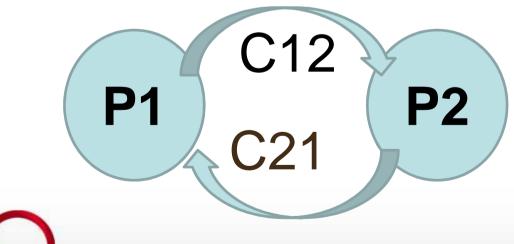


Chandy-Lamport Snapshot Algorithm - Global State

Y. V. Lokeswari, AP/CSE

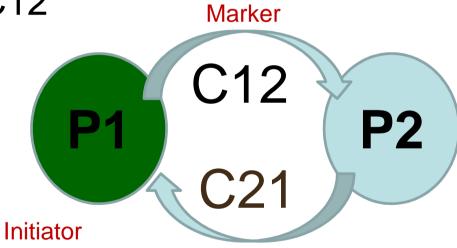
- Consider 2 processes P1 and P2
- Local state of P1 is recorded as LS1
- Local State of P2 is recorded as LS2
- Channel C12 sends message from P1 to P2. Local state is SC12
- Channel C21 sends message from P2 to P1. Local state is SC21





- P1 initiates Marker Sending
 - P1 records its own local state LS1 and
 - Sends marker to all outgoing channels

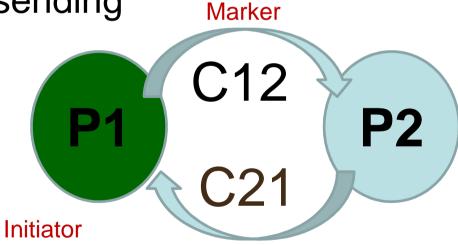
- Here it is Channel C12





- P2 receives Marker
 - P2's local state is not recorded
 - P2 records channel C12 as empty

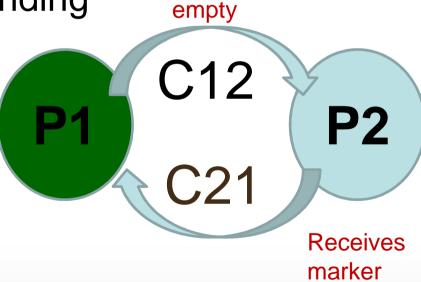
P2 initiates marker sending





- P2 receives Marker
 - P2's local state is not recorded
 - P2 records channel C12 as empty

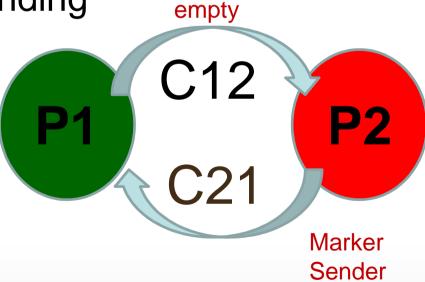
- P2 initiates marker sending





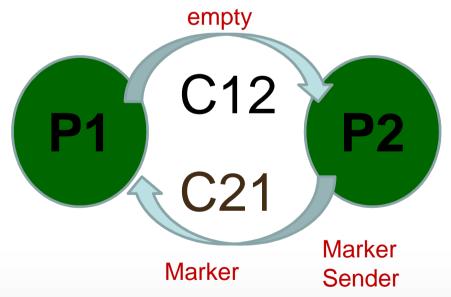
- P2 receives Marker
 - P2's local state is not recorded
 - P2 records channel C12 as empty

P2 initiates marker sending



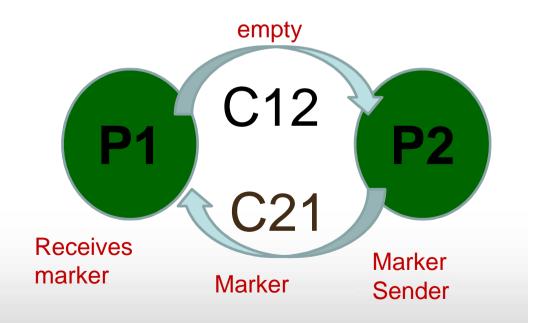


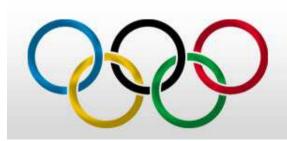
- P2 follows Marker sending rule
 - P2's local state LS2 is now recorded
 - P2 sends marker to all of its outgoing channels
 - Here channel C21





