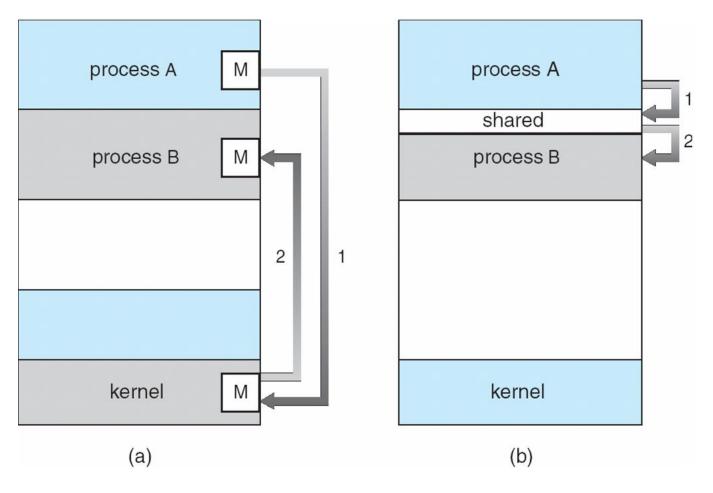
# Process Management (II)

August 30, 2017

## Interprocess Communication

- Processes within a system may be independent or cooperating
- Cooperating process can affect or be affected by other processes, including sharing data
- Reasons for cooperating processes:
  - Information sharing
  - Computation speedup
  - Modularity/Convenience
- Cooperating processes need inter-process communication (IPC)

### Communications Models

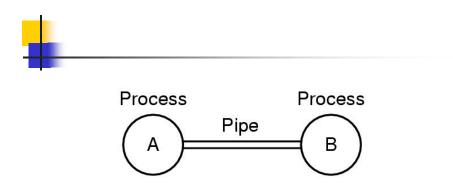


Message Passing

**Shared Memory** 

## Example of Message Passing: Unix Pipes

Pipe sets up communication channel between two (related) processes.



Two processes connected by a pipe

## Unix Pipes

- One process writes to the pipe, the other reads from the pipe.
- Similar to reading from/writing to a file.
- System call:

int fd[2];

pipe(fd);

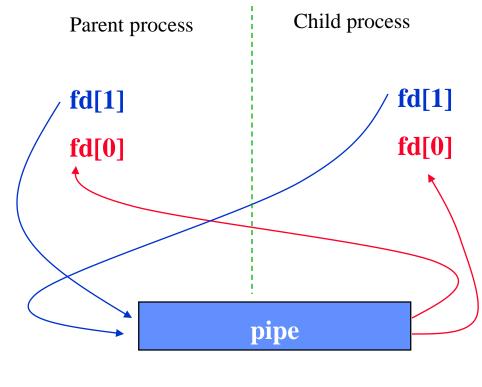
fd[0] now holds descriptor to read from pipe fd[1] now holds descriptor to write into pipe **fd[0]** 

fd[1]

pipe

#### A Simple Example

```
#include <unistd.h>
#include <fcntl.h>
#include <stdio.h>
char *message = "This is a message!!!";
main()
{ char buf[1024];
  int fd[2];
  pipe(fd); /*create pipe*/
  if (fork() != 0) { /* I am the parent */
    write(fd[1], message, strlen (message) + 1);
  else { /*Child code */
    read(fd[0], buf, 1024);
    printf("Got this from parent!!: %s\n", buf);
```



## Synchronization

- Message passing may be either blocking or non-blocking
- Blocking is considered synchronous
  - Blocking send has the sender block until the message is received
  - Blocking receive has the receiver block until a message is available
- Non-blocking is considered asynchronous
  - Non-blocking send has the sender send the message and continue
  - Non-blocking receive has the receiver receive a valid message or null
- Read in pipe is blocking

## Limitations of Pipes

- Processes using a pipe must come from a common ancestor
  - e.g., parent and child; siblings
- Pipes are not permanent
  - disappearing when the process terminates

## Named Pipes (Called FIFO in UNIX)

- Similar to pipes, but with some advantages
  - A FIFO can be created separately from the processes that will use it
  - FIFOs look like files
    - Having an owner, size, access permissions
    - Permanent until deleted with rm

# Creating a FIFO

Using unix command line (calling system program): \$mkfifo <FIFO-NAME> e.g., \$mkfifo fifo1

Using C program
#include <sys/types.h>
#include <sys/stat.h>
int mkfifo(const chair \*pathname, mode\_t mode);

- pathname: path-name of the FIFO
- mode: permission. It could be any combination of S\_IRUSR (owner read), S\_IWUSR (owner write), S\_IRGRP (owner's group read), S\_IWGRP (group write), S\_IROTH (other users read), S\_IWOTH (other write)

e.g., mkfifo("fifo1", S\_IRUSR | S\_IWUSR); //create FIFO fifo1 for owner's processes' communication

