



Chapter 3: Introduction to SQL

Database System Concepts, 7th Ed.

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Outline

- Overview of The SQL Query Language
- SQL Data Definition
- Basic Query Structure of SQL Queries
- Additional Basic Operations
- Set Operations
- Null Values
- Aggregate Functions
- Nested Subqueries
- Modification of the Database



History

- IBM Sequel language developed as part of System R project at the IBM San Jose Research Laboratory
- Renamed Structured Query Language (SQL)
- ANSI and ISO standard SQL:
 - SQL-86
 - SQL-89
 - SQL-92
 - SQL:1999 (language name became Y2K compliant!)
 - SQL:2003
- Commercial systems offer most, if not all, SQL-92 features, plus varying feature sets from later standards and special proprietary features.
 - Not all examples here may work on your particular system.



Data Definition Language

The SQL data-definition language (DDL) allows the specification of information about relations, including:

- The schema for each relation.
- The type of values associated with each attribute.
- The Integrity constraints
- The set of indices to be maintained for each relation.
- Security and authorization information for each relation.
- The physical storage structure of each relation on disk.



Domain Types in SQL

- **char(n).** Fixed length character string, with user-specified length n .
- **varchar(n).** Variable length character strings, with user-specified maximum length n .
- **int.** Integer (a finite subset of the integers that is machine-dependent).
- **smallint.** Small integer (a machine-dependent subset of the integer domain type).
- **numeric(p,d).** Fixed point number, with user-specified precision of p digits, with d digits to the right of decimal point. (ex., **numeric(3,1)**, allows 44.5 to be stored exactly, but not 444.5 or 0.32)
- **real, double precision.** Floating point and double-precision floating point numbers, with machine-dependent precision.
- **float(n).** Floating point number, with user-specified precision of at least n digits.
- More are covered in Chapter 4.



Create Table Construct

- An SQL relation is defined using the **create table** command:

create table *r*

$(A_1 D_1, A_2 D_2, \dots, A_n D_n,$
 $\text{(integrity-constraint}_1\text{)},$
 $\dots,$
 $\text{(integrity-constraint}_k\text{)})$

- *r* is the name of the relation
 - each A_i is an attribute name in the schema of relation *r*
 - D_i is the data type of values in the domain of attribute A_i
- Example:

```
create table instructor (  
    ID          char(5),  
    name        varchar(20),  
    dept_name   varchar(20),  
    salary      numeric(8,2))
```



Integrity Constraints in Create Table

- Types of integrity constraints
 - **primary key** (A_1, \dots, A_n)
 - **foreign key** (A_m, \dots, A_n) **references** r
 - **not null**
- SQL prevents any update to the database that violates an integrity constraint.
- Example:

```
create table instructor (
    ID          char(5),
    name        varchar(20) not null,
    dept_name   varchar(20),
    salary      numeric(8,2),
    primary key (ID),
    foreign key (dept_name) references department);
```



And a Few More Relation Definitions

- **create table student (**
ID **varchar(5),**
name **varchar(20) not null,**
dept_name **varchar(20),**
tot_cred **numeric(3,0),**
primary key (ID),
foreign key (dept_name) references department);

- **create table takes (**
ID **varchar(5),**
course_id **varchar(8),**
sec_id **varchar(8),**
semester **varchar(6),**
year **numeric(4,0),**
grade **varchar(2),**
primary key (ID, course_id, sec_id, semester, year) ,
foreign key (ID) references student,
foreign key (course_id, sec_id, semester, year) references section);



And more still

- **create table course (**
course_id **varchar(8),**
title **varchar(50),**
dept_name **varchar(20),**
credits **numeric(2,0),**
primary key (course_id),
foreign key (dept_name) references department);



Updates to tables

- **Insert**
 - **insert into *instructor* values ('10211', 'Smith', 'Biology', 66000);**
- **Delete**
 - Remove all tuples from the *student* relation
 - **delete from *student***
- **Drop Table**
 - **drop table *r***
- **Alter**
 - **alter table *r* add *A D***
 - where *A* is the name of the attribute to be added to relation *r* and *D* is the domain of *A*.
 - All existing tuples in the relation are assigned *null* as the value for the new attribute.
 - **alter table *r* drop *A***
 - where *A* is the name of an attribute of relation *r*
 - Dropping of attributes not supported by many databases.



