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1
    Math
1.1 卡塔蘭數 Catalan number
// 1, 1, 2, 5, 14, 42
11 c[N];
c[0] = 1;
for(int i=1;i<n;i++)</pre>
   c[i]=2*(2*(i-1)+1)*c[i-1]/(i+1);
1.2 錯排公式 Derangements
// 1, 0, 1, 2, 9, 44, 265
ll d[n]; //d[0]=1
d[1]=0, d[2]=1;
for(int i=3;i<n;i++){</pre>
   d[i]=(i-1)*(d[i-1]+d[i-2]);
1.3 Euler function
11 phi(11 n){ // 計算小於n的數中與n互質的有幾個
   ll res = n, a=n; // O(sqrtN)
   for(ll i=2;i*i<=a;i++){</pre>
       if(a%i==0){
           res = res/i*(i-1);
           while(a%i==0) a/=i;
   if(a>1) res = res/a*(a-1);
   return res;
}
1.4 Euler function(建表)
void phi_table(int n,int *phi){
   memset(phi,0,sizeof(phi)); //初始化
   phi[1] = 1;
   for(int i=2;i<=n;i++){</pre>
       if(phi[i]==0){
           for(int j=i;j<=n;j+=i){</pre>
              if(phi[j]==0) phi[j] = j;
              phi[j] = phi[j] / i * (i - 1);
           }
       }
   }
1.5 Sieve Prime
```

```
const int N = 20000000; // 質數表大小
bool sieve[N];
vector<int> prime;
void linear_sieve(){
    for (int i = 2; i < N; i++){</pre>
         if (!sieve[i]) prime.push_back(i);
         for (int p : prime){
             if (i * p >= N) break;
sieve[i * p] = true;
             if (i % p == 0) break;
         }
    }
```

#### 1.6 快速冪

# 1.7 Cn 取 m 模反, Combination, inverse

```
| #define MXN 1'000'005
| #define N 1'000'000
| long long fac[MXN], inv[MXN];
| fac[0] = 1; // 0! = 1
| for(long long i = 1; i <= N; i++)
| fac[i] = fac[i-1] * i % MOD; // 階乘
| inv[N] = FastPow(fac[N], MOD-2); // 模逆元
| for(ll i = N-1; i >=0; i--)
| inv[i] = inv[i+1] * (i+1) % MOD;
| ll c(ll n, ll m){return fac[n]*inv[m]*inv[n-m];}
```

### 1.8 質因數分解 Factorizationn

#### 1.9 Miller Rabinn

```
/*Miller_Rabin 質數判定*/
// n < 4,759,123,141
                            3 : 2, 7, 61
// n < 1,122,004,669,633
                            4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383
                                  6 : pirmes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
11 magic[N] = {};
bool witness(ll a, ll n, ll u, int t) {
   if (!a) return 0;
   ll x = mypow(a, u, n); //快速冪
    for (int i = 0; i < t; i++) {</pre>
       ll nx = mul(x, x, n); //快速乘
        if (nx == 1 && x != 1 && x != n - 1) return 1;
        x = nx:
   return x != 1;
bool miller_rabin(ll n) {
    int s = (magic number size);
        // iterate s times of witness on n
        if (n < 2) return 0;
```

```
if (!(n & 1)) return n == 2;
ll u = n - 1; int t = 0;
// n-1 = u*2^t
while (!(u & 1)) u >>= 1, t++;
while (s--) {
    ll a = magic[s] % n;
    if (witness(a, n, u, t)) return 0;
}
return 1;
}
```

### 1.10 乘法取餘 Multiplication

```
/*大數乘法取餘數*/
ll mul(ll x, ll y, ll mod) {
    ll ret = x * y - (ll)((long double)x / mod * y) *
        mod;
    // LL ret=x*y-(LL)((long double)x*y/mod+0.5)*mod;
    return ret < 0 ? ret + mod : ret;
}
```

#### 1.11 快速乘法 karatsuba

```
/*karatsuba 快速乘法*/
// Get size of the numbers
int getSize(ll num){
    int count = 0;
    while (num > 0)
        count++;
        num /= 10;
    return count;
11 karatsuba(11 X, 11 Y){
   // Base Case
    if (X < 10 && Y < 10)
        return X * Y;
    // determine the size of X and Y
    int size = fmax(getSize(X), getSize(Y));
    // Split X and Y
    int n = (int)ceil(size / 2.0);
    11 p = (11)pow(10, n);
    11 a = (11)floor(X / (double)p);
    11 b = X \% p;
    11 c = (11)floor(Y / (double)p);
    11 d = Y \% p;
    // Recur until base case
    11 ac = karatsuba(a, c);
    11 bd = karatsuba(b, d);
    11 e = karatsuba(a + b, c + d) - ac - bd;
    // return the equation
    return (11)(pow(10 * 1L, 2 * n) * ac + pow(10 * 1L,
         n) * e + bd);
```

### 1.12 ax+by=gcd(a,b)

```
| /*ax+by=gcd(a,b) 一組解*/
| ll a, b, x, y;
| ll exgcd(ll a, ll b, ll& x, ll& y) {
| if (b) {
| ll d = exgcd(b, a % b, y, x);
| return y -= a / b * x, d;
| }
| return x = 1, y = 0, a;
| }
```

#### 1.13 josephusl

```
| /*約瑟夫問題:n個人圍成一桌,數到m的人出列*/
| int josephus(int n, int m) { //n人每m次
| int ans = 0;
| for (int i = 1; i <= n; ++i)
| ans = (ans + m) % i;
| return ans;
| }
```

# 1.14 大數 Big number

```
/*大數(Big Number)*/
void add(int a[100], int b[100], int c[100]){
    int i = 0, carry = 0;
    for (i = 0; i < 100; ++i) {</pre>
        c[i] = a[i] + b[i] + carry;
        carry = c[i] / 10;
        c[i] %= 10;
void sub(int a[100], int b[100], int c[100]){
    int i = 0, borrow = 0;
    for (i = 0; i < 100; ++i) {
        c[i] = a[i] - b[i] - borrow;
        if (c[i] < 0) {</pre>
            borrow = 1;
            c[i] += 10;
        }
        else
            borrow = 0;
    }
void mul(int a[100], int b[100], int c[100]){
    int i = 0, j = 0, carry = 0;
    for (i = 0; i < 100; ++i) {
        if (a[i] == 0) continue;
        for (j = 0; j < MAX; ++j)
            c[i + j] += a[i] * b[i];
    for (i = 0; i < MAX; ++i) {</pre>
        carry = c[i] / 10;
        c[i] %= 10;
void div(int a[100], int b[100], int c[100]){
    int t[100];
    for (i = 100 - 1; i >= 0; i--) {
        for (int k = 9; k > 0; k--) // 嘗試商數
            mul(b + i, k, t);
            if (largerthan(a + i, t))
                 sub(a + i, t, c + i);
                break:
            }
        }
    }
}
```

### 1.15 矩陣快速冪

```
struct Martix_fast_pow{ //O(len^3 lg k)
     LL init(int _len,LL m=9223372036854775783LL){
         len=_len, mod=m;
         // mfp.solve(k,{0, 1}, {1, 1}) k'th fib {值,係
         數} // 0-base
     LL solve(LL n,vector<vector<LL>> poly){
         if(n<len) return poly[n][0];</pre>
         vector<vector<LL>> mar(len,vector<LL>(len,0)),x
             (len,vector<LL>(len,0));
                                     mar[i][i]=1;
         for(int i=0;i<len;i++)</pre>
         for(int i=0;i+1<len;i++) x[i][i+1]=1;</pre>
         for(int i=0;i<len;i++)</pre>
                                    x[len-1][i]=poly[i
             ][1];
         while(n){
             if(n&1) mar=mar*x;
             n >>=1, x=x*x;
         LL ans=0;
         for(int i=0;i<len;i++)</pre>
                                   ans=(ans+mar[len-1][i
             ]*poly[i][0]%mod)%mod;
         return ans;
}mfp;
```

#### 1.16 GaussElimination

```
/*GaussElimination*/
// by bcw_codebook
const int MAXN = 300;
const double EPS = 1e-8;
int n;
double A[MAXN][MAXN];
void Gauss() {
  for(int i = 0; i < n; i++) {</pre>
    bool ok = 0;
    for(int j = i; j < n; j++) {</pre>
       if(fabs(A[j][i]) > EPS) {
         swap(A[j], A[i]);
         ok = 1;
         break;
      }
    if(!ok) continue;
    double fs = A[i][i];
    for(int j = i+1; j < n; j++) {</pre>
       double r = A[j][i] / fs;
       for(int k = i; k < n; k++) {</pre>
         A[j][k] -= A[i][k] * r;
    }
  }
}
```

### 1.17 Epsilon

```
| /*精準度(Epsilon)*/
float eps = 1e-8;
bool Equal(float a, float b)
    return fabs(a - b) < eps
bool NEqual(float a, float b)
    return fabs(a - b) > eps
bool Less(float a, float b)
    return (a - b) < -eps
bool Greater(float a, float b)
    return (a - b) > eps
```

#### 1.18 Floor-Ceil

```
/*floor向下取整,ceil向上取整*/
int floor(int a,int b){return a/b - (a%b and a<0^b<0);}
int ceil (int a,int b){return a/b + (a%b and a<0^b>0);}
```

# 2 Geometry

#### 2.1 Struct

