

CMPE 152: Compiler Design

September 7 Class Meeting

Department of Computer Engineering
San Jose State University

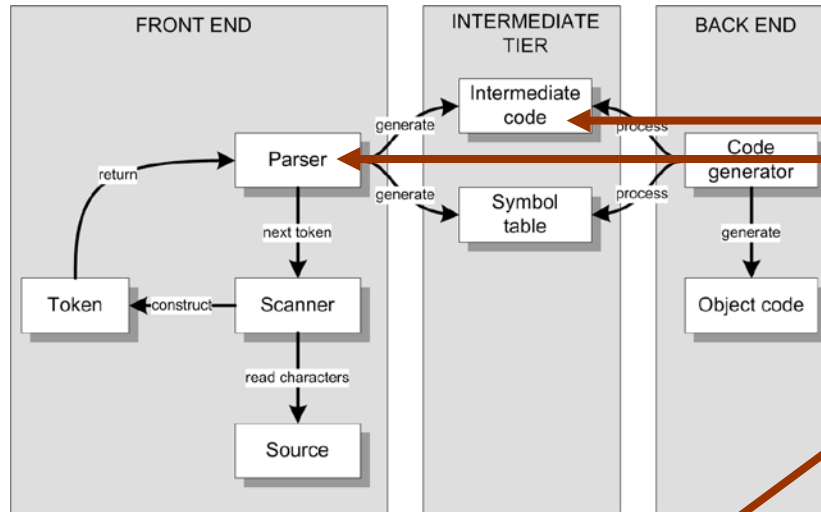


Fall 2017
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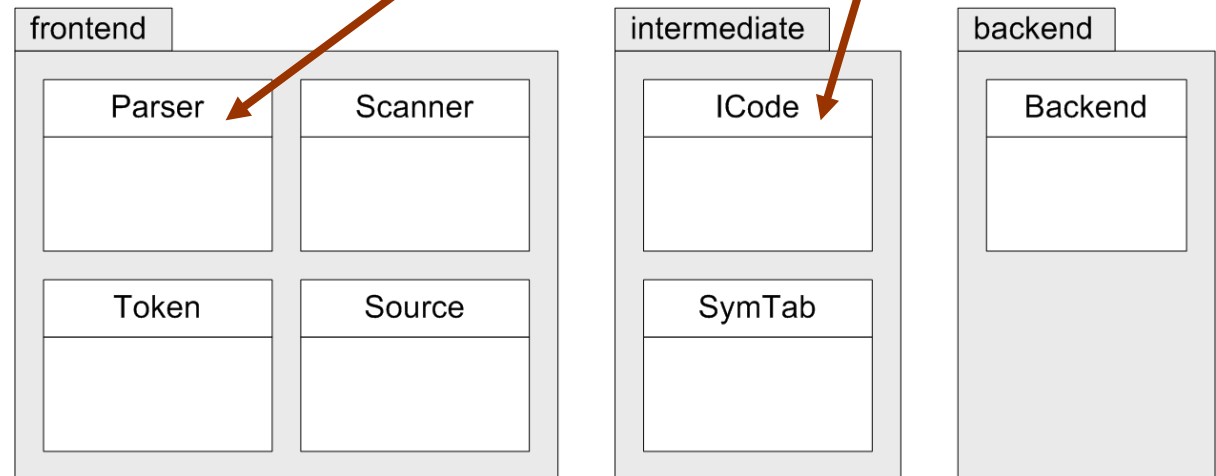
Quick Review of the Framework

FROM:



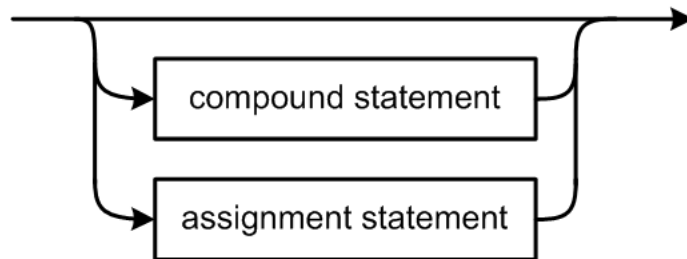
Next topic:
Parsing assignment statements
and expressions, and
generating parse trees.

TO:

