Jeffrey Cho

Professor Charles Palmer

COSC 55

September 27, 2020

Lab 1: Lab Report

Task 1 Frequency Analysis: Given the provided online resources (1.0), I have created an excel document (1.1) that highlights all the letter, bigram, trigram, and cipher text frequencies.

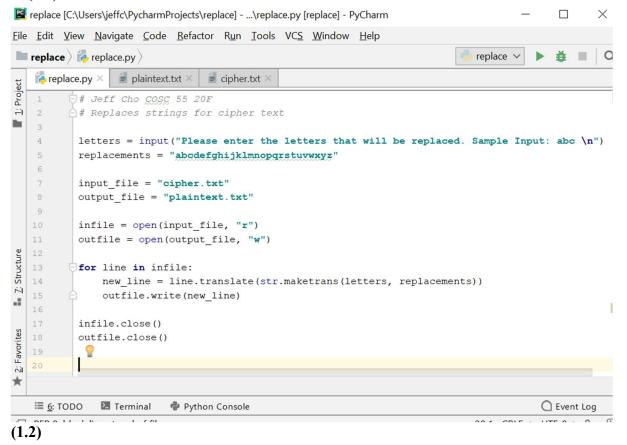
- http://www.richkni.co.uk/php/crypta/freq.php: This website can produce the statistics fro a ciphertext, including the single-letter frequencies, bigram frequencies (2-letter sequence), and trigram frequencies (3-letter sequence), etc.
- https://en.wikipedia.org/wiki/Frequency_analysis: This Wikipedia page provides frequencies for a typical English plaintext.
- https://en.wikipedia.org/wiki/Bigram: Bigram frequency.
- https://en.wikipedia.org/wiki/Trigram: Trigram frequency.

d				D		F G					N	0	P	Q	R		T
	traditional alphabet				subtrep		xt percentage	bigram freq		cipher text bigram freq				percentage		cipher text trigram	
2		12.7		9 (k	12.2	th	1.52	ok	78		the	1.81		hok	36
3		9.1	E) (u	h	10.5	he	1.28	ho	65		and	0.73		tok	25
4		8.2	C	i		P	8.2	in	0.94	ke	52		ing	0.72		oke	24
5		7.5	C	d c	С	q	8	er	0.94	wm	40		ent	0.42		wmr	21
6		7	e		k	0	6.9	an	0.82	wh	37		ion	0.42		qmc	15
7	n	6.7	f	t	Ь	W	6.4	re	0.68	hp	36		for	0.34		ake	13
8		6.3	<u>c</u>	3 г	r	m	6.3	nd	0.63	to	29		nde	0.34		zqt	13
9	h	6.1	r	1 0	0	t	5.9	at	0.59	qh	25		has	0.34		how	11
10	г	6	i		w	е	5.4	on	0.57	me	23		nce	0.34		who	10
11	d	4.3	i	r	n	С	4.4	nt	0.56	mr	23		edt	0.34		pzm	9
12	I	4	k	: f	f	×	4.2	ha	0.56	qm	23		tis	0.34		ppf	9
13	С	2.8	- 1)	X	z	3.1	es	0.56	qt	23		tha	0.33		plh	9
14	u	2.8	г	n i		1	2.7	st	0.55	kx	22		tio	0.31		keq	9
15	m	2.4	r	1	m	ь	2.3	en	0.55	km	20		oft	0.22		wik	9
16	W	2.4	C) [D	i	1.9	ed	0.53	oq	20		sth	0.21		kxx	9
17		2.2) 8	s	u	1.9	to	0.52	hw	20		men	0.21		owm	9
18	a	2		۱ ۱	v	г	1.9	it	0.5	kc	20						
19		2				s	1.8	ou	0.5	zq	20						
20		1.9		t t		ф	1.7	ea	0.47	pl pl	19						
21		1.5		i		ī	1.4	hi	0.46	xx	19						
22		1		ı İ		f	1.3	is	0.46	kz	18						
23		0.8			a	a	0.9	or	0.43	kh	18						
24		0.15		v 2		n	0.1	ti	0.34	ak	18						
25		0.15				v	0	as	0.33	hq	17						
26		0.1		, ,		y	Ö	te	0.27	gx	17						
27		0.1				q	0	et	0.19	th	17						
28	-	0.01	-		9	9		ng	0.18	De	17						
29								of	0.16	mh	17						
30								al	0.09	kt	16						
31								de	0.03	pm	16						
32								se	0.03	ck	16						
33								le .	0.08	ob	16						
34									0.06	ek	16						
34 35								sa .:	0.05		16						
20								si	0.05	eq Ih	15						
36 37								ar	0.04		15						
37								ve		qe				-			
38								га	0.04	kq	14						
39								ld	0.02	wi	14						
40								ur	0.02	xw	13						
41																	

