第十三讲: 准备阶段 Rewrite 和 Prepare 设计与实现

知春路遇上八里桥

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1 前情提要

② Rewrite 过程

③ Prepare 过程

4 代码分析









前情提要

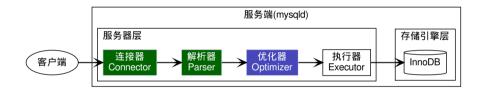








执行流程











本节内容

• 连接器

▶ ☑ 连接管理器 Connection Manager

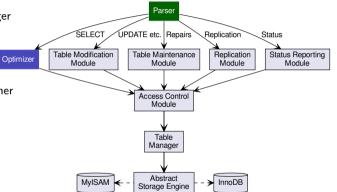
- ▶ ☑ 线程管理器 Thread Manager
- ▶ ☑ 用户模块 User Module

● 解析器

- ▶ ☑ 网络模块 Net Module
- ▶ ☑ 派发模块 Commander Dispatcher
- ▶ ☑ 词法分析 Lexical Analysis
- ▶ ☑ 语法分析 Syntax Analysis

• 优化器

▶ □ 准备模块 Prepare Module











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Rewrite 过程







Rewrite 功能介绍

- MySQL 支持查询 Rewrite ¹ 插件,可以在执行之前检查并修改接收的语句
 - ▶ 该插件允许在服务器处理之前重写到达服务器的 SQL 语句
 - ▶ 该插件接收一个语句字符串,并可能返回一个不同的字符串
- ② 这个设计可以满足 DBA 在不修改业务代码前提下在 Query 中添加 hits
- ③ Rewrite 的安装和卸载通过自带的脚本完成
 - ▶ 脚本位于 \$MYSQL_HOME/share 目录下
 - ▶ 安装脚本 install_rewriter.sql
 - ▶ 卸载脚本 uninstall_rewriter.sql
- Rewrite 的安装和卸载过程

```
mysql> source /opt/mysql/share/install_rewriter.sql Query OK, 1 row affected (0.01 sec) ... mysql> source /opt/mysql/share/uninstall_rewriter.sql Query OK, 1 row affected (0.05 sec) ...
```









添加 Rewrite 规则

• 可以通过 Rewrite 将发送的查询语句添加一个 LIMIT 1 限制只返回一行









规则生效示例

```
mvsal> select * from querv rewrite.rewrite rules\G
id: 1
        pattern: SELECT * FROM employees WHERE emp no > ?
 pattern_database: employees
     replacement: SELECT * FROM employees WHERE emp no > ? LIMIT 1
        enabled: YES
        message: NULL
   pattern digest: 193b6eb2e5a0f8cc9351110a142f3bfa196f09b17cfb34bea295bd0bc77ec92a
normalized pattern: select `*` from `employees`.`employees` where (`emp no` > ?)
1 row in set (0.00 sec)
mysal> select count(*) from employees where emp no > 11100:
+----+
 count(*) |
+-----
   298924 I
+----
1 row in set (2.24 sec)
mysql> select * from employees where emp_no > 11100;
+----+
emp no | birth date | first name | last name | gender | hire date
+----+
  11101 | 1961-10-31 | Manibrata | Horswill | M
+-----
1 row in set. 1 warning (0.01 sec)
mysql>
```









Rewrite 实现示例

• Rewrite 的实现包含两种模式: pre-parse 和 post-parse 2



Rewrite 的调用点在 dispatch_sql_command() 函数中, 精简代码如下
void dispatch_sql_command(THD *thd, Parser_state *parser_state) {
 thd->reset_rewritten_query();

 thd->m_parser_state = parser_state;
 invoke_pre_parse_rewrite_plugins(thd);
 thd->m_parser_state = nullptr;

if (!err) {
 err = parse_sql(thd, parser_state, nullptr);
 if (!err) err = invoke_post_parse_rewrite_plugins(thd, false);

- Rewrite 实现见文件 ☞ sql/sql_query_rewrite.cc
 - ▶ invoke_pre_parse_rewrite_plugins() 前置解析 Rewrite 的实现
 - ▶ invoke_post_parse_rewrite_plugins() 后置解析 Rewrite 的实现

My<mark>SQ</mark>L







重写 Query 断点调试过程

Breakpoint 8 at 0x555558c4959c: THD::set query. (2 locations)

