Scala

Lightweight Modular Staging

Steven Both, Toby Rufinus, Jaspreet Singh, Daniël Stekelenburg Now that we are familiar with Scala...



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lets look at an awesome library implemented in Scala.
Lightweight Modular Staging (LMS)

LMS is..

'A library-based multi-stage programming approach that uses types to distinguish between binding time.'

Outline

- A gentle introduction to LMS
- Generative Programming
- ▶ How to write a multi-staged program
- Language virtualization
- Intermediate representation
- ► How do we program in LMS?



A gentle introduction to LMS

Power function in Scala

```
def power(b: Double, x: Int): Double =
  if (x == 0) 1.0 else b * power(b, x - 1)
```

Power function in Scala LMS

```
trait PowerA { this: Arith =>
  def power(b: Rep[Double], x: Int): Rep[Double] =
   if (x == 0) 1.0 else b * power(b, x - 1)
}
```

A gentle introduction to LMS

What did we just see?

- ▶ T versus Rep[T]
- ▶ def Versus trait

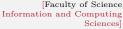
Productivity vs Performance

Software performance depends more on programmer productivity

- ▶ Processor clock speed doesn't double every 18 months
- High-level programming is hard to translate to efficient code
- Shift towards big data workloads

Result: Hand-optimized programs (BAD IDEA)

- abandoning all best practices and benefits of high-level programming
- programs become hard to read, maintain, verify...
- this attracts bugs and security vulnerabilities





Solution: Generative Programming

Write a program generator

- Produces the code of a program as output
- Reorganization of a programs' execution into stages, also called multi-stage programming

How do we write MSP-programs?

- A single-stage program is developed, implemented and tested
- ► Ensure the the program can be used in a staged manner. Otherwise "refactor"
- Introduce staging annotations

Another stage!

From

- Compilation-based program execution:
- ► Compile-time, run-time.

To

- Generated program execution:
- ► Generation-time, Compile-time, run-time.

Generative Programming

What is generative programming?

Generative Programming

Commonalities of LMS and generative programming

Generative Programming

Differences of LMS and generative programming



Drawbacks of LMS

- ► The Scala LMS library only implements staged operations for a subset of Scala.
- Debugging is painful LMS can give obscure errors, run into an infinite loop, or generate wrong code.
- Documentation is lacking.

Intermediate representation



How to LMS:

- Staging
- Generating code
- Data types



Power function in Scala:

```
def power(b: Double, p: Int): Double = {
  if (p == 0)
    1.0
  else
    b * power(b, p - 1)
}
```

Staged power function in LMS:

```
def power(b: Rep[Double], p: Int): Rep[Double] = {
  if (p == 0)
    1.0
  else
    b * power(b, p - 1)
}
```

Staged power function in LMS:

```
def power(b: Rep[Double], p: Int): Rep[Double] = {
  if (p == 0)
    1.0
  else
    b * power(b, p - 1)
}
power(b, 3)
```

```
power(b, 3)

Generated code:

def apply(x3: Double): Double = {
  val x4 = x3 * x3
  val x5 = x3 * x4
  x5
}
```

Simple factorial function

```
def fac(n: Rep[Int]): Rep[Int] = {
  if (n == 0) 1
  else n * fac(n - 1)
}
```

Simple factorial function

```
def fac(n: Rep[Int]): Rep[Int] = {
  if (n == 0) 1
   else n * fac(n - 1)
}
fac(n)
```

Simple factorial function

```
def fac(n: Rep[Int]): Rep[Int] = {
  if (n == 0) 1
  else n * fac(n - 1)
fac(n)
[error] (run-main) java.lang.StackOverflowError
[error] (compile:run) Nonzero exit code: 1
```

power(b, 3) vs. fac(n)

```
power(b, 3) vs. fac(n)
Make use of a lambda function:
def fac: Rep[Int => Int] = doLambda { n =>
  if (n == 0) 1
  else n * fac(n-1)
Now we can try it again:
fac(n)
```