(1.1)

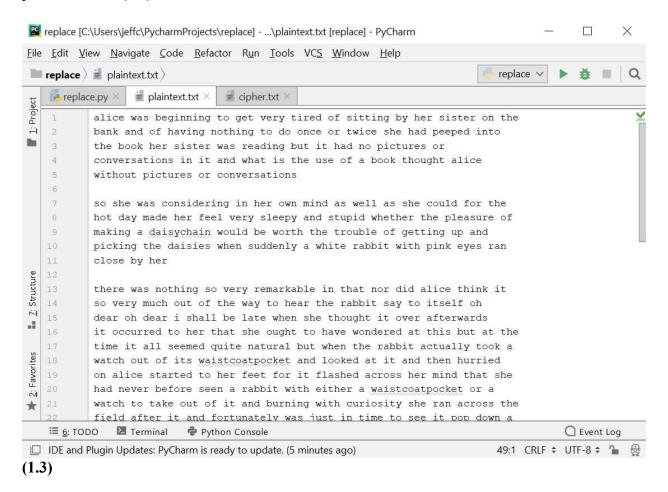
The process to figure out the plaintext and encryption key involved numerous testing. However, to create starting data, that will eventually be modified, the frequencies of letters in the cipher text and the traditional alphabet were compared and sorted in descending order next to each

other. From there, the corresponding frequencies of letters in the traditional alphabet and the cipher text were determined, which can be seen in **col 1** and **col 3** in **(1.1)**. The testing portion would have to include trial and error of replacing such letters, and to minimize the time taken for testing, a python testing script has been written to make the testing much easier, which is shown in **(1.2)**.



The python script simply takes in the input as a string concatenated together with no spaces between. It also assumes that the first letter inputted is the sub (substitution) for a and so on. Then the python program takes in the cipher.txt as input and creates a new output file which tests the inputted encryption key. The encryption key is not fully complete of course until there are no spelling errors in the output file and the story makes sense. Soon after testing, the encryption key was recovered, which is in **col. 2** in **(1.1)**. The **'key'** part is the traditional alphabet, and the

'sub/rep' is what the cipher text replaced the traditional alphabet letters with. The plaintext is pictured below (1.3).



Task 2 Encryption using Different Ciphers and Modes: First in order to try out different ciphers, a plain.txt file in the VM was created. The contents strictly in the file are "hello my name is Jeff. I am in cs 55." The file is strictly 38 bytes. The three ciphers used were "aes-128-cbc, aes-128-cfb, aes-128-ctr", which can be seen in the command line arguments (2.0).

```
[09/26/20]seed@VM:~$ openssl enc -aes-128-cfb -e -in pl ain.txt -out output_cfb.txt -K 00112233445566778889aabb ccddeeff -iv 0102030405060708 (2.0) [09/26/20]seed@VM:~$ openssl enc -aes-128-cbc -e -in pl ain.txt -out output_cbc.txt -K 00112233445566778889aabb ccddeeff -iv 0102030405060708 [09/26/20]seed@VM:~$ openssl enc -aes-128-ctr -e -in pl ain.txt -out output_ctr.txt -K 00112233445566778889aabb ccddeeff -iv 0102030405060708
```

Because of how small the file is, there was not noticeable time difference in speed for the commands. However, the encrypted contents were noticeably different, which can be seen in (2.1).

```
[09/26/20]seed@VM:~$ cat output_cbc.txt
`$\tilde{\partial} \tilde{\partial} \tilde{\partia
```

There is a clear difference in regard to spacing in addition to the wide variety of symbols used to encrypt the file. However, one other noticeable different was the sizes of the resulting encryptions. The file sizes respectively for "output_cbc.txt, output_cfb.txt, and output_ctr.txt" are 48 bytes, 38 bytes, and 38 bytes. For some reason it seems as though the aes-128-cbc cipher specifically has a different encrypted file size for its content.

