

File-System Interface

Unit-IV

Session Objectives

File-System
Interface

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File Concept

- To explain the function of file systems
- To describe the interfaces to file systems
- To discuss file-system design tradeoffs, including access methods, file sharing, file locking, and directory structures
- To explore file-system protection

Session Outcomes

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File Concept

At the end of this session, participants will be able to

- Discuss File Concepts, Access Methods
- Disk and Directory Structure
- File-System Mounting
- File Sharing Protection

Agenda

File-System
Interface

Unit-IV

File Concept

1 File Concept

Presentation Outline

File-System
Interface

Unit-IV

File Concept

1 File Concept

File Concept

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File Concept

- Contiguous logical address space
- Types:
 - Data
 - ▶ numeric
 - ▶ character
 - ▶ binary
 - Program
- Contents defined by file's creator
 - Many types
 - ▶ Consider **text file, source file, executable file**

File Attributes

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File Concept

- ❑ **Name** – only information kept in human-readable form
- ❑ **Identifier** – unique tag (number) identifies file within file system
- ❑ **Type** – needed for systems that support different types
- ❑ **Location** – pointer to file location on device
- ❑ **Size** – current file size
- ❑ **Protection** – controls who can do reading, writing, executing
- ❑ **Time, date, and user identification** – data for protection, security, and usage monitoring
- ❑ Information about files are kept in the directory structure, which is maintained on the disk
- ❑ Many variations, including extended file attributes such as file checksum
- ❑ Information kept in the directory structure

File Operations

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File Concept

- File is an **abstract data type**
- **Create**
- **Write** – at **write pointer** location
- **Read** – at **read pointer** location
- **Reposition within file - seek**
- **Delete**
- **Truncate**
- ***Open(F_i)*** – search the directory structure on disk for entry F_i , and move the content of entry to memory
- ***Close (F_i)*** – move the content of entry F_i in memory to directory structure on disk

Open Files

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File Concept

Several pieces of data are needed to manage open files:

- **Open-file table:** tracks open files
- **File pointer:** pointer to last read/write location, per process that has the file open
- **File-open count:** counter of number of times a file is open – to allow removal of data from open-file table when last processes closes it
- Disk location of the file: cache of data access information
- Access rights: per-process access mode information

Open File Locking

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File Concept

- Provided by some operating systems and file systems
- **Shared lock** similar to reader lock – several processes can acquire concurrently
- **Exclusive lock** similar to writer lock
- Mediates access to a file
- **Mandatory or advisory:**
 - Mandatory – access is denied depending on locks held and requested
 - Advisory – processes can find status of locks and decide what to do

File Structure

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File Concept

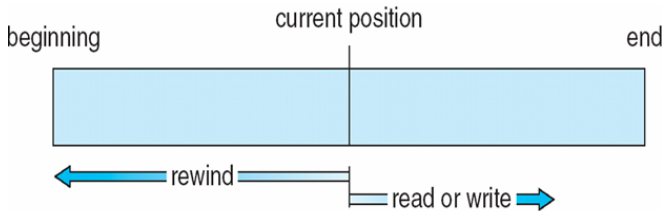
- None - sequence of words, bytes
- Simple record structure
- Lines
- Fixed length
- Variable length
- Who decides:
- Operating system

Sequential-access File

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Access Methods

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□ Sequential Access

```
read next
write next
reset
no read after last write
(rewrite)
```

□ Direct Access – file is fixed length **logical records**

```
read n
write n
position to n
    read next
    write next
rewrite n
```

n = **relative block number**

- Relative block numbers allow OS to decide where file should be placed

Simulation of Sequential Access on Direct-access File

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sequential access	implementation for direct access
<i>reset</i>	<i>cp = 0;</i>
<i>read next</i>	<i>read cp;</i> <i>cp = cp + 1;</i>
<i>write next</i>	<i>write cp;</i> <i>cp = cp + 1;</i>

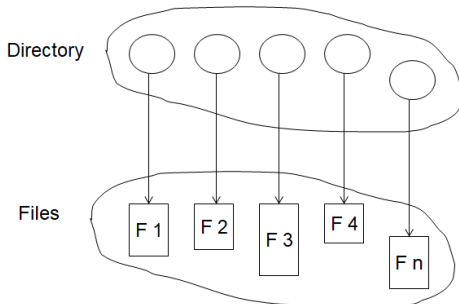
Directory Structure

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File Concept

- A collection of nodes containing information about all files



Both the directory structure and the files reside on disk

Disk Structure

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File Concept

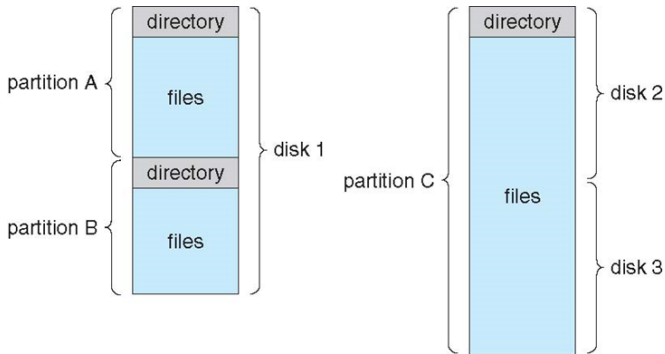
- Disk can be subdivided into **partitions**
- Disks or partitions can be **RAID** protected against failure
- Disk or partition can be used **raw** – without a file system, or **formatted** with a file system
- Partitions also known as minidisks, slices
- Entity containing file system known as a **volume**
- Each volume containing file system also tracks that file system's info in **device directory or volume table of contents**