```
struct Pt{
    ld x, y;
};
struct Line{
    Pt st, ed;
struct Circle{
    Pt o; // 圓心
    ld r; // 半徑
};
struct poly{
    int n; // n 邊形
    vector<Pt> pts;
};
struct Pt {
  11 x, y; // ll or ld
  Pt(){}
  Pt(ll _x,ll _y){
    x=_x, y=_y;
  Pt operator+(const Pt &a) {
    return Pt(x+a.x, y+a.y);
  Pt operator-(const Pt &a) {
    return Pt(x-a.x, y-a.y);
  Pt operator*(const ld &a) {
    return Pt(x*a, y*a);
  Pt operator/(const ld &a) {
    return Pt(x/a, y/a);
  ll operator*(const Pt &a){ //內積
    return x*a.x + y*a.y;
  ll operator^(const Pt &a){ //外積
    return x*a.y - y*a.x;
  bool operator<(const Pt &a) const {</pre>
    return x < a.x || (x == a.x && y < a.y);
  friend int cross(const Pt& o,const Pt& a,const Pt& b)
      { //向量外積
    Pt lhs = o-a, rhs = o-b;
    return lhs.x*rhs.y - lhs.y*rhs.x;
  friend bool operator ==(const Pt& lhs,const Pt& rhs){
    return (rhs.x==lhs.x && rhs.y==lhs.y);
};
```

#### 2.2 Function

```
//是否三點共線
bool collinearity(const Pt& a, const Pt& b, const Pt& c
){
return (b-a)^(c-a) < EPS; //外積為0: ==0
}
//判斷點是否在線段上
bool inLine(const Pt& p, const Line& li){
return collinearity(li.st, li.ed, p) && (li.st-p)*(
li.ed-p) < EPS; //內積為負: <=0
}
//平行四邊形面積
ll areaPt(Pt i, Pt j, Pt k){
return abs(cross(i, j, k));
}
//兩點距離
ld distance(Pt i,Pt j){
```

```
Pt vt = i - j;
    return sqrt((ld)(vt.x)*(vt.x)+(ld)(vt.y)*(vt.y));
}
ld distance(int i,int j){
    Pt vt = p[i]-p[j];
    return sqrt((ld)(vt.x)*(vt.x)+(ld)(vt.y)*(vt.y));
}

// 輸出
round(double x); //四捨五入至整數
cout << fixed << setprecision(11); //精度
void print(ll x){ //兩倍面積輸出判斷
    if(x&1)
        cout << (x>>1) << ".5\n";
else
        cout << (x>>1) << "\n";
}
```

#### 2.3 凸包 Convex hull

```
vector<Pt> convex_hull(vector<Pt> hull){ //凸包
    sort(hull.begin(),hull.end());
    int top=0:
    vector<Pt> stk;
    for(int i=0;i<hull.size();i++){</pre>
        while(top>=2&&cross(stk[top-2],stk[top-1],hull[
            i])<=0) //或<0
            stk.pop_back(),top--;
        stk.push_back(hull[i]);
        top++;
    for(int i=hull.size()-2,t=top+1;i>=0;i--){
        while(top>=t&&cross(stk[top-2],stk[top-1],hull[
            i])<=0) //或<0
            stk.pop back(),top--;
        stk.push_back(hull[i]);
        top++;
    stk.pop_back();
    return stk;
}
```

### 2.4 旋轉卡尺

```
// 最遠兩點配對
double FarthestPair(vector<Pt> arr){
  double ret = 0;
  for (int i = 0, j = i + 1; i < arr.size(); i++){</pre>
    while (distance(i, j) < distance(i, (j + 1) % arr.</pre>
        size())){
      j = (j + 1) \% arr.size();
   ret = max(ret, distance(i, j));
  }
  return ret;
}
// 點是否凸包內
bool in_hull(vector<Pt> &hull, Pt p){
  int flag = 0;
  for (int i = 0; i < hull.size(); i++){</pre>
    if (cross(p, hull[i], hull[(i + 1) % hull.size()])
      //與凸包上前後兩點作外積,如果夾角在0~180度內,則
          在內側。
      flag++;
   }
  if (flag == hull.size())
    return true;
  else
    return false;
}
```

```
// 線找點最短距離
ld rot(vector<Pt> hull){
  1d mn = 2.0 * r;
  int n = hull.size();
  hull.pb(hull[0]);
  for (int i = 0, j = i + 2; i < n; i++){
    //平行四邊形面積 除以 底 = 高
    ld area = abs(cross(hull[i], hull[i + 1], hull[j]))
    ld base = abs(dis(hull[i], hull[i + 1]));
    while (area <= abs(cross(hull[i], hull[i + 1], hull</pre>
        [(j + 1) % n])){
      j = (j + 1) \% n;
      area = abs(cross(hull[i], hull[i + 1], hull[j]));
    mn = min(mn, area / base);
  return mn;
}
// 最大三角形
ld solve(vector<Pt> arr){
  1d ret = 0;
  for (int i = 0; i < arr.size(); i++)</pre>
    for (int j = i + 1; j < arr.size(); j++)</pre>
      for (int k = j + 1; k < arr.size(); k++)</pre>
        ret = max(ret, area(i, j, k));
  return ret;
}
// 最大四邊形
ll solve(){//兩倍面積
  11 \text{ ret} = 0;
  for (int i = 0; i < hull.size(); i++){</pre>
    int k = (i + 1) \% hull.size(), k2 = (i + 3) \% hull.
        size();
    for (int j = i + 2; j <= i + hull.size() - 2; j++){</pre>
      while (area(i, j % hull.size(), k) < area(i, j %</pre>
          hull.size(), (k + 1) % hull.size())){
        k = (k + 1) % hull.size();
      while (area(i, j % hull.size(), k2) < area(i, j %</pre>
           hull.size(), (k2 + 1) % hull.size())){
        k2 = (k2 + 1) \% hull.size();
      ret = max(ret, area(i, j % hull.size(), k) + area
          (i, j % hull.size(), k2));
    }
  }
  return ret;
```

# 3 Algorithm

#### 3.1 Binary search

#### 3.2 DFS

```
|/*n queen*/
|//k為第幾行,a[k]為第幾列,n個皇后
int a[100], n, count;
|void DFS(int k) {
| if (k > n) {//當k=n+1時找到解
```

```
count++;
        print("anser");
    else {
        for (int i = 1; i <= n; i++) {
            a[k] = i; //存入皇后
            if (check(a, k))DFS(k + 1);//放入,求下一行
/*Traveling Knight Problem*/
#define X 5 //棋盤
#define Y 5
int dir[8][2] =
\{\{1,2\},\{2,1\},\{1,-2\},\{-2,1\},\{-1,2\},\{2,-1\},\{-1,-2\},\{-2,-1\}\};
int board[X][Y]={0},tot=0,_x,_y;
void dfs(int x,int y,int t) {
 if (t>X*Y) {
    print("anser");
    return;
  for (int i=0;i<8;i++) {</pre>
    int xx=x+dir[i][0];
    int yy=y+dir[i][1];
    if ((xx>=X)||(xx<0)||(yy>=Y)||(yy<0)||(board[xx][yy</pre>
        ])) continue;
    board[xx][yy]=t;
    dfs(xx,yy,t+1);
    board[xx][yy]=0; //回溯
signed main(){
  board[_x][_y]=1; //starts
  dfs(_x,_y,2);
```

#### 3.3 Brute Force

```
#define MAXN 1<<18+5 //雙倍空間
/*折半枚舉 與 二進制枚舉*/
int main() {
    int n, m, i, temp;
    11 \mod, \mod \max = 0;
    vector<11> arr, ans(MAXN,0), ans2(MAXN,0);
    cin >> n >> m;
    for(i=0;i<n;i++){</pre>
        cin >> temp;
        arr.push_back(temp%m);
    //折半枚舉
    for(int i=0; i<(1<<(n/2)); i++){ //2^{(n/2)}}
        for(int j=0;j<n/2;j++){</pre>
            if(i>>j&1) //二進制枚舉(選或不選)
                ans[i] = (ans[i] + arr[j]) % m; //前半
                    枚舉
        }
    for(int i=0;i<(1<<(n-n/2));i++){ //2^(n-n/2)
        for(int j=0;j<(n-n/2);j++){</pre>
            if(i>>j&1) ans2[i] = (ans2[i] + arr[n/2+j])
                 % m; //後半枚舉
        }
    //二分維護
    temp = 1 << (n-n/2);
    sort(ans2.begin(), ans2.begin() + temp);
    for(auto i:ans){
        mod_max = max(mod_max, i + *(upper_bound(ans2.
            begin(), ans2.begin() + temp, m-1-i)-1));
        //mod 最大為m-1,配對另一半最優解
    cout << mod_max <<"\n";</pre>
```

#### 3.4 Divide and Conquer

```
/*最近點對*/
double dist(pair<double, double> a, pair<double, double</pre>
    > b)
  return sqrt(pow((a.first - b.first), 2) + pow((a.
      second - b.second), 2));
double closest(int 1, int r) {
  if (1 >= r)
    return 10000;
  int mid = (1 + r) / 2;
  double radl = closest(1, mid);
  double radr = closest(mid + 1, r);
  double range = min(radl, radr), middle = (point[mid].
      first + point[mid + 1].first) / 2, minimum;
  minimum = range;
  for (int i = mid + 1; i <= r && point[i].first <</pre>
      middle + range; i++) {
    for (int j = mid; j >= 1 && point[j].first > middle
          - range; j--) {
      minimum = min(minimum, dist(point[i], point[j]));
    }
  }
  return minimum;
void quicksort(int arr[], int left, int right){
    if (left >= right) return;
    int i = left, j = right;
    int k = (left + right) / 2;
    int pivot = arr[k];
    while (true){
        while (arr[i] < pivot) i++;</pre>
        while (arr[j] > pivot) j--;
        if (i >= j) return;
        swap(arr[i], arr[j]);
        i++; j--;
    quicksort(arr, left, j);
    quicksort(arr, i, right);
void Merge(vector<int> &Arr, int front, int mid, int
    end) {
    vector<int> Left(Arr.begin() + front, Arr.begin() +
         mid + 1);
    vector<int> Right(Arr.begin() + mid + 1, Arr.begin
        () + end + 1);
    int idxLeft = 0, idxRight = 0;
    Left.insert(Left.end(), INT_MAX);
    Right.insert(Right.end(), INT_MAX);
    for (int i = front; i <= end; i++) {</pre>
        if (Left[idxLeft] < Right[idxRight]) {</pre>
            Arr[i] = Left[idxLeft];
             idxLeft++;
        } else {
            Arr[i] = Right[idxRight];
             idxRight++;
        }
    }
void MergeSort(vector<int> &Arr, int front, int end) {
    if (front >= end)
        return;
    int mid = front + (end - front) / 2;
    MergeSort(Arr, front, mid);
    MergeSort(Arr, mid + 1, end);
    Merge(Arr, front, mid, end);
}
```

#### 4 Data Structure

#### **4.1 SQRT**