Task 3 Encryption Mode – ECB vs. CBC:

1) In order to encrypt the pic_original.bmp file, the following commands were used. The first command depicts 128-bit AES with CBC mode (3.0), and the second command depicts 128-bit AES with ECB mode (3.1).

```
[09/26/20]seed@VM:~/lab 1$ openssl enc -aes-128-cbc -e
-in pic_original.bmp -out pic_cbc.bmp -K 00112233445566
778889aabbccddeeff -iv 0102030405060708
(3.0)
```

```
[09/26/20]seed@VM:~/lab 1$ openssl enc -aes-128-ecb -e
-in pic_original.bmp -out pic_ecb.bmp -K 00112233445566
778889aabbccddeeff
```

(3.1)

As you can see the **CBC** encrypted file was saved as **pic_cbc.bmp** file and the **ECB** encrypted file was saved as **pic_ecb.bmp** file. Using the following commands, the first 54 bytes of the encrypted pictures were replaced with the 54-byte header of the original picture (3.2).

```
[09/27/20]seed@VM:~/lab 1$ head -c 54 pic_original.bmp
> header
[09/27/20]seed@VM:~/lab 1$ tail -c +55 pic_cbc.bmp > bo
dy
[09/27/20]seed@VM:~/lab 1$ cat header body > new_file.b
mp
(3.2)
```

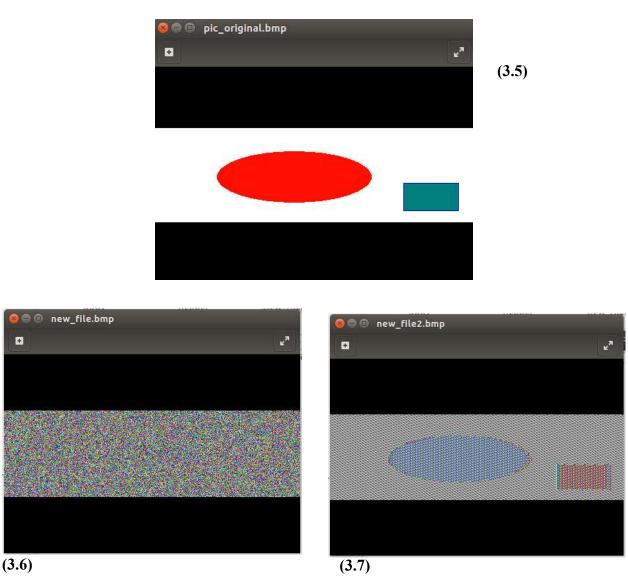
Here the first 54 bytes of the original picture was used as the header while the remaining bytes starting with 55 from the encrypted picture was used as the body. This led to the creation of the new_file.bmp, which is composed of the header and body.

2) The encrypted photos are depicted in (3.3) and (3.4). Here it is quite clear that there is no information that can be gained from the photos when comparing to the original photo in (3.5).



The only information we have for both the CBC and ECB encrypted photos is that there are bogus headers. However, when viewing the encrypted photos where the first 54 bytes from the original picture (3.5) replaced the bogus header, the ECB encrypted photo begins to appear (3.7).

However, in the CBC encrypted photo (3.6), there is still no clear indication of what the photo is; however, we can see a somewhat blurry version of the original picture in the ECB encrypted photo. Based on (3.7), we can get a basic shape and short preview of what the original picture could look like; however, for (3.6), it is still quite unclear as to what the original picture is.



3) Repeat using another picture

Another .bmp image was encrypted, and the image can be seen in (3.8). The terminal command lines are similar as seen in (3.0) and (3.1) with the only difference being image.bmp and not "pic_original.bmp."