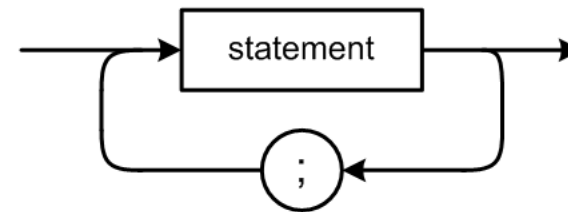


Pascal Statement Syntax Diagrams

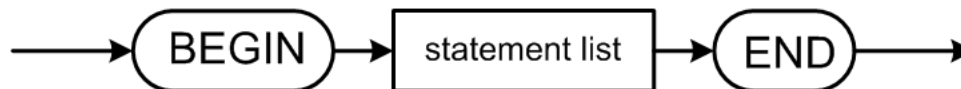
statement



statement list

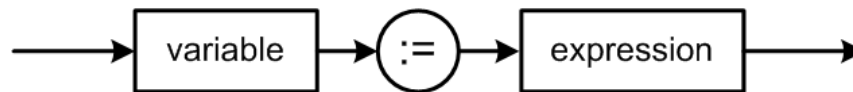


compound statement

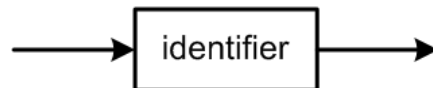


Pascal Statement Syntax Diagrams, *cont'd*

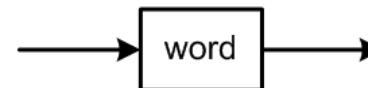
assignment statement



variable

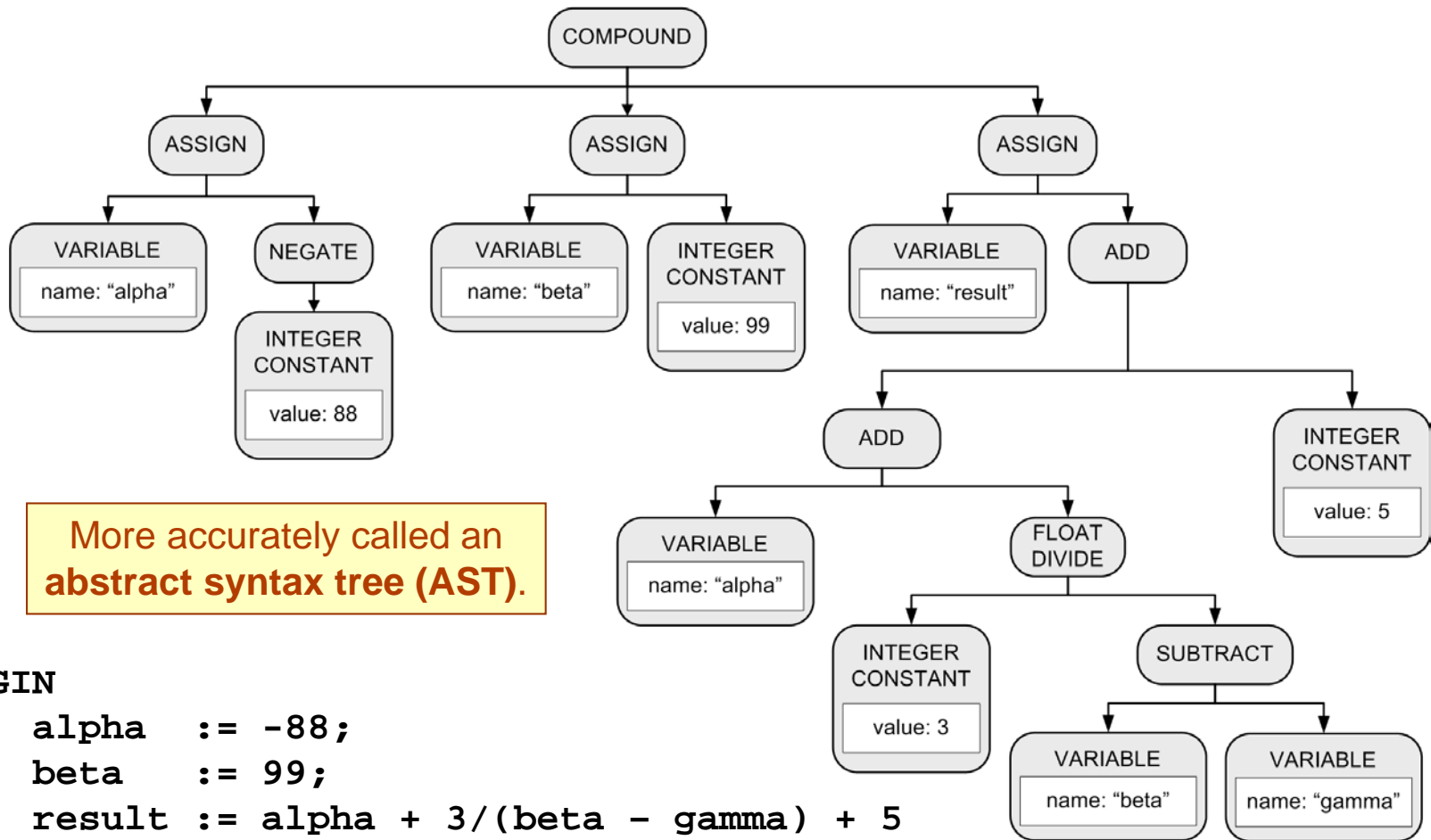


identifier



For now,
greatly simplified!

Parse Tree: Conceptual Design



BEGIN

alpha := -88;

beta := 99;