## Open an Existing FIFO

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
int open( const chair *pathname, mode_t mode );
```

- pathname: path-name of the FIFO
- open a FIFO to write: open (FIFOName, O\_WRONLY);
- open a FIFO to read: open (FIFOName, O\_RDONLY);
- After a FIFO has been open, read/write operations can be performed on it like with the pipe.

## POSIX Shared Memory

Process first creates shared memory segment segment id = shmget(IPC PRIVATE, size, S IRUSR S IWUSR); //a new shared memory segment is created segment id = shmget(key, size, S IRUSR|S IWUSR|IPC CREAT); //if a SM segment associated with key exists, its ID is returned; otherwise, a new SM segment is created Process wanting access to that shared memory must attach to it Shared memory = (char \*) shmat(id, NULL, 0); Now the process could write to the shared memory sprintf(shared memory, "Writing to shared memory"); When done a process can detach the shared memory from its address space shmdt(shared memory);

## A simple program

```
#include <stdio.h>
#include <sys/shm.h>
#include <sys/stat.h>
int main(){
   int segment_id;
   char *shared_memory;
   const int size = 4096;
   segment_id = shmget(IPC_PRIVATE, size, S_IRUSR | S_IWUSR);
   shared_memory=(char *)shmat(segment_id, NULL,0);
   sprintf(shared_memory, "Hi there!");
   shmdt(shared_memory);
   shmctl(segment_id,IPC_RMID,NULL); //remove the shared memory block
```

# A simple program (cont.): Sharing a memory between parent and child processes

```
#include <stdio.h>
#include <sys/shm.h>
#include <sys/stat.h>
int main(){
   int segment_id;
   char *shared_memory;
   const int size = 4096;
   segment_id = shmget(IPC_PRIVATE, size, S_IRUSR | S_IWUSR);
   shared_memory=(char *)shmat(segment_id, NULL,0);
   if (fork() == 0)
         sprintf(shared_memory, "Hi, this is the child!");
         shmdt(shared_memory);
         exit(0);
   }else{
         wait(NULL);
         printf("%\n", shared_memory);
         shmdt(shared_memory);
         shmctl(segment_id,IPC_RMID,NULL); //remove the shared memory
```

# Another example: Sharing a memory between two processes

```
#define KEY 9876
int main(){
  int segment_id;
  char *shared_memory;
  const int size = 4096;
  segment_id = shmget(KEY, size, S_IRUSR|S_IWUSR|IPC_CREAT);
  shared_memory=(char *)shmat(segment_id, NULL,0);
  sprintf(shared_memory, "Hi, this is the child!");
  shmdt(shared_memory); }
```

```
#define KEY 9876
int main(){
  int segment_id;
  char *shared_memory;
  const int size = 4096;
  segment_id = shmget(KEY, size, S_IRUSR|S_IWUSR|IPC_CREAT);
  shared_memory=(char *)shmat(segment_id, NULL,0);
  printf("%s\n", shared_memory);
  shmdt(shared_memory);}
15
```

Signals are software interrupts

• Examples: A user hitting ctrl+c; a process sending a signal to kill another process.

In Linux, signals are defined at /usr/include/bits/signum.h, and every signal has a name that begins with characters SIG:

- SIGINT: when a user presses ctrl+c.
- SIGKILL
- SIGUSR1, SIGUSR2: you are free to define what they mean

How to send a signal to a process?

• System program kill

```
example: >kill -9 1000 //kill process with ID 1000
```

• System call kill

```
#include <sys/types.h>
#include <signal.h>
int kill(pid_t pid, int sig);
```

How to handle a signal in a process?

- Define a handler function
- Make system call signal to link the function with a signal

```
#include <stdio.h>
#include <signal.h>
#include <unistd.h>
void sig_handler(int signo) {
  if(signo == SIGINT) printf("received SIGINT\n");
int main(void) {
  if(signal(SIGINT, sig_handler)==SIG_ERR)
       printf("Cannot catch SIGINT\n");
  while (1) sleep (1);
```

#### Quiz

```
#include <stdio.h>
#include <sys/shm.h>
#include <sys/stat.h>
int main(){
  int i;
  for(i=0;i<2;i++) fork();
}</pre>
```

How many processes are created while the above program is executed?

Thread (I)
Kernel Threads

September 1, 2017

#### Example of Multi-Process Collaboration

- Finding the lines containing key word "Linux" in all .txt files in the current directory
- Linux command line:
  grep "Linux" \*.txt
- Using multi-process collaboration to do the job?