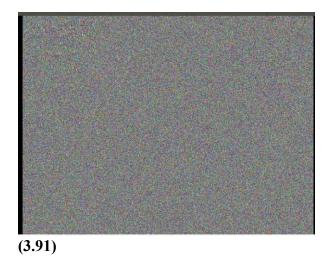


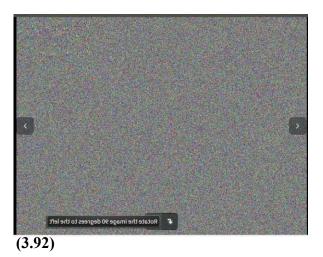
The similar methodology was used to encrypt the images with the first 54 bytes from the "**image.bmp**" and the remaining bytes from the encrypted images. We can see the command line arguments below **(3.9).** They are exactly the same as used previously.

```
[09/27/20]seed@VM:~/lab 1$ head -c 54 image.bmp > heade
r1
[09/27/20]seed@VM:~/lab 1$ tail -c +55 image_cbc.bmp >
body1
[09/27/20]seed@VM:~/lab 1$ cat header1 body1 > new_imag
e.bmp
(3.9)
```

Below, we have the encrypted images (with the byte modifications). Unlike the previous trial with the "pic_original.bmp", here we see that even the "image_ecb_encrypted.bmp" is not as revealing as before for "pic_original.bmp" (3.92). However, similar to the cbc image for

"pic_original.bmp," the "image_cbc_encrypted.bmp" also shares the same encrypted screen where we get little to no information about the original image (3.91).





Task 4 Padding: Before encrypting the files to report for padding, a file was created. The file "pad.txt" was created and its contents are "hello world today is a wonderful day with lots of sunshine." The size of the file is 58 bytes.

1) For pad.txt, **128-bit AES with CBC, CFB, OFB, and ECB modes** were used. We can see the terminal commands in **(4.0)**. For example, on output files, if **ECB** mode

```
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -e -in pa
d.txt -out pad_cbc.txt -K 00112233445566778899aabbccdde
eff -iv 1234567890123457
[09/27/20]seed@VM:~$ openssl enc -aes-128-cfb -e -in pa
d.txt -out pad_cfb.txt -K 00112233445566778899aabbccdde
eff -iv 1234567890123457
[09/27/20]seed@VM:~$ openssl enc -aes-128-ofb -e -in pa
d.txt -out pad_ofb.txt -K 00112233445566778899aabbccdde
eff -iv 1234567890123457
[09/27/20]seed@VM:~$ openssl enc -aes-128-ecb -e -in pa
d.txt -out pad_ecb.txt -K 00112233445566778899aabbccdde
eff
(4.0)
```

was used, then the output file would be pad_ecb.txt. This goes for all the modes that were tested. After each output file was created, the sizes for the resulting documents for pad_cbc.txt, pad_cfb.txt, pad_ofb.txt, and pad_ecb.txt were 64 bytes, 58 bytes, 58 bytes, and 64 bytes respectively. Through the differences of the size for the original file (58 bytes), it is quite noticeable that the CBC and ECB modes do indeed use padding in contrast to the CFB and OFB encryption modes.

2) Three files were created. The files for the 5 bytes, 10 bytes, and 16 bytes are "five.txt, ten.txt, and sixteen.txt" respectively and were created through terminal command lines as shown in (4.1).

```
[09/27/20]seed@VM:~$ echo -n "hello" > five.txt
[09/27/20]seed@VM:~$ echo -n "hello worl" > ten.txt
[09/27/20]seed@VM:~$ echo -n "hello world what" > sixte
en.txt
(4.1)
```

The three files were then encrypted using 128-bit AES with CBC mode through the terminal command lines as shown in (4.2).

```
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -e -in fi
ve.txt -out five_out.txt -K 00112233445566778899aabbccd
deeff -iv 1234567890123457
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -e -in te
n.txt -out ten_out.txt -K 00112233445566778899aabbccdde
eff -iv 1234567890123457
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -e -in si
xteen.txt -out sixteen_out.txt -K 00112233445566778899a
abbccddeeff -iv 1234567890123457
(4.2)
```