Generated code

```
def apply(x12:Int): Int = {
 var x1 = {x2: (Int) => }
   val x3 = x2 == 0
   val x8 = if (x3) { 1 }
    else {
     val x4 = x2 - 1
     val x5 = x1(x4) // recursion
     val x6 = x2 * x5
     x6 }
   x8: Int }
 val x13 = x1(x12)
                           // recursion
 x13 }
```

```
power(b, 3)

Generated code:

def apply(x3: Double): Double = {
  val x4 = x3 * x3
  val x5 = x3 * x4
  x5
}
```

Optimized power function

```
def powerOpt(b: Rep[Double], p: Int): Rep[Double] = {
  def loop(x: Rep[Double], ac: Rep[Double],
                            y: Int): Rep[Double] =
    if (y == 0)
      ac
    else if (y \% 2 == 0)
      loop(x * x, ac, y / 2)
    else
      loop(x, ac * x, y - 1)
  loop(b, 1.0, p)
```

```
powerOpt(b, 3)

Generated code:

def apply(x3: Double): Double = {
  val x4 = x3 * x3
  val x5 = x3 * x4
  x5
}
```

```
powerOpt(b, 3)

Generated code:

def apply(x3: Double): Double = {
  val x4 = x3 * x3
  val x5 = x3 * x4
  x5
}
```

LMS can generate the same code from different staged codes.



But not per se.



But not per se.

For example:

powerOpt(b, 6)

```
But not per se.

For example:

powerOpt(b, 6)

Generated code of optimized version
```

```
def apply(x4:Double): Double = {
  val x5 = x4 * x4
  val x6 = x5 * x5
  val x7 = x5 * x6
  x7
```

Trivial regular expression

```
checkRegex(".", "Hello world")
```

The char '.' is a wildcard in Scala.

Partial regex code

```
def matchStar(...): Rep[Boolean] = { ... }
def matchBegin(...): Rep[Boolean] = { ... }
def matchEnd(...): Rep[Boolean] = { ... }

def matchChar(c: Char, t: Rep[Char]): Rep[Boolean] = { c == '.' || c == t }
```

Partial regex code

```
def matchStar(...): Rep[Boolean] = { ... }
def matchBegin(...): Rep[Boolean] = { ... }
def matchEnd(...): Rep[Boolean] = { ... }

def matchChar(c: Char, t: Rep[Char]): Rep[Boolean] = { c == '.' || c == t }

Now do:
checkRegex(".", s)
```

```
val x47 = while ({ val x28 = x27
  val x34 = if (x28) \{ false \}
  else \{ val x30 = x26 \}
    val x32 = x30 < x31
    x32 } x34}) {
  val x36 = x26 += 1
  val x37 = x26
  val x38 = x37 < x31
  val x42 = if (x38) {
    val x39 = x25(x37)
    val x40 = '.' == x39 // matchChar(...)
    val x41 = true | x40 // c == '.' | c == t
    x41 }
  else { false }
```



Generated code is not meant to be human-readable

The previous examples only considered:

- ► Rep[Int]
- ► Rep[Double]
- Rep[Char]
- ► Rep[Boolean]

The previous examples only considered:

- ► Rep[Int]
- Rep[Double]
- Rep[Char]
- Rep[Boolean]

But what to do for datatypes that are of your own making?

- Scala LMS only has pre-defined operations for standard library types.
- ➤ You add Rep[T] types to your own functions to stage them, allowing them to work on staged values.

Unstaged:

```
case class Vec3(x: Double, y: Double, z: Double) {
  def +(that: Vec3): Vec3 =
    Vec3(this.x + that.x,
        this.y + that.y, this.z + that.z)
}
```

```
Unstaged:
case class Vec3(x: Double, y: Double, z: Double) {
  def +(that: Vec3): Vec3 =
    Vec3(this.x + that.x,
      this.y + that.y, this.z + that.z)
}
Staged:
case class Vec3(x: Rep[Double], y: Rep[Double],
    z: Rep[Double]) {
  def +(that: Vec3): Vec3 =
    Vec3(this.x + that.x,
      this.y + that.y, this.z + that.z)
```

What about this function?