```
// build O(n)
// update O(√n)
```

```
// query 0(√n)
//分塊結構
//假設要求區間總和
struct blk{
                         //每塊的全部元素
    vector<int> local;
                         //儲存每塊的總和
    int global;
    int tag;
                         //儲存整塊一起更新的值
    blk(){
                         //初始化
        local.clear();
                        //清空區間元素
        tag = global = 0; //將區間總和先設為0
};
vector<blk> b;
void build(){
    int len=sqrt(n),num=(n+len-1)/len;
    for(int i=0;i<n;i++){</pre>
                          //第i個元素分在第 i/len 塊
        cin>>x;
        //存入區間中
        b[i/len].local.push_back(x);
        //更新區間總和
        b[i/len].global += x;
void update(int ql,int qr,int v){
    int blk_l=ql/len,blk_r=qr/len,ret=0;
    if(blk_l == blk_r){
        //如果都在同一塊直接一個一個跑過去就好
        for(int i=ql;i<=qr;i++)</pre>
            b[blk_1].local[i%len]+=v;
        b[blk_1].global+=(qr-ql+1)*v;
        return;
    for(int i=ql;i<(blk_l+1)*len;i++){ //最左的那一塊
        b[blk_1].local[i%len]+=v;
        b[blk_1].global+=v;
    for(int i=blk_l+1;i<blk_r;i++){ //中間每塊
        b[i].tag+=v;
        b[i].global+=v*len;
    for(int i=blk_r*len;i<=qr;i++){ //最右的那一塊
        b[blk_r].local[i%len]+=v;
        b[blk_r].global+=v;
    }
int query(int ql,int qr){
    int blk_l=q1/len,blk_r=qr/len,ret=0;
    if(blk_l == blk_r){
        //如果都在同一塊直接一個一個跑過去就好
        for(int i=ql;i<=qr;i++)</pre>
            ret+=b[blk_1].local[i%len]+b[blk_1].tag;
        return ret;
    for(int i=ql;i<(blk_l+1)*len;i++)</pre>
                                       //最左的那一塊
        ret+=b[blk_1].local[i%len]+b[blk_1].tag;
                                    //中間每塊的總和
    for(int i=blk_l+1;i<blk_r;i++)</pre>
        ret+=b[i].global;
    for(int i=blk_r*len;i<=qr;i++)</pre>
                                    //最右的那一塊
        ret+=b[blk_r].local[i%len]+b[blk_r].tag;
    return ret;
| }
```

### 4.2 Mo's Algorithm

```
//先判斷是不是在同一塊
        //不同塊的話就比較塊的順序,否則比較右界r
};
int num = 0;
int cnt[1000005], ans[30005];
vector<query> q;
void add(int x){
                  //新增元素到區間內
   ++cnt[x];
   if(cnt[x] == 1)
                    ++num:
void sub(int x){
                 //從區間內移除元素
    --cnt[x];
   if(cnt[x] == 0)
                     --num:
void solve(){
   sort(q.begin(),q.end());
   for(int i=0,l=-1,r=0;i<n;i++){</pre>
       while(1>q[i].1) add(--1);
       while(r<q[i].r) add(++r);//記得要先做新增元素的
       while(1<q[i].1) sub(1++);//再做移除元素的
       while(r>q[i].r) sub(r--);
                             //移到區間後儲存答案
       ans[q[i].id] = num;
}
```

#### 4.3 BIT

```
/*BIT 樹狀數組(動態前綴和)*/
//BIT and Array start at 1
#define MAXN 100005 //最大區間<MAXN
vector <int> arr(MAXN); //原始陣列
vector <int> bit(MAXN); //BIT數組
//前綴和查詢
11 query(int i) { //index
    11 \text{ ret} = 0;
    while(i > 0) ret += bit[i], i -= i & -i; // 1-base
        i-lowbit(i)
    return ret;
}
// 單點增值
void modify(int i, int val) { //index,value
    while(i <= MAXN) bit[i] += val, i += i & -i; // i+</pre>
        Lowbit(i)
}
```

### 4.4 Segment tree

```
/*Segment tree 線段樹(區間問題)*/
//segment tree and Array start at 1
// [l,r] 最大區間設為[1,n]
// [ql,qr] 目標區間
// pos,val 修改位置,修改值
#define MAXN 100005*4 //tree大小為4n
#define cl(x) (x*2)
                   //左子節點index
#define cr(x) (x*2+1) //右子節點index
#define NO_TAG 0 //懶惰記號
vector <int> tag(MAXN);
vector <int> arr(MAXN);
vector <int> tree(MAXN);
void build(int i, int l, int r){ //i為當前節點index, l,r
   為當前遞迴區間
   if(1 == r){ // 遞迴到區間大小為1
       tree[i] = arr[l];
       return;
   int mid=(1+r)/2; //往兩邊遞迴
   build(cl(i),l,mid);
   build(cr(i),mid+1,r);
```

```
tree[i] = max(tree[cl(i)], tree[cr(i)]); //<-可修改
   //將節點的值設成左右子節點的最大值
}
// i 為當前節點index, l, r當前區間左右界, ql, qr詢問左
int query(int i,int l,int r,int ql,int qr){
   if(q1 <= 1 && r <= qr){ //若當前區間在詢問區間內,
       直接回傳區間最大值
       return tree[i];
   int mid=(1+r)/2, ret=0; //<-可修改條件
   if(ql<=mid) // 如果左子區間在詢問區間內
       ret = max(ret, query(cl(i),1,mid,ql,qr));
           <-可修改條件
   if(qr> mid) // 如果右子區間在詢問區間內
       ret = max(ret, query(cr(i),mid+1,r,ql,qr)); //
           <-可修改條件
   return ret;
/*單點修改*/
void update(int i,int l,int r,int pos,int val){
   if(1 == r){ // 修改 a[pos] 的值為 val
       tree[i] = val;
       return;
   int mid=(1+r)/2;
   if(pos <= mid) // 如果修改位置在左子節點,往左遞迴
       update(cl(i),1,mid,pos,val);
   else // 否則往右遞迴
       update(cr(i),mid+1,r,pos,val);
   tree[i] = max(tree[cl(i)], tree[cr(i)]);
                                           //<-可
       修改條件
/*區間修改*/
//將區間 [l, r] 的值都加 v
void push(int i,int l,int r){
   if(tag[i] != NO_TAG){ // 判斷是否有打標記,NO_TAG=0
       tree[i] += tag[i]; // 有的話就更新當前節點的值
       if(1 != r){ // 如果有左右子節點把標記往下打
           tag[cl(i)] += tag[i];
          tag[cr(i)] += tag[i];
       tag[i] = NO_TAG; // 更新後把標記消掉
void pull(int i,int l,int r){
   int mid = (1+r)/2;
   push(cl(i),l,mid); push(cr(i),mid+1,r);
   tree[i] = max(tree[cl(i)], tree[cr(i)]);
void update(int i,int l,int r,int ql,int qr,int v){
   push(i,l,r);
   if(q1<=1 && r<=qr){
       tag[i] += v; //將區間 [l, r] 的值都加 v
       return:
   int mid=(1+r)/2;
   if(ql<=mid) update(cl(i),l,mid,ql,qr,v);</pre>
   if(qr> mid) update(cr(i),mid+1,r,ql,qr,v);
   pull(i,l,r);
/*動態開點*/
struct node{
   node *1, *r;
   int val, tag;
void update(node *x, int 1, int r, int q1, int qr, int
   v){
   push(x, 1, r);
```

```
if(q1 <= 1 && r <= qr){
    x->tag += v;
    return;
}
int mid=(l+r)>>1;
if(q1 <= mid){
    if(x->1 == nullptr)//判斷是否有節點
        x->1 = new node();
    update(x->1, 1 , mid, q1, qr, v);
}
if(qr > mid){
    if(x->r == nullptr)//判斷是否有節點
        x->r = new node();
    update(x->r, mid+1, r, q1, qr, v);
}
pull(x, 1, r);
}
```

## 4.5 Persistent Segment tree and DSU

```
#define push_back emplace_back
struct node{
   ll val;
   node *1, *r;
   node()\{val = 0;\}
11 n,idx=0;
vector<node *> version;
//用一個vector紀錄全部版本的根節點
node mem[MAXN*25];
node *newNode(){
   return &mem[idx++];
node *build(int 1, int r){
   node *x = newNode();
   if(1 == r) return x;
   int mid = (l+r)>>1;
   x->1 = build(1, mid);
   x->r = build(mid+1, r);
   return x;
node *update_version(node *pre, 11 1, 11 r, 11 pos, 11
   node *x = newNode(); //當前位置建立新節點
   if(1 == r){
       x \rightarrow val = v;
       return x;
   int mid = (l+r)>>1;
   if(pos <= mid){ //更新左邊
    //左邊節點連向新節點
       x->l = update_version(pre->l, l, mid, pos, v);
                      //右邊連到原本的右邊
       x->r = pre->r;
   }
   else{ //更新右邊
   //右邊節點連向新節點
                      //左邊連到原本的左邊
       x->1 = pre->1;
       x->r = update_version(pre->r, mid+1, r, pos, v)
   x-val = min(x->l->val, x->r->val); //<-修改
   return x;
11 query(node *x,int ql,int qr,int v){ //bin search
   if(ql == qr) return qr;
   int mid=(ql+qr)>>1;
   if(x->1->val<v) // 如果左子區間在詢問區間內
       return query(x->1,q1,mid,v);
    else// 如果右子區間在詢問區間內
       return query(x->r,mid+1,qr,v);
void add_version(int x,int v){ //修改位置 x 的值為 v
   version.push_back(update_version(version.back(), 0,
        n-1, x, v));
```

