

# COMP2610/COMP6261 – Information Theory

## Tutorial 6

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### Question 1. Entropy

Let  $p = (p_1, p_2, \dots, p_m)$  be a probability distribution on  $m$  elements, i.e.,  $p_i \geq 0$ , and  $\sum_{i=1}^m p_i = 1$ . Define a new distribution  $q$  on  $m-1$  elements as  $q_1 = p_1, q_2 = p_2, \dots, q_{m-2} = p_{m-2}$ , and  $q_{m-1} = p_{m-1} + p_m$ , i.e., the distribution  $q$  is the same as  $p$  on any  $i \in \{1, 2, \dots, m-2\}$ , and the probability of the last element in  $q$  is the sum of the last two probabilities of  $p$ . Show that

$$H(p) = H(q) + (p_{m-1} + p_m)H\left(\frac{p_{m-1}}{p_{m-1} + p_m}, \frac{p_m}{p_{m-1} + p_m}\right).$$

**Solution:**

$$\begin{aligned} H(p) &= -\sum_{i=1}^m p_i \log p_i \\ &= -\sum_{i=1}^{m-2} p_i \log p_i - p_{m-1} \log p_{m-1} - p_m \log p_m \\ &= -\sum_{i=1}^{m-2} p_i \log p_i - p_{m-1} \log \frac{p_{m-1}}{p_{m-1} + p_m} - p_m \log \frac{p_m}{p_{m-1} + p_m} - (p_{m-1} + p_m) \log(p_{m-1} + p_m) \\ &= H(q) - p_{m-1} \log \frac{p_{m-1}}{p_{m-1} + p_m} - p_m \log \frac{p_m}{p_{m-1} + p_m} \\ &= H(q) - (p_{m-1} + p_m) \left( \frac{p_{m-1}}{p_{m-1} + p_m} \log \frac{p_{m-1}}{p_{m-1} + p_m} + \frac{p_m}{p_{m-1} + p_m} \log \frac{p_m}{p_{m-1} + p_m} \right) \\ &= H(q) + (p_{m-1} + p_m)H\left(\frac{p_{m-1}}{p_{m-1} + p_m}, \frac{p_m}{p_{m-1} + p_m}\right). \end{aligned}$$

## Question 2. Mutual Information and Relative Entropy

Let  $X, Y, Z$  be three random variables with a joint probability mass function  $p(X, Y, Z)$ .

(a). Show that

$$I(X, Y; Z) - I(Y, Z; X) = I(Y; Z) - I(X; Y).$$

(b). The relative entropy between the joint distribution and the product of the marginals is  $D(p(x, y, z) || p(x)p(y)p(z))$ . Show that

$$D(p(x, y, z) || p(x)p(y)p(z)) = I(X; Y) + I(X, Y; Z).$$

**Solution:**

(a). Note that

$$I(X, Y; Z) = H(Z) - H(Z|X, Y),$$

$$I(Y, Z; X) = H(X) - H(X|Y, Z),$$

$$I(Y; Z) = H(Z) - H(Z|Y),$$

$$I(X; Y) = H(X) - H(X|Y).$$

We obtain

$$I(X, Y; Z) + I(X; Y) = H(Z) - H(Z|X, Y) + H(X) - H(X|Y)$$

$$= H(Z) + H(X) - H(X, Z|Y)$$

$$= H(X) + H(Z) + H(Y) - H(X, Y, Z).$$

$$I(Y, Z; X) + I(Y; Z) = H(X) - H(X|Y, Z) + H(Z) - H(Z|Y)$$

$$= H(Z) + H(X) - H(X, Z|Y) = I(X, Y; Z) + I(X; Y).$$

Hence, it can be proved that

$$I(X, Y; Z) - I(Y, Z; X) = I(Y; Z) - I(X; Y).$$

(b).