The sizes of the resulting encryption files for "five_out.txt, ten_out.txt, and sixteen_out.txt" are 16 bytes, 16 bytes, and 32 bytes respectively. This provides evidence for padding in the CBC mode. Now the files are to be decrypted while using the "-nopad" option, which will not remove padded data in the decrypted data. We can see how this is achieved through the terminal command lines as shown in (4.3).

```
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -d -in fi
ve out.txt -out five decrypt.txt -K 0011223344556677889
9aabbccddeeff -iv 1234567890123457
[09/27/20]seed@VM:~$ hd five decrypt.txt
00000000
         68 65 6c 6c 6f
     |hello|
00000005
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -d -nopad
 -in five out.txt -out five decrypt nopad.txt -K 001122
33445566778899aabbccddeeff -iv 1234567890123457
[09/27/20]seed@VM:~$ hd five decrypt nopad.txt
00000000 68 65 6c 6c 6f 0b 0b 0b 0b 0b 0b 0b 0b 0b
 0b
     |hello......|
00000010
(4.3)
```

We can see the padding in the "five_decrypt_nopad.txt" when comparing it with "five_decrypt.txt." The difference in hex is the padding which is "0b 0b 0b" as shown in (4.3). We can also see the padding when decrypting "ten.txt" shown in (4.4).

We can see the padding in the "five decrypt nopad.txt" when comparing it with

"five_decrypt.txt." The difference in hex is the padding which is "0b 0b 0b" as

shown in (4.3). We can also see the padding when decrypting "ten.txt" shown in (4.4).

```
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -d -in te
n out.txt -out ten decrypt.txt -K 00112233445566778899a
abbccddeeff -iv 1234567890123457
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -d -nopad
 -in ten out.txt -out ten decrypt nopad.txt -K 00112233
445566778899aabbccddeeff -iv 1234567890123457
[09/27/20]seed@VM:~$ hd ten decrypt.txt
00000000
         68 65 6c 6c 6f 20 77 6f 72 6c
     Thello worll
0000000a
[09/27/20]seed@VM:~$ hd ten decrypt nopad.txt
         68 65 6c 6c 6f 20 77 6f 72 6c 06 06 06 06 06
    |hello worl....|
00000010
 (4.4)
```

It is quite clear that for "ten_decrypt_nopad.txt" that the padding is hex "06 06 ... 06" as shown in (4.4). For "sixteen_out.txt" we also see hex padding in

"sixteen_decrypt_nopad.txt," as shown in (4.5) along with terminal command lines.

```
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -d -in si
xteen out.txt -out sixteen decrypt.txt -K 0011223344556
6778899aabbccddeeff -iv 1234567890123457
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -d -nopad
 -in sixteen out.txt -out sixteen decrypt nopad.txt -K
00112233445566778899aabbccddeeff -iv 1234567890123457
[09/27/20]seed@VM:~$ hd sixteen decrypt.txt
        68 65 6c 6c 6f 20 77 6f 72 6c 64 20 77 68 61
0000000
74
    [hello world what]
00000010
[09/27/20]seed@VM:~$ hd sixteen decrypt nopad.txt
00000000 68 65 6c 6c 6f 20 77 6f 72 6c 64 20 77 68 61
    |hello world what|
        00000010
10
00000020
 (4.5)
```

We see that the hex padding is "10 10 10 ... 10" when we compare

"sixteen_decrypt_nopad.txt" with "sixteen_decrypt.txt."

Task 5 Error Propagation – Corrupted Cipher Text: A text file of at least 1000 bytes was created and it was named "baby.txt." The contents of this file are ironically enough the song lyrics to the hit song by Justin Bieber "Baby."

