

CS 642: Computer Security and Privacy

Cryptography [Intro]

Spring 2020

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Thanks to Franz Roesner, Dan Boneh, Dieter Gollmann, Dan Halperin, Yoshi Kohno, Ada Lerner, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

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There, I changed it

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Announcements

- You should've received an email from the mailing list with a link to course website – email me if you haven't received this (check spam)
- Office hours today: 1.30 to 2.30pm in my office CS 7387.
- Office hours Friday: Zijun Ma, 2 to 3pm, CS 4217

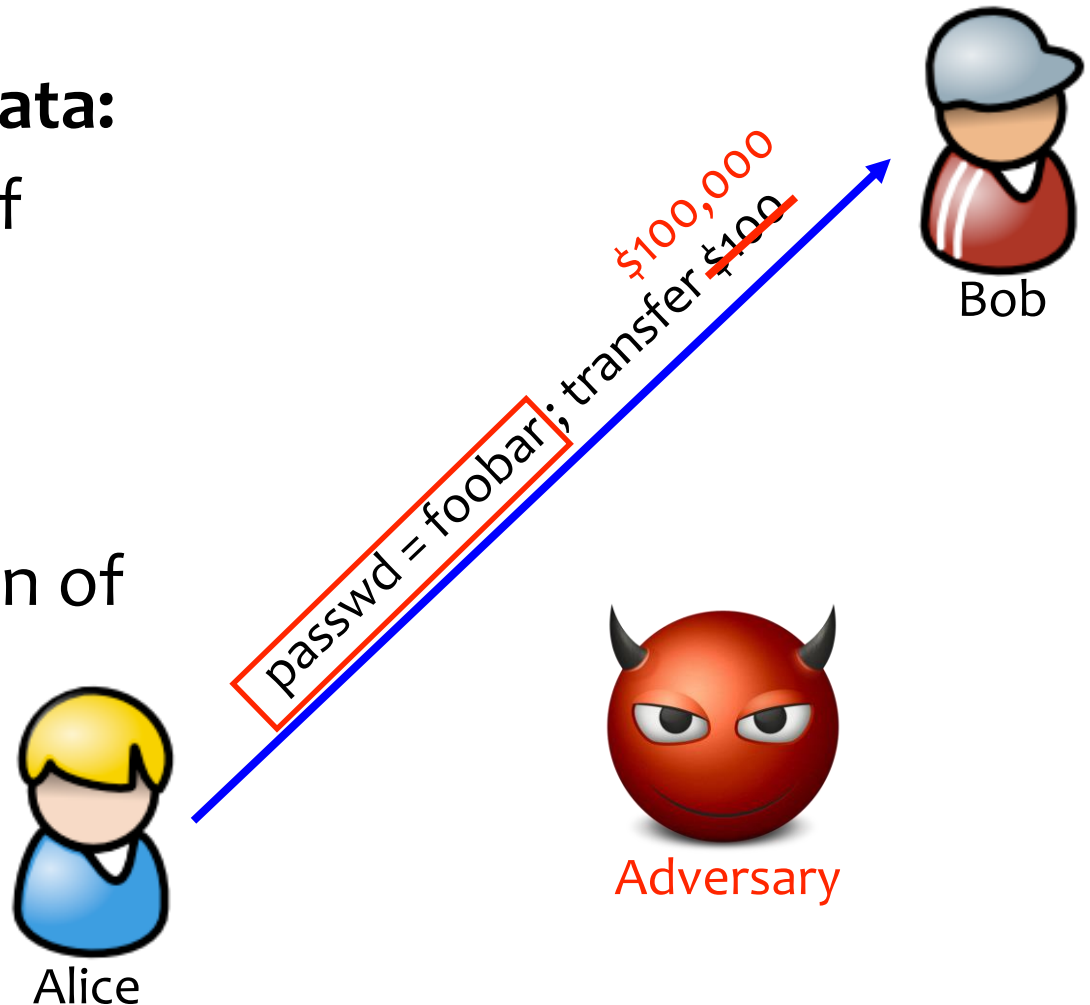
Common Communication Security Goals

Confidentiality of data:

Prevent exposure of information

Integrity of data:

Prevent modification of information



Recall Bigger Picture

- Cryptography only one small piece of a larger system
- Must protect entire system
 - Physical security
 - Operating system security
 - Network security
 - Users
 - Cryptography (following slides)
- Recall the weakest link
- Still, cryptography is a crucial part of our toolbox



Kerckhoff's Principle

- Security of a cryptographic object should depend only on the secrecy of the secret (private) key.
- Security should not depend on the secrecy of the algorithm itself.

