# CMPE 152: Compiler Design

September 26 Lab

Department of Computer Engineering San Jose State University



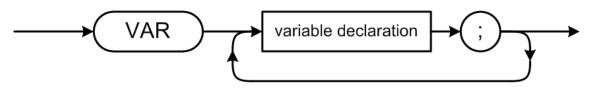
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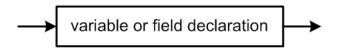


## Pascal Variable Declarations

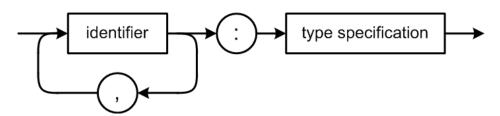
#### variable declarations



#### variable declaration

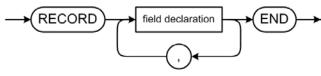


#### variable or field declaration



#### Compare to:

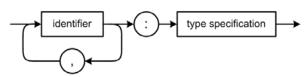




#### field declaration



#### variable or field declaration





### VariableDeclarationsParser.parse()

- Repeatedly call parse\_identifier\_sublist()
  to parse the semicolon-separated
  sublists of variables.
  - Loop to parse the <u>comma</u>-separated list of variable names.

```
i, j, k : integer;

x, y, z : real;

comma-separated list
```



### VariableDeclarationsParser.parse()

- Repeatedly call parse\_identifier\_sublist()
  to parse the semicolon-separated
  sublists of variables (cont'd).
  - Call parse\_identifier() to parse each variable name.
    - Enter each identifier into the current symbol table (the one at the top of the symbol table stack).
    - Set each identifier's definition to VARIABLE

```
VAR
i, j, k : integer;
x, y, z : real;
```



### VariableDeclarationsParser.parse()

- Repeatedly call parse\_identifier\_sublist() to parse the <u>semicolon</u>-separated sublists of variables (cont'd).
  - Call parseTypeSpec() to parse the type specification.
    - Consume the token.
    - Call TypeSpecificationParser::parse() to parse the type specification.
  - Assign the <u>type specification</u> to each variable in the list.

```
VAR
    i, j, k : integer;
    x, y, z : real;
```



### Demo

- Pascal Cross-Referencer II
  - Parse declarations
  - Generate a detailed cross-reference listing
  - Syntax check declarations



## Now that we can parse declarations ...

We can parse variables that have subscripts and fields.

Chapter 10

Example:

```
var9.rec.flda[b][0,'m'].flda[d] := 'p'
```

- We can perform type checking.
  - A semantic action.



# Type Checking

- Ensure that the types of the operands are type-compatible with their operator.
  - Example: You can only perform an integer division with the DIV operator and integer operands.
  - Example: The relational operators AND and OR can only be used with boolean operands.
- Ensure that a value being assigned to a variable is assignment-compatible with the variable.
  - Example: You cannot assign a string value to an integer variable.



# Type Specifications and the Parse Tree

- Every Pascal expression has a data type.
- Add a type specification to every parse tree node.
  - "Decorate" the parse tree with type information.
  - In interface ICodeNode.h:

```
void set_typespec(TypeSpec *spec) = 0;
TypeSpec *get_typespec() const = 0;
```

In class ICodeNodeImpl.h:

```
private:
    TypeSpec *typespec; // data type specification

public:
    void set_typespec(TypeSpec typeSpec) { ... }
    TypeSpec *get_typespec() { ... }
```



## Class TypeChecker

Static boolean methods for type checking:

```
is_integer()
                    In namespace
are_both_integer()
                    intermediate::typeimpl
is_real()
is_integer_or_real()
is_at_least_one_real()
is_boolean()
are both boolean()
is char()
are_assignment_compatible()
are_comparison_cCompatible()
equal_length_strings()
```



