

1125. Smallest Sufficient Team II DP II Bit Manipulation

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1125. Smallest Sufficient Team

Hard 1K 20 Companies

In a project, you have a list of required skills `req_skills`, and a list of people. The i^{th} person `people[i]` contains a list of skills that the person has.

Ultimately you want to gather all "req_skills". ↗ have list of skills
→ smallest size

Consider a sufficient team: a set of people such that for every required skill in `req_skills`, there is at least one person in the team who has that skill. We can represent these teams by the index of each person.

- For example, `team = [0, 1, 3]` represents the people with skills `people[0]`, `people[1]`, and `people[3]`.

Return any sufficient team of the smallest possible size, represented by the index of each person. You may return the answer in any order.

It is **guaranteed** an answer exists.

Example 1:

Input: `req_skills = ["java", "nodejs", "reactjs"]`, `people = [["java"], ["nodejs"], ["nodejs", "reactjs"]]`
Output: `[0, 2]` ↗ take Max ↗ rare scenario for skills. ↗ thinking of Greedy ↗ Always prove it or NOT WRONG

Example 2:

Input: `req_skills = ["algorithms", "math", "java", "reactjs", "csharp", "aws"]`,
`people = [[["algorithms", "math", "java"], ["algorithms", "math", "reactjs"],`
`["java", "csharp", "aws"], ["reactjs", "csharp"], ["csharp", "math"], ["aws", "java"]]]`
Output: `[1, 2]`

Constraints:

- `1 <= req_skills.length <= 16`
- `1 <= req_skills[i].length <= 16`
- `req_skills[i]` consists of lowercase English letters.
- All the strings of `req_skills` are **unique**.
- `1 <= people.length <= 60`
- `0 <= people[i].length <= 16`
- `1 <= people[i][j].length <= 16`
- `people[i][j]` consists of lowercase English letters.
- All the strings of `people[i]` are **unique**.
- Every skill in `people[i]` is a skill in `req_skills`.
- It is guaranteed a sufficient team exists.

↑ taking the person with max skills & then compensating other skills with other people.

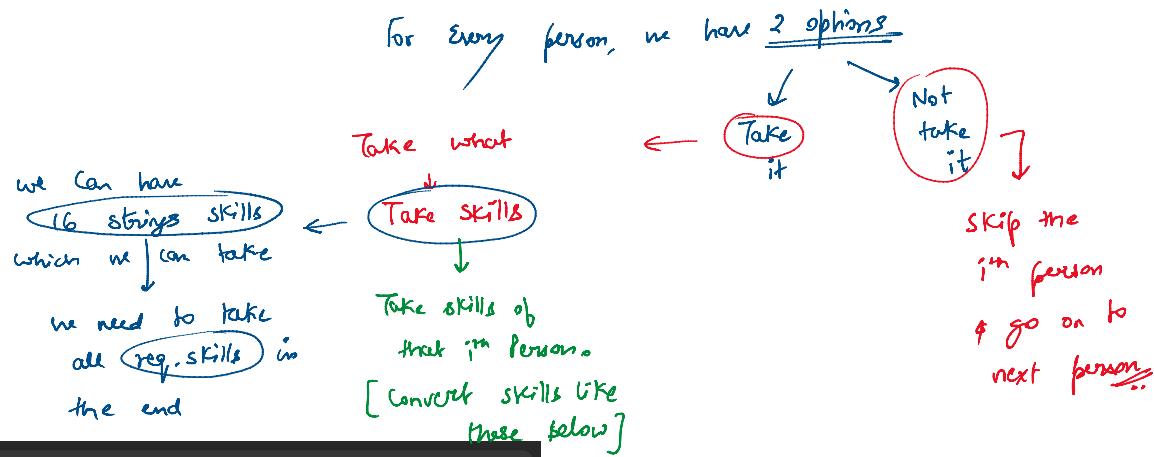
Greedy won't work.
people : [[a, b, c, d], [e], [f]] } → Greedy → 3 people
[a, b, e] } → optimal → 2 people
[c, d, f] }

Thus, try for ALL POSSIBILITIES



For every person, we have 2 options

↙ ↘ Next



Input: req_skills = ["java", "nodejs", "reactjs"], people = [[["java"], ["nodejs"], ["reactjs"]]]
Output: [0,2] ↴ 011 ↴ 100 ↴ 010

req_skills = | | |

↓

Just making a mask

of skills → so, that at every person, we would know if we want that skill or not

How to make Mask of skills of people :

Example 1:

→ Give indices to req_skills

Input: req_skills = ["java", "nodejs", "reactjs"], people = [[["java"], ["nodejs"], ["reactjs"]]]
Output: [0,2] ↴ 0 | 1 | 2
 ↴ index 1 ↴ index 2 ↴ index 3

```

12
13     unordered_map<string, int> skills;
14     for (int i = 0; i < n; ++i)
            skills[req_skills[i]] = i;
    
```

110 [OR operation]

```

16
17     vector<int> people;
18     for (auto &v: people_skills) {
19         int h = 0;
20         for (string &skill: v)
21             h |= 1 << skills[skill];
22     }
23     people.push_back(h);
    
```

left shift
OR
operation

```

25     req_mask = (1 << n) - 1;
26
27     memset(dp, -1, 1 << n + 8);
28     memset(choice, 0, 1 << n + 6);
29
30     solve(people, 0, 0);