```
//前一個版本
int find(int x) {
  int fa = query(version.back(), 0, n - 1, x);
  if (fa == x)return x;
  return find(fa);
void merge(int a, int b) {
  int fa = find(a), fb = find(b);
  if (sz[fa] < sz[fb])</pre>
    swap(fa, fb);
  sz[fa] += sz[fb];
  add_version(fb, fa);
signed main(){
    io
    11 q,temp,i,l,r;
    cin >> n >> q;
    version.push_back(build(0,n-1));
    for(i=1;i<=n;i++){</pre>
        cin >> temp;
        add_version(temp,i);
    for(i=0;i<q;i++){</pre>
        cin >> 1 >> r;
        cout << query(version[r],0,n-1,1) <<"\n";</pre>
}
```

### 4.6 Treap

```
struct Treap{
   int key,pri,sz;
                      //key,priority,size
   Treap *1, *r;
                      //左右子樹
   Treap(){}
   Treap(int _key){
       key = _key;
                      //隨機的數維持樹的平衡
       pri = rand();
       sz = 1;
       l = r = nullptr;
   }
};
Treap *root;
int Size(Treap* x){ return x ? x->sz : 0 ; }
void pull(Treap *x){ x->sz = Size(x->l) + Size(x->r) +
   1;}
Treap* merge(Treap *a,Treap *b){
   //其中一個子樹為空則回傳另一個
   if(!a || !b)
                  return a ? a : b;
   if(a->pri > b->pri){//如果a的pri比較大則a比較上面
       a->r = merge(a->r,b);//將a的右子樹跟b合併
       pull(a);
       return a;
   }
   else{} //如果b的pri比較大則b比較上面
       b->1 = merge(a,b->1);//將b的左子樹根a合併
       pull(b);
       return b;
void splitByKth(Treap *x,int k,Treap*& a,Treap*& b){
   if(!x){ a = b = nullptr; }
   else if(Size(x->1) + 1 <= k){
       splitByKth(x->r, k - Size(x->l) - 1, a->r, b);
       pull(a);
   }
   else{
       splitByKth(x->1, k, a, b->1);
       pull(b);
void splitByKey(Treap *x,int k,Treap*& a,Treap*& b){
   if(!x){ a = b = nullptr; }
```

```
else if(x->key<=k){</pre>
        a = x:
        splitByKey(x->r, k, a->r, b);
        pull(a);
    else{
        b = x;
        splitByKey(x->1, k, a, b->1);
        pull(b);
}
                                //新增一個值為val的元素
void insert(int val){
    Treap *x = new Treap(val); //設一個treap節點
    Treap *1,*r;
    splitByKey(root, val, l, r);//找到新節點要放的位置
    root = merge(merge(1,x),r); //合併到原本的treap裡
void erase(int val){
                                 //移除所有值為val的元素
    Treap *1,*mid,*r;
    splitByKey(root, val, 1, r);//把小於等於val的丟到L
    splitByKey(l, val-1, l, mid);
    //小於val的丟到l,等於val的就會在mid裡
                                //將除了val以外的值合併
    root = merge(1,r);
int findVal(int val){ //小於等於val的size
    int size = -1;
    Treap *1, *r;
    splitByKey(root, val, l, r); //把小於等於val的丟到l
    size = Size(1):
    root = merge(1,r);
    return size;
void interval(Treap *&o, int l, int r) {// [l,r]區間
    Treap *a, *b, *c;
    splitByKey(o, l - 1, a, b), splitByKey(b, r, b, c);
    // operate
    o = merge(a, merge(b, c));
void inOrderTraverse(Treap* o, int print) {// 中序
    if (o != NULL){
        push(o);
        inOrderTraverse(o->1, print);
        // print
        if(print) cout << o->val <<"</pre>
        inOrderTraverse(o->r, print);
    }
// Rank Tree
// Kth(k): 查找第k小的元素
// Rank(x):x的名次,即x是第幾小的元素
int kth(Treap* o, int k){
    if(o == NULL || k > o -> sz || k <= 0)    return 0;</pre>
    int s = (o \rightarrow 1 == NULL ? 0 : o \rightarrow 1 \rightarrow sz);
    if(k == s + 1) return o -> key;
    else if(k <= s) return kth(o -> 1, k);
                    return kth(o \rightarrow r, k - s - 1);
int rank(Node* o, int x){
    if(o == NULL) return 0;
    int res = 0;
    int s = (o \rightarrow 1 == NULL ? 0 : o \rightarrow 1 \rightarrow sz);
    if(x \le o \rightarrow key){
        res += rank(o -> 1, x);
        res += x == o \rightarrow key;
    else{
        res += s + 1;
        res += rank(o \rightarrow r, x);
    return res;
}
```

# 5 Graph(path)

## 5.1 Disjoint Set(Union-Find)

```
// 宣告父節點陣列 f
int f[N];
int sz[N]; // 子樹大小
void init(int n) {
    for (int i = 0; i < n; i++){</pre>
        f[i] = i;
        sz[i] = 1;
    }
int find(int x) {
    return f[x] == x ? x : f[x] = find(f[x]);
void merge(int x, int y) {
    x = find(x), y = find(y);
    if (x != y){
        sz[x] += sz[y];
        f[y] = x;
    }
| }
```

# 5.2 Kruskal's algorithm 最小生成樹

```
/*Kruskal's algorithm 最小生成樹*/
//搭配 Disjoint Set(Union-Find)
struct Edge {
    int u, v, w; // 點 u 連到點 v 並且邊權為 w
    friend bool operator<(const Edge& lhs, const Edge&</pre>
        rhs) {
        return lhs.w > rhs.w;//兩條邊比較大小用邊權比較
    }
}:
priority queue < Edge > graph();// 宣告邊型態的陣列 graph
int kruskal(int m){
    int tot = 0;
    for (int i = 0; i < m ; i++) {</pre>
        if (find(graph.top().u) != find(graph.top().v))
             { // 如果兩點未聯通
           merge(graph.top().u, graph.top().v);
              // 將兩點設成同一個集合
           tot += graph.top().w; // 權重加進答案
        graph.pop();
    return tot;
signed main() {
    cin >> n >> m; //node,edge
    init(n);
    for (int i = 0; i < m; i++) {</pre>
        cin >> u >> v >> w;
        graph.push(Edge{u,v,w});
    cout << kruskal(m) << "\n";</pre>
```

# 5.3 Minimum Steiner Tree.cpp 最小斯坦因樹

```
// Minimum Steiner Tree 重要點的mst

// O(V 3^T + V^2 2^T)

struct SteinerTree{

#define V 33

#define T 8

#define INF 1023456789

int n , dst[V][V] , dp[1 << T][V] , tdst[V];

void init( int _n ){

    n = _n;

    for( int i = 0 ; i < n ; i ++ ){

        for( int j = 0 ; j < n ; j ++ )
```

```
dst[ i ][ j ] = INF;
       dst[ i ][ i ] = 0;
  void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
    dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  void shortest_path(){ // using spfa may faster
    for( int k = 0 ; k < n ; k ++ )</pre>
       for( int i = 0 ; i < n ; i ++ )</pre>
         for( int j = 0 ; j < n ; j ++ )</pre>
            dst[ i ][ j ] = min( dst[ i ][ j ],
                  dst[ i ][ k ] + dst[ k ][ j ] );
  }// call shorest_path before solve
  int solve( const vector<int>& ter ){
    int t = (int)ter.size();
    for( int i = 0 ; i < ( 1 << t ) ; i ++ )</pre>
       for( int j = 0 ; j < n ; j ++ )</pre>
         dp[ i ][ j ] = INF;
     for( int i = 0 ; i < n ; i ++ )</pre>
       dp[0][i] = 0;
     for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
       if( msk == ( msk & (-msk) ) ){
         int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )</pre>
           dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
         continue;
       }
       // Transfer starts at i=0, These points will be
            used as relay points
       for( int i = 0 ; i < n ; i ++ )</pre>
         for( int submsk = ( msk - 1 ) & msk ; submsk ;
                    submsk = (submsk - 1) \& msk)
              dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                dp[ submsk ][ i ] +
                                dp[ msk ^ submsk ][ i ] );
       for( int i = 0 ; i < n ; i ++ ){</pre>
         tdst[ i ] = INF;
         for( int j = 0 ; j < n ; j ++ )
  tdst[ i ] = min( tdst[ i ],</pre>
                        dp[ msk ][ j ] + dst[ j ][ i ] );
       for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = tdst[ i ];</pre>
    int ans = INF;
    for( int i = 0 ; i < n ; i ++ )</pre>
       ans = min(ans, dp[(1 << t) - 1][i]);
    return ans;
} }solver;
```

# 5.4 Dijkstra's algorithm

```
/*Dijkstra's algorithm 單源最短路徑*/
#define MAX_V 100
#define INF 10000
struct Edge {
  int idx,w;
bool operator>(const Edge& a, const Edge& b) {
 return a.w > b.w;
int dist[MAX_V];
vector<vector<Edge> > adj(MAX_V);
void dijkstra(int vn, int s) {
 vector <bool> vis(vn, false);
  fill(dist, dist + vn, INF); dist[s] = 0;
  priority_queue <Edge, vector<Edge>, greater<Edge> >
      pq;
  Edge node;
  node.idx = s; node.w = 0;
  pq.emplace(node);
  while (!pq.empty()) {
    int u = pq.top().idx; pq.pop();
    if (vis[u])continue;
    vis[u] = true;
```