result := alpha + 3/(beta - gamma) + 5

END

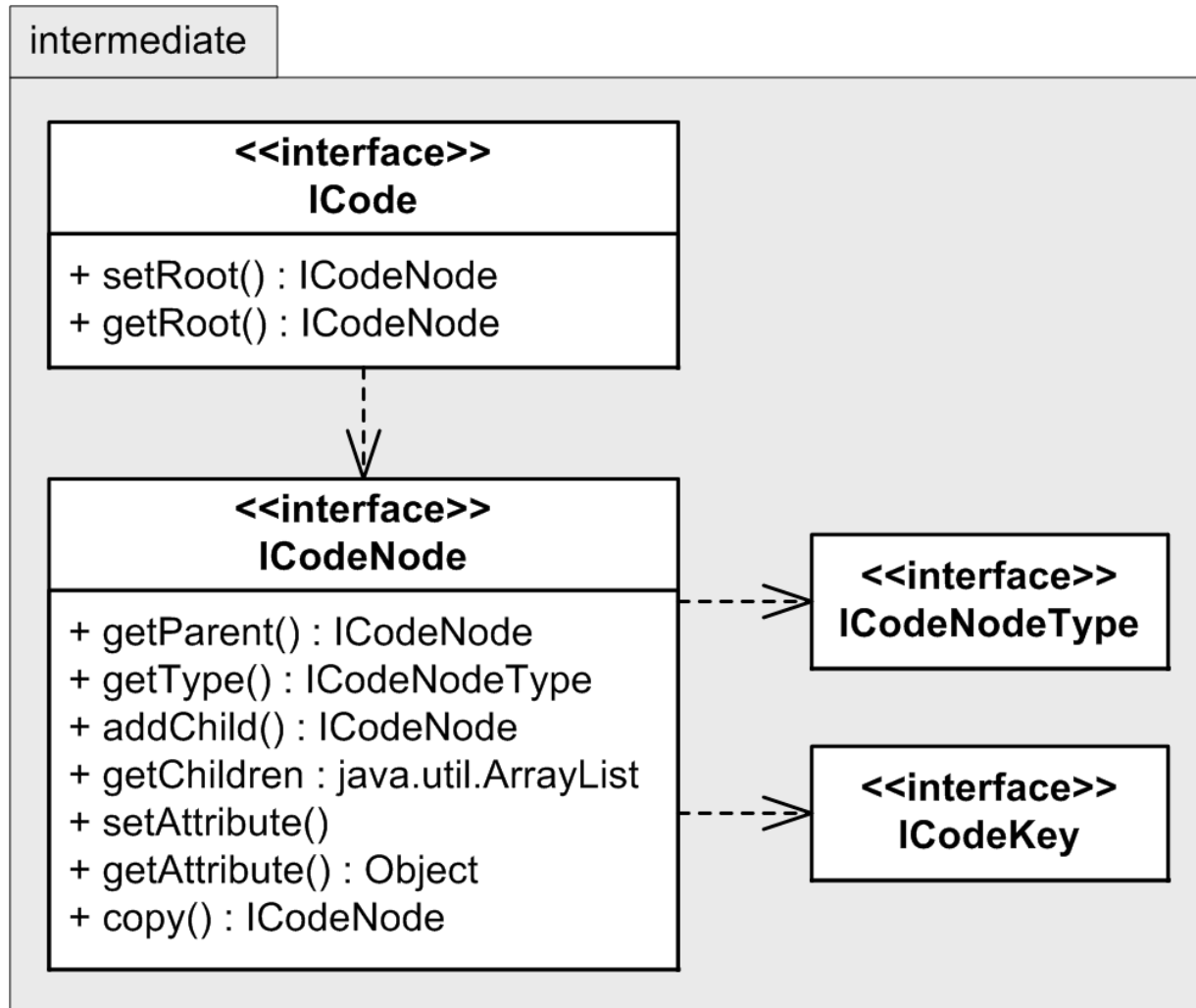
Parse Tree: Conceptual Design

- At the conceptual design level,
we don't care how we implement the tree.
- This should remind you of how
we first designed the symbol table.

Parse Tree: Basic Tree Operations

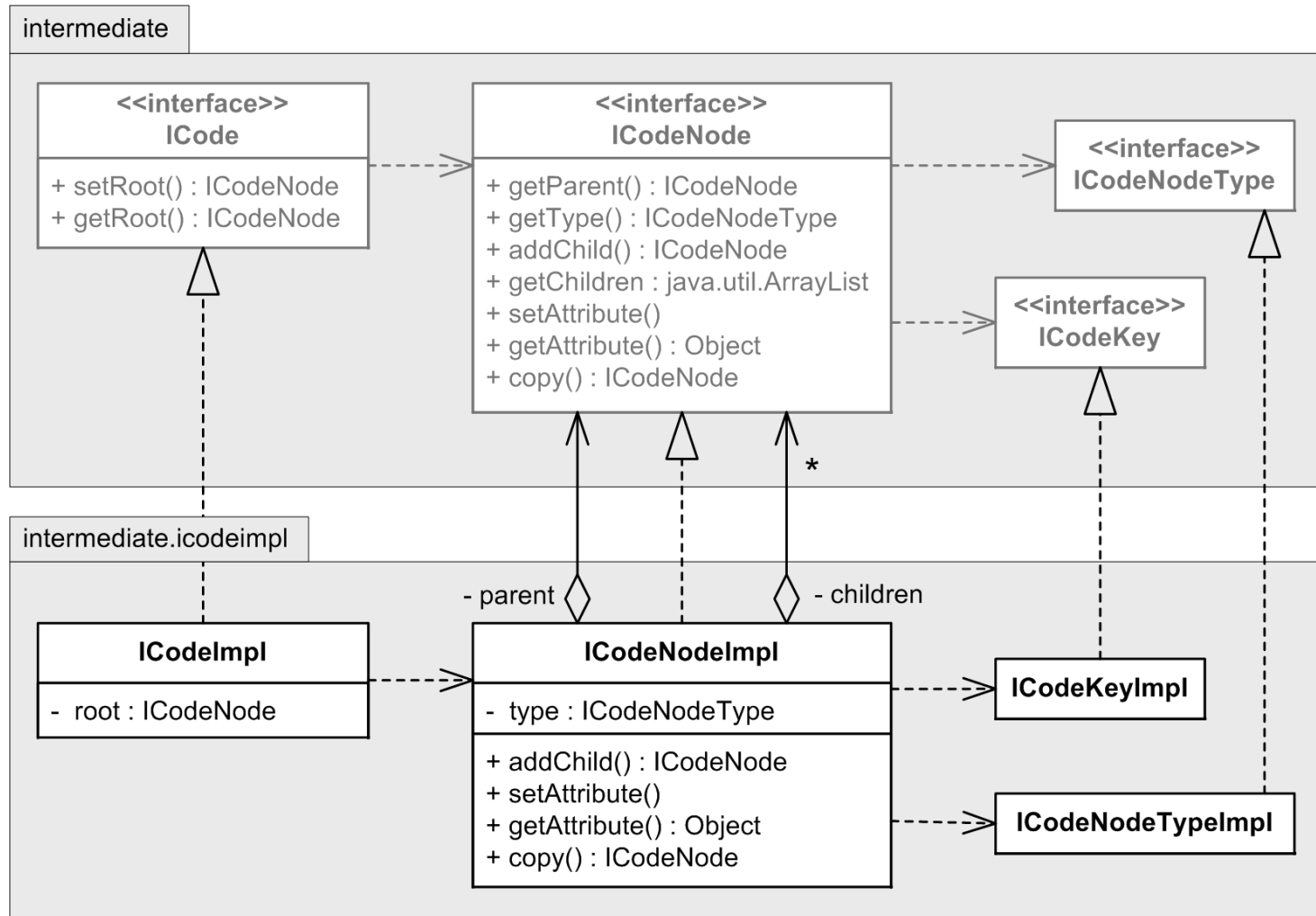
- ❑ Create a new node.
- ❑ Create a copy of a node.
- ❑ Set and get the root node of a parse tree.
- ❑ Set and get an attribute value in a node.
- ❑ Add a child node to a node.
- ❑ Get the list of a node's child nodes.
- ❑ Get a node's parent node.

