

Synchronization (Absence of global clock)

$N_3 \rightarrow N_1, N_2 : 01$

$N_1 \rightarrow N_3 : 01$

$N_2 \rightarrow N_3 : 46$

$N_2 \rightarrow N_1, N_3 : 47$

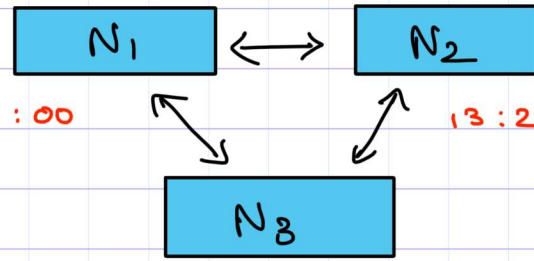
$N_1 \rightarrow N_2 : 03$

$N_3 \rightarrow N_2 : 04$

13:25:00

13:25:45

13:26:01



Ordering of msgs w/ clock order (of N_1 , say) leads to wrong ordering!

- Clock syncing isn't practical.
- Physical clocks aren't possible.
- Geographically distributed nodes \Rightarrow different timezones
- Cannot have a common clock.

Assumptions:

\rightarrow Comm. channel is error free.

\rightarrow Comm. channel is FIFO.

Clock syncing with message passing:

[rcvd. timestamp + $RTT/2$]

based on assumptions

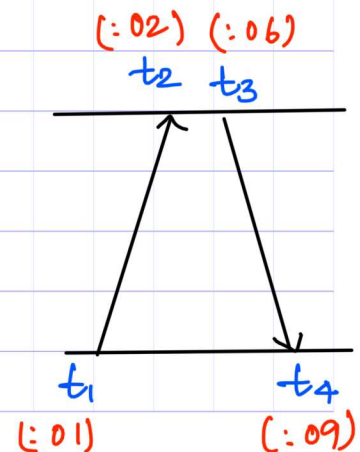
In case of delay & not FIFO:

$$(t_4 - t_1) - (t_3 - t_2) = 8 - 4 = 4$$

\downarrow
RTT

\downarrow
processing overhead

\downarrow
true RTT



$$\underbrace{\text{True RTT}/2}_{\downarrow} = 4/2 = 2$$

almost accurate.

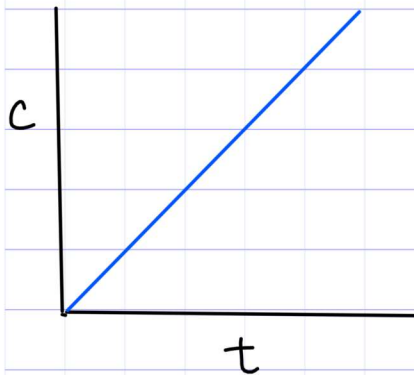
Delay is unpredictable.

\therefore Cannot sync 2 physical clocks.

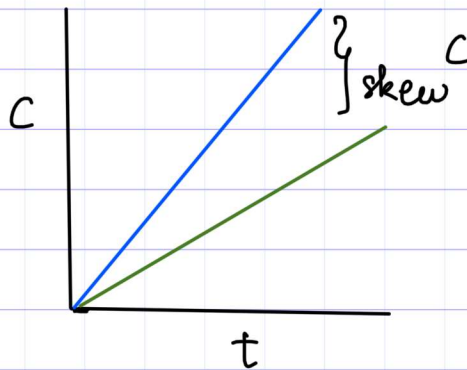
Clock Synchronization

Clock drift : clocks tick @ diff. rates.
 quartz clock : 10^{-6} seconds
 diff. of 1 sec / 11.6 days

clock skew : diff. between 2 clocks @
 one point in time

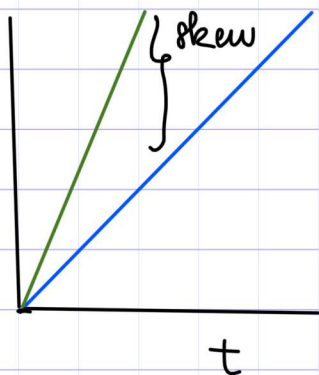


$$\frac{dC}{dt} = 1$$



$$\frac{dC}{dt} < 1$$

make clock
 run faster
 till sync



$$\frac{dC}{dt} > 1$$

make clock
 run slower
 till sync

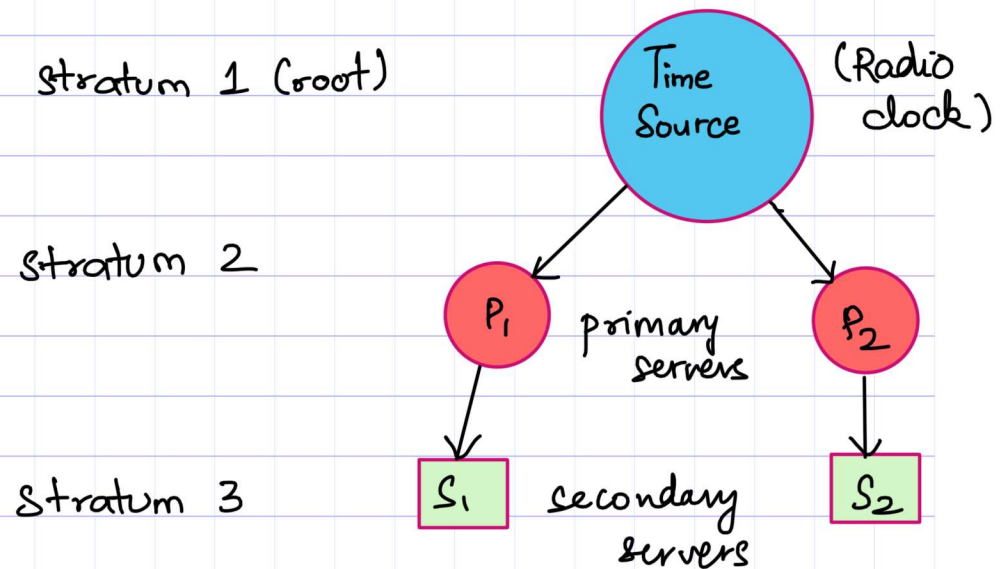
Mechanisms to sync:

Cristian's Algorithm
Berkeley Algorithm
Network Time Protocol

$(T_{new} = T_{server} + RTT/2)$
(Fault tolerant avg.
timestamping)

Network Time Protocol:

Clients across Internet sync accurately
to UTC despite message delays



Modes:

multicast mode - sends time to all other nodes.

procedure call mode - similar to Cristian's algorithm

symmetric mode - for master servers pair of servers exchange messages.

UDP protocol.