

Dependency Analysis

(2)

- ① Dependence: Constraints of executⁿ bcz of each other.
- ② Dependence Analysis: To identify which factors creating ~~obstacles~~ ^{dependency}.
- ③ Dependency Types: Control Dependency (Graph Based) (Task Based)
Data Dependency: (Interactⁿ Based)

Dependency Analysis

- ↳ Two memory access are involved.
- ↳ Data may refer to a same memory locatⁿ.
- ↳ Ensure data is produced / consumed in right order.

Type of Dependency Analysis

① → Flow Dependency

- ↳ Instructⁿ depends on final outcome of previous instructions.

- ↳ also called: Write or write/read dependency.

Add, R₁, R₂

MOV, R₄, R₁

② → Anti Dependency

- ↳ One Instructⁿ depends upon the data that could be destroyed by another instructⁿ.

- ↳ Occurance: One ^{inst.} ~~register~~ reads a register, & subsequent inst. writes value to same locatⁿ.

Condition: Two inst. have Anti-D only if swapping their order would result in true dependency.

$$\begin{bmatrix} A = B + C \\ C = D + E \end{bmatrix}$$

④ O/P dependency

- ↳ Occur: When two ins both write a result.

- ↳ Write write dependency.

- ↳ Exist due to limited no. of arch. registers.

$$\begin{array}{l} A = B + C \\ d = C + D \end{array} \quad \left. \vphantom{\begin{array}{l} A = B + C \\ d = C + D \end{array}} \right\} \text{O/P dependency}$$

GB23: HPC: 3:45 - 4:45 : 15/03/24 : Tuesday

Roll_NO: 18, 19, 20, 21, 23, 24, 25, 26, 27, 30, 31, 32, 33, 34, 35
36, 37, 38, 41, 42, 44, 45, 46, 47.

PF : 1, 2, 3, 13.

Mapping Parallel Algorithm onto Parallel Architecture.

→ Mapping requirement

Analysis of Parallel Algo. + definition of logical configuratⁿ of platforms + Mapping of algo to logic. Haffor.

→ Issues of mapping arises when

↳ No. of process > No. of units processing

↳ Communication structure of algo. differs from interconnect structure of parallel machine : Topological Variation

↳ No. of process in PA > No. of units in Parallel M/c : Cardinality Ratio

Solutⁿ : Contractⁿ : Embedding : Multiplexing

Threads Basics. —

Pthread

↳ creatⁿ / Terminatⁿ

↳ Synchronizatⁿ

↳ controlling Threads & Synchronization Attributes.

Attribute of Objects for Threads.

" ——— for Mutexes

Composite Synchronization Constructs.

OpenMP.

↳ Programming Model.

↳ Reduction Clause

↳ Program : Example

↳ specifying con. Task

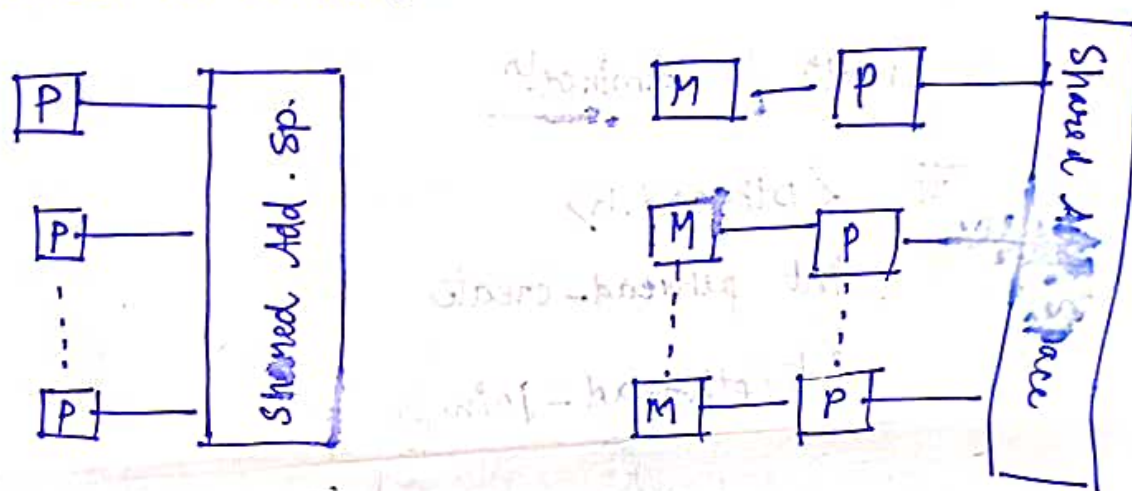
Take that after Thread / Pthreads



Thread Basis

↳ Memory in the logical m/c model of thread is globally accessible to every thread.

