

# Process Synchronization (III)

September 20, 2017

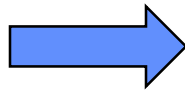
# Reader-Writer Problem

- ❏ A data set is shared among a number of concurrent processes
  - ❏ Readers: only read the data set; they do **not** perform any updates
  - ❏ Writers: can both read and write
- ❏ Problem – allow multiple readers to read at the same time. Only one single writer can access the shared data at a time.

# Reader-Writer Problem: Analogies



It is fine for a group of people to watch (**read**) a TV simultaneously.



But changing (**write**) programs must be done sequentially (no concurrency)!

# Reader-Writer Problem: Observations

- ❏ A writer is allowed to access the data set only when there is no other writer or reader accessing
- ❏ Readers can simultaneously access the data set
  - When a reader arrives, it is allowed to access the data set if (i) no one is accessing the data set, or (ii) some other reader is reading; otherwise (some writer is writing) , it should wait.
  - For a waiting reader, it is allowed to read if the writing writer completes.

# Reader-Writer Problem: Observations → Solution

- ❏ A writer is allowed to access the data set only when there is no other writer or reader accessing
- ❏ Data Structure
  - ❏ Semaphore `wrt` initialized to 1 → to ensure at most one writer (no 2+ writers or one writer and reader(s)) allowed to access the data set (To guard access to data set)
- ❏ Code: The structure of a writer process

```
wait (wrt) ;  
.... writing is performed ...  
signal (wrt) ;
```

# Reader-Writer Problem:


## Observations → Solution

- ❏ Readers can simultaneously access the data set
  - When a reader arrives, it is allowed to access the data set if (i) no one is accessing the data set, or (ii) some other reader is reading; otherwise (some writer is writing) , it should wait.
  - For a waiting reader, it is allowed to read if the writing writer completes.

## Additional Data Structures:

- ❏ Integer **readcount** initialized to 0 → to keep track of number of readers that are currently reading the data set
- ❏ Semaphore **mutex** initialized to 1 → to ensure mutual exclusive modification of readcount (To guard access to readcount)


# Readers-Writers Problem: Solution


 The structure of a reader process


```
wait (mutex) ;  
if (readcount == 0) wait (wrt) ;  
    //check if there is already reader being reading  
    //if yes, go ahead to read;  
    //otherwise, go on only if no writer writing  
    //note: if this reader has to wait, follow-up readers should wait too; so block them  
readcount ++ ;  
signal (mutex);  
    ... reading is performed...  
wait (mutex) ;  
readcount -- ;  
if (readcount == 0) signal (wrt) ; //if this is the last reader  
signal (mutex) ;
```

# Semaphore Implementation

 Must guarantee:

 No two processes can actively execute `wait()` on the same semaphore at the same time (but when one is blocked by wait, another is allowed to execute it)

 No two processes can modify the value of semaphore at the same time (i.e., `signal()` and “S- -” of `wait()` must be atomic)

 Thus, implementation becomes the critical section problem where the wait and signal code are placed in the critical section.



# Semaphore Implementation

All competing processes share a variable lock initialized to false

```
Signal(S) {  
    while(testAndSet(&lock));  
    S++;  
    lock=False;  
}
```

# Semaphore Implementation

All competing processes share a variable lock initialized to false

```
wait(S) {  
    while(testAndSet(&lock));  
    while (S <= 0) {  
        lock=False;  
        while(testAndSet(&lock);  
    }  
    S --;  
    lock=False;  
}
```

# Semaphore Implementation

- ☐ Ideally, will not have **busy waiting** in critical section
  - ☐ Otherwise, applications may spend lots of time in critical sections.

# Semaphore Implementation with no Busy waiting


- ❏ With each semaphore, there are
  - ❏ value (of type integer)
  - ❏ list: an associated waiting queue. Each entry in a waiting queue has two items: (1) pointer to a process/thread; (2) pointer to next record in the list
- ❏ Two (atomic) operations:
  - ❏ **block** – place the process invoking the operation on a waiting queue.
  - ❏ **wakeup** – remove one of processes on a waiting queue and place it in the ready queue.

# Semaphore Implementation with no Busy waiting

 (Rewritten) Definition of wait:

```
wait(semaphore *S) {  
    S->value--;  
    if (S->value < 0) {  
        add this process to S->list;  
        block();  
    }  
}
```

# Semaphore Implementation with no Busy waiting

 Implementation of wait (example):


```
wait(semaphore *S) {  
    while(testAndSet(&lock));  
    S->value--;  
    if (S->value < 0) {  
        add this process to S->list;  
        lock=False;  
        block();  
    }  
    else lock=False;  
}
```

# Semaphore Implementation with no Busy waiting

 Definition of signal:

```
signal(semaphore *S) {  
    S->value++;  
    if (S->value <= 0) {  
        remove a process P from S->list;  
        wakeup(P);  
    }  
}
```

# Semaphore Implementation with no Busy waiting

 Implementation of signal (example):

```
signal(semaphore *S) {  
    while(testAndSet(&lock));  
    S->value++;  
    if (S->value <= 0) {  
        remove a process P from S->list;  
        wakeup(P);  
    }  
    lock=False;  
}
```