## HIGHER ORDER FUNCTIONS

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#### LECTURE OUTLINE

• FUNCTIONS AS VALUES

2 Closures

3 Declarative programming

#### COMMON PATTERNS

We often write functions that are nearly identical. For example, computing

- ★ Sums of squares, cubes, ...
- \* Arithmetic series
- ★ Geometric series
- ★ Taylor series of various functions

all have the same basic pattern.

#### GENERALISING FUNCTIONS

$$\sum_{n=1}^{N} a + (n-1)d$$

$$\sum_{n=1}^{N} ar^n - 1$$

```
let rec arith_sum a d n =
   if n=1 then a
   else
   let t = a +. float (n-1)
    *.d in
   t +. arith_sum a d (n-1)
```

```
let rec geom_sum a r n =
  if n=1 then a
  else
  let t = a *. r ** float (n
    -1) in
  t +. geom_sum a r (n-1)
```

Is there a way to avoid such repetition?

#### HIGHER ORDER FUNCTIONS

$$\sum_{n=1}^{N} t_{n-1}$$

We need a way to abstract out the calculation of the terms

```
let rec series_sum term n =
  if n=1 then term 1
  else term n +. series_sum term (n-1)
```

Note that the parameter term is a function.

```
let aterm i = 1. +. float (i-1) *. 2. in
series_sum aterm 10
```

series\_sum is an example of a higher order function.

## EXERCISE

```
let rec series_sum term n =
   if n=1 then term 1
   else term n +. series_sum term (n-1)
;;
let aterm i = 1. +. float (i-1) *. 2.in
series_sum aterm 2
```

- 1) Show the evaluation of the expression above.
- 2 Write an expression to calculate the sum of the first ten terms of the geometric series with a=5, r=2.

# Anonymous functions (Lambdas)

Ocaml provides a shorthand notation to define functions without a name.

```
(* square function *)
fun x -> x *x

let rec series_sum term n =
   if n=1 then term 1
   else term n +. series_sum term (n-1)
   ;;
series_sum (fun i -> 1. +. float (i-1) *. 2.) 2
```

Can you define a recursive function using a lambda?

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#### Specialising an HOF

Our "general solution" is a lot less convenient to use than a specialised function. e.g. compare the following equivalant expressions:

```
arith_sum 1. 2. 2
series_sum (fun i -> 1. +. float (i-1) *. 2.) 2
```

We can rewirte arith\_sum in terms of series sum so we have to provide the lambda ourselves

```
let arith_sum a d n =
  series_sum (fun i -> a +. float (i-1) *. d) n
```

## EXERCISE

```
1 let arith sum a d n =
     series_sum (fun i \rightarrow a +. float (i-1) *. d) n
  let rec series_sum term n =
    if n=1 then term 1
    else term n +. series sum term (n-1)
```

2 Write an HOF compose f g which returns f(g(x)).

Show the evaluation of arith sum 1 1 1.

#### LEXICAL CLOSURES

```
arith_sum 1 1 1
series sum (fun i -> 1. +. float(i-1) *. 1.) 1
term 1
1. +. float(1-1) *. 1.
```

Note how the values for a and d of arith\_sum are *bound* within the lambda.

They are available at the point of use in series\_sum.

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DECLARATIVE PROGRAMMING

#### Abstracting control flow

- $\star$  Recursion is difficult to reason about<sup>1</sup>.
- ★ Writing an HOF for a particular task lets us avoid direct use of recursion.

```
let arith_sum a d n =
  series_sum (fun i -> a +. float (i-1) *. d) n
```

Once we have written series\_sum we can ignore the "low level details" of how the summation happens.

 $\star$  Can we do the same for other kinds of operations as well?

<sup>&</sup>lt;sup>1</sup>Loops are no easier!

#### COMMON LIST OPERATIONS

- MAP Perform an operation on individual elements of a list, e.g. scale each element by a constant.
- FOLD <sup>2</sup> Combine the elements of a list using an operation, e.g. sum of elements.
- FILTER Remove elements that do not meet a particular condition, e.g. remove all negative elements.
  - ZIP Combine pairs of elements in two lists together using an operation.

Have you seen these functions elsewhere?

#### MAP

```
(* Apply the function f to each element of list I *)
let map f l =
  match l with
  hd::tl -> f hd :: map f tl
  [] -> []
```

What is the type of this function?

# FOLD (LEFT)

```
(* Compose element of list I using
  function f and identity e *)
let fold f e l =
  match l with
  hd::tl -> f hd (fold f e tl)
  [] -> _____
Example: fold (+) 0 [1; 2; 3]
```

### EXERCISES

- ★ Multiply the numbers in the list [1; 2; 3] together.
- ★ Show the evaluation of the expression map (fun  $\times -\times \times$ ) [2].
- ★ Define the function filter p l that removes the elements that fail the predicate<sup>3</sup> p from list l.

<sup>&</sup>lt;sup>3</sup>boolean function

## WHY HOFS?

- ★ Abstracts common patterns leading to less code.
- ★ Makes it easier to reason about programs by abstracting control flow.
- ★ Easy to parallelise. e.g. Map-reduce computing on clusters and CUDA kernels.

#### ALGEBRAIC DATA TYPES

Ocaml provides three basic ways to structure data

Tuples and records represent data combinations that are a product sets  $\alpha \times \beta$ 

Variants represent data combinations that are a disjoint union of sets  $\alpha \sqcup \beta$ 

Functions represent mappings from one data type to another  $\alpha \to \beta$ . This corresponds to the *power set*<sup>4</sup>  $b^a$ .

<sup>&</sup>lt;sup>4</sup>Functions can be represented as a (possibly infinite) set of pairs.