```
for (auto v : adj[u]) {
    if (dist[v.idx] > dist[u] + v.w) {
        dist[v.idx] = dist[u] + v.w;
        node.w = dist[v.idx];
        node.idx = v.idx;
        pq.emplace(node);
    }
}

signed main() {
    //從start連接到end的最短路徑
    cin >> start >> end;
    dijkstra(n, start);
    if(dist[end]==INF) cout << "NO\n";
    else cout << dist[end] <<"\n";
}
```

## 5.5 Floyd-Warshall

#### 5.6 SPFA

```
/*SPFA 單源最短路徑(negative cycle)*/
struct Edge {
   int idx, w;
};
vector<Edge> adj[MAX_V]; //adjacency list
vector<bool> inp(MAX_V);
int dist[MAX_V];
//return true if negative cycle exists
bool spfa(int vn, int s) {
    fill(dist, dist + vn, INF); dist[s] = 0;
    vector<int> cnt(vn, 0);
    vector<bool> inq(vn, 0);
    queue<int> q; q.push(s); inq[s] = true;
    while (!q.empty()) {
        int u = q.front(); q.pop();
        inq[u] = false;
        for (auto v : adj[u]) {
            if (dist[v.idx] > dist[u] + v.w) {
                if (++cnt[v.idx] >= vn)return true;
                dist[v.idx] = dist[u] + v.w;
                if (!inq[v.idx]) inq[v.idx] = true, q.
                    push(v.idx);
            }
        }
    return false;
```

# 6 Graph(Tree)

#### 6.1 DFS and BFS

```
//DFS
void dfs(int x){
    vis[x]=1;
    for(int i:adj[x]){
        if(!vis[i])
             dfs(i);
    }
}
//BFS
void bfs(int s){
    queue<int> q;
    q.push(s);
    vis[s]=1;
    while(!q.empty()){
        int x=q.front();q.pop();
        for(int i:adj[x]){
             if(!vis[i])
                 q.push(i), vis[i]=1;
        }
void init(int N){
    for(int i=0;i<N;i++){</pre>
        if(!adj[i].empty()) adj[i].clear();
}
```

### 6.2 Eulerian Path and Circuit

```
// O(M)
            歐拉迴路
                            歐拉路徑
//
// 無向圖/所有點的度數為偶數/度數為奇數的點數量不超過2
// 有向圖/所有點入度等於出度/全部點的入度出度一樣
//或剛好一個點出度-1=入度 另一點入度-1=出度,其他點入度
   等於出度
vector<int> path;
void dfs(int x){
   while(!edge[x].empty()){
      int u = edge[x].back();
       edge[x].pop_back();
      dfs(u);
   path.push_back(x);
int main(){
   build_graph();
   dfs(st); // 如果剛好一個點出度-1=入度 則為起點
   reverse(path.begin(),path.end());
}
```

### 6.3 Topological Sort

```
// O(N+M)
for(int i=0;i<m;i++){</pre>
                      //點 u 連到點 v
    cin >> u >> v;
    adj[u].push_back(v);
    ++deg[v];
// can solve DAG
void bfs(int n){ // Topological Sort
    queue<int> q;
    for(int i=0;i<n;i++)</pre>
        if(!deg[i]) q.push(i); //入度0先出
    while(!q.empty()){
        int u = q.front(); q.pop();
        for(int i:adj[u]){
             --deg[i];
            deg_sum[i] += deg_sum[u];
            if(deg[i] == 0) q.push(i);
    }
}
```

#### 6.4 LCA

```
// pre O(NLgN)
// query O(LqN)
// 最近共同祖先
// 兩點間距離 / 兩點間最大邊 / 兩點間重合長度
// 時間戳記,判斷祖先關係
int ti = 0; // 當前時間
int tin[MAXN+5],tout[MAXN+5];
int dis[MAXN+5]; // 計算距離深度
int query[MAXN+5][1gN+5]; //點N的2^LgN祖先的最大邊
void dfs(int x,int f,int deep){
    fa[x] = f;
    tin[x] = ti++;
    dis[x] = deep;
    for(auto i:edge[x]){
        if(i.v == f){
            //query[x][0] = i.w;
            continue;
        dfs(i.v, x, deep+i.w);
    tout[x] = ti++;
bool isAncestor(int u, int v){
    return tin[u]<=tin[v] && tout[u] >= tout[v];
}
// LCA
int n,lgn;
int anc[MAXN+5][lgN+5]; //點N的2^LgN祖先
int getLca(int u, int v){
    if(isAncestor(u, v))
                           return u;
    if(isAncestor(v, u))
                           return v;
    for(int i=lgn;i>=0;i--){
        // 判斷 2^lgN, 2^(lgN-1),...2^1, 2^0 倍祖先
        if(!isAncestor(anc[u][i], v))
            u = anc[u][i]; // 則往上移動
    return anc[u][0];
}
// 找出路徑最大邊
int max_cost(int u, int v){
    int max_cost = 0;
    if(u == v) return max_cost;
    for(int i=lgn;i>=0;i--){
        if(!isAncestor(anc[u][i], v)){
           max_cost = max(max_cost,query[u][i]);
            u = anc[u][i]; // 則往上移動
    return max(max cost,query[u][0]);
} // max(max_cost(u,nodeLca), max_cost(v,nodeLca))
// 兩點距離
int dist(int u, int v){
    //depth[X] + depth[Y] - 2 * depth[ancestor]
    return dis[u] + dis[v]-2*dis[find(v)];
// init 建表
for(s=1;s<=n;s++) anc[s][0] = fa[s];</pre>
for(i=1;i<=lgn;i++)</pre>
    for(s=1;s<=n;s++)</pre>
        anc[s][i] = anc[anc[s][i-1]][i-1];
        // query[s][i] = max(query[s][i-1], query[anc[s
            ][i-1]][i-1]);
```

# 6.5 樹上差分

```
dep[rt] = dep[f] + 1;
  for (int i = 1; i <= 20; i++) {
    fa[rt][i] = fa[fa[rt][i - 1]][i - 1];
  for (auto i : edge[rt]) {
    if (i == f)continue;
    dfs(i, rt);
  }
int lca(int a, int b) {
 if (dep[a] < dep[b]) {</pre>
    swap(a, b);
  for (int i = 20; i >= 0; i--) {
    if (dep[fa[a][i]] >= dep[b]) {
      a = fa[a][i]; //上跳
 if (a == b)
    return a;
  for (int i = 20; i >= 0; i--) {
    if (fa[a][i] != fa[b][i]) {
      a = fa[a][i];
      b = fa[b][i];
    }
  return fa[a][0];
void dfssum(int rt,int f) {
  for (auto i : edge[rt]) {
    if (i == f)continue;
    dfssum(i, rt);
    cnt[rt] += cnt[i];
void solve() {
 int u, v, cmnlca;
  for (int i = 0; i < n; i++) {
    cin >> a[i];
  for (int i = 0; i < n - 1; i++) {</pre>
    cin >> u >> v;
    edge[u].push_back(v);
    edge[v].push_back(u);
  dfs(1, 0);
  for (int i = 0; i < n - 1; i++) {
    cmnlca = lca(a[i], a[i + 1]);
    cnt[fa[cmnlca][0]]--; //父節點 -v
    cnt[cmnlca]--;
                         //Lca -v
                       //兩端點 +v
    cnt[a[i]]++;
    cnt[a[i + 1]]++;
  dfssum(1, 0);
  for (int i = 1; i <= n; i++) { //多加的減回去
    cnt[a[i]]--;
  for (int i = 1; i <= n; i++) {
    cout << cnt[i] << "\n";
int main() {
  cin >> n;
  solve();
```

#### 6.6 HLD with Segment tree

```
| #define MXN 10005
| #define cl(x) (x<<1)
| #define cr(x) (x<<1|1)
| #define INF 1e9+5
| int n;
| int sz[MXN], fa[MXN], heavy[MXN], dep[MXN];
| int root[MXN]; //鍊的根節點
```

```
int len[MXN]; //鍊長度
struct Edge {int u, v;};
struct node {int v, w;};
vector<Edge> edge;
vector<node> graph[MXN];
vector<int> tree[MXN]; // 第i個節點為根的線段樹
                      // 第i個節點為根的序列
vector<int> val[MXN];
// 子樹大小
void dfs_sz(int u, int f, int d){
  sz[u] = 1, fa[u] = f, dep[u] = d;
    for(auto v : graph[u]){
        if(v.v != f){
            dfs_sz(v.v, u, d+1);
            sz[u] += sz[v.v];
            if(sz[v.v] > sz[heavy[u]]) heavy[u] = v.v;
                //重兒子
        }
    }
}
// 樹鍊剖分
void dfs_hld(int u,int f){
    for(auto v : graph[u]){
        if(v.v != f){
            if(v.v == heavy[u])root[v.v] = root[u];
            //重兒子的根,重鍊的頭
            else root[v.v] = v.v; //輕兒子的根
           val[root[v.v]].push_back(v.w); //點權
           dfs_hld(v.v, u);
        }
    }
    len[root[u]]++; //鍊長度
// LCA
int getLca(int x,int y){
    while(root[x] != root[y]){
        if(dep[root[x]] > dep[root[y]])
           x = fa[root[x]]; //跳鍊
        else
           y = fa[root[y]];
    return (dep[x] \leftarrow dep[y] ? x : y);
// 線段樹
void build(int ver,int i,int l,int r){
   if(1 == r){
        tree[ver][i] = val[ver][1];
        return;
    int mid=(l+r)>>1;
    build(ver, cl(i), l, mid);
    build(ver, cr(i), mid+1, r);
    tree[ver][i] = max(tree[ver][cl(i)], tree[ver][cr(i
        )]); //最大邊
void update(int ver,int i,int l,int r,int pos,int val){
    if(1 == r){ // 修改 a[pos] 的值為 val
       tree[ver][i] = val; return;
    int mid=(l+r)>>1;
    if(pos <= mid) update(ver,cl(i),l,mid,pos,val);</pre>
                   update(ver,cr(i),mid+1,r,pos,val);
    else
    tree[ver][i] = max(tree[ver][cl(i)], tree[ver][cr(i
}
// i 為當前節點index, l, r當前區間左右界, ql, qr詢問左
     右界
int query(int ver,int i,int l,int r,int ql,int qr){
    if(ql <= 1 && r <= qr){
       return tree[ver][i];
    int mid=(l+r)>>1, ret=-INF;
    if(ql<=mid) ret =</pre>
        max(ret, query(ver,cl(i),l,mid,ql,qr));
    if(qr) mid) ret =
        max(ret, query(ver,cr(i),mid+1,r,ql,qr));
    return ret;
```