Intermediate Code Interfaces



Goal: Keep it source language-independent.

Intermediate Code Implementations



An Intermediate Code Factory Class

```
ICode *ICodeFactory::create_icode()  
{  
    return new ICodeImpl();  
}  
  
ICodeNode *ICodeFactory::create_icode_node(const ICodeNodeType type)  
{  
    return new ICodeNodeImpl(type);  
}
```

Coding to the Interfaces (Again)

```
// Create the ASSIGN node.
ICodeNode *assign_node =
    ICodeFactory::create_icode_node((ICodeNodeType) NT_ASSIGN);

...

// Create the variable node and set its name attribute.
ICodeNode *variable_node =
    ICodeFactory::create_icode_node(
        (ICodeNodeType) NT_VARIABLE);
NodeValue *node_value = new NodeValue();
node_value->id = target_id;
variable_node->set_attribute((ICodeKey) ID, node_value);

// The ASSIGN node adopts the variable node as its first child.
assign_node->add_child(variable_node);
```

Intermediate Code (ICode) Node Types

```
enum class ICodeNodeTypeImpl
{
    // Program structure
    PROGRAM, PROCEDURE, FUNCTION,

    // Statements
    COMPOUND, ASSIGN, LOOP, TEST, CALL, PARAMETERS,
    IF, SELECT, SELECT_BRANCH, SELECT_CONSTANTS, NO_OP,

    // Relational operators
    EQ, NE, LT, LE, GT, GE, NOT,

    // Additive operators
    ADD, SUBTRACT, OR, NEGATE,

    // Multiplicative operators
    MULTIPLY, INTEGER_DIVIDE, FLOAT_DIVIDE, MOD, AND,

    // Operands
    VARIABLE, SUBSCRIPTS, FIELD,
    INTEGER_CONSTANT, REAL_CONSTANT,
    STRING_CONSTANT, BOOLEAN_CONSTANT,

    // WRITE parameter
    WRITE_PARM,
};
```

Do not confuse
node types (ASSIGN, ADD, etc.)
with **data types** (integer, real, etc.).

We use the enumerated type
ICodeNodeTypeImpl for node types
which is different from the enumerated
type **PascalTokenType** to help maintain
source language independence.

Intermediate Code Node Implementation

```
class ICodeNodeImpl : public ICodeNode
{
public:
    ICodeNodeImpl(const ICodeNodeType type);
    ...

protected:
    map<ICodeKey, NodeValue *> contents;
    ICodeNodeType type;           // node type
    ICodeNode *parent;           // node's parent
    vector<ICodeNode *> children; // node's children
}
```

```
ICodeNodeImpl::ICodeNodeImpl(const ICodeNodeType type)
    : type(type), parent(nullptr)
{
}
```

- ❑ Each node has a `map<ICodeKey, NodeValue *>`.
- ❑ Each node has an `vector<ICodeNode *>` of child nodes.

A Parent Node Adopts a Child Node

```
ICodeNode *ICodeNodeImpl::add_child(ICodeNode *node)
{
    if (node != nullptr)
    {
        children.push_back(node);
        node->parent = this;
    }
    return node;
}
```

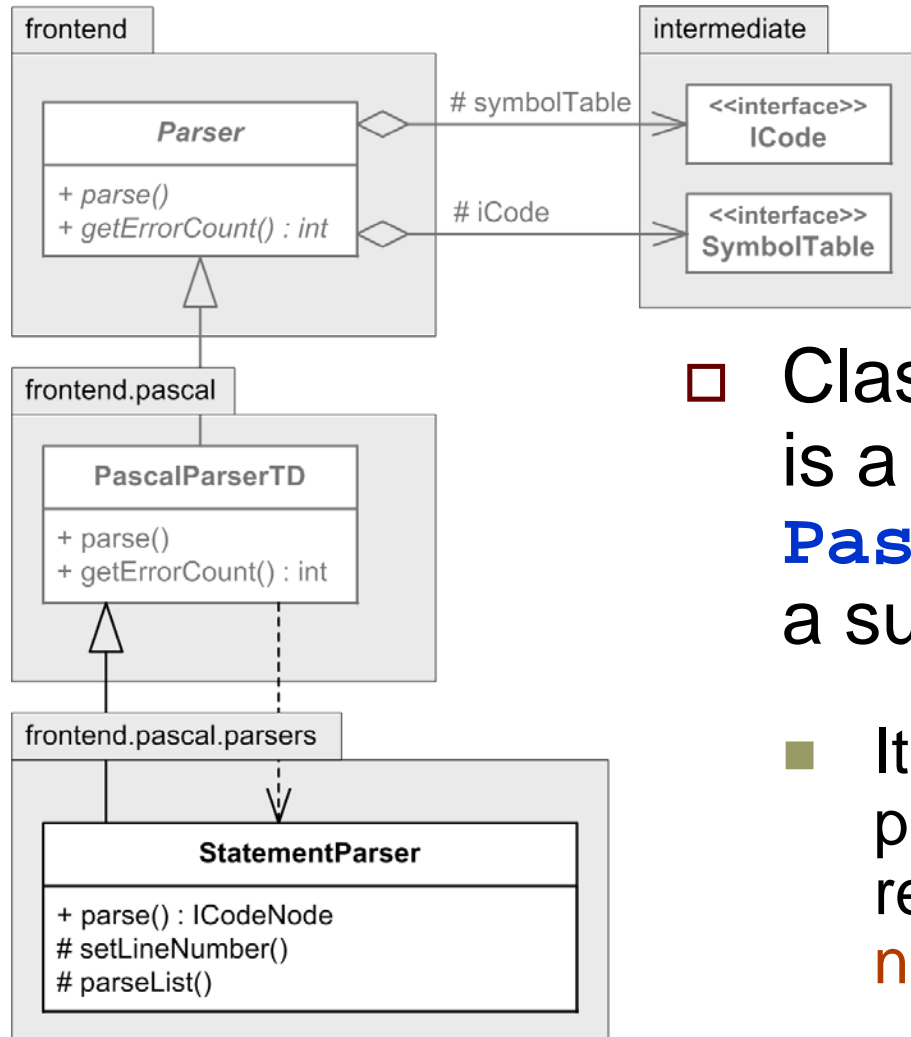
- ❑ When a parent node adds a child node, we can say that the parent node “**adopts**” the child node.
- ❑ Keep the parse tree implementation simple!

