

Homework Assignment 1

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1 Exercise 1 - [Edit Distance/Levenshtein Distance]

1.1 Question 1

Proposition 1.1.1 (Complexity and Correction). *If we denote C_f the complexity of f , this algorithm has time complexity $\mathcal{O}\left(C_f \left(\frac{n}{t}\right)^2\right)$. This algorithm is correct.*

Démonstration. — Moreover, it is clear this algorithm is correct as it only just applies the dynamic programming algorithm for the Levenshtein distance by steps.

— This algorithm complexity comes from the fact it has two while loops for which the commands are executed at most n/t times. The commands in both *while* loops are executed in $\mathcal{O}(C_f)$. The *for* loops inside the *left while* loop are equivalent to loops for i between left and left + $t - 1$ and thus are disjoint. The sum of their complexity over the *left while* loop is then n . The number of operations inside the *up while* loop is then in $\mathcal{O}\left(C_f \frac{n}{t}\right)$ and thus the total complexity is, as announced, in $\mathcal{O}\left(C_f \left(\frac{n}{t}\right)^2\right)$ ■

1.2 Question 2

By the recurrence formula : $\mathbf{D}[i][j] = \max \begin{cases} \mathbf{D}[i-1][j] + 1 \\ \mathbf{D}[i][j-1] + 1 \\ \mathbf{D}[i-1][j-1] + 1 \text{ if } S[i] \neq T[j] \text{ else } 0 \end{cases}$ we see

that $\mathbf{D}[i][j]$ is at most 1 plus one of its left neighbour, upper neighbour or upper left corner neighbour.

Algorithm 1 Question 1 - Levenshtein Distance with f

Input S, T, f, t \triangleright Two Strings, the function f computing the values and the step t

$\mathbf{D} = \text{zeros}(n+1, n+1)$ $\triangleright \text{len}(S) = \text{len}(T) = n$

for $i \leftarrow 0$ **to** $n+1$ **do**
 $\mathbf{D}[i][0] \leftarrow i$
end for

for $j \leftarrow 0$ **to** $n+1$ **do**
 $\mathbf{D}[0][j] \leftarrow j$
end for

$\text{up}, \text{left} \leftarrow 0, 0$
while $\text{up} < n$ **do**
 $\text{left} \leftarrow 0$
 while $\text{left} < n$ **do**
 $\text{down} \leftarrow \min(n - \text{up}, t)$
 $\text{right} \leftarrow \min(n - \text{left}, t)$

 $b \leftarrow \mathbf{D}[\text{up}][\text{left}]$
 $a \leftarrow \mathbf{D}[\text{up} + 1 \rightarrow \text{up} + 1 + \text{down}][\text{left}]$
 $c \leftarrow \mathbf{D}[\text{up}][\text{left} + 1 \rightarrow \text{left} + 1 + \text{right}]$

 $f(a, b, c, d, e)$ \triangleright We can suppose here that f modifies only the last line and column of F in \mathbf{D} with side-effect.
 $\text{left} \leftarrow \text{left} + \text{right}$
 for $i \leftarrow 1$ **to** $\text{down} - 1$ **do**
 $\mathbf{D}[\text{up} + i][\text{left}] \leftarrow \min \begin{cases} \mathbf{D}[\text{up} + i][\text{left} - 1] + 1 \\ \mathbf{D}[\text{up} + i - 1][\text{left}] + 1 \\ \mathbf{D}[\text{up} + i - 1][\text{left} - 1] + \mathbb{1}_{\{S[\text{up}+i]=T[\text{left}]\}} \end{cases}$
 \triangleright We update the first Column of the block we consider.
 end for
 end while
 $\text{up} \leftarrow \text{up} + d$
 for $i \leftarrow 1$ **to** n **do**
 $\mathbf{D}[\text{up}][i] \leftarrow \min \begin{cases} \mathbf{D}[\text{up}][i - 1] + 1 \\ \mathbf{D}[\text{up} - 1][i] + 1 \\ \mathbf{D}[\text{up} - 1][i - 1] + \mathbb{1}_{\{S[\text{up}+i]=T[i]\}} \end{cases}$
 \triangleright We update the first line of the blocks we will consider.
 end for
end while
return $\mathbf{D}[n][n]$

1.3 Question 3

By recurrence formula, if we subtract from all the values in A, B, C a certain integer k , then we get for the new matrix F , the one we would have gotten with A, B, C with k subtracted to all values. Thus, if the values in A, A', B, B' and C, C' all differ from a common integer, the resulting values after applying f will differ from this same value. Thus, $F' = F + (A' - A)$ and $A' - A$ is a matrix with all values equal.