↳ Subprocess is thread.



↳ Properties :

↳ Provide Software Portability : through threaded API

↳ Inherent support for latency hiding.

↳ Scheduling & load balancing.

↳ Ease of programming & wide spread use.

The POSIX thread API

↳ Referred to Pthread.

↳ Standard thread APIs.

↳ There can be used for programming with other thread APIs. (NT threads, Solaris thread, Java threads, as well).

↳ Create / Termination

<pthread.h>

int pthread_create

int pthread_join



P-threads (POSIX)

- ↳ P-threads are standard threads API
- ↳ can be used for programming with other thread APIs.
like NT threads, solaris threads, java threads.

↳ include `<pthread.h>`

↳ Header file contains function declarations & defines no. of constant used by these functions

↳ `pthread_create (pthread_t* __restrict __thread, const pthread_attr_t *attr, void *(*start_routine) (void *arg), void * __restrict __arg);`

↳ `pthread_create` : creates a new thread within the process with

attr: is the thread attribute object.

→ `pthread_t` : Datatype to uniquely identify a thread.

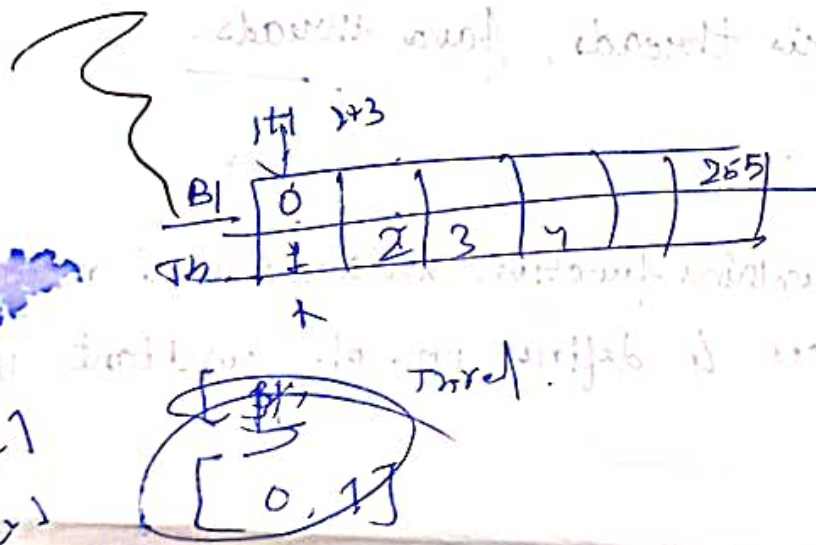
→ `start_routine`: pointer to sub routine that is executed by thread.

→ `Arg`: pointer to void that contains the argument to the function defined in the previous argument.

↳ pthread_join (pthread_t thread, void ** ptr)

: Used to wait for the termination of thread.

thread_t: thread id of the thread for which current thread waits.



1 + 2
A = [1]
B = [1]



Synchronization primitives in Pthread

- ↳ Multiple threads attempt to manipulate same data item
- ↳ Results can be ~~not~~ Incoherent

Example :

Each thread tries to update best_cost as follows

If (my_cost < best_cost)

best_cost = my_cost ;

best_cost = 100;

t1 = 50, t2 = 75.

Depend upon the schedule of the threads,

the best_cost could be 50 or 75.

↳ The value 75 does not correspond to any serializⁿ. of threads.

Controlling Thread & Synchronization Attributes


↳ Pthreads allows programmer to change the default attr using attributes objects.

↳ Attribute object: Data structure that describe entity

Entities (thread, mutex, condition variable)

↳ Enhances: Modularity, Reusability

Attributes objects for Thread.

 pthread_attr_t: Create attributes objects.

- Setdetachstate:

- setguardsize_np

- setstacksize

- setinheritsched.