1. The text file creation in the terminal using command lines is shown in (5.0).

```
[09/27/20]seed@VM:~$ echo -n "Ooh whoa, ooh whoa
> You know you love me, I know you care
> Just shout whenever and I'll be there
> You are my love, you are my heart
> And we will never, ever, ever be apart
> Are we an item? Girl quit playin'
> We're just friends, what are you sayin'
> Yeah, yeah, yeah (now I'm all gone)
> Yeah, yeah, yeah
> Yeah, yeah, yeah (now I'm all gone)
> Yeah, yeah, yeah
> Yeah, yeah, yeah
> Yeah, yeah, yeah (now I'm all gone)
> Gone, gone, gone, I'm gone" > baby.txt
(5.0)
```

2. Using the commands shown in (5.1), the file "baby.txt" was encrypted using 128-bit

AES with ECB, CBC, CFB, and OFB modes, and their respective files were created.

```
[09/27/20]seed@VM:~$ openssl enc -aes-128-cbc -e -in ba by.txt -out baby_cbc.txt -K 00112233445566778899aabbccd deeff -iv 1234567890123457 [09/27/20]seed@VM:~$ openssl enc -aes-128-ofb -e -in ba by.txt -out baby_ofb.txt -K 00112233445566778899aabbccd deeff -iv 1234567890123457 [09/27/20]seed@VM:~$ openssl enc -aes-128-cfb -e -in ba by.txt -out baby_cfb.txt -K 00112233445566778899aabbccd deeff -iv 1234567890123457 [09/27/20]seed@VM:~$ openssl enc -aes-128-ecb -e -in ba by.txt -out baby_ecb.txt -K 00112233445566778899aabbccd deeff
```

- 3. Use the bless hex editor to simulate the corruption ("flip") of a single bit of the 55th byte in each encrypted file. Flip the same bit in each file. Using bless hex editor to simulate corruption, a single bit of the 55th byte in each encrypted file was "flipped." For "baby_cbc.txt" the 8A hex was changed to 8B. For "baby_cfb.txt" the 37 hex was changed to 36 hex. In "baby_ofb.txt" the 1F hex was changed to 1E hex, and lastly, in "baby_ecb.txt" the F8 hex was changed to F9.
- 4. Using the same key and IV when encrypting the files, the files were decrypted after they were "corrupted." The terminal commands can be seen in (5.2).

We now can take a look at how the corrupted files changed after they were decrypted as shown in (5.3, 5.4, 5.5, and 5.6).

```
[09/27/20]seed@VM:~$ cat baby_ecb_decrypt.txt

Ooh whoa, ooh whoa, ooh whoa

You know you love m@@@@@@

zT@[@7@Ere

Just shout whenever and I'll be there

(5.3)
```

We can see here how the **ECB** mode specifically decrypted the corrupted document (5.3). There is clearly a chunk of "baby.txt" that is still not decrypted, which can be explained

by the corrupt "flipping." However, we know that the corruption affects the while block the corrupted bit is in, meaning that if the whole block is unable to be decrypted, then by the result we can see that the whole block will appear decrypted even if some parts of the block may not be "corrupted."

```
[09/27/20]seed@VM:~$ cat baby_cbc_decrypt.txt

Ooh whoa, ooh whoa

You know you love m/���]7 t}��ſre

Jusu shout whenever and I'll be there

(5.4)
```

Here, we see how the **CBC** mode specifically decrypted the corrupted file **(5.4)**. There seems to be some plaintext missing, which is quite similar to how the ECB mode decryption turned out. **CBC** mode also involves block ciphering, and since the corrupted bit was in a block, it is quite like the **ECB** in that the block cipher is not wholly decrypted because of the "*flipped*" bit.

```
[09/27/20]seed@VM:~$ cat baby_cfb_decrypt.txt

0oh whoa, ooh whoa, ooh whoa

10u know you love me, I koow you calo@@@mWW00

@enever and I'll be there

You are my love, you are my heart

(5.5)
```

We can see how the **CFB** mode specifically decrypts the file with the "*flipped*" bit as shown above **(5.5).** However, unlike the **CBC** or even the **ECB** there are some words that are not correct and are simply out of place and missing. **CFB** mode however is still quite similar to **CBC**, but **CFB** mode is quite particular in that one corrupted bit could throw off the whole block decryption, which is evident when "*know*" is spelled incorrectly and parts of the verse "*Just shout whenever*" are still encrypted.

```
[09/27/20]seed@VM:~$ cat baby_ofb_decrypt.txt
Ooh whoa, ooh whoa
You know you love me, I koow you care
Just shout whenever and I'll be there
(5.6)
```

Based on the decryption using **OFB** mode, even with the single corrupted bit, we see that there are minimal errors here **(5.6)**. The only error lies in the 55th byte, which is where we "flipped" the bit. This show how **OFB** functions with errors, which is that it should be able to continue to correctly decrypt the rest of the block(s) even with the one bit "flipped" in the corrupted file.