Basic Query Structure

- A typical SQL query has the form:

```
select A1, A2, ..., An  
from r1, r2, ..., rm  
where P
```

- A_i represents an attribute
- R_j represents a relation
- P is a predicate.
- The result of an SQL query is a relation.



The select Clause

- The **select** clause lists the attributes desired in the result of a query
 - corresponds to the projection operation of the relational algebra
- Example: find the names of all instructors:

```
select name  
      from instructor
```
- NOTE: SQL names are case insensitive (i.e., you may use upper- or lower-case letters.)
 - E.g., *Name* \equiv *NAME* \equiv *name*
 - Some people use upper case wherever we use bold font.



The select Clause (Cont.)

- SQL allows duplicates in relations as well as in query results.
- To force the elimination of duplicates, insert the keyword **distinct** after **select**.
- Find the department names of all instructors, and remove duplicates

```
select distinct dept_name  
from instructor
```

- The keyword **all** specifies that duplicates should not be removed.

```
select all dept_name  
from instructor
```



The select Clause (Cont.)

- An asterisk in the select clause denotes “all attributes”

```
select *
  from instructor
```

- An attribute can be a literal with no **from** clause

```
select '437'
```

- Results is a table with one column and a single row with value “437”
- Can give the column a name using:

```
select '437' as FOO
```

- An attribute can be a literal with **from** clause

```
select 'A'
  from instructor
```

- Result is a table with one column and N rows (number of tuples in the *instructors* table), each row with value “A”



The select Clause (Cont.)

- The **select** clause can contain arithmetic expressions involving the operation, +, -, *, and /, and operating on constants or attributes of tuples.
 - The query:

```
select ID, name, salary/12  
from instructor
```

would return a relation that is the same as the *instructor* relation, except that the value of the attribute *salary* is divided by 12.

- Can rename “*salary/12*” using the **as** clause:

```
select ID, name, salary/12 as monthly_salary
```



The where Clause

- The **where** clause specifies conditions that the result must satisfy
 - Corresponds to the selection predicate of the relational algebra.
- To find all instructors in Comp. Sci. dept

```
select name  
from instructor  
where dept_name = 'Comp. Sci.'
```

- SQL allows the use of the logical connectives **and**, **or**, and **not**
- The operands of the logical connectives can be expressions involving the comparison operators <, <=, >, >=, =, and <>.
- Comparisons can be applied to results of arithmetic expressions
- To find all instructors in Comp. Sci. dept with salary > 80000

```
select name  
from instructor  
where dept_name = 'Comp. Sci.' and salary > 80000
```



The from Clause

- The **from** clause lists the relations involved in the query
 - Corresponds to the Cartesian product operation of the relational algebra.
- Find the Cartesian product *instructor X teaches*

```
select *
  from instructor, teaches
```

 - generates every possible instructor – teaches pair, with all attributes from both relations.
 - For common attributes (e.g., *ID*), the attributes in the resulting table are renamed using the relation name (e.g., *instructor.ID*)
- Cartesian product not very useful directly, but useful combined with where-clause condition (selection operation in relational algebra).