$$\begin{aligned} D(p(x, y, z) || p(x)p(y)p(z)) &= \sum_{x,y,z} p(x, y, z) \log \frac{p(x, y, z)}{p(x)p(y)p(z)} \\ &= \sum_{x,y,z} p(x, y, z) \log p(x, y, z) - \sum_{x,y,z} p(x, y, z) \log p(x) - \sum_{x,y,z} p(x, y, z) \log p(y) - \sum_{x,y,z} p(x, y, z) \log p(z) \\ &= \sum_{x,y,z} p(x, y, z) \log p(x, y, z) - \sum_x p(x) \log p(x) - \sum_y p(y) \log p(y) - \sum_z p(z) \log p(z) \\ &= -H(X, Y, Z) + H(X) + H(Y) + H(Z) \\ &= I(X, Y; Z) + I(X; Y). \end{aligned}$$

### Question 3. Markov Chain

Suppose a Markov chain  $X_1 \rightarrow X_2 \rightarrow X_3$ , starts in one of  $n$  states, i.e.,  $X_1 \in \{1, 2, \dots, n\}$ . Suppose  $X_2$  will go down to  $k < n$  states, i.e.,  $X_2 \in \{1, 2, \dots, k\}$ . Then  $X_3$  go back to  $m > k$  states, i.e.,  $X_3 \in \{1, 2, \dots, m\}$ .

(a). What is the upper bound of  $I(X_1; X_3)$ ?

(b). Evaluate  $I(X_1; X_3)$  for  $k = 1$ , and conclude that no dependence can survive.

**Solution:**

(a). From the data processing inequality, and the fact that entropy is maximum for a uniform distribution, we get

$$\begin{aligned} I(X_1; X_3) &\leq I(X_1; X_2) \\ &= H(X_2) - H(X_2|X_1) \leq H(X_2) \\ &\leq \log |\mathcal{X}_2| = \log k, \end{aligned} \tag{1}$$

$$I(X_1; X_3) \leq H(X_1) \leq \log |\mathcal{X}_1| = \log n, \tag{2}$$

$$I(X_1; X_3) \leq H(X_3) \leq \log |\mathcal{X}_3| = \log m. \tag{3}$$

Hence,  $I(X_1; X_3) \leq \log k$ .

(b). For  $k = 1$ ,  $I(X_1; X_3) \leq \log 1 = 0$ . Since  $I(X_1; X_3) \geq 0$ , we obtain  $I(X_1; X_3) = 0$ . Thus, for  $k = 1$ ,  $X_1$  and  $X_3$  are independent.

## Question 4. Inequalities

A coin is known to land heads with probability  $\frac{1}{5}$ . The coin is flipped  $N$  times for some even integer  $N$ .

- (a). Using Markov's inequality, provide a bound on the probability of observing  $\frac{N}{2}$  or more heads.  
(b). Using Chebyshev's inequality, provide a bound on the probability of observing  $\frac{N}{2}$  or more heads. Express your answer in terms of  $N$ .

**Solution:**

(a). Let's denote the flip result of a coin as  $X$  with head as 1 and tail as 0, i.e.,  $P(X = 1) = \frac{1}{5}$  and  $P(X = 0) = \frac{4}{5}$ . Then we obtain

$$\mathbb{E}[X] = P(X = 1) \times 1 + P(X = 0) \times 0 = \frac{1}{5}.$$

Based on Markov inequality, we obtain

$$p\left(X^N \geq \frac{N}{2}\right) \leq \frac{\mathbb{E}[X^N]}{N/2} = \frac{2}{5}.$$

(b). Based on Chebyshev's inequality, we obtain

$$p\left(|X^N - \mathbb{E}[X^N]| \geq \frac{3N}{10}\right) \leq \frac{V[X^N]}{(3N/10)^2} = \frac{NV[X]}{(3N/10)^2} = \frac{16}{9N}.$$

Note that  $p\left(|X^N - \mathbb{E}[X^N]| \geq \frac{3N}{10}\right) = p\left(X^N - \mathbb{E}[X^N] \geq \frac{3N}{10}\right) + p\left(\mathbb{E}[X^N] - X^N \geq \frac{3N}{10}\right) = p\left(X^N - \mathbb{E}[X^N] \geq \frac{3N}{10}\right)$ . Hence, we obtain

$$p\left(X^N \geq \frac{N}{2}\right) \leq \frac{16}{9N}$$