Ingredient: Randomness

- Many applications (especially security ones) require randomness
- Explicit uses:
 - Generate secret cryptographic keys
 - Generate random initialization vectors for encryption
- Other “non-obvious” uses:
 - Generate passwords for new users
 - Shuffle the order of votes (in an electronic voting machine)
 - Shuffle cards (for an online gambling site)

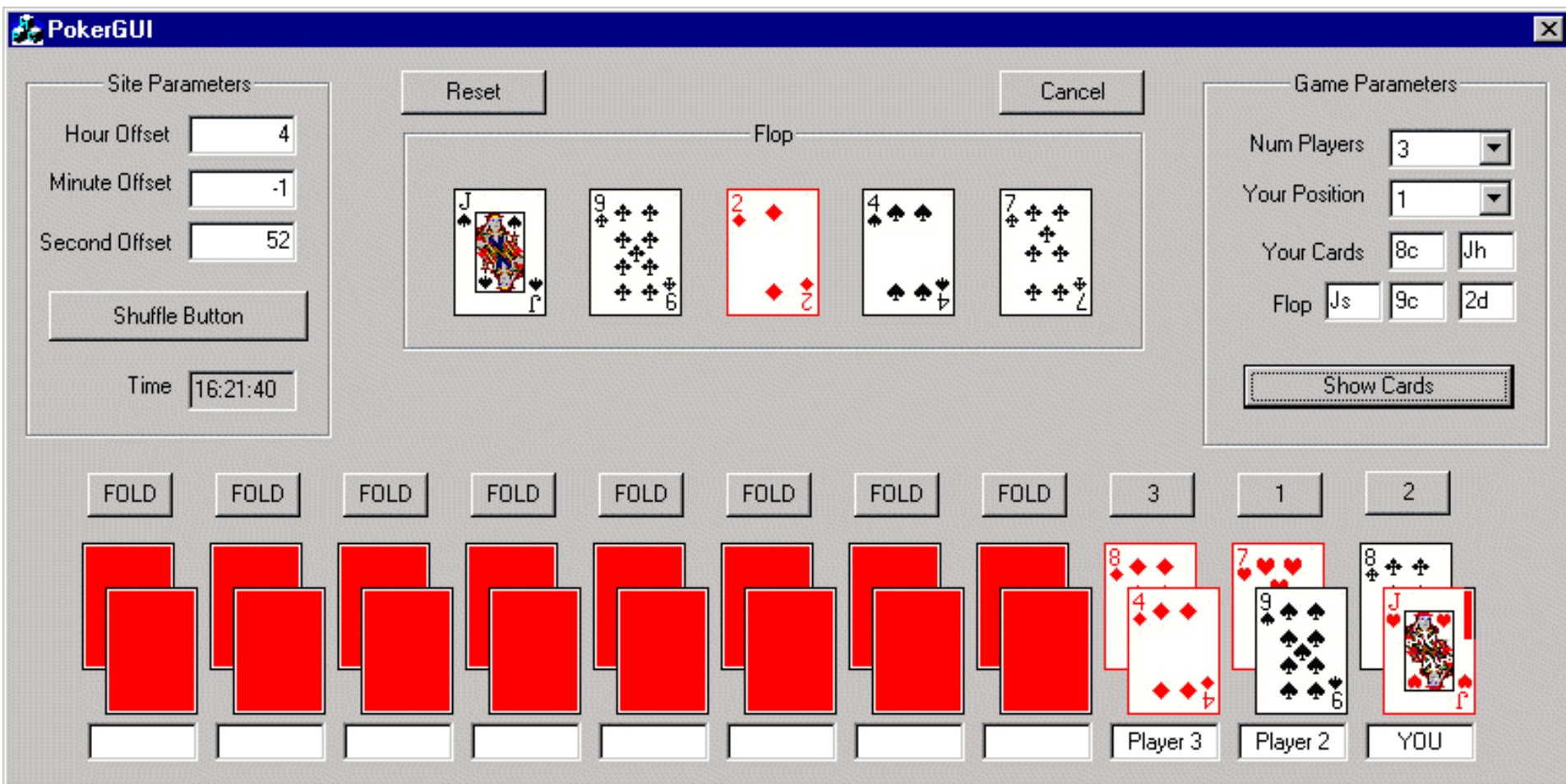
C's rand() Function

- C has a built-in random function: rand()

```
unsigned long int next = 1;
/* rand:  return pseudo-random integer on 0..32767 */
int rand(void) {
    next = next * 1103515245 + 12345;
    return (unsigned int)(next/65536) % 32768;
}
/* srand:  set seed for rand() */
void srand(unsigned int seed) {
    next = seed;
}
```

- Problem: don't use rand() for security-critical applications!