Examples

- Find the names of all instructors who have taught some course and the course_id
 - **select name, course_id
from instructor , teaches
where instructor.ID = teaches.ID**
- Find the names of all instructors in the Art department who have taught some course and the course_id
 - **select name, course_id
from instructor , teaches
where instructor.ID = teaches.ID and instructor.dept_name = 'Art'**



The Rename Operation

- The SQL allows renaming relations and attributes using the **as** clause:
 $old-name \text{ as } new-name$
- Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci'.
 - **select distinct T.name
from instructor as T, instructor as S
where T.salary > S.salary and S.dept_name = 'Comp. Sci.'**
- Keyword **as** is optional and may be omitted
 $instructor \text{ as } T \equiv instructor T$



String Operations

- SQL includes a string-matching operator for comparisons on character strings. The operator **like** uses patterns that are described using two special characters:
 - percent (%). The % character matches any substring.
 - underscore (_). The _ character matches any character.
- Find the names of all instructors whose name includes the substring “dar”.

```
select name  
from instructor  
where name like '%dar%'
```

- Match the string “100%”

```
like '100 \%' escape '\'
```

in that above we use backslash (\) as the escape character.



String Operations (Cont.)

- Patterns are case sensitive.
- Pattern matching examples:
 - 'Intro%' matches any string beginning with “Intro”.
 - '%Comp%' matches any string containing “Comp” as a substring.
 - '___' matches any string of exactly three characters.
 - '___ %' matches any string of at least three characters.
- SQL supports a variety of string operations such as
 - concatenation (using “||”)
 - converting from upper to lower case (and vice versa)
 - finding string length, extracting substrings, etc.



Ordering the Display of Tuples

- List in alphabetic order the names of all instructors

```
select distinct name  
from instructor  
order by name
```

- We may specify **desc** for descending order or **asc** for ascending order, for each attribute; ascending order is the default.
 - Example: **order by name desc**
- Can sort on multiple attributes
 - Example: **order by dept_name, name**



Where Clause Predicates

- SQL includes a **between** comparison operator
- Example: Find the names of all instructors with salary between \$90,000 and \$100,000 (that is, $\geq \$90,000$ and $\leq \$100,000$)
 - **select name
from instructor
where salary between 90000 and 100000**
- Tuple comparison
 - **select name, course_id
from instructor, teaches
where (instructor.ID, dept_name) = (teaches.ID, 'Biology');**



Set Operations

- Find courses that ran in Fall 2017 or in Spring 2018

```
(select course_id from section where sem = 'Fall' and year = 2017)
union
(select course_id from section where sem = 'Spring' and year = 2018)
```

- Find courses that ran in Fall 2017 and in Spring 2018

```
(select course_id from section where sem = 'Fall' and year = 2017)
intersect
(select course_id from section where sem = 'Spring' and year = 2018)
```

- Find courses that ran in Fall 2017 but not in Spring 2018

```
(select course_id from section where sem = 'Fall' and year = 2017)
except
(select course_id from section where sem = 'Spring' and year = 2018)
```



Null Values

- It is possible for tuples to have a null value, denoted by **null**, for some of their attributes
- **null** signifies an unknown value or that a value does not exist.
- The result of any arithmetic expression involving **null** is **null**
 - Example: $5 + \text{null}$ returns **null**
- The predicate **is null** can be used to check for null values.
 - Example: Find all instructors whose salary is null.
**select name
from instructor
where salary is null**
- The predicate **is not null** succeeds if the value on which it is applied is not null.



Null Values (Cont.)

- SQL treats as **unknown** the result of any comparison involving a null value (other than predicates **is null** and **is not null**).
 - Example: $5 < \text{null}$ or $\text{null} <> \text{null}$ or $\text{null} = \text{null}$
- The predicate in a **where** clause can involve Boolean operations (**and**, **or**, **not**); thus the definitions of the Boolean operations need to be extended to deal with the value **unknown**.
 - **and** : $(\text{true and unknown}) = \text{unknown}$,
 $(\text{false and unknown}) = \text{false}$,
 $(\text{unknown and unknown}) = \text{unknown}$
 - **or**: $(\text{unknown or true}) = \text{true}$,
 $(\text{unknown or false}) = \text{unknown}$
 $(\text{unknown or unknown}) = \text{unknown}$
- Result of **where** clause predicate is treated as *false* if it evaluates to *unknown*



