# Modules

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- OCAML MODULE LANGUAGE
- 2 STRUCTURES: MODULE IMPLEMENTATION
- 3 Signatures: module interfaces
- ADTs: TYPE ABSTRACTION
- **10** Building and executing Ocaml programs

# **MODULARITY**

We can make a software system *modular* by structuring it as a set of units that are *loosly coupled*, that is, relatively independent of each other.

#### This aids

- $\star$  understanding large systems.
- $\star$  making modifications without affecting the rest of the system.
- $\star$  reuse of existing code.
- $\star$  independent development of modules in parallel.

### Language support

Language support for modularity must address several aspects:

- NAMESPACES Prevent clashes due to use of the same name in different parts of the program or in different libraries. Supported via *packages* in OO languages.
- INTERFACE A specification containing the *operations* supported by a particular module.

  Module signatures in Ocaml and interfaces in OO

languages provide interface specification.

IMPLEMENTATION The actual code to implement the operations specified by an interface.

Module structures in Ocaml and classes in OO languages provide implementation specification.

C and Javascript lack some of these features, making it difficult to structure large programs.

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### STRUCTURES

Structures define the implementation details of a module. It can contain types, functions and, nested modules.

#### STRUCTURE SYNTAX

#### Example:

```
module Circle = struct
  let pi = 3.14159
  let circum r = 2.0 *. pi *. r
  let sq x = x *. x
  let area r = pi *. (sq r)
end
```

# QUALIFIED NAMES

- ★ To reference functions and types in another module, prefix them with the module name separated by a dot. e.g. Circle.area 1.0.
- ★ Alternatively, you may open a module, and use unqualified names.

```
let open Circle in
sq pi
```

Opening two modules defining the same names causes confusion, so always use the *local open* shown here.

# NESTED MODULES

content...

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# IMPLEMENTATION HIDING

- ★ The key to constructing larger systems is defining clear interfaces between modules.
- ★ A signature (module type) can be used to restrict the visibility of implementation details.
- ★ A module can implement many signatures. Conversely, a signature may be implemented by many modules.

#### SIGNATURE SYNTAX

module type Name = sig interface definitions

#### DEFAULT SIGNATURE

A module has a *default signature* that exposes all implementation details.

UTop echos back this default signature when you enter a module definitions.

```
module Circle: sig
  val pi : float
  val circum : float -> float
  val sq : float -> float
  val area : float -> float
end
```

Note the similarity with C function prototypes.

#### Implementation hiding

To *hide* the function such as sq, we simply omit it from the signature.

```
module type CIRCLE = sig
  val pi : float
  val circum : float -> float
  val area : float -> float
end

module Circle :CIRCLE = struct
...
end
```

The sq function is no longer visible.

```
Circle.sq 6. ;;
Error: Unbound value Circle.sq
```

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# Abstract data types

- ★ We have seen several examples of data structures, e.g. lists and trees.
- ★ Data structures *concretely* specify how they are represented using a type definition.
- ★ In contrast, an abstarct data type (ADT) hides its representation and specifies just the available operations on its type.
- ★ We perform *type abstraction* in module signatures when specifying ADTs in Ocaml.

# SET ADT SIGNATURE

Sets are mathematical objects that support various operations.

```
module type SET = sig
  type 'a set
  (** the empty set **)
  val empty : 'a set
  (** insert element into given set **)
  val insert : 'a -> 'a set -> 'a set
  (** is element a member of given set? **)
  val member : 'a -> 'a set -> bool
end
```

Note that type set is *abstract*, we have not specified how it is represented.

This gives us the flexibility to implement Set using any suitable data structure.

```
(* Set ADT implemented as a sorted list of unique elements *)
module ListSet: SFT = struct
  type 'a set = 'a list
  let empty = []
  let rec member x l =
    match l with
    | [] -> false
    | hd::tl -> (x = hd) || ((x > hd) && (member x tl))
  let rec insert x l =
    match l with
    | | | -> | | |
      hd::tl -> if x < hd then x::l else if x = y then l
                else hd::(insert hd tl)
```

#### BENEFITS OF TYPE ABSTRACTION

\* We can only manipulate sets using the operations in ListSet.

```
# ListSet.insert 1 ListSet.empty ;;
- : int ListSet.set = <abstr>
```

The internal use of lists to represent sets is completely hidden.

- ★ Invalid manipulations like violating the sorted order, or uniqueness of a list that represents a set, are automatically prevented.
- \* We can define multiple implementations of the Set interface, e.g. a more efficient BSTset based on a binary search tree. The *client code* that uses Set remains unchanged even when changing implementations.

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★ An ADT implementation and its clients can be coded independently open an interface has been agreed on ZIYAN MARAIKAR

# PITFALLS OF DEPENDING ON CONCRETE TYPES

Suppose we implement ListSet without bothering to define a signature for the SET ADT.

```
module ListSet
   ...
end
```

We lose all the benefits of type abstraction with this implementation.

```
# ListSet.insert 1 ListSet.empty ;;
- : int list = [1]
```

# PITFALLS OF DEPENDING ON CONCRETE TYPES

We are prone to errors due to list being visible when dealing with sets.

```
# let s1 = ListSet.insert 1 ListSet.empty ;;
- : int list = [1]
# let s2 = 2::s1 ;;
- : int list = [2; 1]
# ListSet.insert 1 s2;;
- : int list = [1; 2; 1]
```

This is the hallmark of a lack of modularity!

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# SOURCE FILES, STRUCTURES AND SIGNATURES

An Ocaml program is structured as a collection of source files.

Module implementations A source file with a .ml extension defines a structure. Its name is the same as the file, with the first character uppercased.

e.g. the file binarytree.ml defines a module called

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Module interfaces A source file with a .mli extension defines a signature. Its name is the same as the file, with the first character uppercased.

e.g. the file set.mli defines a signature called Set.

### Building code

Building Ocaml programs is a two-step process, similar to how C code is built.

COMPILATION Each source file is a separate *compilation unit*, compiled individually using the -c switch to ocamlc,

ocamlc -c binarytree.ml ocamlc -c set.mli

This produces two *object files* named binarytree.cmo and set.cmi.

LINKING The resulting object files are then linked together using ocamlc again.

ocamlc — o example binarytree.cmo set.cmi

This entire process can be automated using the ocambuild tool.

# PROGRAM EXECUTION

- ★ At runtime definitions (types, functions etc.) in the entire program are evaluated.
- ★ There is no unique main function. Any function application (call) at the top level is evaluated.
- ★ To call a function producing a *side-effect* like printing we precede the call with let () =. This pattern matches the function result against unit type.

```
let () = Printf.printf "hello world\n" ;;
```