CMPE 152: Compiler Design

September 5 Class Meeting

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Teams

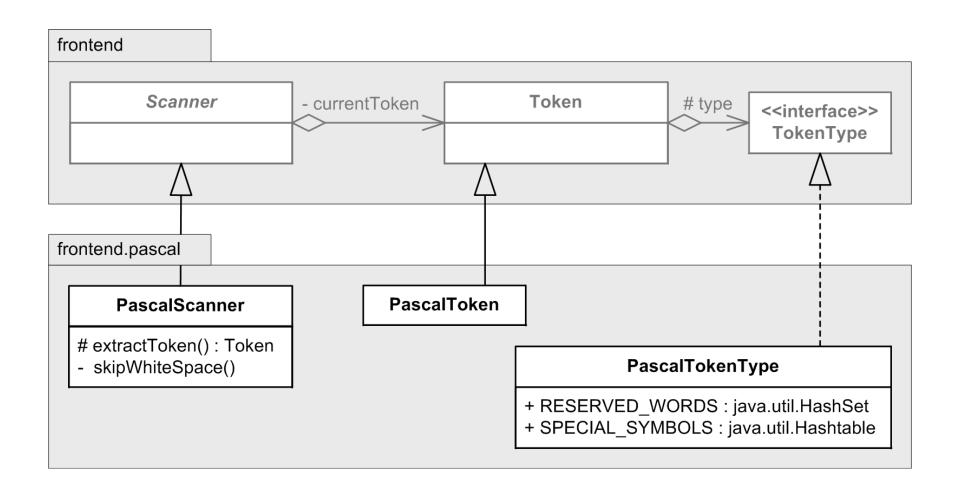


Basic Scanning Algorithm

- Skip any blanks until the current character is nonblank.
 - In Pascal, a comment and the end-of-line character each should be treated as a blank.
- The current (nonblank) character determines what the next token is and becomes that token's first character.
- Extract the rest of the next token by copying successive characters up to but not including the first character that does not belong to that token.
- Extracting a token consumes all the source characters that constitute the token.
 - After extracting a token, the current character is the first character after the last character of that token.



Pascal-Specific Subclasses



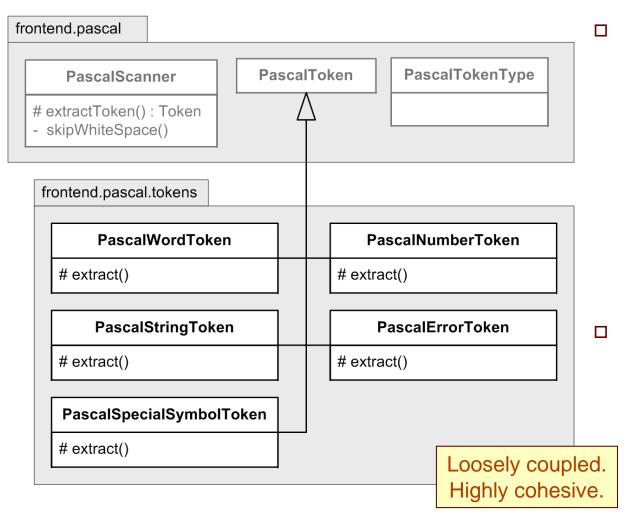


Class PascalScanner

```
Token *PascalScanner::extract_token() throw (string)
   skip_white_space();
    Token *token;
   char current_ch = current_char();
    string string_ch = " ";
   string_ch[0] = current_ch;
    // Construct the next token. The current character determines the
    // token type.
    if (current ch == Source::END OF FILE)
                                                   The first character
        token = nullptr;
                                                   determines the type
                                                   of the next token.
    else if (isalpha(current_ch))
        token = new PascalWordToken(source);
    else if (isdigit(current_ch))
        token = new PascalNumberToken(source);
   return token;
```



Pascal-Specific Token Classes



- Each class
 PascalWordToken,
 PascalNumberToken,
 PascalStringToken,
 PascalSpecialSymbolToken, and
 PascalErrorToken is
 is a subclass of class
 PascalToken.
- PascalToken is a subclass of class Token.
- Each Pascal token subclass overrides the default extract() method of class Token.
- The default method could only create single-character tokens.