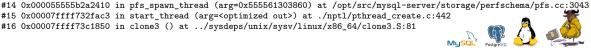
(gdb) b THD::set query

(gdb) c Continuing.

```
Thread 51 "connection" hit Breakpoint 8, THD::set_query (this=0x7fff3807b2e0, query_arg=0x7fff381e7a08 "SELECT * FROM en
                      void set querv(const char *querv arg, size t querv length arg) {
4142
(gdb) bt
#0 THD::set query (this=0x7fff3807b2e0, query arg=0x7fff381e7a08 "SELECT * FROM employees WHERE emp no > 11100 LIMIT 1"
        0x0000555558d8cd94 in alloc query (thd=0x7fff3807b2e0, packet=0x7fff381ae130 "SELECT * FROM employees WHERE emp no >
        0x00005555594b2997 in mysql parser parse (thd=0x7fff3807b2e0, guery=..., is prepared=0 '\000', handle condition=0x0.
        0x00007fffe899e481 in services::parse (thd=0x7fff3807b2e0, query="SELECT * FROM employees WHERE emp no > 11100 LIMIT
        0x00007fffe899a153 in rewrite query notify (thd=0x7fff3807b2e0, event class=MYSQL AUDIT PARSE CLASS, event=0x7fffe01
#5
        0x0000555558ff0428 in plugins dispatch (thd=0x7fff3807b2e0, plugin=0x7fff380a1820, arg=0x7fffe01ed620) at /opt/src/m
        0x0000555558ff0530 in event class dispatch (thd=0x7fff3807b2e0, event class=MYSQL AUDIT PARSE CLASS, event=0x7fffe01
        0x0000555558ff05fc in event class dispatch error (thd=0x7fff3807b2e0, event class=MYSQL AUDIT PARSE CLASS, event named to the class and the control of the c
#8
        0x0000555558fee63d in mysql audit notify (thd=0x7fff3807b2e0, subclass=MYSQL AUDIT PARSE POSTPARSE, subclass name=0x
        0x0000555558de54dd in invoke_post_parse_rewrite_plugins (thd=0x7fff3807b2e0, is_prepared=false) at /opt/src/mysql-se
#10 0x0000555558d95048 in dispatch sql command (thd=0x7fff3807b2e0, parser state=0x7fffe01ed970) at /opt/src/mysql-serve
#11 0x0000555558d8aa69 in dispatch command (thd=0x7fff3807b2e0, com data=0x7fffe01ee2c0, command=COM QUERY) at /opt/src/
```







#13 0x0000555558fd61c5 in handle_connection (arg=0x5555610a30a0) at /opt/src/mysql-server/sql/conn_handler/connection_ha

#12 0x0000555558d888d3 in do command (thd=0x7fff3807b2e0) at /opt/src/mysgl-server/sgl/sgl parse.cc:1440

#15 0x00007fffff732fac3 in start thread (arg=<optimized out>) at ./nptl/pthread create.c:442 #16 0x00007fffff73c1850 in clone3 () at ../sysdeps/unix/sysv/linux/x86 64/clone3.S:81

Prepare 过程









准备阶段功能概述

- 准备阶段用于为整个查询做准备工作,具体划分成以下操作函数
 - ▶ Setup/Fix 设置和修复工作,也包含一部分语义检查
 - △ 准备临时表
 - ② 打开表 检查表是否可以访问
 - ◎ 检查字段
 - 检查非 group 函数
 - ▶ Transformation 变换工作,通过恒等变换来降低后序优化分析代码的复杂度
 - 代数变换
 - ② 逻辑变换
- 准备阶段的大部分代码位于 🔊 sql/sql_resolver.cc 文件中
 - ▶ 历史遗留的变换方法,例如 bool setup_order(...)
 - ▶ Query_block 类是成员函数实现代码









Setup/Fix 部分举例

• Setup 设置函数

- ▶ setup_tables() 设置 Query Block 所有表中的叶子节点的表
- ▶ setup_wild() 将 '*' 设置成表中对应的所有列名
- 1 select * from employees; -- 将 * 通配符替换成 employees 表的列名
- ▶ setup_fields() 检查表中是否存在所有给定的列,并将查找的列填充进去
- ▶ setup_conds() 设置 WHERE 子句中的条件和 JOIN 子句中的条件
- ▶ setup_group() 设置 GROUP BY 子句中的字段列表
- ▶ setup_order() 设置 ORDER BY 子句
- ▶ setup_order_final() 在 Query Block 全部解析完成后,再一次设置 ORDER BY 子句
- select first_name, last_name
- from employees where emp_no <10010
- 3 order by 2 -- 这里的 ORDER BY 必须等 Query Block 全部解析完成才能确定
- ▶ setup_ftfuncs() 在解析 HAVING 子句后设置全文函数 (full-text functions)

• Fix 修复函数

- ▶ m_having_cond->fix_fields() 设置 HAVING 子句
- ▶ resolve_limits() 设置 OFFSET 和 LIMIT 子句
- ▶ resolve_rollup() 汇总 SELECT 和 ORDER BY 子句中的 Item 对象









Transformation 变换举例(壹)