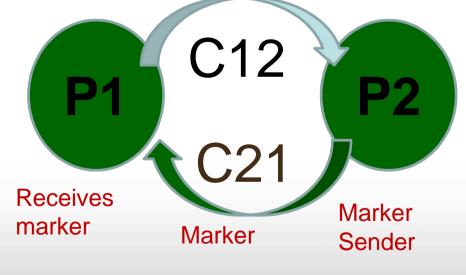
- P1 receives marker
 - P1's local state LS1 is already recorded
 - Records all the messages in channel C21
 right from the point P1 recorded its local state





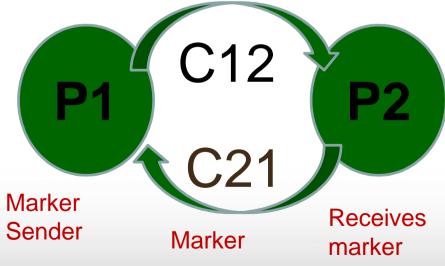
- P1 receives marker
 - P1's local state LS1 is already recorded
 - Records all the messages in channel C21 right from the point P1 recorded its local state till P1 receives second marker in all incoming channels

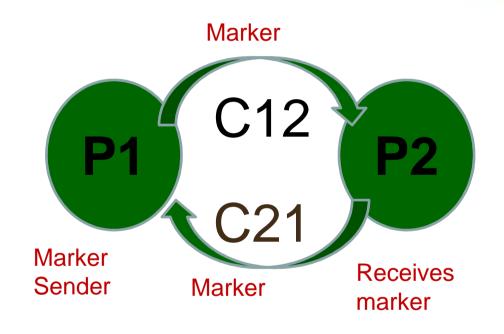




- P1 follows marker sending and P2 receives it
 - P2's local state LS2 is already recorded
 - Records all the messages in channel C12 right from the point P2 recorded its local state till P2receives second marker in all incoming channels

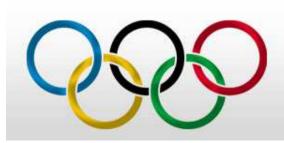


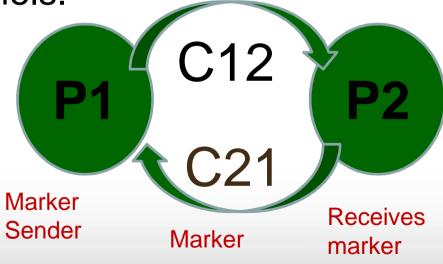


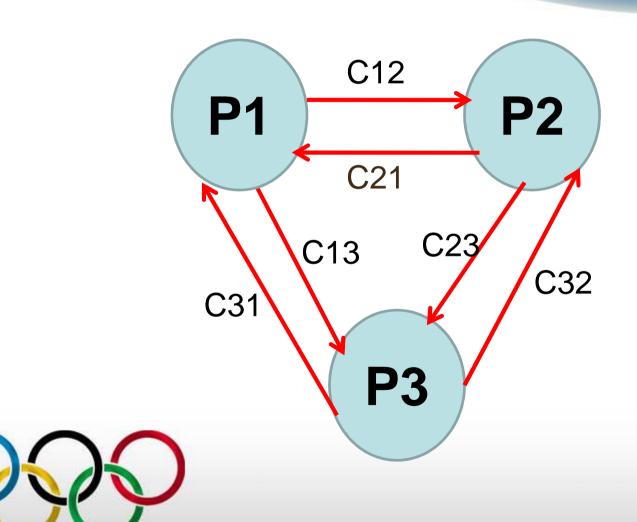




- When does algorithm terminates:
 - The local state is disseminated to all other nodes in the form of second marker message. When all of the processes in the Distributed systems have received second marker in all of its incoming channels.



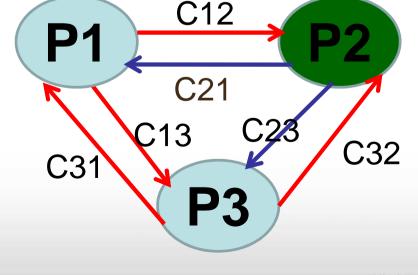




P2 is initiator.

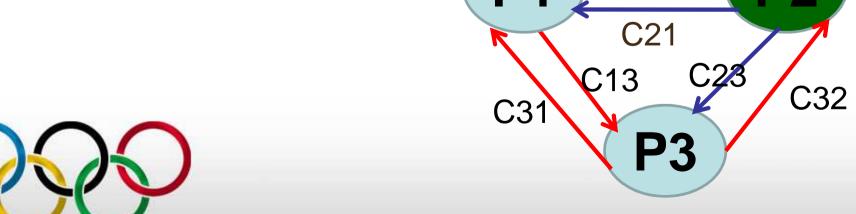
P2 records its local state LS2 and P2 sends marker as follows

P1 and P3 receives marker and their states are not recorded.





