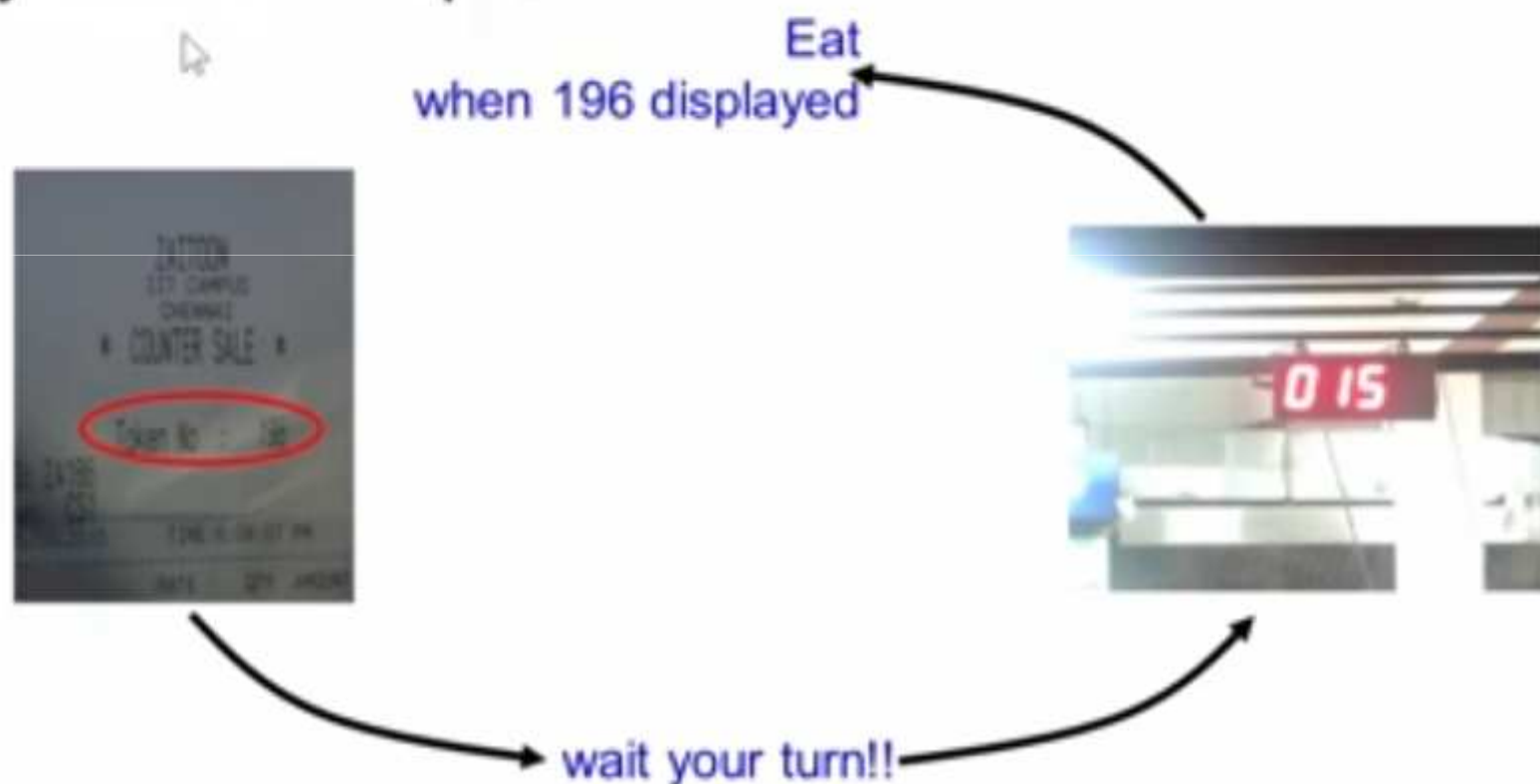


# Lamport's Bakery Algorithm for DSM

Bakery Algorithm for Mutually Exclusive access to the shared variables stored in DSM

# Bakery Algorithm

- Synchronization between  $N > 2$  processes
- By Leslie Lamport



# Simplified Bakery Algorithm

- Processes numbered 0 to N-1
- num is an array N integers (initially 0).
  - Each entry corresponds to a process

```
lock(i){  
    num[i] = MAX(num[0], num[1], ....., num[N-1]) + 1  
    for(p = 0; p < N; ++p){  
        while (num[p] != 0 and num[p] < num[i]);  
    }  
}
```

critical section

```
unlock(i){  
    num[i] = 0;  
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# Simplified Bakery Algorithm (example)

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P1	P2	P3	P4	P5
0	4	1	2	3



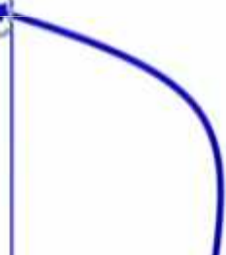
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```




This is at the doorway!!!  
It has to be atomic  
to ensure two processes  
do not get the same token

# Simplified Bakery Algorithm (why atomic doorway?)

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This is at the doorway!!!  
Assume it is not atomic



critical section

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- Without atomic operation assumptions
- Introduce an array of N Booleans: *choosing*, initially all values False.

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critical section

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doorway

Choosing ensures that a process  
is not at the doorway  
i.e., the process is not 'choosing'  
a value for num

$(a, b) < (c, d)$  which is equivalent to:  $(a < c)$  or  $((a == c) \text{ and } (b < d))$



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check  $p < i$

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# References

- Leslie Lamport's Bakery Algorithm Paper.  
<http://lamport.azurewebsites.net/pubs/bakery.pdf>
- NPTEL Video Lecture:  
<https://www.youtube.com/watch?v=YHQxp-XduS0>