```
void init(){
    edge.clear(); edge.resize(n-1);
    for(int i=1;i<=n;i++){</pre>
        graph[i].clear();
        tree[i].clear();
        val[i].clear();
        heavy[i]=len[i]=0;
signed main(){
    int i,t,a,b,w,ti;
    string op;
    cin >> n;
    init();
    for(i=0; i+1<n; i++){</pre>
        cin >> a >> b >> w;
        graph[a].push_back(node{b,w});
        graph[b].push_back(node{a,w});
        edge[i] = Edge{a,b};
    val[1].push_back(-INF);
    root[1] = 1;
    dfs_sz(1, 1, 0);
    dfs_hld(1, 1);
    // build tree
    for(i=1;i<=n;i++){ // 第i個節點為根的線段樹
        if(root[i] == i){
            tree[i].resize(len[i]*4,0);
            build(i, 1, 0, len[i]-1);
        }
    }
    // query
    while(cin >> op){
        if(op == "DONE") break;
        else if(op == "CHANGE"){
            cin >> i >> ti; i--;
            if(dep[edge[i].u] < dep[edge[i].v])</pre>
                swap(edge[i].u,edge[i].v);
            i = edge[i].u;
            update(root[i], 1, 0, len[root[i]]-1,
            dep[i]-dep[root[i]], ti);
        else if(op == "QUERY"){
            cin >> a >> b;
            int ans = -INF;
            while(root[a] != root[b]){ //不同鍊
                if(dep[root[a]] < dep[root[b]])</pre>
                     swap(a, b);
                // 深鍊的最大邊
                ans = max(ans, query(root[a], 1, 0, len
                     [root[a]]-1, 0, dep[a]-dep[root[a
                     11));
                a = fa[root[a]]; //跳鍊
            if(a != b){ //不同節點
                int mn =
                    min(dep[a],dep[b])-dep[root[a]]+1;
                int mx =
                    max(dep[a],dep[b])-dep[root[a]];
                //所在節點區間 mn,mx
                ans = max(ans, query(root[a], 1, 0, len
                     [root[a]]-1, mn, mx));
            cout << ans <<"\n";
}
   }
       }
```

### 6.7 DSU on Tree

```
void add(int v, int p, int x){
   cnt[ col[v] ] += x;
   // now you can insert test
   for(auto u: g[v])
      if(u != p && !big[u])
      add(u, v, x);
```

```
void dfs(int v, int p, bool keep){
    int mx = -1, bigChild = -1;
    for(auto u : g[v])
       if(u != p \&\& sz[u] > mx)
          mx = sz[u], bigChild = u;
    for(auto u : g[v])
        if(u != p && u != bigChild)
        // run a dfs on small childs and clear them
            dfs(u, v, 0);
    if(bigChild != -1)
    // bigChild marked as big and not cleared from cnt
         dfs(bigChild, v, 1), big[bigChild] = 1;
    add(v, p, 1);
    ans[v] = sum;
    //now cnt[c] is the number of vertices in subtree
         of vertex v that has color c. You can answer
         the queries easily.
    if(bigChild != -1)
        big[bigChild] = 0;
    if(keep == 0){
        add(v, p, -1);
         // now you can init to 0
}
```

# 6.8 Scc 強連通分量 (轉 DAG)

```
#define MXN 200005
#define FZ(x) memset(x, 0, sizeof(x))
// 用於 DAG 2-SAT
struct Scc{ // 強連通分量 0-base
  int n, nScc, vst[MXN], bln[MXN]; // 最後每個點所屬的
      連通分量存在bln陣列
  vector<int> E[MXN], rE[MXN], vec;
  void init(int _n){ //先初始化點的數量
    n = n;
    for (int i=0; i<MXN; i++)</pre>
      E[i].clear(), rE[i].clear();
  void addEdge(int u, int v){ // 加有向邊
    E[u].PB(v); rE[v].PB(u);
  void DFS(int u){
    vst[u]=1;
    for (auto v : E[u]) if (!vst[v]) DFS(v);
    vec.PB(u);
  void rDFS(int u){
    vst[u] = 1; bln[u] = nScc;
    for (auto v : rE[u]) if (!vst[v]) rDFS(v);
  void solve(){ // 跑 kosaraju
    nScc = 0;
    vec.clear();
    FZ(vst);
    for (int i=0; i<n; i++)</pre>
      if (!vst[i]) DFS(i);
    reverse(vec.begin(),vec.end());
    FZ(vst);
    for (auto v : vec)
      if (!vst[v]){
        rDFS(v); nScc++;
}scc;
```

#### 6.9 2-SAT

```
// build
scc.init(2*m);
```

```
for(i=0;i<m;i++){ //所有點比較
  for(j=i+1;j<m;j++){</pre>
    if(graph[i].x == graph[j].x){ //同一行
      // m中的2個控制
      // (x \text{ or } y) x -> !y, y -> !x -> !x = x + m
       if(abs(graph[i].y - graph[j].y)<=2*r){ // light</pre>
           // x or y
           scc.addEdge(i,j+m);
           scc.addEdge(j,i+m);
      }
    }
    else if(graph[i].y == graph[j].y){ //同一列
      if(abs(graph[i].x - graph[j].x)<=2*r){ // light</pre>
           // !x or !y
           scc.addEdge(i+m,j);
           scc.addEdge(j+m,i);
       }
    }
  }
// solve
scc.solve();
for(i=0;i<m;i++){</pre>
  if(scc.bln[i] == scc.bln[i+m]){
       cout << "NO\n";</pre>
       return 0;
  }
cout << "YES\n";</pre>
```

# 6.10 BccVertex 點雙連通分量 (關節點、橋)

```
#define MXN 100005
#define FZ(x) memset(x, 0, sizeof(x))
#define REP(i, b) for(int i = 0; i < (b); ++i)
// 找關節點(分量重疊) 與 橋(只有兩點)
struct BccVertex { // 點雙連通分量(無向) 0-base
 int n,nScc,step,dfn[MXN],low[MXN];
  vector<int> E[MXN],sccv[MXN];
  int top,stk[MXN];
 void init(int _n) { //初始化點的數量
   n = _n; nScc = step = 0;
    for (int i=0; i<n; i++) E[i].clear();</pre>
  void addEdge(int u, int v) // 加無向邊
  { E[u].PB(v); E[v].PB(u); }
  void DFS(int u, int f) {
   dfn[u] = low[u] = step++;
    stk[top++] = u;
    for (auto v:E[u]) {
      if (v == f) continue;
      if (dfn[v] == -1) {
        DFS(v,u);
        low[u] = min(low[u], low[v]);
        if (low[v] >= dfn[u]) {
          int z;
          sccv[nScc].clear();
          do {
           z = stk[--top];
            sccv[nScc].PB(z);
          } while (z != v);
         sccv[nScc++].PB(u);
      }else
        low[u] = min(low[u],dfn[v]);
  } }
 vector<vector<int>>> solve() { // 跑 Tarjan
    vector<vector<int>> res;
    for (int i=0; i<n; i++)</pre>
      dfn[i] = low[i] = -1;
    for (int i=0; i<n; i++)</pre>
      if (dfn[i] == -1) {
        top = 0;
        DFS(i,i);
```

```
REP(i,nScc) res.PB(sccv[i]);
  return res;
}
}graph;
```

#### 7 DP

### 7.1 LCS 與 DP 回溯

```
n=str.size(); m=str2.size();
vector<vector<int>> dp(n+1, vector<int>(m+1,0));
// LCS
for(i=1;i<=n;i++){</pre>
    for(j=1;j<=m;j++){</pre>
        if(str[i-1]==str2[j-1]){
            dp[i][j] = dp[i-1][j-1] + 1;
        }else{
            dp[i][j] = max(dp[i][j-1],dp[i-1][j]);
    }
// DP回溯
string ans="";
i=n; j=m;
while(i>0 && j>0){
    if(str[i-1]==str2[j-1]){
        ans = str[i-1] + ans;
        i--; j--;
    else if(dp[i][j-1]>dp[i-1][j]) j--;
    else i--;
cout << dp[n][m] <<"\n";</pre>
cout << ans <<"\n";
```

#### 7.2 LIS

```
| /*LIS 最長遞增子序列 */
int LIS(int n, int s[], int length[]) {
    for (int i = 0; i < n; i++) length[i] = 1;
    for (int j = 0; j < n; j++)
        for (int i = j + 1; i < n; i++)
        if (s[j] < s[i]) length[i] =
            max(length[i], length[j] + 1);
    return *max_element(length, length+n);
```

#### 7.3 背包問題

```
|// n: 第0種到第n種物品要放進背包內。
// w: 背包耐重限制。
// c(n, w): 只有第0種到第n種物品
// 耐重限制為w,此時的背包問題答案。
// weight[n]: 第n種物品的重量。
// cost[n]: 第n種物品的價值。
// number[n]: 第n種物品的數量。
// 0/1背包滾動
// 每種物品只會放進背包零個或一個。
const int N = 500, W = 2000000; //N個物品,耐重W
int cost[N], weight[N];
int c[W + 1];
void knapsack(int n, int w){
  c[0] = 0;
  for (int i = 0; i < n; ++i)</pre>
    for (int j = w; j - weight[i] >= 0; --j)
     c[j] = max(c[j], c[j - weight[i]] + cost[i]);
  cout << c[w];
// Subset Sum Problem
```