- setschedpolicy.

- setschedparam.



Attributes Objects for Mutexes.

↳ Mutexes (Mutual Exclusion)

: to save the processes from being in deadlock

: If thread is using resources, the value is set = 0;

If thread is available, then value = 1

Functions

↳ pthread_mutex_init :

↳ for the initialization of attribute object :

Types of Mutexes.

→ pthread_mutex_t Normal.NP.

→ pthread_mutex_t RECURSIVE.NP.

→ pthread_mutex_t ERRORCHECK.NP.

Threads

Directive Based

Standard thread.



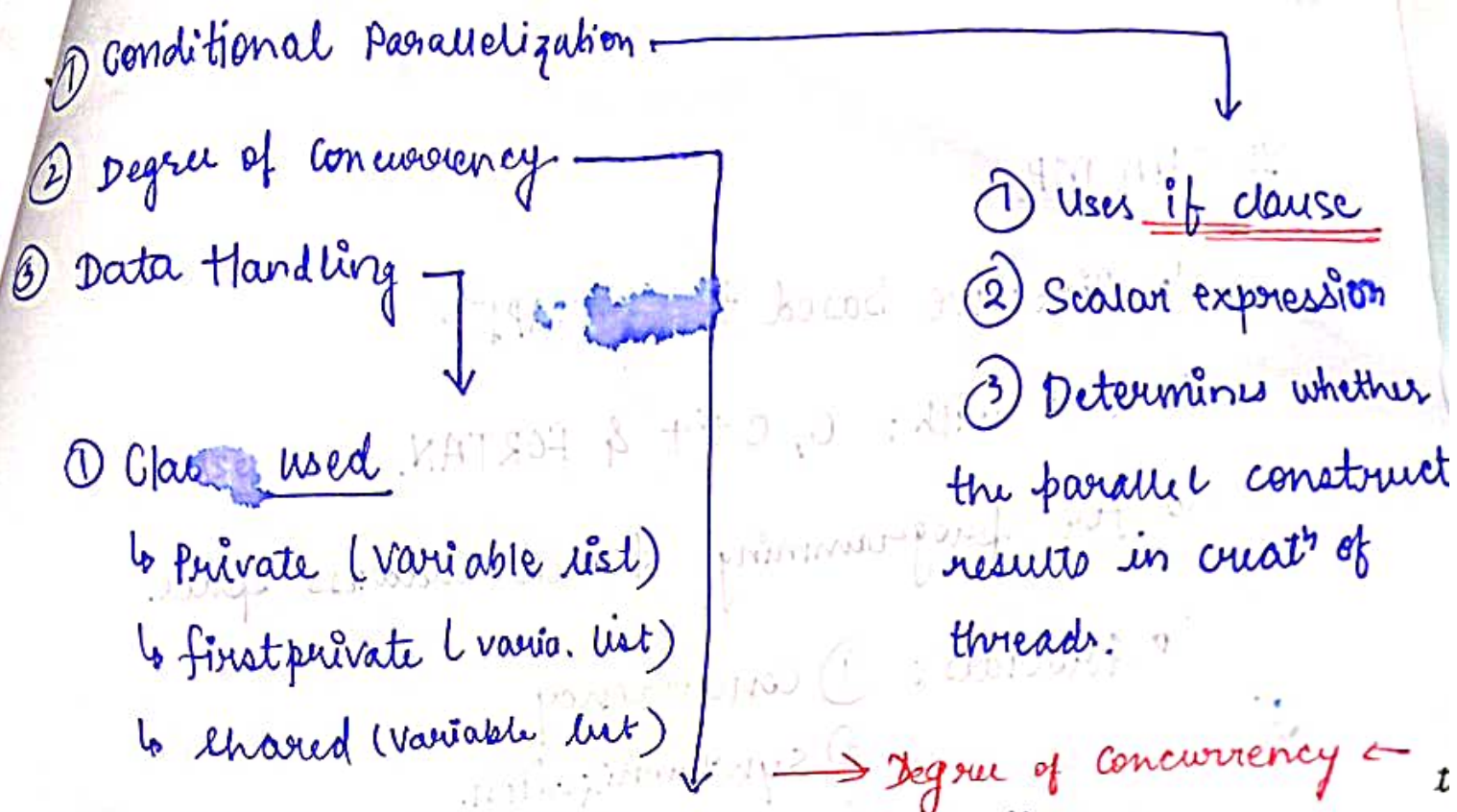
① ROPEN MP.