```
int numFiles; //total number of files
char fileNames[MAX_NUM_FILES][255]; //array storing the file names
int i;
for(i=0; i<numFiles; i++)
  if (fork() == 0){
     FILE *fp = fopen(fileNames[i], "r");
     ... //search for lines containing "Linux" in the file
     exit(0);
}</pre>
```

#### Multi-Process

Stack

Heap

Data

Code (Text)

PCB

Register Values (including PC)

Information about memory, files,

. . .

Stack

Heap

Data

Code (Text)

**PCB** 

Register Values (including PC)

Information about memory, files, Stack

Heap

Data

Code (Text)

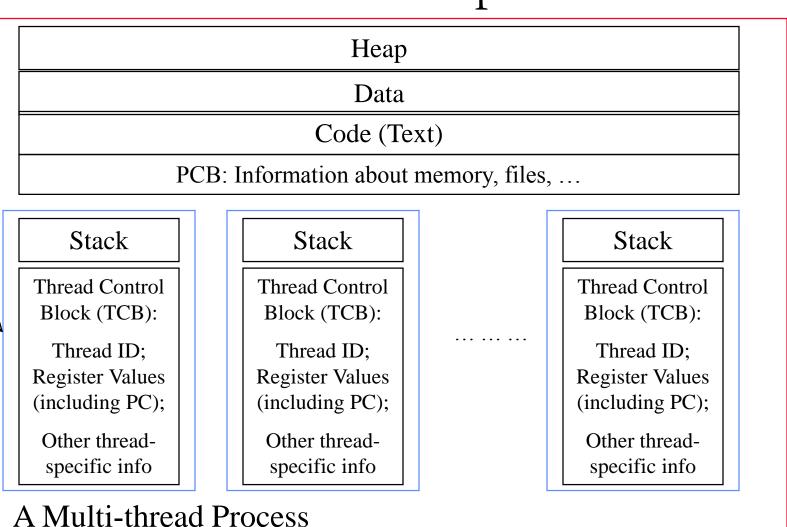
**PCB** 

Register Values (including PC)

Information about memory, files,

. . .

# A Process with Multiple Threads

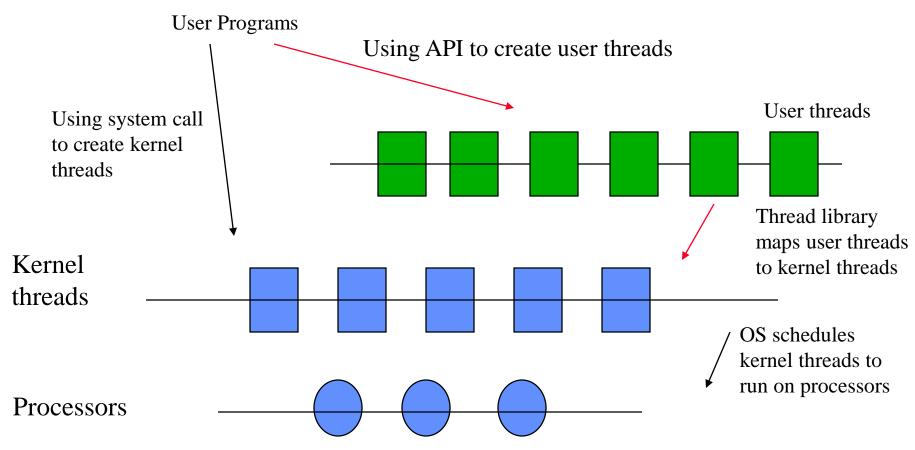


Thread

#### Benefits

- Responsiveness (similar to multi-process)
  - Allowing a program to continue running even if part of it is blocked
  - Allowing a program to perform multiple tasks concurrently
- Resource sharing
  - Threads of the same process share the memory and resources
- Economy
  - Allocating memory and resources for process creation is costly.
- Scalability (similar to multi-process)
  - A single-thread process can only run on one processor regardless how many are available; multi-threaded process can increase parallelism on a multi-processor machine.

# Threads in a Computer System



#### Kernel Threads

- Directly created/managed/scheduled by the OS kernel
- Virtually all contemporary OSes support kernel threads

#### Linux Kernel Threads

- Linux refers to them as tasks rather than threads
- Thread creation is done through clone() system call

```
#define _GNU_SOURCE
```

#include <sched.h>

int clone(int (\*fn)(void \*), void \*child\_stack, int flags, void \*arg);

flag	meaning
CLONE_FS	File-system information is shared.
CLONE_VM	The same memory space is shared.
CLONE_SIGHAND	Signal handlers are shared.
CLONE_FILES	The set of open files is shared.

```
#include _GNU_SOURCE
#include <stdio.h>
#include <sched.h>
#include <unistd.h>
#include <stdlib.h>
#include <fcntl.h>

int variable, fd;
int do_something(); //function to be executed by the new task (thread)
```

```
int main() {
void *child_stack;
char tempch;
variable = 9;
fd = open("test.file", O_RDONLY);
child_stack = (void *)malloc(16384); child_stack += 16383;
printf("The variable was %d\n", variable);
clone(do_something, child_stack, CLONE_VM | CLONE_FILES, NULL); /*A*/
sleep(1);
printf("The variable is now %d\n", variable);
if (read(fd, &tempch, 1) \leq 1) {
      perror("File Read Error"); exit(1);}
      printf("Parent could read from the file\n");}
```

```
int do_something() {
  char tempch;
  variable = 42;
 if (read(fd, &tempch, 1) \leq 1) {
    perror("File Read Error");
   _exit(1);
  printf("Child could read from the file\n");
  close(fd);
 _exit(0);
```

1. What is the output of the program?

```
2. What if Line A is changed to: clone(do_something, child_stack, CLONE_VM, NULL); or clone(do_something, child_stack, CLONE_FILES, NULL); or Clone(do_something, child_stack, 0, NULL);
```