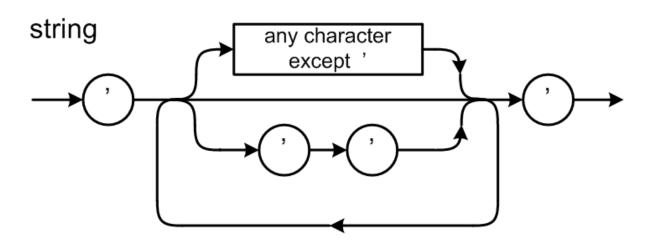


Class PascalWordToken

```
void PascalWordToken::extract() throw (string)
    char current ch = current char();
    // Get the word characters (letter or digit). The scanner has
    // already determined that the first character is a letter.
    while (isalnum(current ch))
        text += current_ch;
        current_ch = next_char(); // consume character
    // Is it a reserved word or an identifier?
    string upper_case(text);
    transform(upper_case.begin(), upper_case.end(),
              upper_case.begin(), ::toupper);
    if (PascalToken::RESERVED WORDS.find(upper case)
            != PascalToken::RESERVED WORDS.end())
        // Reserved word.
        type = (TokenType) PascalToken::RESERVED WORDS[upper case];
        value = new DataValue(upper case);
    else
        // Identifier.
        type = (TokenType) PT_IDENTIFIER;
```



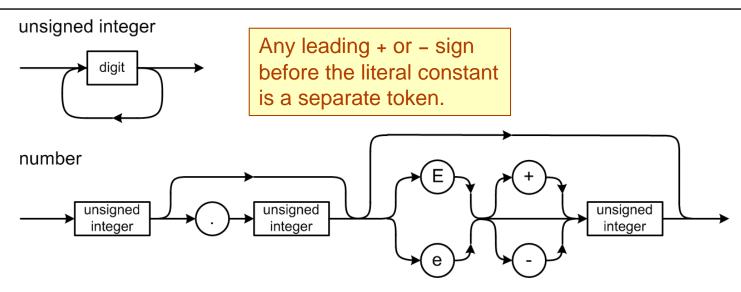
Pascal String Tokens



- A Pascal string literal constant uses single quotes.
- Two consecutive single quotes represents a single quote character inside a string.
 - 'Don''t' is the string consisting of the characters Don't.
- A Pascal character literal constant is simply a string with only a single character, such as 'a'.
- Pascal token subclass PascalStringToken.



Pascal Number Tokens



- A Pascal integer literal constant is an unsigned integer.
- A Pascal real literal constant starts with an unsigned integer (the whole part) followed by either
 - A decimal point followed by another unsigned integer (the fraction part), or
 - An E or e, optionally followed by + or -, followed by an unsigned integer (the exponent part), or
 - A whole part followed by an exponent part.



Class Pascal Number Token

For the token string "31415.926e-4", method extractNumber() passes the following parameter values to method computeFloatValue():

wholeDigits	"31415"
fractionDigits	"926"
exponentDigits	"4"
exponentSign	1 _ 1

Compute variable exponentValue:

4	as computed by computeIntegerValue()
-4	after negation since exponentSign is '-'
-7	after subtracting fractionDigits.length()

 \square Compute 31415926 x 10⁻⁷ = 3.1415926



Syntax Error Handling

- Error handling is a three-step process:
 - Detect the presence of a syntax error.
 - Flag the error by pointing it out or highlighting it, and display a descriptive error message.
 - Recover by moving past the error and resume parsing.
 - For now, we'll just move on, starting with the current character, and attempt to extract the next token.
- □ SYNTAX_ERROR message
 - source line number
 - beginning source position
 - token text
 - syntax error message



Class PascalParserTD

```
void PascalParserTD::parse() throw (string)
    // Loop over each token until the end of file.
    while ((token = next_token(token)) != nullptr)
        TokenType token type = token->get type();
        last_line number = token->get_line number();
        string type_str;
        string value_str;
        switch ((PascalTokenType) token_type)
            case PT STRING:
                type str = "STRING";
                value_str = token->get_value()->s;
                break:
```



Class PascalParserTD, cont'd

```
case PT_IDENTIFIER:
    type_str = "IDENTIFIER";
    value str = "";
    break;
case PT_INTEGER:
    type_str = "INTEGER";
    value str = token->get value()->display();
    break;
case PT_REAL:
    type_str = "REAL";
    value_str = token->get_value()->display();
    break;
case PT_ERROR: break;
```



Class PascalParserTD, cont'd

```
default: // reserved word or special character
    DataValue *token value = token->get value();
    // Reserved word
    if (token_value != nullptr)
        value_str = token_value->s;
        type_str = value_str;
    // Special symbol
    else
        type_str =
            PascalToken::SPECIAL SYMBOL NAMES[
                               (PascalTokenType) token_type];
    break;
```



Class PascalParserTD, cont'd

```
if (token_type != (TokenType) PT_ERROR)
    // Format and send a message about each token.
    Message message (TOKEN,
                    LINE NUMBER, to string(token->get line number()),
                    POSITION, to string(token->get position()),
                    TOKEN TYPE, type str,
                    TOKEN TEXT, token->get text(),
                    TOKEN VALUE, value str);
    send message(message);
else
    PascalErrorCode error code =
                        (PascalErrorCode) token->get value()->i;
    error handler.flag(token, error code, this);
```