② for : C, C++,
fortran.

① POSIX - pthreads

② can be used with NT threads,
solaris threads, java threads.

openMP programming model



Private: Indicates variable local to the each thread

firstprivate: Similar to private

Only difference is values of variables are initialized already.

Shared: Indicates that variables are shared across the threads

Sample program

```
#pragma omp parallel if (is-parallel == 1) num-threads(8)
```

(a) Private (a) (b) first

(c) private

(d) shared.

(e) first private.

Open MP

↳ Directive based ~~the~~ **API**.

↳ Used with: C, C++ & FORTRAN.

↳ for programming shared address space.

↳ Provides: ① concurrency

② Synchronization.

③ Data Handling

with the need to use Mutexes explicitly.

↳ Use #pragma compiler directives.

↳ ① #pragma omp directive. [clause]

② #pragma omp parallel. [clause]

① OpenMP program executes serially until they encounter the parallel directives which creates a group of threads.

② The main thread that encounters with parallel directive is set as 0. Within the group and called as master

If $is_parallel = 1$

then

- ① ~~8~~ threads are created.
- ② Each thread gets private copies of variable (a) and (c).
- ③ ~~8~~ keep a single value of (b).
- ④ The value of each copy of (c) is initialized to the value of C before the parallel directive.
- ⑤ The default state of variables is specified by the ~~nonse~~ default (shared) or default (none)

Reduction clause

↳ How multiple local copies of a variable at different threads are combined into a single copy at the master when ~~the~~ threads exit.

↳ The usage of the ~~reduction~~ clause is reduction
(operator : variable list)

↳ These ~~the~~ variables are private to threads.

↳ operator: $+$, $*$, $-$, $\&$, $!$, (or addition)

pragma omp parallel reduction($+$: sum) num-

($+$: sum) num-threads (8)

{

* compute local sum here *

}

Specifying concurrent Task

↳ Used with conjunction with other directives to specify concurrency across iterations & task.

↳ Two ~~direction~~ directives of OPEN MP.

↙
for

↓
section

58, 63, 65, 67, 70.

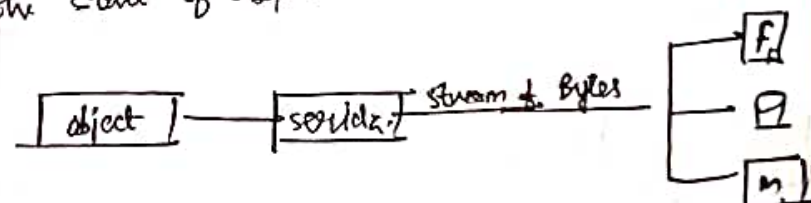
Parallel Processing fundamental Design Issue -

- ① Synchronization
- ② Scheduling
- ③ Job allocation
- ④ ~~job allocation~~
- ⑤ R. job partitioning
- ⑥ Dependency analysis
- ⑦ Mapping parallel algo to Parallel Archi.
- ⑧ Performance Analysis of Parallel Algo.

Synchronization

- ↳ Seq. of work & task
- ↳ Performance factor
- ↳ Requires serializablⁿ

↓
~~serializabl~~
It is a process of converting a data object : code/data represented in a storage into a series of bytes that save the state of object into transmissible form.



↳ Types

- ↳ Barrier
- ↳ Lock / Semaphore
- ↳ Synchronous com. Operatⁿ.
- ↳ Acknowledgment Based
- ↳ To look up the coordinatⁿ b/w the s-process & x-process.

↳ to protect access to global data.
↳ one data can use one lock at a time.

Job scheduling Policies -

FCFS

Look ahead Optimizing Schedulers

- ↳ Backtracking Algo.
- ↳ Branches the variable to evaluate one of its value.

Task gr! → Gang scheduling

- Multiple gangs run for t time
- One stops then another gang executes
- ↳ Concurrent Gang
- ↳ SHARE scheduling
- ↳ Gangs with some resource utility

↳ Used to ensure that if two or more process / threads will be ready to communicate at the same time.

→