape nondeterministic Turing machine \mathcal{M}_3 for the above decision problem.

This means that the machine should start with

- a block of a s, b s and $\#$ s representing the input $w_0\#w_1\#\dots\#w_k$ on an otherwise blank **Main** tape with the head on the first \sqcup to the left of w_0

- and a blank *Aux* tape.

It should accept (return `True`) if the question is true for this input and reject (return `False`) otherwise. The tape contents and head positions at the end do not matter.

Give the machine \mathcal{M}_3 and briefly explain your solution.

[3 marks]

(Do not use more than 5 states.)

4. *Explain briefly why the problem is in NP.*

[3 marks]

(Note: You may assume that any polytime two-tape nondeterministic Turing machine can be converted into a polytime one-tape nondeterministic machine with the same language. Just as we learnt in lectures for deterministic machines.)