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10. ALGORITHM computeRange ( root )

// Input: root node

// Output: range ( difference between the largest & smallest numbers in the tree )

node ← root

while ( node.left ≠ null ) do

node ← node.left

end while

smallest ← node.value

node ← root

while ( node.right ≠ null ) do

node ← node.right

end while

largest ← node.value

range ← largest - smallest

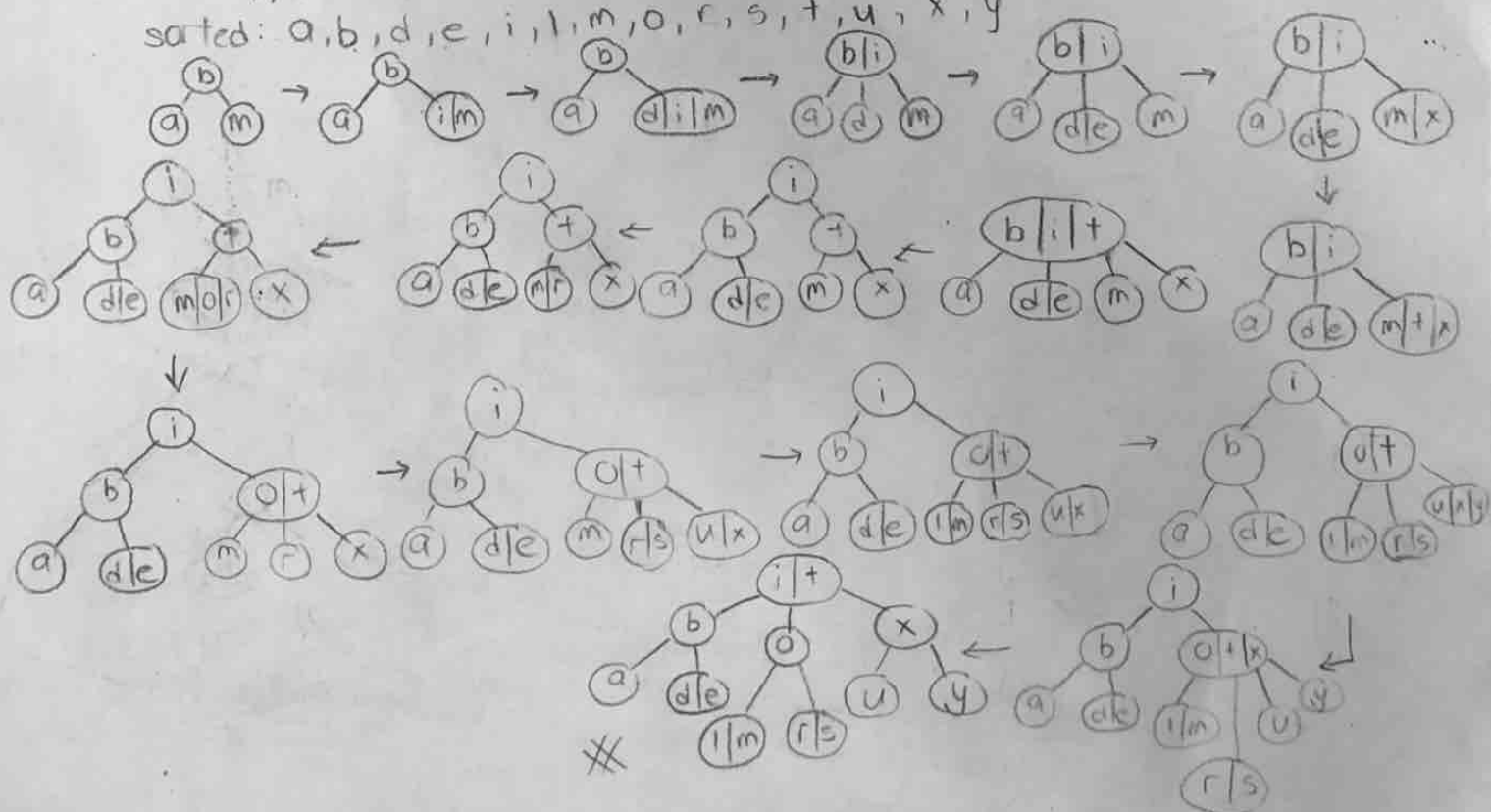
return range

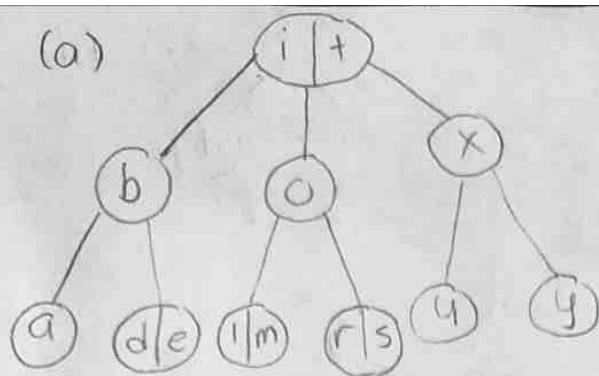
Worst case:  $O(\log n) + O(\log n) \approx O(\log n)$