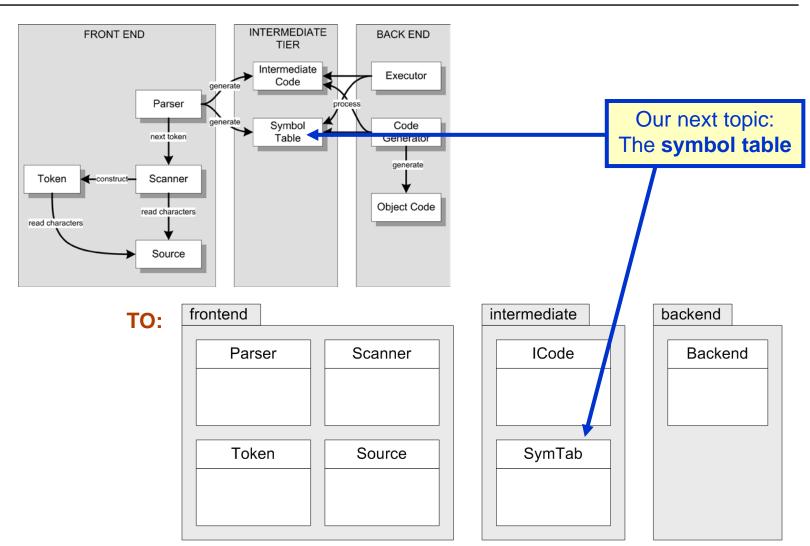
Program: Pascal Tokenizer

- Verify the correctness of the Pascal token subclasses.
- Verify the correctness of the Pascal scanner.
- Demo (Chapter 3)



Quick Review of the Framework

FROM:





The Symbol Table: Basic Concepts

Purpose

- To store information about certain tokens during the translation process (i.e., parsing and scanning)
- What information to store?
 - Anything that's useful!
 - For an identifier:
 - name
 - data type
 - how it's defined (as a variable, type, function name, etc.)

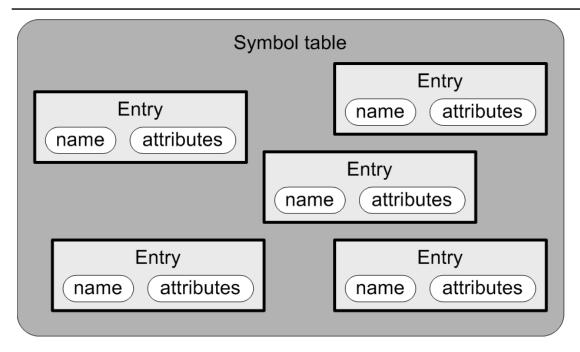


The Symbol Table: Basic Operations

- Enter new information.
- Look up existing information.
- Update existing information.



The Symbol Table: Conceptual Design



Goal: The symbol table should be source language independent.

- Each entry in the symbol table has
 - a name
 - attributes
- At the conceptual level, we don't worry about implementation.



What Needs a Symbol Table?

A Pascal program

Identifiers for constant, type, variable, procedure, and function names.

A Pascal procedure or function

- Identifiers for constant, type, variable, procedure, and function names.
- Identifiers for formal parameter (argument) names.

A Pascal record type

Identifiers for field names.

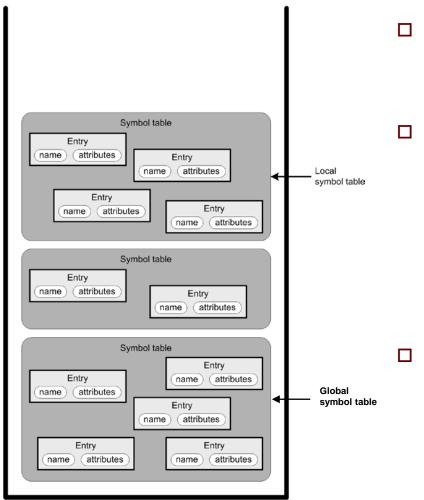


The Symbol Table Stack

- Language constructs can be nested.
 - Procedures and functions are nested inside a program.
 - Procedures and functions can be nested inside of each other.
 - Record types are defined within programs, procedures, and functions.
 - Record types can be nested inside of each other.
- Therefore, symbol tables need to be kept on a symbol table stack.



The Symbol Table Stack, cont'd



- Whichever symbol table is on top of the stack is the local symbol table.
 - The first symbol table created (the one at the bottom of the stack) is the global symbol table.
 - It stores the predefined information, such as entries for the names of the standard types integer, real, char, and boolean.
 - During the translation process, symbol tables are pushed onto and popped off the stack ...
 - ... as the parser enters and exits nested procedures, functions, record types, etc.





The Symbol Table Stack, cont'd

- Symbol table Entry Entry name attributes attributes Local Entry symbol table attributes Global Entry Entry symbol table attributes
 - For now, we'll have only have a <u>single</u> symbol table.
 - Therefore, the local symbol table is the global symbol table.
 - We won't need multiple symbol tables until we start to parse declarations.
 - Implementing the symbol table stack now will make things easier for us later.



Symbol table stack