Task 6 Initial Vector (IV) and Common Mistakes:

Task 6.1: A file "testing.txt" was created that contains one world "Yes."

A. We can see below the terminal commands to encrypt "testing.txt" using the iv:

0102030405060708, which printed out the cipher text underneath the command (6.0).

B. Below, we see the terminal commands to encrypt "testing.txt" using a different iv:

1020304050607080, which prints out a different cipher text seen in (6.1).

C. We can see below the terminal commands to encrypt "testing.txt" using the iv used in part A. Using iv: 0102030405060708, we see that the cipher text is identical to the cipher text from part A (6.2).

Overall, based on parts A, B, and C, we can see that using the same IV or even a predictable IV can make it easier to figure out the cipher text and the message behind it. For example, if someone knows plaintext 1, and for some reason a plaintext 2 exists that the person does not know the message. It is quite easy to figure out the message if the same IV is used and the cipher texts are exactly the same. Therefore, not using a unique IV would make it much easier to possibly decrypt the cypher text as opposed to a unique IV, which we see in part B creates a whole different cypher text than the ones from part A and C.

Task 6.2 Common Mistake: Use the Same IV: If an attacker gets hold of a plaintext (P1) and the corresponding ciphertext (C1), he/she should be able to decrypt other encrypted messages if the IV is always the same. For example: Given the following information (6.3), we can figure out P2.

```
Plaintext (P1): This is a known message!
Ciphertext (C1): a469b1c502c1cab966965e50425438e1bb1b5f9037a4c159

Plaintext (P2): (unknown to you)
Ciphertext (C2): bf73bcd3509299d566c35b5d450337e1bb175f903fafc159

(6.3)
```

To get the IV, we can XOR the plaintext 1 and cipher text 1. (P1 XOR C1). \rightarrow We get IV to be f001d8b622a8b99907b6353e2d2356c1d67e2ce356c3a478 in hex. We then can use this IV and

XOR it with C2 to get the value of the message in P2. We then can get in ASCII for P2 the value to be *bf73bcd3509299d566c35b5d450337e1bb175f903fafc159*. If that is converted to text, we can get the message P2, which is quite alarming: "*Order: Launch a missile!*"

After "Yes" XOR IV1, we get the hex "68574039383b3a35343d3c3f3e39383b," which is then used to XOR with IV2. When we finish that computation, we find out that the resulting value is very similar to "Yes" in that the hex is "5965730d0d0d0d0d0d0d0d0d0d0d0d0d0d0d0d0d0c" with minor padding differences. Using the bless editor in VM, we edited a blank text file and proceeded to put in the hex value of "5965730d0d0d0d0d0d0d0d0d0d0d0d0d0d0d0d0d0c" that we found earlier (6.5). We do this so we can send this message to Bob who will then encrypt it for us.



We can see the terminal command line from Bob's encryption in **(6.6).** We also see that when we receive the ciphertext from Bob that it very much so resembles the ciphertext provided in C1, meaning that Bob's **P1** message was indeed "**Yes**."

Overall Thoughts: I was particularly interested and surprised by Task 6.2 and 6.3. Both tasks for me were quite aggravating, resulting in numerous trial and error attempts, but the overall aspect of being able to deduce and figure out messages given only so much information is something that I find hard to wrap my head around, and in the end those tasks were the most rewarding. However, the idea that I can find out information with only a certain amount of info is something that very much interesting and something I would never know or learn if not for these challenges. Even though these tasks are primarily exposition to the overall subject of Secret-Key Encryption, it makes me think of what I could actually do if I learned much more on the subject (not in a harmful or terrible way of course). It is quite amazing what can be done with only so much information and very specific information that makes me rethink all my saved passwords and secure messages/emails while also making me double check security measure and what I can do to not let some future issues arise.