- b. The smallest key is always located at the leftmost node, while the largest key is always located at the rightmost node.

2a. list: a, m, b, i, d, e, x, t, r, o, u, s, l, y

sorted: a, b, d, e, i, l, m, o, r, s, t, u, x, y





(b) - There are 14 keys in the AVL tree.

- The largest number of key comparisons in a successful search will be in the searches for e, m, and s will be equal to 5.

- Average number of comparisons:

$$\begin{aligned}
 & \frac{1}{14} C(i) + \frac{1}{14} C(+) + \frac{1}{14} C(b) + \frac{1}{14} C(o) + \frac{1}{14} C(x) + \frac{1}{14} C(a) + \frac{1}{14} C(d) + \frac{1}{14} C(e) \\
 & + \frac{1}{14} C(l) + \frac{1}{14} C(m) + \frac{1}{14} C(r) + \frac{1}{14} C(s) + \frac{1}{14} C(u) + \frac{1}{14} C(y) \\
 & = (1 + 2 + 3 + 3 + 3 + 4 + 4 + 5 + 4 + 5 + 4 + 5 + 4 + 4) \frac{1}{14} \\
 & = \frac{51}{14}
 \end{aligned}$$

3 ALGORITHM computeRange (root)

// Input : root node

// Output : range (difference between the largest and smallest numbers in the tree)

node  $\leftarrow$  root

while (node.left  $\neq$  null) do

node  $\leftarrow$  node.left

end while

smallest  $\leftarrow$  node.value

node  $\leftarrow$  root

while (node.right  $\neq$  null) do

node  $\leftarrow$  node.right

end while

largest  $\leftarrow$  node.value

range  $\leftarrow$  largest - smallest

return range

Worst case efficiency:  $O(\log n) + O(\log n) \approx O(\log n)$

4. - No, it is not possible.

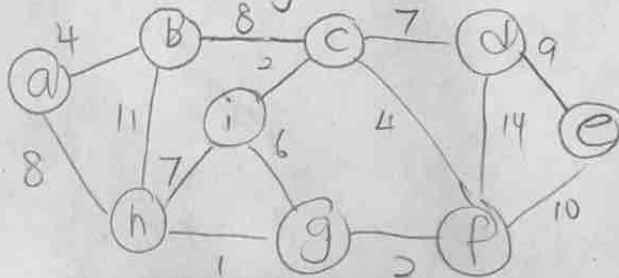
- If we assume that  $T$  is generated with vertices,  
 $V_1, V_2, \dots, V_n$

- If a new edge is added and let  $V_{n+1}$  to be  
the vertex  $(V_1, V_{n+1})$

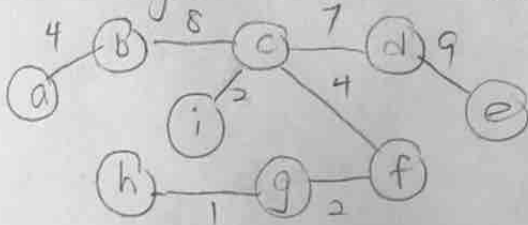
- The diagram will change as  $(V_1, V_{n+1})$  will replace  
 $(V_1, V_2)$  if the weight of  $(V_1, V_{n+1})$  is smaller than  
 $(V_1, V_2)$

- Therefore, the statement is wrong.

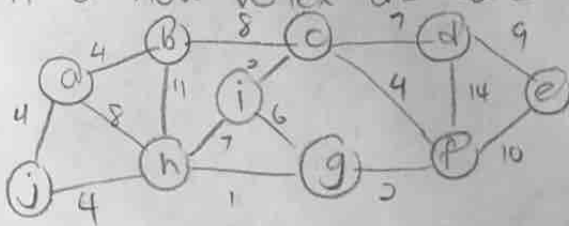
- set the diagram as



-  $T$  is generated as shown below



- If a new vertex and some edges are added at the beginning



-  $T_{new}$  is generated as below which proves the statement is wrong