Aggregate Functions

- These functions operate on the multiset of values of a column of a relation, and return a value

avg: average value

min: minimum value

max: maximum value

sum: sum of values

count: number of values



Aggregate Functions Examples

- Find the average salary of instructors in the Computer Science department
 - **select avg (salary)
from instructor
where dept_name= 'Comp. Sci.');**
- Find the total number of instructors who teach a course in the Spring 2018 semester
 - **select count (distinct ID)
from teaches
where semester = 'Spring' and year = 2018;**
- Find the number of tuples in the *course* relation
 - **select count (*)
from course;**



Aggregate Functions – Group By

- Find the average salary of instructors in each department
 - `select dept_name, avg (salary) as avg_salary
from instructor
group by dept_name;`

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000

<i>dept_name</i>	<i>avg_salary</i>
Biology	72000
Comp. Sci.	77333
Elec. Eng.	80000
Finance	85000
History	61000
Music	40000
Physics	91000



Aggregation (Cont.)

- Attributes in **select** clause outside of aggregate functions must appear in **group by** list

- /* erroneous query */
select dept_name, ID, avg (salary)
from instructor
group by dept_name;



Aggregate Functions – Having Clause

- Find the names and average salaries of all departments whose average salary is greater than 42000

```
select dept_name, avg (salary) as avg_salary  
from instructor  
group by dept_name  
having avg (salary) > 42000;
```

- Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups



Nested Subqueries

- SQL provides a mechanism for the nesting of subqueries. A **subquery** is a **select-from-where** expression that is nested within another query.
- The nesting can be done in the following SQL query

```
select  $A_1, A_2, \dots, A_n$ 
  from  $r_1, r_2, \dots, r_m$ 
  where  $P$ 
```

as follows:

- **From clause:** r_i can be replaced by any valid subquery
- **Where clause:** P can be replaced with an expression of the form:
 $B <\text{operation}> (\text{subquery})$

B is an attribute and $<\text{operation}>$ to be defined later.

- **Select clause:**
 A_i can be replaced by a subquery that generates a single value.



Set Membership



Set Membership

- Find courses offered in Fall 2017 and in Spring 2018

```
select distinct course_id  
from section  
where semester = 'Fall' and year= 2017 and  
course_id in (select course_id  
from section  
where semester = 'Spring' and year= 2018);
```

- Find courses offered in Fall 2017 but not in Spring 2018



Set Membership (Cont.)

- Name all instructors whose name is neither “Mozart” nor Einstein”

```
select distinct name
from instructor
where name not in ('Mozart', 'Einstein')
```

- Find the total number of (distinct) students who have taken course sections taught by the instructor with *ID* 10101

```
select count (distinct ID)
from takes
where (course_id, sec_id, semester, year) in
      (select course_id, sec_id, semester, year
       from teaches
       where teaches.ID= 10101);
```

- Note: Above query can be written in a much simpler manner.
The formulation above is simply to illustrate SQL features



Set Comparison



Set Comparison – “some” Clause

- Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

```
select distinct T.name  
from instructor as T, instructor as S  
where T.salary > S.salary and S.dept name = 'Biology';
```

- Same query using > **some** clause

```
select name  
from instructor  
where salary > some (select salary  
                      from instructor  
                      where dept name = 'Biology');
```



Definition of “some” Clause

- $F \text{ <comp> some } r \Leftrightarrow \exists t \in r \text{ such that } (F \text{ <comp> } t)$
Where <comp> can be: $<$, \leq , $>$, $=$, \neq

(5 < some

0
5
6

) = true (read: 5 < some tuple in the relation)

(5 < some

0
5

) = false

(5 = some

0
5

) = true

(5 \neq some

0
5

) = true (since $0 \neq 5$)

(= some) \equiv in

However, (\neq some) $\not\equiv$ not in



Set Comparison – “all” Clause

- Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

```
select name  
from instructor  
where salary > all (select salary  
                     from instructor  
                     where dept name = 'Biology');
```