What Attributes to Store in a Node?

```
enum class ICodeKeyImpl
{
    LINE, ID, LEVEL, VALUE,
};
```

- ❑ Not much! Not every node will have these attributes.
 - **LINE**: statement line number
 - **ID**: symbol table entry of an identifier
 - **VALUE**: data value
- ❑ Most of the information about what got parsed is encoded in the **node type** and in the **tree structure**.

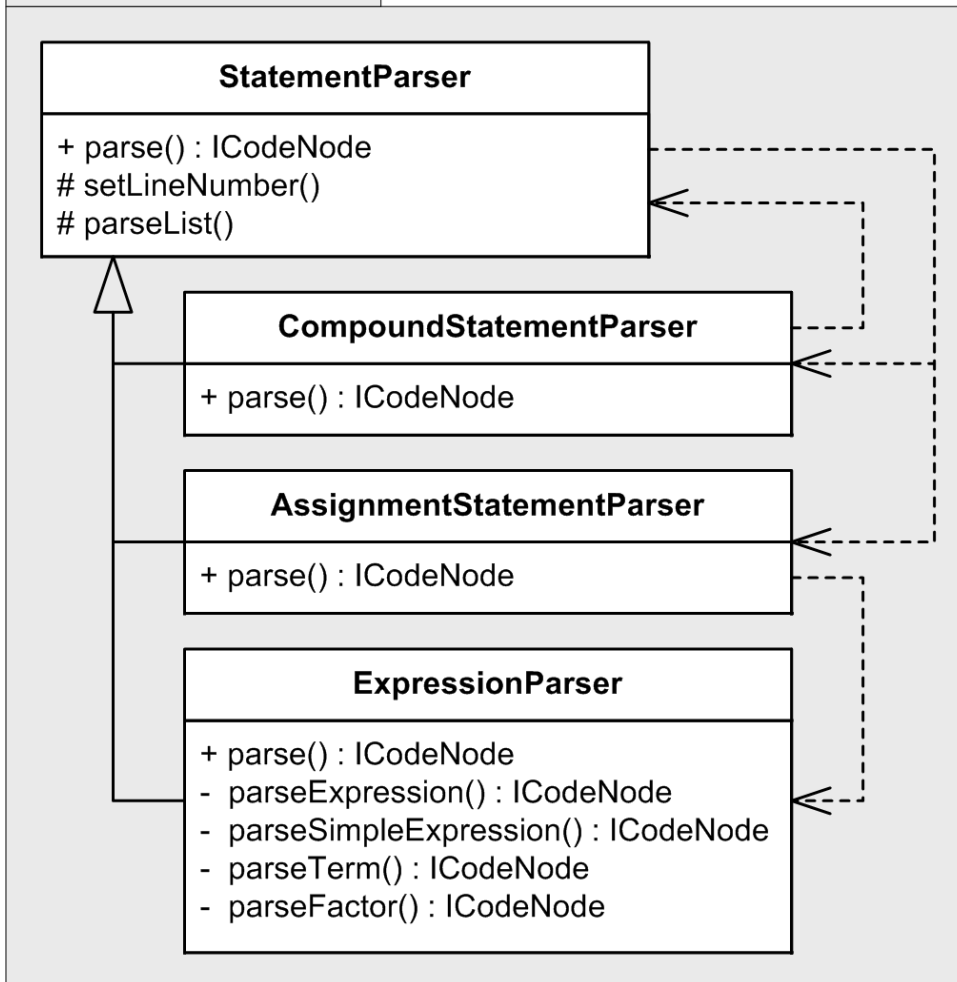
Statement Parser Class



- ❑ Class **StatementParser** is a subclass of **PascalParserTD** which is a subclass of **Parser**.
- Its **parse()** method builds a part of the parse tree and returns the **root node of the newly built subtree**.

Statement Parser Subclasses

frontend.pascal.parsers



- **StatementParser** itself has subclasses:
 - **CompoundStatementParser**
 - **AssignmentStatementParser**
 - **ExpressionParser**
- The **parse()** method of each subclass returns the root node of the subtree that it builds.
- Note the dependency relationships among **StatementParser** and its subclasses.

Building a Parse Tree

- Each `parse()` method builds a subtree and returns the `root node` of the new subtree.
- The caller of the `parse()` method adopts the subtree's root node as a child of the subtree that the caller is building.
 - The caller then returns the root node of its subtree to its caller.
 - This process continues until the entire source has been parsed and we have the entire parse tree.