```
// Partition Problem
// Bin Packing Problem
// 無限背包
// 物品有許多種類,每一種物品都無限量供應的背包問題。
void knapsack(int n, int w){
   memset(c, 0, sizeof(c));
   for (int i=0; i<n; ++i)</pre>
       for (int j = weight[i]; j <= w; ++j)</pre>
           c[j] = max(c[j], c[j - weight[i]] + cost[i]
   cout << "最高的價值為" << c[w];
}
// 有限背包
// 物品有許多種類,每一種物品都是限量供應的背包問題。
int cost[N], weight[N], number[N];
// number[n]: 第n種物品的數量。
void knapsack(int n, int w){
   for (int i = 0; i < n; ++i)</pre>
       int num = min(number[i], w / weight[i]);
       for (int k = 1; num > 0; k *= 2)
           if (k > num) k = num;
           num -= k;
           for (int j = w; j >= weight[i] * k; --j)
               c[j] = max(c[j], c[j - weight[i] * k] +
                    cost[i] * k);
   }
   cout << "最高的價值為" << c[w];
```

# 7.4 找零問題

```
// n:用第0種到第n種錢幣來湊得價位。
// m: 欲湊得的價位值。
// c(n, m):用第0種到第n種錢幣湊得價位m的湊法數目。
// price[n]: 第n種錢幣的面額大小。
// 能否湊得某個價位,無限
int price[5] = {5, 2, 6, 11, 17}; //錢幣面額
bool c[1000+1];
// 這些面額湊不湊得到價位 m
void change(int m){
   memset(c, false, sizeof(c));
   c[0] = true;
   // 依序加入各種面額
   for (int i = 0; i < 5; ++i)</pre>
       // 由低價位逐步到高價位
       for (int j = price[i]; j <= m; ++j)</pre>
          c[j] |= c[j-price[i]];
   if (c[m]) cout << "yes";</pre>
   else cout << "no";</pre>
// 湊得某個價位的湊法總共幾種
void change(int m){
   memset(c, 0, sizeof(c));
   c[0] = 1;
   for (int i = 0; i < 5; ++i)</pre>
       for (int j = price[i]; j <= m; ++j)</pre>
          c[j] += c[j-price[i]];
   cout << "湊得價位" << m;
   cout << "湊法總共" << c[m] << "種";
// 湊得某個價位的最少錢幣用量
// c(n, m): 用第0種到第n種錢幣湊得價位m,最少所需要的錢
    幣數量。
void change(int m){
   memset(c, 0x7f, sizeof(c));
```

```
c[0] = 0;
    for (int i = 0; i < 5; ++i)</pre>
       for (int j = price[i]; j <= m; ++j)</pre>
           c[j] = min(c[j], c[j-price[i]] + 1);
    cout << "湊得價位" << m;
    cout << "最少需(只)要" << c[m] << "個錢幣";
// 湊得某個價位的錢幣用量,有哪幾種可能性。
void change(int m){
   memset(c, 0, sizeof(c));
    c[0] = 1;
    for (int i = 0; i < 5; ++i)
       for (int j = price[i]; j <= m; ++j)</pre>
           // 錢幣數量加一,每一種可能性都加一。
           c[j] |= c[j-price[i]] << 1;
    for (int i = 1; i <= 63; ++i)
       if (c[m] & (1 << i))</pre>
           cout << "用" << i << "個錢幣可湊得價位" <<
// 能否湊得某個價位,限量
int price[5] = {5, 2, 6, 11, 17};
int number[5] = {4, 5, 5, 3, 2}; // 各種錢幣的供應數量
bool c[1000+1];
void change(int m){
    memset(c, 0, sizeof(c));
    c[0] = true;
    for (int i = 0; i < 5; ++i)
       // 各種餘數分開處理
       for (int k = 0; k < price[i]; ++k){</pre>
           int left = number[i]; // 補充彈藥
           // 由低價位到高價位
           for (int j = k; j <= m; j += price[i])</pre>
               // 先前的面額已能湊得,當前面額可以省著
               if (c[j])
                   left = number[i]; // 補充彈藥
               // 過去都無法湊得,一定要用目前面額硬
               else if (left > 0){
                  left--; // 用掉一個錢幣
                  c[j] = true;
    if (c[m]) cout << "yes";</pre>
   else cout << "no";</pre>
// 買東西找回最少硬幣。
int price[5] = {50, 20, 10, 4, 2}; // 面額由大到小排列
void cashier(int n){ // n 是總共要找的錢。
    int c = 0;
    for (int i=0; i<5; ++i)</pre>
       while (n >= price[i]){
           n -= price[i]; // 找了 price[i] 元
           c++;
    if(n != 0)cout << "找不出來";
    else cout << "找了" << c << "個錢幣";
}
```

#### 8 DP on tree

#### 8.1 全點對距離 Tree Distance

```
| int dp[MAXN]={0};
| void dfs_sz(int x,int f){
| sz[x] = 1, fa[x] = f;
| for(int i:edge[x]){
| if(i == f) continue;
| dfs1(i, x); // 先計算完子節點的答案再算自己的
| sz[x]+=sz[i];
| dp[x]+=(dp[i]+sz[i]);
```

### 8.2 最大獨立集 Independent set

# 8.3 最小點覆蓋 Vertex Cover

# 8.4 最小支配集 Dominating Set

# 9 String

### 9.1 Trie

```
set<int> cnt_idx, sz_idx;
   trie():cnt(0),sz(0){
       memset(nxt,0,sizeof(nxt));
trie *root = new trie();
void insert(string& s, int idx){
   trie *now = root; // 每次從根結點出發
   for(auto i:s){
       now->sz++; now->sz_idx.emplace(idx); //被誰經過
       if(now->nxt[i-'a'] == NULL){
           now->nxt[i-'a'] = new trie();
       now = now->nxt[i-'a']; //走到下一個字母
   now->cnt++; now->cnt_idx.emplace(idx);//以此點結尾
   now->sz++; now->sz_idx.emplace(idx); //被誰經過
//query
int query_prefix(string& s){ //查詢有多少前綴為 s
                      // 每次從根結點出發
   trie *now = root;
   for(auto i:s){
       if(now->nxt[i-'a'] == NULL){
           return 0;
       now = now -> nxt[i-'a'];
   return now->sz;
int query_count(string& s){ //查詢字串 s 出現次數
                      // 每次從根結點出發
   trie *now = root;
   for(auto i:s){
       if(now->nxt[i-'a'] == NULL){
           return 0;
       now = now->nxt[i-'a'];
   return now->cnt;
//str有沒有在[l,r]的前綴中
bool query_ArrPrefix(string& s,int 1,int r){
                      // 每次從根結點出發
   trie *now = root;
   for(auto i:s){
       if(now->nxt[i-'a'] != NULL && now->nxt[i - 'a'
           ]->sz > 0 ){ //存在
           now = now - > nxt[i - 'a'];
       }else return false; //不存在,無解
   // 這個s的節點, [l,r]有沒有經過
   auto L = now->sz_idx.lower_bound(1);
   if(1<=*L && *L<=r)return true;</pre>
   else return false;}
//[l,r]有沒有存在於str的前綴中
bool query_StrPrefix(string& s,int l,int r){
   trie *now = root;
                       // 每次從根結點出發
   for(auto i:s){
       if(now->nxt[i-'a'] != NULL && now->nxt[i - 'a'
           ]->sz > 0 ){ //存在
           now = now->nxt[i-'a'];
       }else return false; //不存在,無解
   // [l,r]存在於str的前綴中,代表有字串以str為結尾
       auto L = now->cnt_idx.lower_bound(1);
       if(1<=*L && *L<=r)return true;</pre>
   } return false;
```

#### 9.2 01Trie

```
|// insert O(lgx)
|// query O(lgx)
|// 處理XOR問題
|// struct
| struct trie{
| trie *nxt[2]; // 差別
```

```
//紀錄有多少個數字以此節點結尾
    int cnt;
               //有多少數字的前綴包括此節點
    int sz;
    trie():cnt(0),sz(0){
       memset(nxt,0,sizeof(nxt));
};
//創建新的字典樹
trie *root = new trie();
void insert(int x){
    trie *now = root; // 每次從根結點出發
    for(int i=30;i>=0;i--)
       now->sz++;
       if(now->nxt[x>>i&1] == NULL){
           now->nxt[x>>i&1] = new trie();
       now = now->nxt[x>>i&1]; //走到下一個字母
   }
   now->cnt++:
   now->sz++;
// in this set, the maximum value of bitwise XOR x
int query(int x){
    trie *now = root;
    int ans=0;
    for(int i=30;i>=0;i--){// 不等於為1(0xr1=1,1xr0=1)
       if (now->nxt[!(x>>i&1)] != NULL && now->nxt[!(x
           >>i&1)]-> sz > 0){ //下一個存在
           ans += 1<<i;
           now = now->nxt[!(x>>i&1)];
       else now = now->nxt[x>>i&1];
    }
    return ans;
}
```

#### 9.3 Hash

```
// build O(n)
// query 0(1)
// double hash
// P = 53,97,193,49157,805306457,1610612741,1e9+9,1e9+7
const 11 P1 = 75577;
const 11 P2 = 12721;
                        // 多一個質數 p2
const 11 MOD = 998244353;
pair<ll,ll> Hash[MXN];//Hash[i] 為字串 [0,i] 的hash值
void build(const string& s){
    pair<11,11> val = make_pair(0,0);
    for(int i=0; i<s.size(); i++){</pre>
        val.first = (val.first * P1 + s[i]) % MOD;
        val.second = (val.second * P2 + s[i]) % MOD;
        Hash[i] = val;
    }
// query:
// H[l,r] = Hr - H(l-1) * p^{(r-l+1)} %MOD + MOD )%MOD
```

#### **10 STL**

#### 10.1 常用 tool