Definition of “all” Clause

- $F \text{ <comp> all } r \Leftrightarrow \forall t \in r (F \text{ <comp> } t)$

$(5 < \text{all} \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$

$(5 < \text{all} \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$

$(5 = \text{all} \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$

$(5 \neq \text{all} \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true} \text{ (since } 5 \neq 4 \text{ and } 5 \neq 6\text{)}$

$(\neq \text{all}) \equiv \text{not in}$

However, $(= \text{all}) \neq \text{in}$



Use of “exists” Clause

- Yet another way of specifying the query “Find all courses taught in both the Fall 2017 semester and in the Spring 2018 semester”

```
select course_id  
from section as S  
where semester = 'Fall' and year = 2017 and  
exists (select *  
        from section as T  
        where semester = 'Spring' and year= 2018  
          and S.course_id = T.course_id);
```

- Correlation name** – variable S in the outer query
- Correlated subquery** – the inner query



Use of “not exists” Clause

- Find all students who have taken all courses offered in the Biology department.

```
select distinct S.ID, S.name
from student as S
where not exists ( (select course_id
                     from course
                     where dept_name = 'Biology')
                   except
                   (select T.course_id
                     from takes as T
                     where S.ID = T.ID));
```

- First nested query lists all courses offered in Biology
- Second nested query lists all courses a particular student took
- Note that $X - Y = \emptyset \Leftrightarrow X \subseteq Y$
- Note: Cannot write this query using = all and its variants



Subqueries in the From Clause



Subqueries in the Form Clause

- SQL allows a subquery expression to be used in the **from** clause
- Find the average instructors' salaries of those departments where the average salary is greater than \$42,000.”

```
select dept_name, avg_salary
from ( select dept_name, avg (salary) as avg_salary
       from instructor
       group by dept_name)
      where avg_salary > 42000;
```

- Note that we do not need to use the **having** clause
- Another way to write above query

```
select dept_name, avg_salary
from ( select dept_name, avg (salary)
       from instructor
       group by dept_name)
      as dept_avg (dept_name, avg_salary)
      where avg_salary > 42000;
```



Deletion

- Delete all instructors

delete from *instructor*

- Delete all instructors from the Finance department

delete from *instructor*

where *dept_name*= 'Finance';

- *Delete all tuples in the instructor relation for those instructors associated with a department located in the Watson building.*

delete from *instructor*

where *dept name* in (**select** *dept name*

from *department*

where *building* = 'Watson');



Insertion

- Add a new tuple to *course*

```
insert into course
```

```
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- or equivalently

```
insert into course (course_id, title, dept_name, credits)
```

```
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- Add a new tuple to *student* with *tot_creds* set to null

```
insert into student
```

```
values ('3003', 'Green', 'Finance', null);
```



Updates

- Give a 5% salary raise to all instructors
update *instructor*
 set salary = salary * 1.05
- Give a 5% salary raise to those instructors who earn less than 70000
update *instructor*
 set salary = salary * 1.05
 where salary < 70000;
- Give a 5% salary raise to instructors whose salary is less than average
update *instructor*
 set salary = salary * 1.05
 where salary < (select avg (salary)
 from *instructor*);



Case Statement for Conditional Updates

- Same query as before but with case statement

```
update instructor
  set salary = case
    when salary <= 100000 then salary * 1.05
    else salary * 1.03
  end
```



Updates with Scalar Subqueries

- Recompute and update tot_creds value for all students

```
update student S
set tot_cred = (select sum(credits)
                 from takes, course
                where takes.course_id = course.course_id and
                      S.ID= takes.ID.and
                      takes.grade <> 'F' and
                      takes.grade is not null);
```

- Sets *tot_creds* to null for students who have not taken any course
- Instead of **sum(credits)**, use:

```
case
  when sum(credits) is not null then sum(credits)
  else 0
end
```



End of Chapter 3