 - Given a few sample outputs, you can predict subsequent ones





More details: “How We Learned to Cheat at Online Poker: A Study in Software Security”
<https://www.developer.com/tech/article.php/616221/How-We-Learned-to-Cheat-at-Online-Poker-A-Study-in-Software-Security.htm>

PS3 and Randomness

Hackers obtain PS3 private cryptography key due to epic programming fail? (update)

<http://www.engadget.com/2010/12/29/hackers-obtain-ps3-private-cryptography-key-due-to-epic-programm/>

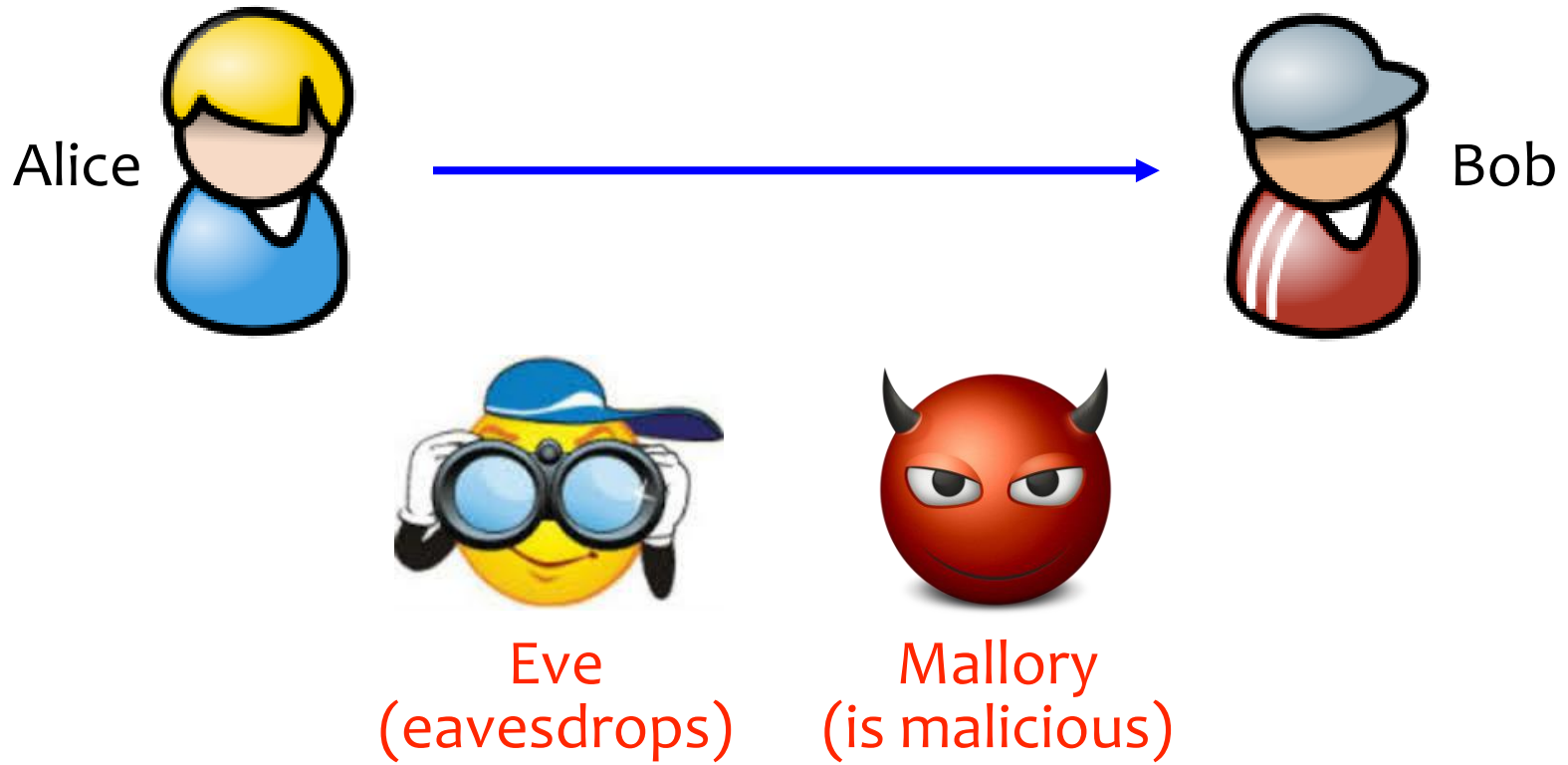
- 2010/2011: Hackers found/released private root key for Sony's PS3
- Key used to sign software – now can load any software on PS3 and it will execute as “trusted”
- Due to bad random number: same “random” value used to sign all system updates