A Typical File-system Organization

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Types of File Systems

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- But systems frequently have many file systems, some general- and some special- purpose
- Consider Solaris has
 - tmpfs – memory-based volatile FS for fast, temporary I/O
 - objfs – interface into kernel memory to get kernel symbols for debugging
 - ctfs – contract file system for managing daemons
 - lofs – loopback file system allows one FS to be accessed in place of another
 - procfs – kernel interface to process structures
 - ufs, zfs – general purpose file systems

Operations Performed on Directory

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File Concept

- Search for a file
- Create a file
- Delete a file
- List a directory
- Rename a file
- Traverse the file system

Directory Organization

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File Concept

The directory is organized logically to obtain

- Efficiency - locating a file quickly
- Naming - convenient to users
- Two users can have same name for different files
- The same file can have several different names
- Grouping - logical grouping of files by properties, -

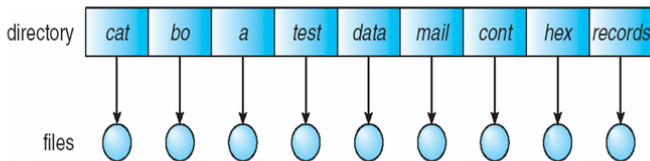
Single-Level Directory

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File Concept

- A single directory for all users



- Naming problem
- Grouping problem

Two-Level Directory

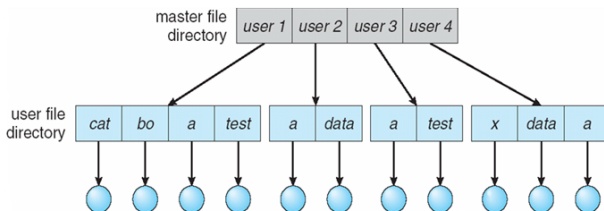
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Two-Level Directory

- Separate directory for each user



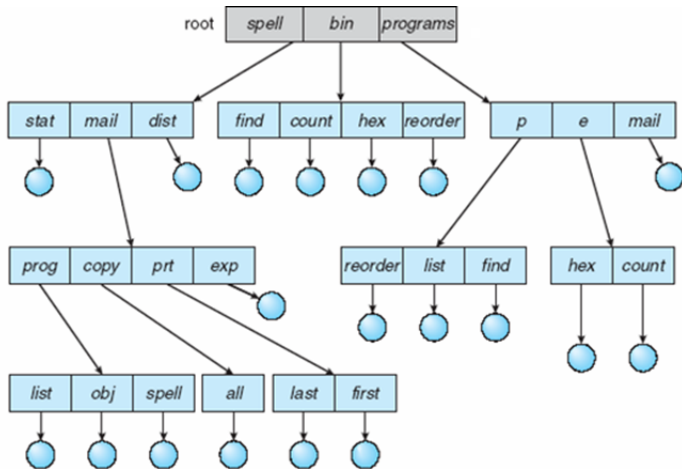
- Path name
- Can have the same file name for different user
- Efficient searching
- No grouping capability

Tree-Structured Directories

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Efficient searching - Grouping Capability

Tree-Structured Directories (Cont)

- ❑ **Absolute** or **relative** path name
- ❑ Creating a new file is done in current directory
- ❑ Delete a file

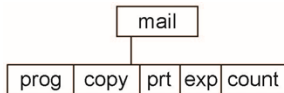
```
rm <file-name>
```

- ❑ Creating a new subdirectory is done in current directory

```
mkdir <dir-name>
```

Example: if in current directory `/mail`

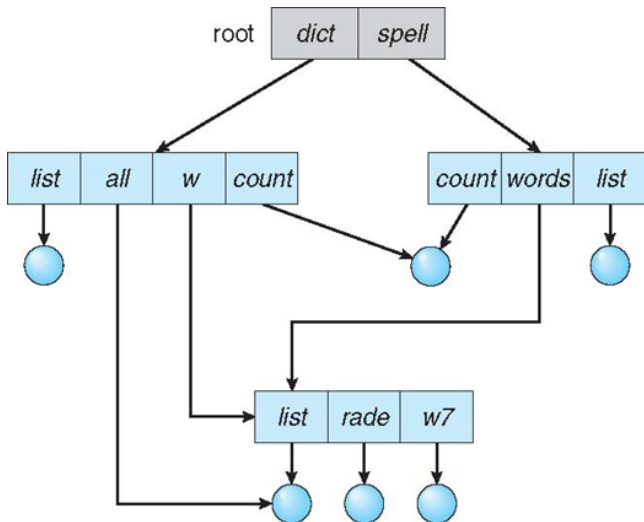
```
mkdir count
```



Deleting “mail” \Rightarrow deleting the entire subtree rooted by “mail”