1.4 Question 4

We will show here that we can pre-compute all $t \times t$ matrices in $\mathcal{O}(3^{2t}\sigma^{2t}t^2)$. First, as $\mathbf{D}[i][j]$ here is the minimal number of elementary operations to go from string $S[0 : i]$ to string $T[0 : j]$. We can thus interpret the submatrix of \mathbf{D} between (i, j) and $(i + t, j + t)$ as the dynamic programming matrix for minimum number of operations to go from string $S[i : i + t + 1]$ to string $T[j : j + t + 1]$ to which we added a first line and a first column.

We thus see that those matrices can be fully determined by their first line, first column and by two words.

As the values in \mathbf{D} are bounded by 0 and $2n$ (we can always remove all letters in S and add all letters in T), and by question 2., as values along a line or a column differ by an integer in $-1, 0, 1$, the number of first lines and columns is $\mathcal{O}(n3^{2t})$. However, from question 3., if we allow negative values for f (which we have no reason not to do), we can always subtract from the first line and column the value in the top left corner, and re-add it in $\mathcal{O}(t^2)$ after pre-processing.

Moreover, there are σ^t words of length t over the alphabet so the number of submatrices is $\mathcal{O}(3^{2t}\sigma^{2t})$.

We can then use the recurrence equation to derive the values on the submatrix in $\mathcal{O}(t^2)$.

Finally, we can pre-compute all $t \times t$ submatrices in $\mathcal{O}(3^{2t}\sigma^{2t}t^2)$.

Then, to access the values of the submatrix from $(up, left)$ to $(up+t, left+t)$, we need to identify the preprocessed corresponding matrix and thus we need to go through $S[up : up+t]$, $T[left : left+t]$, the first column and row of this submatrix to which we subtracted the upper left value, which all are done in $\mathcal{O}(t)$.

1.5 Question 5

From Question 4, we have got an algorithm that allows us to compute the result in $\mathcal{O}(3^{2t}\sigma^{2t}t^2 + \left(\frac{n^2}{t}\right))$.

Indeed, after pre computing, we only need to check in $\mathcal{O}(1)$ the $\mathcal{O}\left(\left(\frac{n}{t}\right)^2\right)$ $t \times t$ submatrices in \mathbf{D} .

Thus, if we take $t = \log_{(3\sigma)}(\sqrt{n})$, we have complexity in $\mathcal{O}\left((3\sigma)^{\log_{3\sigma}(n)} \log_{3\sigma}(\sqrt{n}) + \left(\frac{n^2}{\log_{3\sigma}(\sqrt{n})}\right)\right)$.

As $\log_{3\sigma}(\sqrt{n}) = \Theta(\log(n))$ and $(3\sigma)^{\log_{3\sigma}(n)} \log_{3\sigma}(\sqrt{n}) = \mathcal{O}(n \log(n)^2) = o\left(\frac{n^2}{\log(n)}\right)$, this algorithm has complexity in $\mathcal{O}\left(\frac{n^2}{\log n}\right)$

2 Exercice 2

In this exercise, we will denote by $\lg(n)$ the log in base 2 of n

2.1 Question 1

We will try to precompute arrays containing part of the answer by dividing the bitvector B into multiple blocks, of smaller and smaller size. We will use two intermediate arrays of size $\mathcal{O}\left(\frac{n}{\lg(n)}\right) = o(n)$ bits and a final array (meaning lowest level array) of size $\mathcal{O}(n)$.

We build a first array A_0 which will contain blocks of size $s_0 = \lg^2(n)$ (we will justify this value later). The blocks of this array contain the rank of the first position in B that is compacted in the block, i.e., $A[i] = \text{rank}_1(i \times s_0)$ where rank designates the rank operation in B . As each entry in A takes at most $\lg(n)$ bits to store, A takes $\mathcal{O}\left(\frac{n}{\lg^2(n)} \lg(n)\right) = \mathcal{O}\left(\frac{n}{\lg(n)}\right) = o(n)$ bits to store, thus justifying the value chosen for s_0 .