Unstaged:

```
case class Vec3(x: Double, y: Double, z: Double) {
  def length: Double =
    sqrt(x * x + y * y + z * z)
}
```

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```
What about this function?
Unstaged:
case class Vec3(x: Double, y: Double, z: Double) {
  def length: Double =
    sqrt(x * x + y * y + z * z)
Staged:
case class Vec3(x: Rep[Double], y: Rep[Double],
   z: Rep[Double]) {
  def length: Double =
    sqrt(x * x + y * y + z * z)
```

A problem: undefined operations on staged types

- Scala LMS defines some fundamental operations on staged types, such as integer/floating point arithmetic.
- ► If your staged function only uses those operations (like the + operator in the Vec3 example) you're fine.



How to Stage Your Algorithm

- 1. Add Rep [T] type annotations
- Define an interface for new operations on staged types
- 3. Implement the interface in terms of IR nodes
- (optional) Define optimizations, rewriting certain patterns of IR nodes
- Extend code generator so new IR nodes can be turned into code

You can create a wave-like function by summing up a number of sine waves. The Fourier transform decomposes the waveform back into its sine components. The Fast Fourier Transform (FFT) is a fast numerical algorithm that can do this.

```
def fft(xs: Array[Complex]): Array[Complex]
= xs match {
  case (x :: Nil) => xs
  case =>
    val N = xs.length // assume it's a power of two
    val (even0, odd0) = splitEvenOdd(xs)
    val (even1, odd1) = (fft(even0), fft(odd0))
    val (even2, odd2) = (even1 zip odd1 zipWithIndex)
      case ((x, y), k) \Rightarrow
        val z = omega(k, N) * y
        (x + z, x - z)
    }.unzip;
    even2 ::: odd2
```

```
case class Complex(re: Double, im: Double) {
  def +(that: Complex) = Complex(this.re + that.re,
     this.im + that.im)
  def *(that: Complex) = ...
}

def omega(k: Int, N: Int): Complex = {
  val kth = -2.0 * k * Math.Pi / N
     Complex(cos(kth), sin(kth))
}
```

Note the operations we perform on doubles: arithmetic (addition, multiplication, ...) and trigonometry (sin, cos)



Step 1: Add Rep[T] type annotations

Unstaged:

```
case class Complex(re: Double, im: Double) {
  def +(that: Complex) = Complex(this.re +
        that.re, this.im + that.im)
  def *(that: Complex) = ...
}
def omega ...
def fftt ...
```

Step 1: Add Rep[T] type annotations

```
trait FFT { this: Arith with Trig =>
  case class Complex(re: Rep[Double], im:
    Rep[Double]) {
    def + Complex(this.re + that.re,
        this.im + that.im)
    def * ...
}
def omega ...
def fft ...
}
```

Step 1: Add Rep[T] type annotations

The staged version of complex numbers consists of a pair of Rep[Double] We need to be able to do arithmetic and trigonometric operations on staged doubles. Fortunately, the Scala LMS library happens to define operations for Rep[Double]

Step 1: Add Rep[T] type annotations

```
trait FFT { this: Arith with Trig =>
  case class Complex(re: Rep[Double], im:
    Rep[Double]) {
    def + Complex(this.re + that.re,
        this.im + that.im)
    def * ...
}
def omega ...
def fft ...
}
```

Step 1: Add Rep[T] type annotations

```
trait FFT { this: Arith with Trig =>
  case class Complex(re: Rep[Double],
    im: Rep[Double]) {
    def + Complex(this.re + that.re,
      this.im + that.im)
    def * ...
  def omega ...
  def fft ...
```

- How to define our own staged operations that work on Rep[Double]?
- ▶ We define these operations in traits and mix them in.
- ► The this: Arith with Trig part means: whenever the trait FFT is instantiated, we need to mix in traits that provide arithmetic and trigonometric operations too.

Step 2: Define an interface for new operations on staged types

Scala LMS provides a Base trait; the Rep[T] type is defined there.