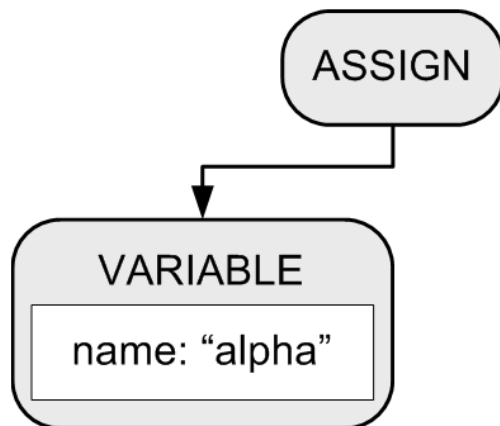
Building a Parse Tree

□ Example:

```
BEGIN
    alpha := 10;
    beta  := 20
END
```

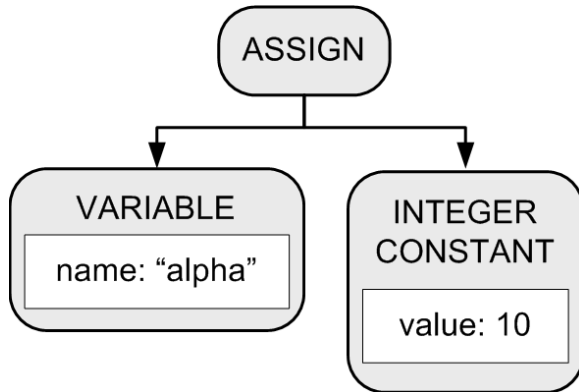
1. **CompoundStatementParser's parse()** method creates a **COMPOUND** node.

COMPOUND



2. **AssignmentStatementParser's parse()** method creates an **ASSIGN** node and a **VARIABLE** node, which the **ASSIGN** node adopts as its first child.

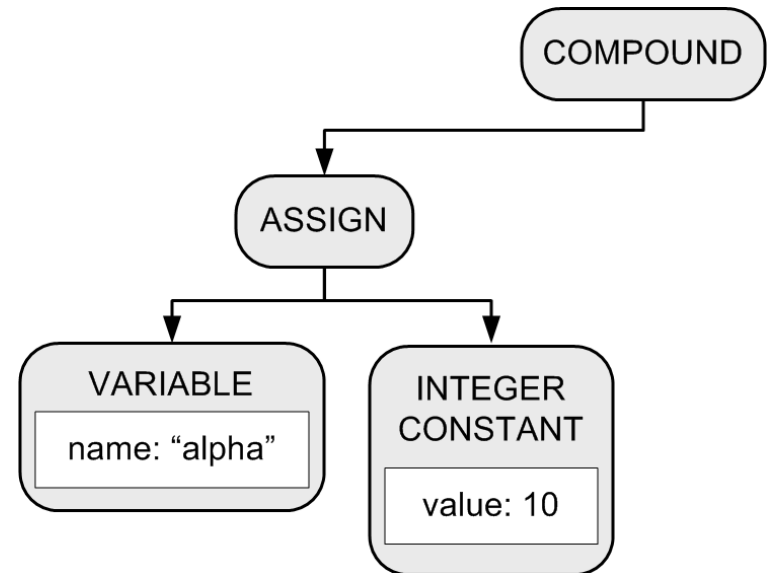
Building a Parse Tree



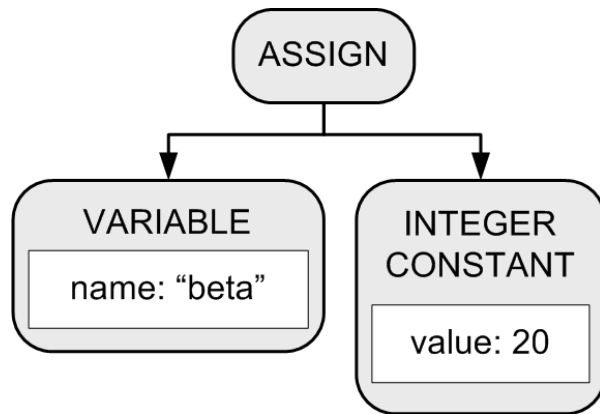
3. **ExpressionParser's parse()** method creates an **INTEGER CONSTANT** node which the **ASSIGN** node adopts as its second child.

```
BEGIN
    alpha := 10;
    beta  := 20
END
```