Then we build a second array A_1 which will store the ranks of smaller blocks of size $s_1 = \lg(n)/2$ (as for s_0 , we will justify this value later). Here, the items in the array only need at most $\lg(\lg^2(n))$ bits to be stored. This table thus takes $\mathcal{O}\left(\frac{n}{s_1} \lg(\lg^2(n))\right) = \mathcal{O}\left(4n \frac{\lg \lg(n)}{\lg(n)}\right) = o(n)$ bits to store, the equality coming from the properties of \lg , thus explaining the general form of s_1 as $\lg * c$ but not yet why $c = 1/2$

To finally answer the query, we need to maintain a third array A_2 which will contain correspondances between every possible $\text{rank}_1(i)$ query on a bitvector of length s_1 with $i < s_1$ and their answers. There are 2^{s_1} such bitvectors, and storing all answers on each can be done in $\mathcal{O}(s_1 \lg(s_1))$. A_2 thus needs $\mathcal{O}(2^{s_1} \lg(s_1) s_1) = \mathcal{O}(\sqrt{n} \lg(n) \lg(\lg(n))) \leq n$, as $\lg(n)/2 = \lg(\sqrt{n})$ and from the properties of \lg , thus explaining the coefficient of \lg in s_1 .

Using these arrays we can compute $\text{rank}_1(i)$ in $\mathcal{O}(1)$ time by searching in blocks i/s_0 , $(i \bmod s_0)/s_1$ and $(i \bmod s_0) \bmod s_1$ in A_0, A_1 and A_2 , and as $\text{rank}_0(i) = i - \text{rank}_1(i)$, we obtain the wanted time complexity. The total number of bits used by these arrays is at most $n + o(n)$ so we have the wanted space complexity, and this solution is one.

2.2 Question 2

The solution proposed before can lead to a $\mathcal{O}(\lg(n))$ time complexity, so it is not sufficient to answer the **select** queries. We will thus refine the latter to get a new structure, by using three levels of intermediate arrays using $o(n)$ bits and two final arrays using $\mathcal{O}(n)$ bits.

We define our first array C_0 which records the positions of the $\lg^2(n)$ -th (again, this value will be explained later) 1 bits. Storing a value in this array costs at most $\lg(n)$ bits and there are at most $\frac{n}{\lg^2(n)}$ values stored in the array. So it only uses $\frac{n}{\lg(n)}$ bits.

Then, we will repeat the operation. Let s_1 be the size of a block in our first array. We want to use at most $\frac{s_1}{\lg(n)} \leq \lg(n)$ bits on this block in our second array C_1 . We can calculate s_1 on the fly in $\mathcal{O}(1)$ during calculation, and using this, we can deduce the location in B of this block. However, there are $\lg^2(n)$ answers in that range, and it so requires $\lg(n)^3$ bits to store. Then, if we have $s_1 \geq \lg(n)^4$, we would have sufficient space to store the values in $o(n)$ bits, but, in the other case, we need to redivide the block in the same way.

We denote by s_2 the size of the considered block in C_1 . We would need $\lg(s_2) \lg(s_1) \lg(n)$ bits to store all the answers. We obtain a similar inequality on s_2 as on s_1 before, by the same methods. If we have $s_2 \geq \lg(s_2) \lg(s_1) \lg(n)$ then we have sufficient space not to need to go further in the

construction and keep $o(n)$ space. Else, we have :

$$\begin{aligned}s_2 &< \lg(s_2)\lg(s_1)\lg(n) \\ s_1 &< (\lg(n))^4\end{aligned}$$

So we have : $\lg(s_1) < 4(\lg(\lg(n)))$. Then, by growing of \lg and as $s_2 < s_1$ we get : $\lg(s_1) < 4\lg(\lg(n))$ and thus :

$$s_2 < 16\lg(n)\lg^2(\lg(n))$$

Using the same idea as in question 1., we answer **select** queries using final arrays, storing the values we need. Indeed, for each of the 2^{s_3} bit pattern of length $s_3 = \frac{\lg(n)}{2}$, we can record the position of the i -th 1 bit in the pattern in an array C_3 , and store the number of 1 in the pattern in an array C_4 . Then to compute the value of $\text{select}_1(i)$ we can scan the range using C_4 to know which subrange contains the answer and use C_3 to get the answer, all in time $\mathcal{O}(1)$.

Finally, $\text{select}_1(i)$ can be computed by finding the block in C_0 in which i is. From here, we compute s_1 and then, based on the case disjunction detailed earlier, we either get the correct answer from C_1 or compute s_2 and start over, using the final arrays. Thus, this data structure supports **select** queries in $\mathcal{O}(1)$ time. For storing the auxiliary directories C_0, C_1 and C_2 , we need at most $\frac{3n}{\lg(n)}$ bits (the simplicity of this expression justifies the values chosen for s_1 and s_2), and for the final arrays we use $\sqrt{n}(\lg(n) + \frac{1}{2}\lg(n)\lg(n))$, by the same calculation as in question 1, thus justifying the choice for the size of the patterns in C_4 . The extra storage is in $\mathcal{O}(n)$, and this data structure is compatible with our requirements.