

## Hash Table Methods

### Quadratic Probing

**Quadratic Probing**

- 1) Set counter  $j=0$
- 2) Get the hash value,  $h(K) = (K+j^2) \bmod \text{table-size}$
- 3) If hash-table  $[h(K)]$  is empty, insert the key and stop. Else the space is occupied, we must find next available space.
  - 3.1 increment  $j$  by 1
  - 3.2 compute a new hash value,  $h(K) = (K+j^2) \bmod \text{table-size}$
  - 3.3 Repeat step 3 till  $j$  is equal to table-size
- 4) The hash table is full
- 5) Stop

### Double Hashing

Exam 2 Practice CSE 331 Spring 2018  $m=11$

(5) (2 points) You are inserting an element  $x$  with  $h_1(x) = 42$  and  $h_2(x) = 7$  into a hash table with 11 cells that uses **double hashing**. What is the index of your **third** attempt to insert  $x$ ?

(a) Index 0  
(b) Index 1  
(c) Index 9  
(d) Index 10  
(e) A different index

$h' = [h_1(x) + i \cdot h_2(x)] \bmod m$   
 $\Rightarrow [42 + 2(7)] \bmod 11$   
 $\Rightarrow [56] \bmod 11 \Rightarrow 1$

(6) (2 points) You are inserting an element  $x$  with  $h_1(x) = 42$  and  $h_2(x) = 7$  into a hash table with 11 cells that uses **quadratic probing**. What is the index of your **third** attempt to insert  $x$ ?

(a) Index 0  
(b) Index 1  
(c) Index 9  
(d) Index 10  
(e) A different index

don't care about  $h_2$   
 $h = [h_1(x) + i^2] \bmod m$   
 $42 + 4 \bmod 11$   
 $46 \bmod 11 = 2$

(7) (2 points) What does `MysteryFunction1` do?

Set  $j$  count = 0  
 + 1 for each attempt  
 reset after insert

### Linear Probing

**Linear Probing**

Linear probing is a strategy for resolving collisions or keys that map to the same index in a hash table.

**Strategy:**

- 1) Use a hash function to find the index for a key.
- 2) If that spot contains a value, use the next available spot "a higher index". If you reach the end of the array, go back to the front.

### Path Finding

#### Dijkstra's Algorithm vs. Floyd's

- Dijkstra's algorithm is  $O(m \log n) = O(n^2 \log n)$  in the worst case.
- Floyd's algorithm is  $O(n^3)$ , but computes all shortest paths.
- Dijkstra's algorithm can compute all shortest paths by starting it  $n$  times, once for each node. Thus, to compute all paths, Dijkstra's algorithm is  $O(nm \log n) = O(n^3 \log n)$  in the worst case.
- Which is better depends on the application.
  - Dijkstra's algorithm is better if only a path from a single node to another or to all others is needed.
  - For all paths, Dijkstra's algorithm is better for sparse graphs, and Floyd's algorithm is better for dense graphs.

#### Two main methods

- Depth-first search (DFS)
  - Go as far as possible, then backtrack
  - Uses stacks
- Breadth-first search (BFS)
  - Check neighborhood first and then spread out
  - Uses queues

## Heaps

### Heap(Min) Insert

(11) (2 points) The result of inserting 1 into the min-heap [2, 3, 4, 5, 6, 7]?

(a) [1, 2, 3, 4, 5, 6, 7]  
 (b) [2, 3, 4, 5, 6, 7, 1]  
 (c) [1, 2, 4, 5, 6, 7, 3]  
 (d) [1, 3, 2, 5, 6, 7, 4]  
 (e) None of the above

(15) (2 points) The result of calling `extract-min` from the min-heap [1, 2, 3, 4, 5, 6, 7]?

(a) [2, 3, 4, 5, 6, 7]  
 (b) [7, 2, 3, 4, 5, 6]  
 (c) [2, 7, 3, 4, 5, 6]  
 (d) [2, 4, 3, 5, 6, 7]  
 (e) None of the above

Parent  $\leq$  Child

insert @ bottom

\* Compare left & right [2 4 3 7 5 6]  
 Child to determine which side of tree to traverse  
 min heap (smaller values)  
 Max heap (larger values)  
 Child =  $2n+1$  &  $2n+2$   
 $n = \text{index}$  L R

### Directed Graphs

#### (17) (10 points) Directed Graphs

$v_a$	$v_b, v_c, v_e$
$v_b$	$v_a, v_d$
$v_c$	$v_d$
$v_d$	$v_c, v_e$
$v_e$	

- (a) (6 points) Draw the directed graph associated with the adjacency list shown above.