```
|// freopen()
| freopen("intel.in","r",stdin); // 讀 file.in 檔
| freopen("intel.out","w",stdout); // 寫入 file.out 檔
|//二進制"1"的個數
    __builtin_popcount(n) -> int
    __builtin_popcountl(n) -> long int
    __builtin_popcountl(n) -> long long
|/|陣列處理
| sort(arr,arr+n);
| reverse(arr,arr+n); //value
```

```
| min_element(arr, arr+n) - arr; //index | *lower_bound(arr, arr+4, c) << '\n'; //第一個大於等於c | *upper_bound(arr, arr+4, c) << '\n'; //第一個大於c | fill(arr, arr+3, 123); | //輸出 | cout << fixed << setprecision(10); | cout << setw(n) << setfill(c); | //迭代器 | T.begin() | T.rend() | T.rend() //逆序迭代器 | T.rend() //逆序迭代器 | T.find() //可用於set,map的earse()。
```

# 10.2 Erase 操作

```
|// erase
| iterator erase (iterator position);//刪除指定元素
| iterator erase (iterator first, iterator last);//刪除指定範圍
|// 刪除指定數值
| auto it = find(p.begin(), p.end(), val);
| p.erase(it);
| // 刪除所有指定數值
| p.erase(remove(p.begin(), p.end(), val), p.end());
```

## 10.3 StringStream

```
| stringstream ss;
| getline(cin, str);
| ss.str("");
| ss.clear();
| ss << oct << s; //以8進制讀入流中
| ss << hex << s; //以16進制讀入流中
| ss >> n; //10進制int型輸出
| ss >> s; //x進制str型輸出
```

#### 10.4 List

```
 push_back()
    pop_back()
    push_front()
    pop_front()
    back()
    front()
    insert(index, obj)
    erase()
//遍歷
for (auto iter = _list.begin(); iter != _list.end();
        iter++)
        cout << *iter << "\n";</pre>
```

#### 10.5 Set

```
unordered_set<int> u_set;
multiset<int> st;
• insert(val)
• erase(1, r) //t與r皆為iterator
• erase(it)
• erase(val)
• count(val) //元素是否存在
it = st.find(val);
it = st.lower_bound(val);
it = st.upper_bound(val);
//遍歷
int mints[] = { 75,23,65,42,13,75,65 };
set<int> myset(myints, myints + 7);
```

12

```
for (auto it = myset.begin(); it != myset.end(); it++)
    cout << ' ' << *it;</pre>
```

#### 10.6 Map

```
// 無序map O(1)
unordered_map<type,type> u_map;
//find
auto iter = mymap.find("a");
if (iter != mapStudent.end())
    cout << "Find, the value is" << iter->second <<
        endl;
else
    cout << "Do not Find" << endl;
//erase
auto iter = mymap.find("a");
mymap.erase(iter);
//map遍歷
for (auto it = mymap.begin(); it != mymap.end(); it++)
cout << it->first << ", " << it->second << endl</pre>
```

## 10.7 Priority Queue

```
|//預設由大排到小
| priority_queue<T> pq
| priority_queue<int, vector<int>, less<int> > pq;
| //改成由小排到大
| priority_queue<T, vector<T>, greater<T> > pq;
| //自行定義 cmp 排序
| priority_queue<T, vector<T>, cmp> pq;
| struct cmp {
| bool operator()(node a, node b) {
| //priority_queue@先判定為!cmp
| //,所以「由大排到小」需「反向」定義
| //實現「最小值優先」
| return a.x < b.x;
| }
| };
```

#### 10.8 Bitset

# 11 Other

```
|/*unroll-loops*/
|#pragma GCC optimize("00")//不優化(預設)
|#pragma GCC optimize("01")//優化一點
|#pragma GCC optimize("02")//優化更多
|#pragma GCC optimize("03")//02優化再加上inline函式優化
|#pragma GCC optimize("unroll-loops")
|/*常數宣告*/
|#define MXN 1'000'005
|#define EPS 1e-6
|#define INF 0x3f3f3f3f
|#define PI acos(-1)
```

# 12.1 MinCostMaxFlow

Flow

```
struct MinCostMaxFlow{
typedef int Tcost;
  static const int MAXV = 20010;
  static const int INFf = 1000000;
  static const Tcost INFc = 1e9;
  struct Edge{
    int v, cap;
    Tcost w;
    int rev;
    Edge(){}
    Edge(int t2, int t3, Tcost t4, int t5)
    : v(t2), cap(t3), w(t4), rev(t5) {}
  };
  int V, s, t;
  vector<Edge> g[MAXV];
  void init(int n, int _s, int _t){
    V = n; s = _s; t = _t;
    for(int i = 0; i <= V; i++) g[i].clear();</pre>
  void addEdge(int a, int b, int cap, Tcost w){
    g[a].push_back(Edge(b, cap, w, (int)g[b].size()));
    g[b].push_back(Edge(a, 0, -w, (int)g[a].size()-1));
  Tcost d[MAXV];
  int id[MAXV], mom[MAXV];
  bool inqu[MAXV];
  queue<int> q;
  pair<int, Tcost> solve(){
    int mxf = 0; Tcost mnc = 0;
    while(1){
      fill(d, d+1+V, INFc);
      fill(inqu, inqu+1+V, 0);
      fill(mom, mom+1+V, -1);
      mom[s] = s;
      d[s] = 0;
      q.push(s); inqu[s] = 1;
      while(q.size()){
        int u = q.front(); q.pop();
        inqu[u] = 0;
        for(int i = 0; i < (int) g[u].size(); i++){</pre>
          Edge &e = g[u][i];
          int v = e.v;
          if(e.cap > 0 && d[v] > d[u]+e.w){
            d[v] = d[u] + e.w;
            mom[v] = u;
            id[v] = i;
            if(!inqu[v]) q.push(v), inqu[v] = 1;
      } } }
      if(mom[t] == -1) break ;
      int df = INFf;
      for(int u = t; u != s; u = mom[u])
        df = min(df, g[mom[u]][id[u]].cap);
      for(int u = t; u != s; u = mom[u]){
        Edge &e = g[mom[u]][id[u]];
        e.cap
        g[e.v][e.rev].cap += df;
      }
      mxf += df;
      mnc += df*d[t];
    return {mxf,mnc};
} }flow;
                                  Hong~Long~Long~Long~
                          )
```

(00)/---/\_

AC | NO BUG /==

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AC

#### 12.2 Dinic

```
struct Dinic{
 struct Edge{ int v,f,re; };
  int n,s,t,level[MXN];
 vector<Edge> E[MXN];
  void init(int _n, int _s, int _t){
   n = _n; s = _s; t = _t;
   for (int i=0; i<n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f){
   E[u].PB({v,f,SZ(E[v])});
   E[v].PB({u,0,SZ(E[u])-1});
  bool BFS(){
   for (int i=0; i<n; i++) level[i] = -1;</pre>
    queue<int> que;
    que.push(s);
    level[s] = 0:
    while (!que.empty()){
     int u = que.front(); que.pop();
      for (auto it : E[u]){
       if (it.f > 0 && level[it.v] == -1){
         level[it.v] = level[u]+1;
          que.push(it.v);
   } } }
    return level[t] != -1;
  int DFS(int u, int nf){
   if (u == t) return nf;
    int res = 0;
    for (auto &it : E[u]){
     if (it.f > 0 && level[it.v] == level[u]+1){
       int tf = DFS(it.v, min(nf,it.f));
       res += tf; nf -= tf; it.f -= tf;
       E[it.v][it.re].f += tf;
       if (nf == 0) return res;
   } }
    if (!res) level[u] = -1;
    return res;
  int flow(int res=0){
   while ( BFS() )
     res += DFS(s,2147483647);
    return res;
} }flow;
   \mathsf{n} \underline{\hspace{0.1cm}} \mathsf{n}
 (。・ω・。)つ-ο・*。
               ' • ° +.
              ° ° + *´")
   レーJ
          (...' (...' * AC
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 (。・ω・。)つ-ο・*。
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  (...´ (...' * AC
...´ (...' *) ...*")
```

#### 12.3 KM

```
struct KM{ // max weight, for min negate the weights
  int n, mx[MXN], my[MXN], pa[MXN];
  11 g[MXN][MXN], 1x[MXN], 1y[MXN], sy[MXN];
  bool vx[MXN], vy[MXN];
  void init(int _n) { // 1-based, N個節點
    n = n:
    for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);</pre>
  }
  void addEdge(int x, int y, ll w) {g[x][y] = w;} //左
       邊的集合節點x連邊右邊集合節點y權重為w
  void augment(int y) {
    for(int x, z; y; y = z)
      x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
  void bfs(int st) {
    for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
    queue<int> q; q.push(st);
    for(;;) {
      while(q.size()) {
        int x=q.front(); q.pop(); vx[x]=1;
        for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
          11 t = 1x[x]+1y[y]-g[x][y];
           if(t==0){
             pa[y]=x;
             if(!my[y]){augment(y);return;}
             vy[y]=1, q.push(my[y]);
          }else if(sy[y]>t) pa[y]=x,sy[y]=t;
      } }
      11 cut = INF;
      for(int y=1; y<=n; ++y)</pre>
        if(!vy[y]&&cut>sy[y]) cut=sy[y];
       for(int j=1; j<=n; ++j){</pre>
        if(vx[j]) lx[j] -= cut;
        if(vy[j]) ly[j] += cut;
         else sy[j] -= cut;
       for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){</pre>
        if(!my[y]){augment(y);return;}
         vy[y]=1, q.push(my[y]);
  } } }
  11 solve(){ // 回傳值為完美匹配下的最大總權重
    fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
    fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
    for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)</pre>
      lx[x] = max(lx[x], g[x][y]);
     for(int x=1; x<=n; ++x) bfs(x);</pre>
    11 \text{ ans} = 0:
    for(int y=1; y<=n; ++y) ans += g[my[y]][y];</pre>
    return ans;
} }graph;
```