## Class TypeChecker, cont'd

```
bool TypeChecker::is integer(TypeSpec *typespec)
             (typespec != nullptr)
    return
           && (typespec->base type() == Predefined::integer type);
}
bool TypeChecker::are both integer(TypeSpec *typespec1,
                                   TypeSpec *typespec2)
   return is integer(typespec1) && is integer(typespec2);
bool TypeChecker::is at least one real(TypeSpec *typespec1,
                                       TypeSpec *typespec2)
    return (is_real(typespec1) && is_real(typespec2)) ||
           (is_real(typespec1) && is_integer(typespec2)) | |
           (is integer(typespec1) && is real(typespec2));
```

# Assignment and Comparison Compatible

- In classic Pascal, a value is assignment-compatible with a target variable if:
  - both have the same type
  - the target is real and the value is integer
  - they are equal-length strings
- Two values are comparison-compatible (they can be compared with relational operators) if:
  - both have the same type
  - one is integer and the other is real
  - they are equal-length strings



# Assignment Compatible

```
bool TypeChecker::are_assignment_compatible(TypeSpec *target_typespec,
                                            TypeSpec *value typespec)
{
    if ((target_typespec == nullptr) || (value_typespec == nullptr))
        return false;
    target_typespec = target_typespec->base_type();
   value typespec = value typespec->base type();
   bool compatible = false;
    if (target_typespec == value_typespec)
                                             Same type
        compatible = true;
    else if (is_real(target_typespec) && is_integer(value_typespec))
        compatible = true;
                                             real := integer
    else {
        compatible =
                        target_typespec->is_pascal_string()
                     && value typespec->is pascal string();
                                             Both are strings
   return compatible;
```



# Type Checking Expressions

- The parser must perform type checking of every expression as part of its semantic actions.
- Add type checking to class ExpressionParser and to each statement parser.
- Flag type errors similarly to syntax errors.



## Method ExpressionParser.parseTerm()

 Now besides doing <u>syntax checking</u>, our expression parser must also do <u>type checking</u> and determine the result type of each operation.

```
case PT STAR:
    if (TypeChecker::are_both_integer(result_typespec,
                                       factor_typespec))
        result typespec = Predefined::integer type;
                                   integer * integer → integer result
    else if (TypeChecker::is at least one real(result typespec,
                                                 factor typespec))
        result_typespec = Predefined::real type;
                  one integer and one real, or both real -> real result
    else
        error handler.flag(expr_token, INCOMPATIBLE TYPES, this);
    break:
```



# Type Checking Control Statements

## Method IfStatementParser.parse()

```
ICodeNode *IfStatementParser::parse_statement(Token *token) throw (string)
    token = next token(token); // consume the IF
    ICodeNode *if node =
            ICodeFactory::create_icode_node((ICodeNodeType) NT_IF);
    Token *expr_token = new Token(*token);
    ExpressionParser expression parser(this);
    ICodeNode *expr_node = expression_parser.parse_statement(token);
    if node->add child(expr node);
    TypeSpec *expr_typespec = expr_node != nullptr
                                  ? expr node->get typespec()
                                  : Predefined::undefined_type;
    if (!TypeChecker::is_boolean(expr_typespec))
        error_handler.flag(expr_token, INCOMPATIBLE TYPES, this);
```

## ExpressionParser.parseFactor()

Now an identifier can be more than just a variable name.

```
private ICodeNode parse_factor(Token token)
    throw (string)
    switch ((PascalTokenType) tokenType)
        case IDENTIFIER:
            return parse_identifier(token);
```



### ExpressionParser.parseIdentifier()

## Constant identifier

```
CONST

pi = 3.14159;
```

- Previously defined in a CONST definition.
- Create an INTEGER\_CONSTANT,REAL\_CONSTANT, or a STRING\_CONSTANT node.
- Set its VALUE attribute.

## Enumeration identifier

```
TYPE
    direction =
        (north, south,
        east, west);
```

- Previously defined in a type specification.
- Create an INTEGER\_CONSTANT node.
- Set its VALUE attribute.



### ExpressionParser.parseIdentifier()

- Variable identifier
  - Call method variableParser::parse().