- (b) (2 points) What is the indegree of  $v_e$ ? 2
- (c) (2 points) What is the out-degree of  $v_b$ ? 2

### Search Methods

#### Degree

- The *degree* of a vertex  $v$  is the number of edges connected to  $v$ .
- In a digraph
  - The out-degree is the number of arcs leaving  $v$
  - The in-degree is the number of arcs entering  $v$

- (11) (2 points) Which search algorithm should you use to find a path through a maze?

- (a) Depth-first search  
 (b) Breadth-first search  
 (c) Dijkstra's Algorithm  
 (d) Floyd-Warshall Algorithm  
 (e) None of the above

- (12) (2 points) How many edges are there in a clique graph with  $n$  vertices? (Hint: In a clique, every vertex shares an edge with every other vertex.)

- (a)  $n$   
 (b)  $2n$   
 (c)  $n^2$   
 (d)  $\frac{n(n+1)}{2}$   
 (e) None of the above

$$\frac{n(n-1)}{2}$$

$n$  vertices each has  $n-1$  edges, each edge connects 2 vertices that share it

## Path Finding examples /Max & Min & Median Heap Returns

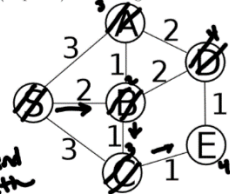
Exam 2 Practice

CSE 331

Spring 2018

(18) (10 points) Pathfinding

only update weights if we find shorter path



(a) (7 points) Find the minimum-weight path from  $s$  to each vertex in the graph above as well as its total weight. Show your work.

	w
A	3
B	2
C	3
D	4
E	4
S	0

(b) (3 points) Construct the shortest path from  $s$  to  $e$ .

$S \rightarrow C \rightarrow E$

(19) (10 points) Generator Functions

```
def getMax():
    # raise exception if list empty.
    error if list.size == 0

    # If maximum unknown, calculate it on demand.
    if maxcount == 0:
        maxval = list[0]
        for each val in list:
            if val == maxval:
                maxcount += 1
            elif val > maxval:
                maxval = val
                maxcount = 1

    # Now it is known, just return it.
    return maxval
```

Problems 1-15: Multiple-choice. Circle the letter of the best response.

(1) (2 points) Below is a Bloom filter with hashes  $h_1(x) = x \bmod 11$  and  $h_2(x) = (3x) \bmod 11$ . Which of the following can be elements of the associated set?

0	1	2	3	4	5	6	7	8	9	10
		X	X		X			X		

- (a) 14  
(b) 27  
(c) 31  
(d) 41  
(e) 42

$h_1$	3	5	9	8
$h_2$	9	4		2

\* both elements must fit in a marked slot

(2) (2 points) What is wrong with the following hash function?

```
def hash_point(pt):
    return pt.x + pt.y
```

- (a) The function is non-deterministic.  
(b) It is not possible to generate all possible hash values.  
(c) The hash values are not uniformly distributed.  
(d) It is not possible to convert a hash function back to the original point.  
(e) Nothing. This is a good hash function.

(3) (2 points) You are inserting an element  $x$  with  $h(x) = 42$  into a hash table with 11 cells that uses linear probing. Unfortunately, the table contains a single element at the spot where you want to insert  $x$ . Where do you insert  $x$ ?

- (a) Index 0  
(b) Index 1  
(c) Index 9  
(d) Index 10  
(e) A different index

$42 \bmod 11 = 9$ . linear probing so we go to the next unfilled slot

(4) (2 points) You are inserting an element  $x$  with  $h(x) = 42$  into a hash table with 11 cells that uses separate chaining. Unfortunately, the table contains a single element at the spot where you want to insert  $x$ . Where do you insert  $x$ ?

- (a) Index 0  
(b) Index 1  
(c) Index 9  
(d) Index 10  
(e) A different index

With chaining we chain the two elements together at that index

```
125 def find_median(seq):
126     """
127     Finds the median (middle) item of the given sequence.
128     Ties are broken arbitrarily.
129     :param seq: an iterable sequence
130     :return: the median element
131     """
132     if not seq:
133         raise IndexError
134     else:
135         min_heap = Heap(lambda a, b: a <= b)
136         max_heap = Heap(lambda a, b: a >= b)
137
138         item = ((len(seq))//2)
139         min_heap.extend(seq[:item])
140         max_heap.extend(seq[item:])
141
142         minpeek = min_heap.peek()
143         maxpeek = max_heap.peek()
144         while minpeek < maxpeek:
145             min_heap.extract()
146             max_heap.extract()
147             min_heap.insert(maxpeek)
148             max_heap.insert(minpeek)
149             minpeek = min_heap.peek()
150             maxpeek = max_heap.peek()
151         return max_heap.peek()
```