4. The **COMPOUND** node adopts the **ASSIGN** node as its first child.

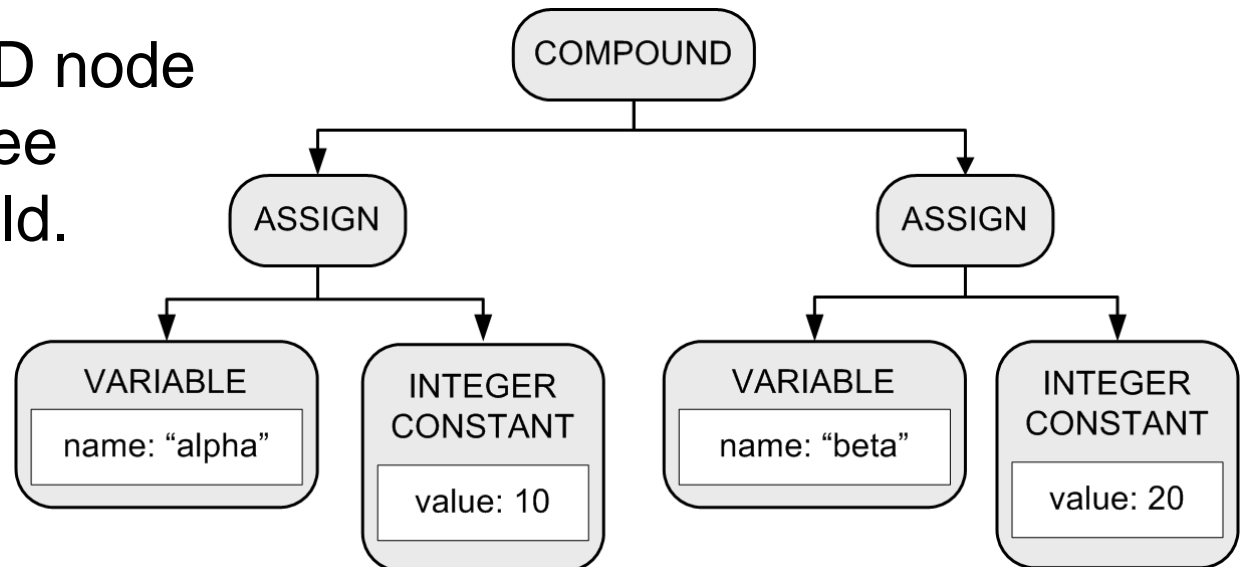


Building a Parse Tree



5. Another set of calls to the `parse()` methods of `AssignmentStatementParser` and `ExpressionParser` builds another assignment statement subtree.

6. The COMPOUND node adopts the subtree as its second child.



BEGIN

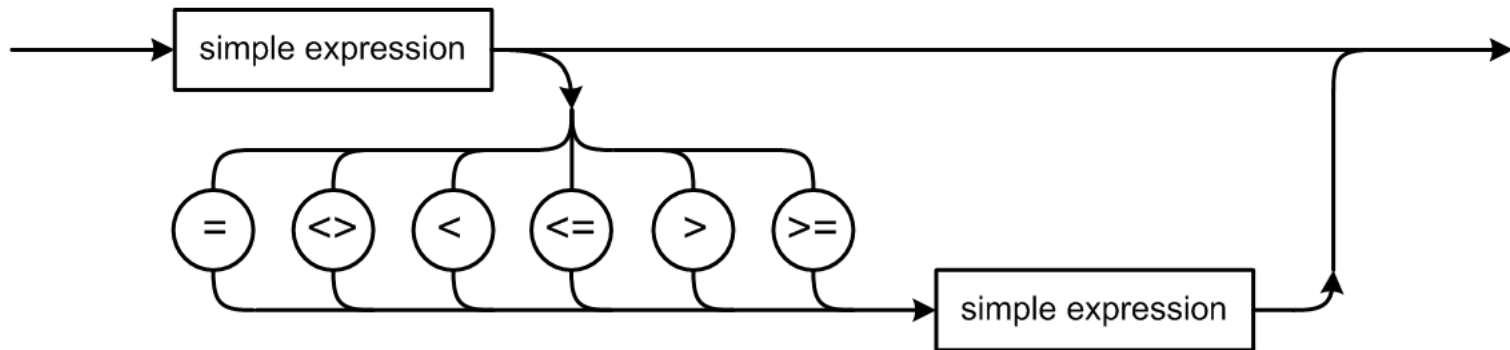
alpha := 10;