Acyclic-Graph Directories

Have shared subdirectories and files



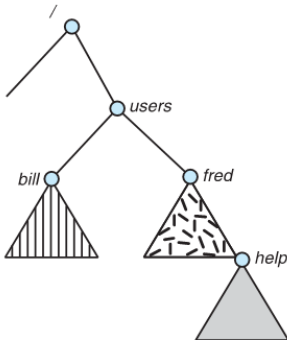
File System Mounting

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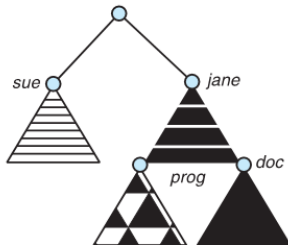
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A file system must be mounted before it can be accessed
A unmounted file system is mounted at a mount point



(a)



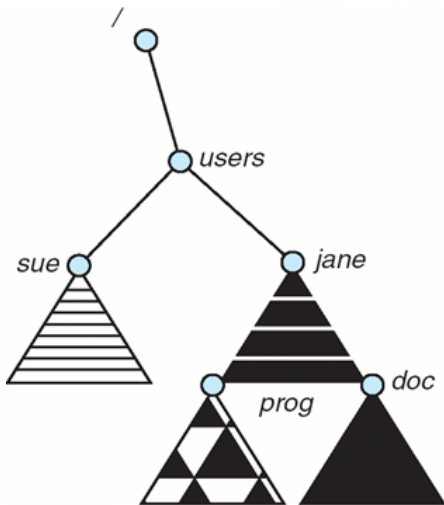
(b)

Mount Point

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File Sharing

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File Concept

- Sharing of files on multi-user systems is desirable
- Sharing may be done through a **protection** scheme
- On distributed systems, files may be shared across a network
- Network File System (NFS) is a common distributed file-sharing method
- In multi-user system
 - **User IDs** identify users, allowing permissions and protections to be per-user **Group IDs** allow users to be in groups, permitting group access rights
 - Owner of a file or directory
 - Group of a file or directory

File Sharing – Remote File Systems

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File Concept

- Uses networking to allow file system access between systems
 - Manually via programs like FTP
 - Automatically, seamlessly using **distributed file systems**
 - Semi automatically via the **world wide web**
- **Client-server** model allows clients to mount remote file systems from servers
 - Server can serve multiple clients
 - Client and user-on-client identification is insecure or complicated
 - **NFS** is standard UNIX client-server file sharing protocol
 - **CIFS** is standard Windows protocol
 - Standard operating system file calls are translated into remote calls
- Distributed Information Systems (**distributed naming services**) such as LDAP, DNS, NIS, Active Directory implement unified access to information needed for remote computing

File Sharing – Failure Modes

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File Concept

- All file systems have failure modes
- For example corruption of directory structures or other non-user data, called metadata
- Remote file systems add new failure modes, due to network failure, server failure
- Recovery from failure can involve state information about status of each remote request
- Stateless protocols such as NFS v3 include all information in each request, allowing easy recovery but less security

Protection

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File Concept

- File owner/creator should be able to control:
 - what can be done
 - by whom
- Types of access
 - Read
 - Write
 - Execute
 - Append
 - Delete
 - List

Access Lists and Groups

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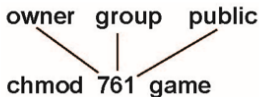
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File Concept

- Mode of access: read, write, execute
- Three classes of users on Unix / Linux

a) owner access	7	⇒	RWX 1 1 1
b) group access	6	⇒	RWX 1 1 0
c) public access	1	⇒	RWX 0 0 1

- Ask manager to create a group (unique name), say G, and add some users to the group.
- For a particular file (say *game*) or subdirectory, define an appropriate access.



Summary

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File Concept

Discussed

- File operations
- File structures
- Disk Structures, types of file systems
- Directory organization

Test Your Understanding

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File Concept

- What is the mount point?
 - a) an empty directory at which the mounted file system will be attached
 - b) a location where every time file systems are mounted
 - c) is the time when the mounting is done
 - d) none of the mentioned
- In distributed file system ————directories are visible from the local machine.
 - a) protected b) local c) private d) remote
- To organise file systems on disk ————
 - a) they are split into one or more partitions
 - b) information about files is added to each partition
 - c) they are made on different storage spaces
 - d) all of the mentioned