

# UCS1524 – Logic Programming

Answer Generation in Logic  
Programming



# Session Meta Data

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Reviewer	
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# Session Objectives

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- Understanding answer generation in logic programming
- Learn about the generation of answers from FOL formulas using resolution

# Session Outcomes

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- At the end of this session, participants will be able to
  - apply the answer generation in logic programming using resolution.

# Agenda

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- Answer Generation
- Generating single answer
- Generating multiple answer

# Answer Generation

- A resolution proof as such shows only that the empty clause is derivable; an answer, in a sense, explains *how it* is obtained.
- Task : how to *generate an answer, a result of the computation*, from the resolution proof.
- Given : Satisfiable set of clauses  $F$
- $F$  consists of predicates, functions, and constants
- Example
  - "Eve likes apples"
  - "Eve likes wine"
  - "Adam likes everybody who likes wine"

Likes: predicate

Eve, apple, wine, Adam: constants

$$F = \{ \{ \text{likes}(\text{Eve}, \text{Apples}) \}, \\ \{ \text{likes}(\text{Eve}, \text{Wine}) \}, \\ \{ \text{likes}(\text{Adam}, x), \neg \text{likes}(x, \text{Wine}) \} \}$$



# Answer Generation

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- Call may be “Is there anybody whom Adam likes?”

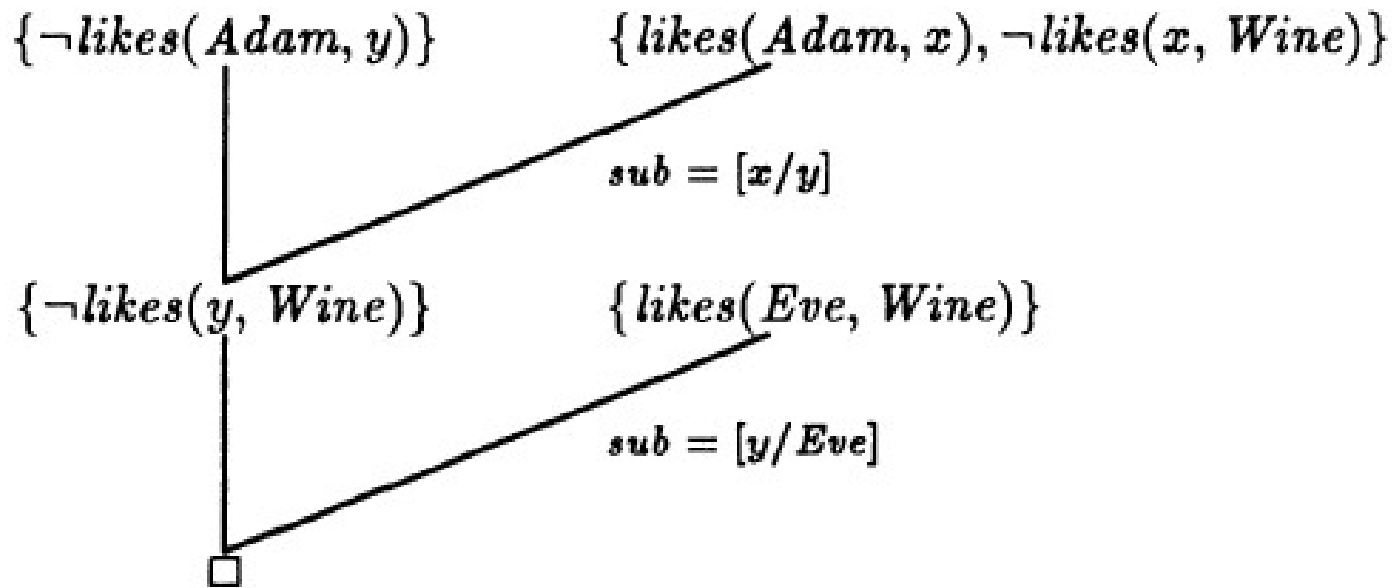
$$G = \exists y \text{ likes}(\text{Adam}, y)$$

- We can verify this using resolution by checking whether there is a resolution refutation of  $F \wedge \neg G$

$$F \wedge \neg G \equiv \{ \{ \text{likes}(\text{Eve}, \text{Apples}) \}, \\ \{ \text{likes}(\text{Eve}, \text{Wine}) \}, \\ \{ \text{likes}(\text{Adam}, x), \neg \text{likes}(x, \text{Wine}) \}, \\ \{ \neg \text{likes}(\text{Adam}, y) \} \}.$$

# Answer Generation

$$F \wedge \neg G \equiv \{ \{ \text{likes}(\text{Eve}, \text{Apples}) \}, \\ \{ \text{likes}(\text{Eve}, \text{Wine}) \}, \\ \{ \text{likes}(\text{Adam}, x), \neg \text{likes}(x, \text{Wine}) \}, \\ \{ \neg \text{likes}(\text{Adam}, y) \} \}.$$