beta := 20

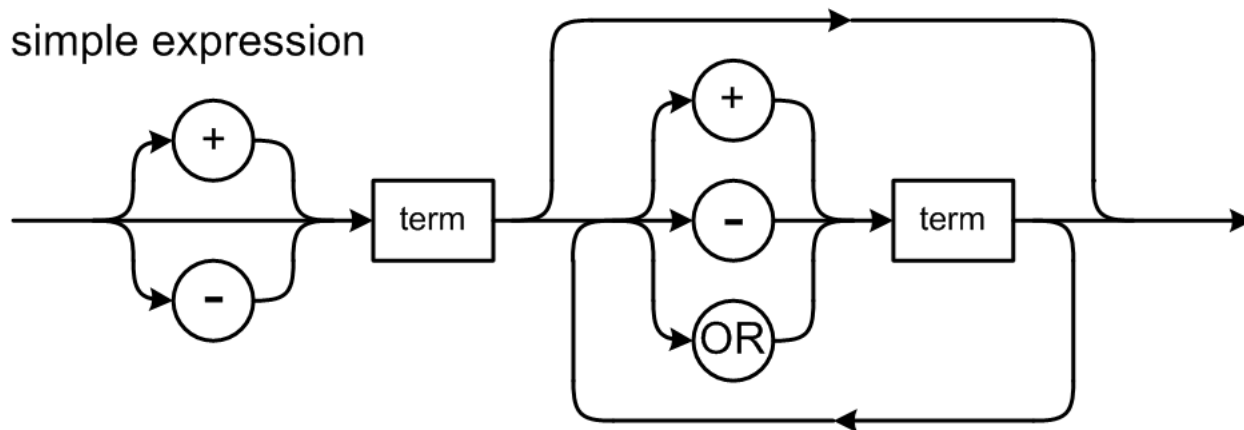
END

Pascal Expression Syntax Diagrams

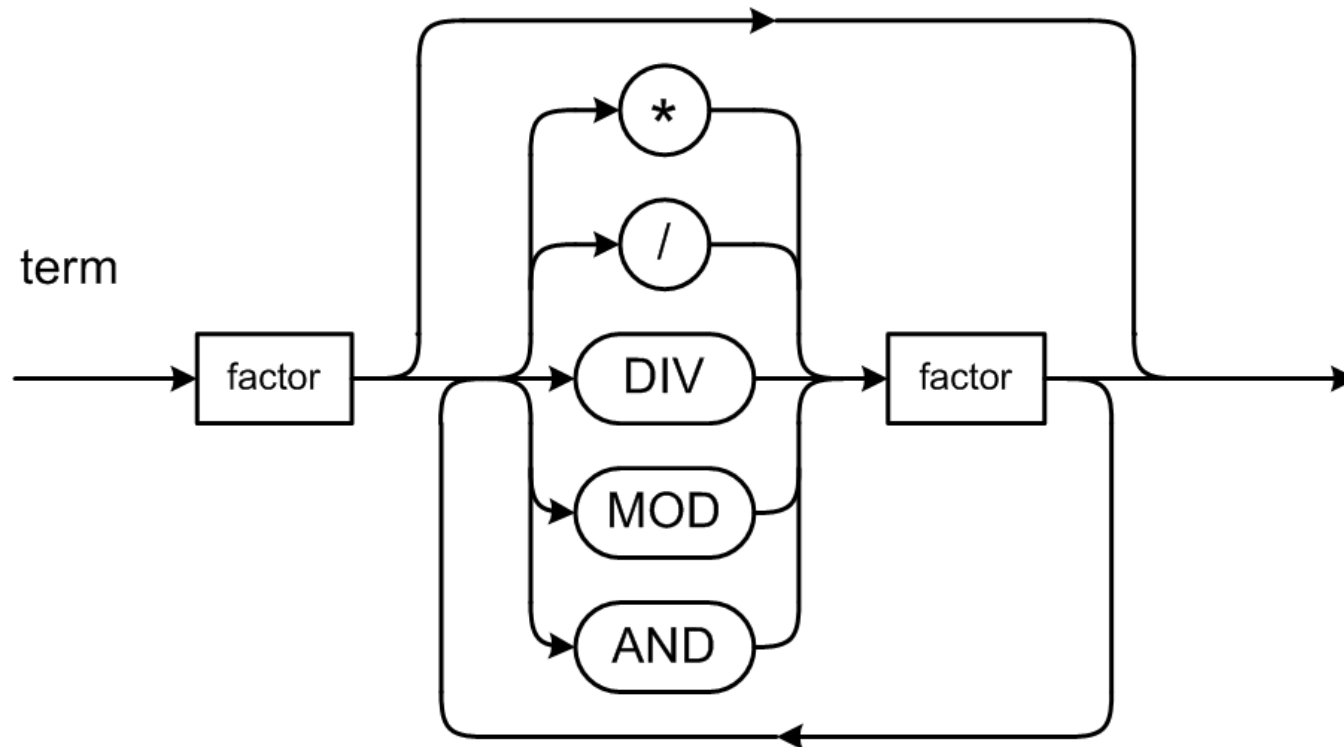
expression



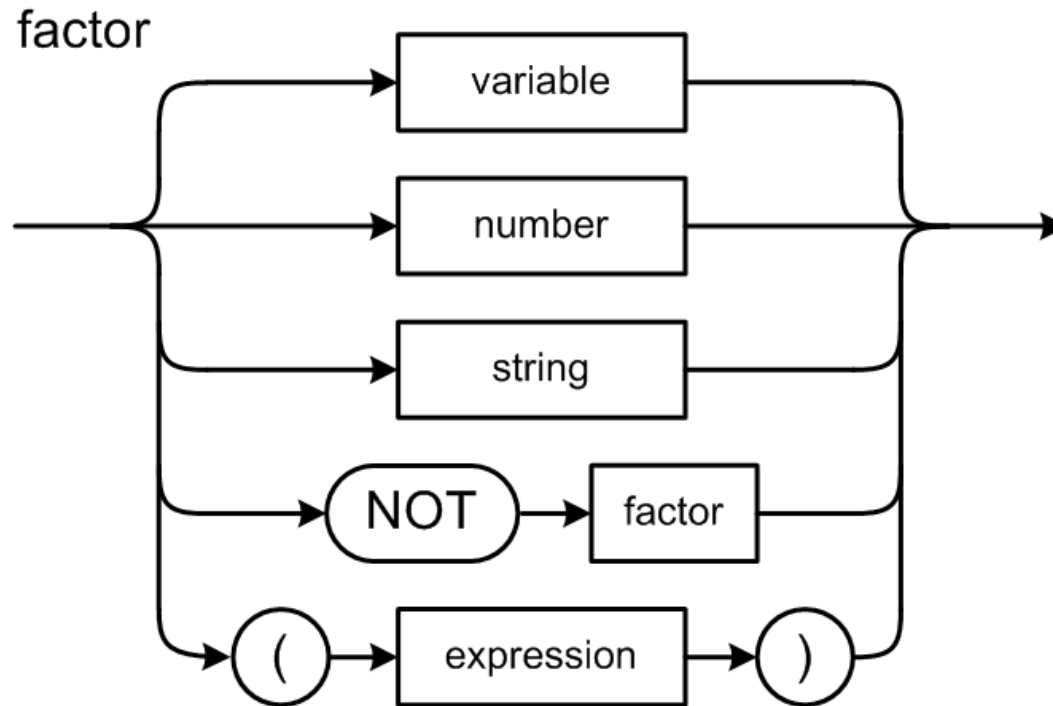
simple expression



Expression Syntax Diagrams, *cont'd*



Expression Syntax Diagrams, *cont'd*



Pascal's Operator Precedence Rules

Level	Operators
1 (factor: <i>highest</i>)	NOT
2 (term)	multiplicative: * / DIV MOD AND
3 (simple expression)	additive: + - OR
4 (expression: <i>lowest</i>)	relational: = <> < <= > >=

- If there are no parentheses:
 - Higher level operators execute before the lower level ones.
 - Operators at the same level execute from left to right.
- Because the **factor** syntax diagram defines parenthesized expressions, parenthesized expressions always execute first, from the most deeply nested outwards.

Example Decomposition

□ $\text{alpha} + 3/(\text{beta} - \text{gamma}) + 5$