```
trait Arith extends Base {
  def infix_+(x: Rep[Double], y:
      Rep[Double]): Rep[Double]
  def infix_*(x: Rep[Double], y:
      Rep[Double]): Rep[Double]
  ...
}
```

trait Trig extends Base {

```
def cos(x: Rep[Double]): Rep[Double]

def sin(x: Rep[Double]): Rep[Double]

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```

These traits contain only abstract members; they are interfaces. We need to create subclasses with concrete implementations.

Step 3: Implement the interface in terms of IR nodes

- Scala LMS uses a node-based intermediate representation for staged expressions.
- ► The BaseExp class from the LMS framework defines some related types.
- Exp[T] is a simple IR expression (a constant or symbol).
- ▶ Def [T] is a composite operation; these operations will be converted to simple expressions.
- ▶ BaseExp also defines Rep[T] = Exp[T] so staged expressions will be converted to IR expressions.



Step 3: Implement the interface in terms of IR nodes

```
trait ArithExp extends Arith with BaseExp {
  // These case classes are IR nodes
  case class Plus(x: Exp[Double], y: Exp[Double])
    extends Def[Double]
  case class Times(x: Exp[Double], y: Exp[Double])
    extends Def[Double]
  // The abstract functions defined in trait Arith ar
  def infix_+(x: Exp[Double], y: Exp[Double]) =
    Plus(x, y)
  def infix_*(x: Exp[Double], y: Exp[Double]) =
    Times(x, y)
```

Step 3: Implement the interface in terms of IR nodes

```
sin(x + 2 * y) + sin(0)
Plus(Sin(Plus(Sym(x),
Times(Const(2),
Sym(y)))),
Sin(Const(0)))
```



Step 4: Define optimizations, rewriting certain patterns of IR nodes

- Scala LMS already contains a number of generic optimizations, such as dead code elimination and reusing identical expressions.
- We can define our own optimizations, both generic optimizations and domain-specific ones.
- ► These are again defined in traits, so you can combine them in a modular way.

```
trait ArithExpOpt extends ArithExp {
  override def infix_*(x: Exp[Double],
    y: Exp[Double]) =
    (x, y) match {
      // Multiplying two constants? We can calculate
      case (Const(x), Const(y)) \Rightarrow Const(x * y)
      // 1 * x = x, and vice versa
      case (x. Const(1)) \Rightarrow x
      case (Const(1), x) => x
      // Base case: apply the regular base function
      case => super.infix *(x, y)
```

Step 4: Define optimizations, rewriting certain patterns of IR nodes

```
trait TrigExpOptFFT extends TrigExpOpt {
  override def cos(x: Exp[Double]) = x match {
    case Const(x)
    if { val z = x / math.Pi / 0.5;
        z != 0 && z == z.toInt } =>
        Const(0.0)
    case _ => super.cos(x)
  }
}
```

Step 5: Extend code generator so new IR nodes can be turned into code

Finally, IR nodes have to be converted to actual code. The LMS framework provides a ScalaGenBase class that we can use. We only have to define what to do when the generator encounters one of the new nodes we added.

Step 5: Extend code generator so new IR nodes can be turned into code

- ▶ It's also possible to generate code for other languages if you define your own generator.
- Scala LMS even comes with a CGenBase trait that can generate C code from your staged Scala functions.

Step 5: Extend code generator so new IR nodes can be turned into code

The CompileScala trait defines a compile function that lets you load the generated code immediately into the running program. Essentially, compile "unstages" your staged function (Rep[A] => Rep[B]) into a regular function (A => B). This function can then be called:

```
val fftCompiled = compile(fft)
// Now we can call fftCompiled with regular values
// Just like any other function in the program
fftCompiled(Array(1.0,0.0, 1.0,0.0, 2.0,0.0, 2.0,0.0))
```

> Array(6.0,0.0, -1.0,1.0, 0.0,0.0, -1.0,-1.0)



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Conclusion

TODO

