Lecture 19: Error Analysis

The last few lectures have studied stability: the sensitivity of a mathematical problem to perturbations in the input. We now turn our attention to the actual error in numerical algorithms. We model this as follows: given a problem $f: Y \to Z$ between two normed vector spaces, consider an algorithm as another problem $\tilde{f}: Y \to Z$ that we hope satisfies $f(x) \approx \tilde{f(x)}$.

There are two types of errors we can look at: *forward* and *backward* error. If these errors are "small" we say that the algorithm is *forward stable* or *backward stable*. The definition of "small" depends on the context.

Forward error

The absolute forward error is defined as

$$\|f(x) - f(x)\|_Z$$

while the relative forward error is

$$\frac{\|f(x) - f(x)\|_{Z}}{\|f(x)\|_{Z}}$$

Backward error

Suppose that there exists a Δx so that $\tilde{f(x)} = f(x + \Delta x)$. Then the absolute backward error is defined as

$$\|\Delta x\|_Y$$

and the relative backward error isdefined as

$$\frac{\|\Delta x\|_Y}{\|x\|_Y}.$$

Warning: Backward error may not always be defined: for example, if f(x) = 1 and $\tilde{f(x)} = 0$ we have $f(x + \Delta x) = 1 \neq 0$ for all Δx . In this case, we can define the backward error as ∞ .

Floating point error analysis

We now consider the forward and backward error in algorithms arising from floating point arithmetic.

Here $\mathrm{fl}(x)$ denotes the operator of rounding a number to floating point. This is always exact up to the last bit in the significand. If there are p bits used to represent the significand, define *machine epsilon* as

$$\epsilon_m \triangleq 2^{-p}$$
.

Since we round to nearest bit, we have the property that

$$fl(x) = x(1 + \delta_x)$$

where $|\delta_x| \le \frac{\epsilon_m}{2}$.

We can verify this numerically. eps(Float32) and eps(Float64) give machine epsilon for Float32 and Float64 respectively. For Float32 we have p=23, so this means it returns 2^{-23} :

In [24]:

```
eps(Float32),2.0f0^(-23)
```

Out[24]:

(1.1920929f-7,1.1920929f-7)

We have $fl(x) = x(1 + \delta_x)$, which implies that

$$|x - \mathrm{fl}(x)| = |x| |\delta_x| \le |x| \frac{\epsilon_m}{2}.$$

We confirm this for a simple example:

In [13]:

```
x=1/3

\tilde{x}=Float32(x)

abs(x-\tilde{x}) \le abs(x)*\epsilon m/2
```

Out[13]:

true

We can explain this using the case x=1. If we add $\epsilon_{\it m}/2$ to one we round back to one:

In [25]:

```
Float32(1.0+2.0^(-24))
```

Out[25]:

1.0f0

Any perturbation above this, no matter how small, rounds up:

In [26]:

Float32(1.0+2.0 $^{-24}$)+2.0 $^{-30}$)

Out[26]:

1.000001f0

Thus the worst case is off by the last bit, but within half:

In [27]:

bits(Float32(1.0+2.0 $^{-24}$)+2.0 $^{-25}$))

Out[27]:

In [28]:

bits(Float32(1.0))

Out[28]:

"00111111100000000000000000000000000"

Example 1: error analysis for rounding a number

We now consider the backward and forward error of rounding. Define f(x) = x and $\tilde{f(x)} = fl(x)$, where $f, \tilde{f}: \mathbb{R} \to \mathbb{R}$ with the absolute value norm attached, which denote as $|\cdot| \to |\cdot|$.

Forward error:

$$\frac{|f(x) - \tilde{f(x)}|}{|f(x)|} = \frac{|x - x(1 + \delta_x)|}{|x|} = |\delta_x| \le \frac{\epsilon_m}{2}.$$

Backward error: Since $\tilde{f(x)} = x(1 + \delta_x) = f(x + x\delta_x) = f(x + \Delta x)$ for $\Delta x = x\delta_x$, we have the (relative) backward error

$$\frac{|\Delta x|}{|x|} = |\delta_x| \le \frac{\epsilon_m}{2}$$

Example 2: error analysis for adding two floating point numbers

Assume x and y are floating point numbers. Consider the problem f(x,y) = x + y calculated via the algorithm $f(x,y) = x \oplus y = fl(x+y) = (x+y)(1+\delta_{x+y})$ where $f,f(x,y) = x \oplus y = fl(x+y) = (x+y)(1+\delta_{x+y})$ where $f,f(x,y) = x \oplus y = fl(x+y) = fl(x+y) = fl(x+y)$ with norms $\|\cdot\|_{\infty} \to \|\cdot\|_{\infty}$.

Forward error:

$$\frac{|f(x) - \tilde{f(x)}|}{|f(x)|} = \frac{|x + y - (x + y)(1 + \delta_{x+y})|}{|x + y|} = |\delta_{x+y}| \le \frac{\epsilon_m}{2}.$$

Backward error: Since $\tilde{f}(x,y) = f(x + x\delta_{x+y}x, y + y\delta_{x+y}) = f(x + \Delta x, y + \Delta y)$, we have the backward error

$$\frac{\|\begin{pmatrix} \Delta x \\ \Delta y \end{pmatrix}\|_{\infty}}{\|\begin{pmatrix} x \\ y \end{pmatrix}\|_{\infty}} = |\delta_{x+y}| \le \frac{\epsilon_m}{2}.$$

Example 3: error analysis for adding two real numbers

Assume x and y are general real numbers. Again consider the problem f(x, y) = x + y, but now calculated via the algorithm

$$f(x,y) = f(x) \oplus f(y) = x(1+\delta_x) \oplus y(1+\delta_y) = (x(1+\delta_x) + y(1+\delta_y))(1+\delta_z) = x+y+x(\delta_x)$$
 where $z = x(1+\delta_x) + y(1+\delta_y)$. As before, $f, f \in \mathbb{R}^2 \to \mathbb{R}$ with norms $\|\cdot\|_{\infty} \to \|\cdot\|$.

Backward error: Since $\tilde{f}(x,y) = f(x + x\delta_{x+y}x, y + y\delta_{x+y}) = f(x + \Delta x, y + \Delta y)$, we have the backward error

$$\frac{\|\begin{pmatrix} \Delta x \\ \Delta y \end{pmatrix}\|_{\infty}}{\|\begin{pmatrix} x \\ y \end{pmatrix}\|_{\infty}} = |\delta_{x+y}| \le \frac{\epsilon_m}{2}.$$

Forward error:

$$\frac{|f(x) - \tilde{f(x)}|}{|f(x)|} = \frac{|x + y - (x + y)(1 + \delta_{x+y})|}{|x + y|} = |\delta_{x+y}| \le \frac{\epsilon_m}{2}.$$