

CS 312 Work Assignment 5

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1. (a) 4

5*6

8/2

(b) $\vdash S_1 : c_1$

$$\frac{\vdash S_1 : c_1 \quad \vdash S_2 : c_2}{\vdash S_1 * S_2 : c_1 * c_2}$$

$\vdash S_2 : c_2$ is not zero

$$\frac{\vdash S_1 : c_1}{\vdash S_1 / S_2 : c_1 / c_2}$$

$$\frac{\vdash e_2[e_1/x] : e}{\vdash \text{let } x = e_1 \text{ in } e_2 : e}$$

(c) $\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 * S_2 : \text{nonzero}}$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{zero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{zero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 / S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 / S_2 : \text{nonzero}}$$

$\Gamma \vdash S_1 : t_1$

$t = t_1$

$$\frac{\Gamma[x \leftarrow t] \vdash S_2 : t_2}{\Gamma \vdash \text{let } x = S_1 \text{ in } S_2 : T_2}$$

(d) **Base:**

nonzero *i*

$$\Gamma \vdash i : \text{nonzero}$$

$$\frac{\text{zero } i}{\Gamma \vdash i : \text{zero}}$$

Inductive Case 1:

$$\frac{\vdash S_1 : c_1 \quad \vdash S_2 : c_2}{\vdash S_1 * S_2 : c_1 * c_2}$$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 * S_2 : \text{nonzero}}$$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{zero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{zero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

Since we know c_1 is a non zero or a zero, and c_2 is a non zero or a zero we know $c_1 * c_2$ is also non zero or zero.

Inductive Case 2:

$$\frac{\vdash S_1 : c_1 \quad \vdash S_2 : c_2}{\vdash S_1 / S_2 : c_1 / c_2}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 / S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 / S_2 : \text{nonzero}}$$

Since we know c_1 is a zero or a non zero, and we know c_2 is a non zero, c_1 / c_2 must be a non zero.

Inductive Case 3:

$$\frac{\vdash e_2[e_1/x] : e}{\vdash \text{let } x = e_1 \text{ in } e_2 : e}$$

$$\begin{array}{l} \Gamma \vdash S_1 : t_1 \\ t = t_1 \\ \hline \Gamma[x \leftarrow t] \vdash S_2 : t_2 \end{array}$$

$$\Gamma \vdash \text{let } x = S_1 \text{ in } S_2 : T_2$$

Since we know c_1 is zero or non zero, we know c_2 must also be zero or non_zero .

(e) **Base:**

$$\frac{\text{nonzero } i}{\Gamma \vdash i : \text{nonzero}}$$

$$\frac{\text{zero } i}{\Gamma \vdash i : \text{zero}}$$

Clearly, if i types as zero or non zero, the corresponding operational semantics rule applies

Inductive Case 1:

$$\frac{\vdash S_1 : c_1 \quad \vdash S_2 : c_2}{\vdash S_1 * S_2 : c_1 * c_2}$$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 * S_2 : \text{nonzero}}$$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{zero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{zero}}{\Gamma \vdash S_1 * S_2 : \text{zero}}$$

We know from the inductive hypothesis that the evaluation of S_1 and S_2 will never get stuck. We also know from preservation that the expressions S_1 and S_2 must evaluate to zero or non zero, therefore the operational semantics rule for multiplication will always apply since the hypotheses only require that c_1 and c_2 are zero or non zero.

Inductive Case 2:

$$\frac{\vdash S_1 : c_1 \quad \vdash S_2 : c_2}{\vdash S_1 / S_2 : c_1 / c_2}$$

$$\frac{\Gamma \vdash S_1 : \text{zero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 / S_2 : \text{zero}}$$

$$\frac{\Gamma \vdash S_1 : \text{nonzero} \quad \Gamma \vdash S_2 : \text{nonzero}}{\Gamma \vdash S_1 / S_2 : \text{nonzero}}$$

We know from the inductive hypothesis that the evaluation of S_1 and S_2 will never get stuck. We also know from preservation that the expression S_1 must evaluate to zero or non zero and S_2 must evaluate to non zero, therefore the operational semantics rule for division will always apply since the hypotheses only require that c_1 is a non zero or zero and c_2 is a non zero.

Inductive Case 3:

$$\frac{\vdash e_2[e_1/x] : e}{\vdash \text{let } x = e_1 \text{ in } e_2 : e}$$

$$\begin{array}{l} \Gamma \vdash S_1 : t_1 \\ t = t_1 \\ \Gamma[x \leftarrow t] \vdash S_2 : t_2 \\ \hline \Gamma \vdash \text{let } x = S_1 \text{ in } S_2 : T_2 \end{array}$$

Here, we know from the inductive hypothesis that S_1 is of type t_1 and $E[x \leftarrow t]S_2 : t_2$ will not get stuck since they are well-typed. Therefore, this rule will also always apply.

2. (a) These are large step Operational Semantics.

$$\begin{array}{l} \text{(b)} \quad TU \text{ 'd' } \rightarrow' D' \\ \quad \quad TU' a' \rightarrow' A' \\ \quad \quad TL' B' \rightarrow' b' \end{array}$$

- (c) 1. A char is defined as a char.
 2. AL returns a lowercase char if it's the correct case.
 3. AU returns an uppercase char if it's the correct case.
 4. TU converts a char to uppercase and returns it.
 5. TL converts a char to lowercase and returns it.
- (d) A runtime error will occur if AL is called on an uppercase char. A runtime error will also occur if AU is called on a lowercase char.
- (e) A type system to prevent runtime errors for this language would consist of the types lowercase and uppercase.
 The concretization of the type lowercase is: $\gamma(\text{lowercase}) = \text{'a'-'b'}$.
 The concretization of the type uppercase is: $\gamma(\text{uppercase}) = \text{'A'-'B'}$.
- (f) lowercase c
 $\Gamma \vdash c : c$
- lowercase c
 $\Gamma \vdash c : c$

$$\frac{\Gamma \vdash P : \textit{lowercase}}{\Gamma \vdash AL\ P : \textit{lowercase}}$$
$$\frac{\Gamma \vdash P : \textit{uppercase}}{\Gamma \vdash AU\ P : \textit{uppercase}}$$
$$\frac{\Gamma \vdash P : \textit{lowercase}}{\Gamma \vdash TL\ P : \textit{lowercase}}$$
$$\frac{\Gamma \vdash P : \textit{uppercase}}{\Gamma \vdash TL\ P : \textit{lowercase}}$$
$$\frac{\Gamma \vdash P : \textit{uppercase}}{\Gamma \vdash TU\ P : \textit{uppercase}}$$
$$\frac{\Gamma \vdash P : \textit{lowercase}}{\Gamma \vdash TU\ P : \textit{uppercase}}$$