- resolve_subquery() 解析子查询
 - ▶ 将子查询转换成半连接 (semi-join)
 - ▶ 标记子查询执行时使用物化策略 (materialization)
 - ▶ 进行 IN → EXISTS 的变换
 - ▶ 进行 ALL/ANY → MIN/MAX 的重写
 - ▶ 使用具体值替换对应标量上下文 (scalar-context) 的子查询
- simplify_joins() 将 OUTER JOIN 尽可能地转化成 INNER JOIN,看代码注释中的例子

```
SELECT * FROM t1 LEFT JOIN t2 ON t2.a=t1.a WHERE t2.b < 5 -- 注意到 t2.b < 5 拒绝 NULL SELECT * FROM t1 INNER JOIN t2 ON t2.a=t1.a WHERE t2.b < 5 -- 首次变换 SELECT * FROM t1, t2 ON t2.a=t1.a WHERE t2.b < 5 AND t2.a=t1.a -- 最终变换
```

- apply_local_transforms() 应用局部变换,例如:
 - ▶ 这里的 "局部" (local) 指转化对象限制在一个 Query Block
 - ▶ join nest simplification (嵌套 JOIN 的简化)
 - ▶ partition pruning (分区剪枝)
 - ▶ condition pushdown to derived tables (对派生表进行条件下推)
 - . .









Transformation 变换举例(贰)

- remove_redundant_subquery_clause() 删除子查询中的多余的字段
 - ▶ 对于包含 LIMIT 的子句,谓词 IN/ANY/ALL/EXISTS 等需要清理
 - ▶ 如果没有聚合函数和 HAVING 子句,清理 GROUP BY 多余的子句
 - ▶ 对于不包含 LIMIT 的标量查询, ORDER BY 是多余的, 因为标量的列数为 1
- flatten_subqueries() 将 semi-join 子查询转换成 semi-join 嵌套 JOIN 子句

```
SELECT ...
FROM ot, ...
WHERE oe IN (SELECT ie FROM it1 ... itN WHERE subq_where) AND outer_where
-- 这里将子查询转换成半连接
SELECT ...
FROM ot SEMI JOIN (it1 ... itN), ...
WHERE outer where AND subg where AND oe=ie
```

- delete_unused_merged_columns() 合并表时删除没有使用的列
 - ▶ 对于嵌套 JOIN 进行递归调用来处理
 - ▶ 如果发现表合并,自动删除没有使用的列









代码分析

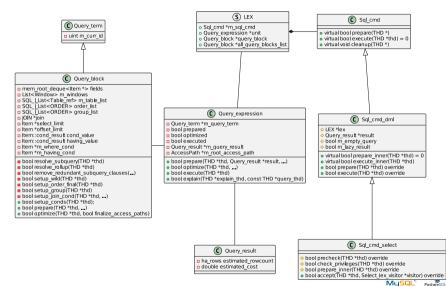








sql_cmd 类关系图







mysgl execute command() 执行命令

● mysql_execute_command()的实现见 ☞ sql/sql_parse.cc , 有超过 2000 行代码 ▶ 见 ☞ sql/sql parse.cc int mysql_execute_command(THD *thd, bool first_level) { return res || thd->is_error(); 5056 5057 ▶ mysgl execute command() 主要功能如下 ● 执行 lex->m_sql_cmd 中的命令 ② 根据 lex->sql_command 来调用不同的处理逻辑 ⑤ 如果是 SQLCOM_SELECT , 根据多态调用 Sql_cmd_dml::execute() 函数 ● Sql_cmd_dml::execute() 执行 SELECT 语句 ▶ 实现见文件 🖙 sql/sql_select.cc bool Sal cmd dml::execute(THD *thd) { DBUG TRACE; 673 return thd->is_error(); 878 879

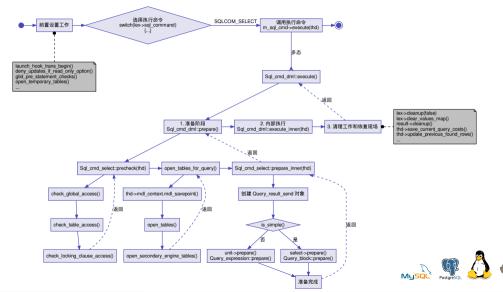




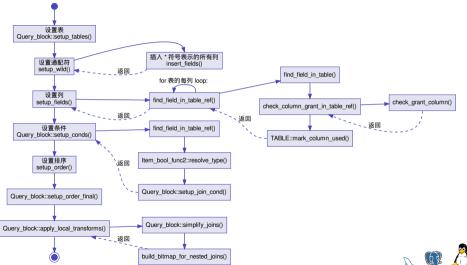




从 mysql_execute_command() 进入准备阶段



查询块的准备过程 - Query_block::prepare() 函数











结束