- P1 and P3 receives.
 - P1 records C21 as empty.
 - P3 records C23 as empty
 - P1 and P3 follows Marker sending rule.





C12

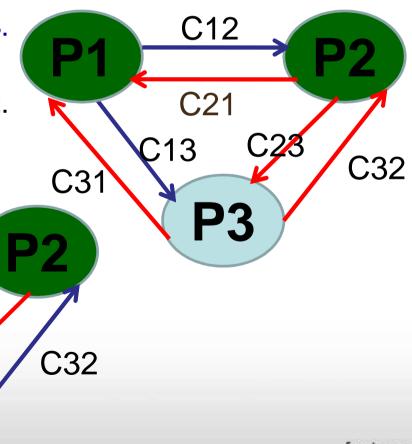
P1 and P3 follows Marker sending rule

C12

C21

C13

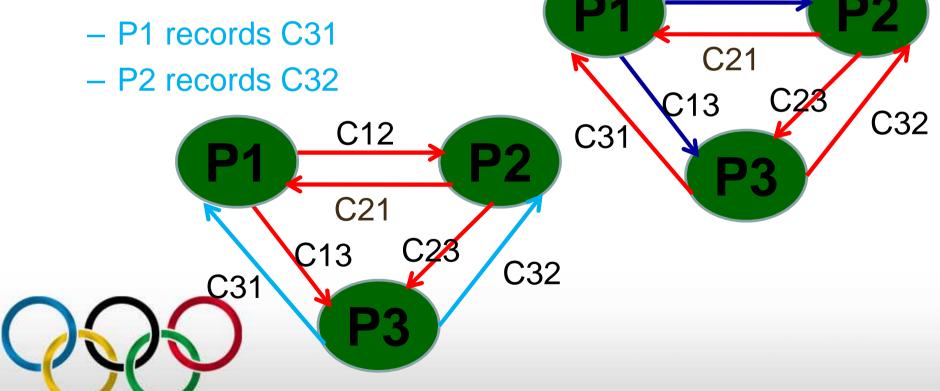
- P1 records its local state as LS1.
- P1 sends marker via C12 and C13.
- P3 records its local state as LS3.
- P3 sends marker via C31 and C32.



P1,P2 & P3 follows Marker Receiving rule



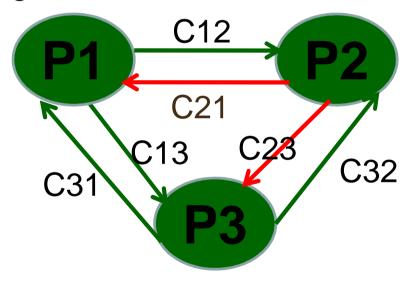
- P3 records C13

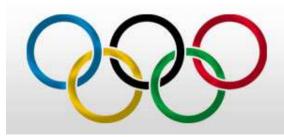


C12

– P1, P2, P3 receives marker andP1, P2, P3 follow marker sending rule

as its states are already recorded, they just record the state of channel right from the point P1, P2 and P3 recorded their states.

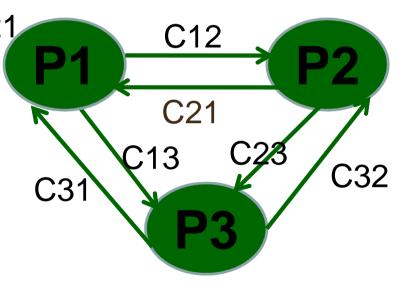


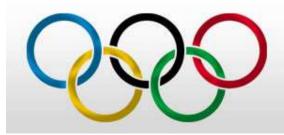


 P3 sends marker to P1 via C31 and P2 via C32 (Already recorded)

 P2 sends marker to P1 via C21 and P3 via C23 (C21 and C23 are recorded)

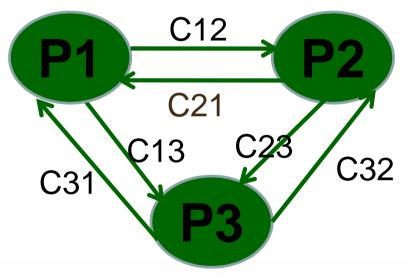
P1 sends marker to P2 via
 C12 and P3 via C13
 (Already recorded)





 All the processes received marker in its all incoming channels.

Chandy-Lamport Snapshot algorithm terminates





- Every process will exchange its recorded local state and channel state to all the other processes.
- All the local states are combined to form a Global State.
- To find whether a global state is consistent or not, introduce cuts and check for Consistency or Inconsistency