Empty clause is derivable thus unsatisfiable

If variable  $y$  is substituted by Eve, we will get the answer Eve



# Answer Generation

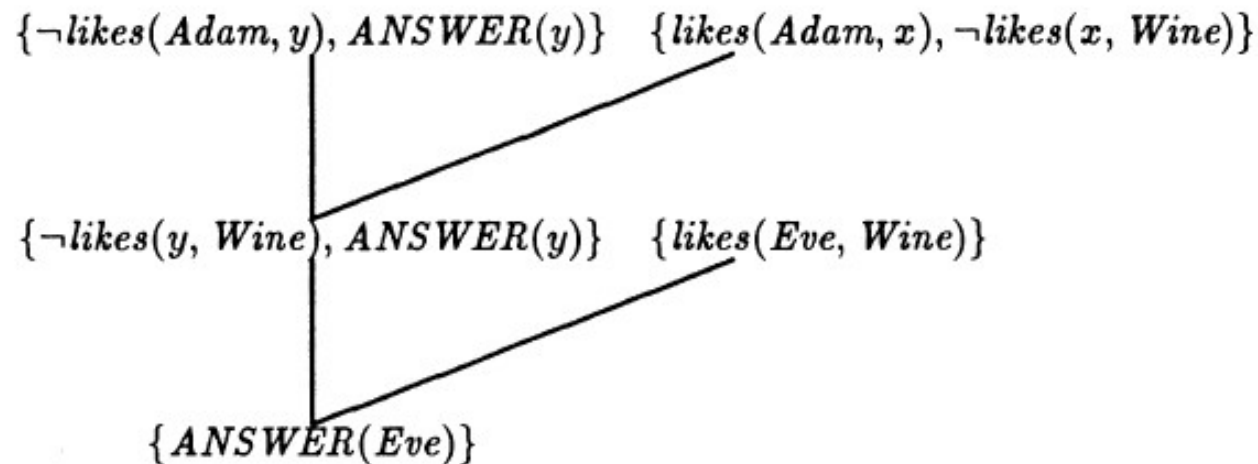
- Instead of

$\{\neg \text{likes}(\text{Adam}, y)\}$

- if we use

$\{\neg \text{likes}(\text{Adam}, y), \text{ANSWER}(y)\}$

- We get



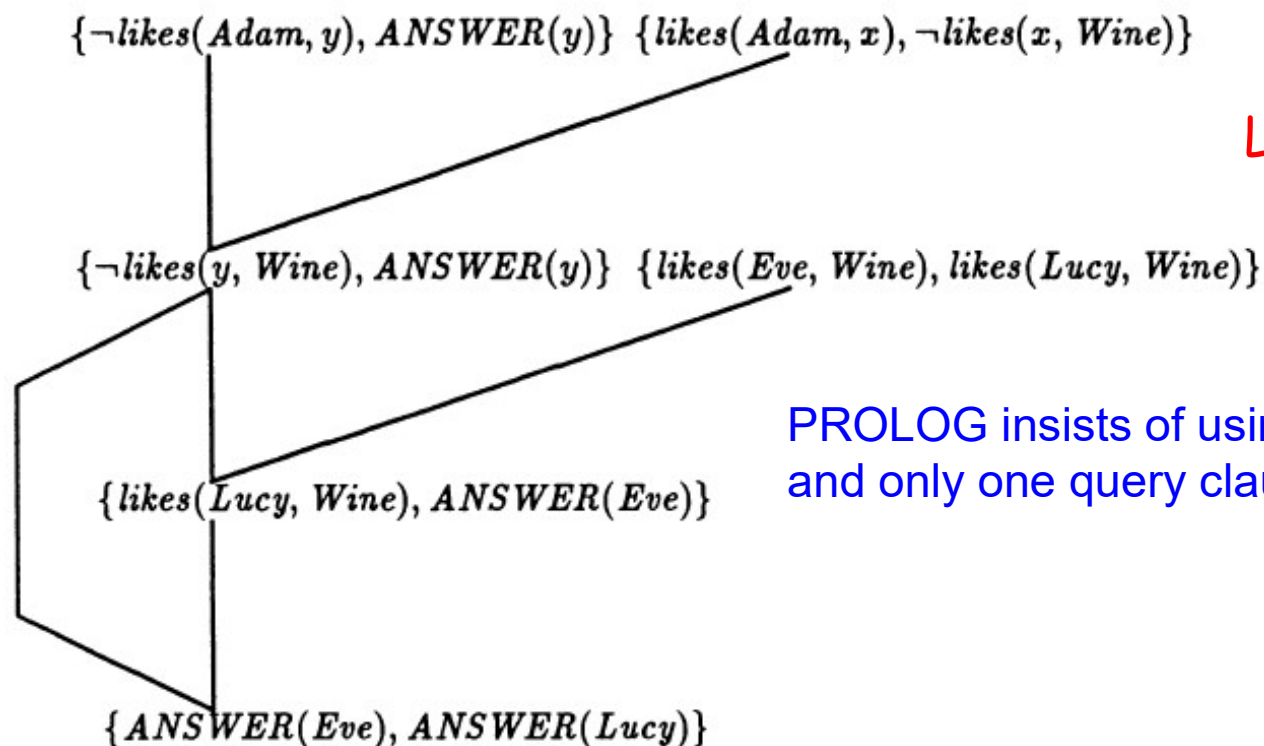
# Answer Generation

- If the clause set is

- "Eve likes apples"
- "Eve or Lucy (or both) like wine"
- "Adam likes everyone who likes wine"

$$\{ \{ \text{likes}(\text{Eve}, \text{Apples}) \}, \\ \{ \text{likes}(\text{Eve}, \text{Wine}), \text{likes}(\text{Lucy}, \text{Wine}) \}, \\ \{ \text{likes}(\text{Adam}, x), \neg \text{likes}(x, \text{Wine}) \}, \\ \{ \neg \text{likes}(\text{Adam}, y), \text{ANSWER}(y) \} \}$$

- Query is "Who does Adam like?"



Linear Resolution

PROLOG insists of using only Horn clauses,  
and only one query clause

# Summary

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- Answer Generation
- Generating single answer
- Generating multiple answer

# Check your understanding

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- Given clauses
  - Laxman is wherever Ram is. Ram is at Ayodhya. Where is Laxman?
- Generate answer