```

↑ index ↑ Mask

ultimately requiring every skill → simple Memoization
To keep track if i^{th} person was chosen or not

```

48     int solve(vector<int> &people, int i, int mask) {
49         if (mask == req_mask)
50             return 0;
51         if (i == m)
52             return 10000;
53
54         int &ans = dp[mask][i];
55         if (ans != -1)
56             return ans;
57
58         int ans1 = solve(people, i+1, mask | people[i]) + 1;
59         int ans2 = solve(people, i+1, mask);
60
61         if (ans1 < ans2)
62             choice[mask][i] = 1;
63
64         return ans = min(ans1, ans2);

```

Memoization ↗ if reached required mask return a 0.
if reached end return high value
add skills in mask. → include i^{th} person
don't include i^{th} person
return min of both ways ↗

if including i^{th} person is beneficial → Put your choice as ①

```

32     // build the answer through the choice matrix (see the solve function first)
33     vector<int> ans;
34     for (int i = 0, mask = 0; i < m; ++i) {
35         // did we include the  $i^{th}$  person?
36         if (choice[mask][i]) {
37             ans.push_back(i);
38             mask |= people[i];
39         }
40
41         if (mask == req_mask)
42             break;
43     }
44
45     return ans;
46 }

```

See if you need to include i^{th} person or not
As you reach all req-skills break

CODE

```
int dp[1 << 16][64];
bool choice[1 << 16][64];
class Solution {
public:
    int req_mask, m;

    vector<int> smallestSufficientTeam(vector<string>& req_skills, vector<vector<string>>& people_skills) {
        int n = req_skills.size();
        m = people_skills.size();

        unordered_map<string, int> skills;
        for (int i = 0; i < n; ++i)
            skills[req_skills[i]] = i;

        vector<int> people;
        for (auto &v: people_skills) {
            int h = 0;
            for (string &skill: v)
                h |= 1 << skills[skill];

            people.push_back(h);
        }

        req_mask = (1 << n) - 1;

        memset(dp, -1, 1 << n + 8);
        memset(choice, 0, 1 << n + 6);

        solve(people, 0, 0);

        // build the answer through the choice matrix (see the solve function first)
        vector<int> ans;
        for (int i = 0, mask = 0; i < m; ++i) {
            // did we include the i'th person?
            if (choice[mask][i]) {
                ans.push_back(i);
                mask |= people[i];
            }

            if (mask == req_mask)
                break;
        }

        return ans;
    }

    int solve(vector<int> &people, int i, int mask) {
        if (mask == req_mask)
            return 0;
        if (i == m)
            return 10000;

        int &ans = dp[mask][i];
        if (ans != -1)
            return ans;

        int ans1 = solve(people, i+1, mask | people[i]) + 1;
        int ans2 = solve(people, i+1, mask);

        if (ans1 < ans2)
            choice[mask][i] = 1;

        return ans = min(ans1, ans2);
    }
};
```

Time: $O(\text{people} * 2^{\text{skills}})$

Space: $O(\text{people} * 2^{\text{skills}})$

Bottom Up Approach would also be same :)

As you were going to every person, same here go to every person & see its contribution in 2^{skills} possible options.

$$\text{if } \text{eq_skills} = 5$$

by total option = 2^5
i.e.

$$[0, 2^5 - 1]$$

$$[00000, \dots, 11111]$$

```

1 class Solution {
2 public:
3     vector<int> smallestSufficientTeam(vector<string>& req_skills, vector<vector<string>>& people) {
4         int n = req_skills.size();
5         unordered_map<int, vector<int>> dp; // initialize dp
6         dp.reserve(1 << n);
7         dp[0] = {};
8
9         unordered_map<string, int> skill_index;
10        for (int i = 0; i < req_skills.size(); ++i) // marking index of req_skills as done previously.
11            skill_index[req_skills[i]] = i;
12
13        for (int i = 0; i < people.size(); ++i) {
14            int cur_skill = 0;
15            for (auto& skill: people[i]) // getting skills of every person.
16                cur_skill |= 1 << skill_index[skill];
17
18            for (auto it = dp.begin(); it != dp.end(); ++it) // going on person
19                for (auto it = dp.begin(); it != dp.end(); ++it) // to all possible skill subset
20                    if (dp.find(comb) == dp.end() || dp[comb].size() > 1 + dp[it->first].size()) // new combination
21                        dp[comb] = it->second; // existing combination will have an impact on existing size :)
22                        dp[comb].push_back(i); // new added
23
24            return dp[(1 << n) - 1]; // if all skills done :)
25        }
26    }
27 };
28 
```

time: $O(\text{people} \cdot 2^{\text{skills}})$

space: $O(2^{\text{skills}})$

CODE

```

class Solution {
public:
    vector<int> smallestSufficientTeam(vector<string>& req_skills, vector<vector<string>>& people) {
        int n = req_skills.size();
        unordered_map<int, vector<int>> dp;
        dp.reserve(1 << n);
        dp[0] = {};
        unordered_map<string, int> skill_index;
        for (int i = 0; i < req_skills.size(); ++i)
            skill_index[req_skills[i]] = i;

        for (int i = 0; i < people.size(); ++i) {
            int cur_skill = 0;
            for (auto& skill: people[i])
                cur_skill |= 1 << skill_index[skill];

            for (auto it = dp.begin(); it != dp.end(); ++it) {
                int comb = it->first | cur_skill;
                if (dp.find(comb) == dp.end() || dp[comb].size() > 1 + dp[it->first].size())
                    dp[comb] = it->second;
                    dp[comb].push_back(i);
            }
        }
        return dp[(1 << n) - 1];
    }
};

```

```
        if (dp.find(comb) == dp.end() || dp[comb].size() > 1 + dp[it->first].size())
    {
        dp[comb] = it->second;
        dp[comb].push_back(i);
    }
}

return dp[(1 << n) - 1];
};
```