Obtaining Pseudorandom Numbers

- For security applications, want “cryptographically secure pseudorandom numbers”
- Libraries include cryptographically secure pseudorandom number generators (CSPRNG)
- Linux:
 - /dev/random
 - /dev/urandom - nonblocking, possibly less entropy
- Internally:
 - Entropy pool gathered from multiple sources
 - e.g., mouse/keyboard timings
- Challenges with embedded systems, saved VMs

Alice and Bob

- Archetypal characters



Received April 4, 1977

A Method for Obtaining Digital Signatures and Public-Key Cryptosystems

R.L. Rivest, A. Shamir, and L. Adleman*

Abstract

An encryption method is presented with the novel property that publicly revealing an encryption key does not thereby reveal the corresponding decryption key. This has two important consequences:

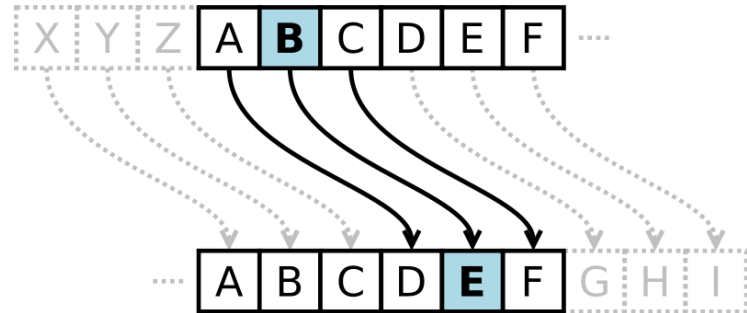
1. Couriers or other secure means are not needed to transmit keys, since a message can be enciphered using an encryption key publicly revealed by the intended recipient. Only he can decipher the message, since only he knows the corresponding decryption key.
2. A message can be "signed" using a privately held decryption key. Anyone can verify this signature using the corresponding publicly revealed encryption key. Signatures cannot be forged, and a signer cannot later deny the validity of his signature. This has obvious applications in "electronic mail" and "electronic funds transfer" systems.

History

- Substitution Ciphers
 - Caesar Cipher
- Transposition Ciphers
- Codebooks
- Machines
- Recommended Reading: **The Codebreakers** by David Kahn and **The Code Book** by Simon Singh.

History: Caesar Cipher (Shift Cipher)

- Plaintext letters are replaced with letters a fixed shift away in the alphabet.



- Example:

– Plaintext: The quick brown fox jumps over the lazy dog

– Key: Shift 3

ABCDEFGHIJKLMNOPQRSTUVWXYZ

DEFGHIJKLMNOPQRSTUVWXYZABC

– Ciphertext: WKHTX LFNEU RZQIR AMXPS VRYHU WKHOD CBGRJ

History: Caesar Cipher (Shift Cipher)

- ROT13: shift 13 (encryption and decryption are symmetric)
- What is the key space?
 - 26 possible shifts.
- How to attack shift ciphers?
 - Brute force.



History: Substitution Cipher

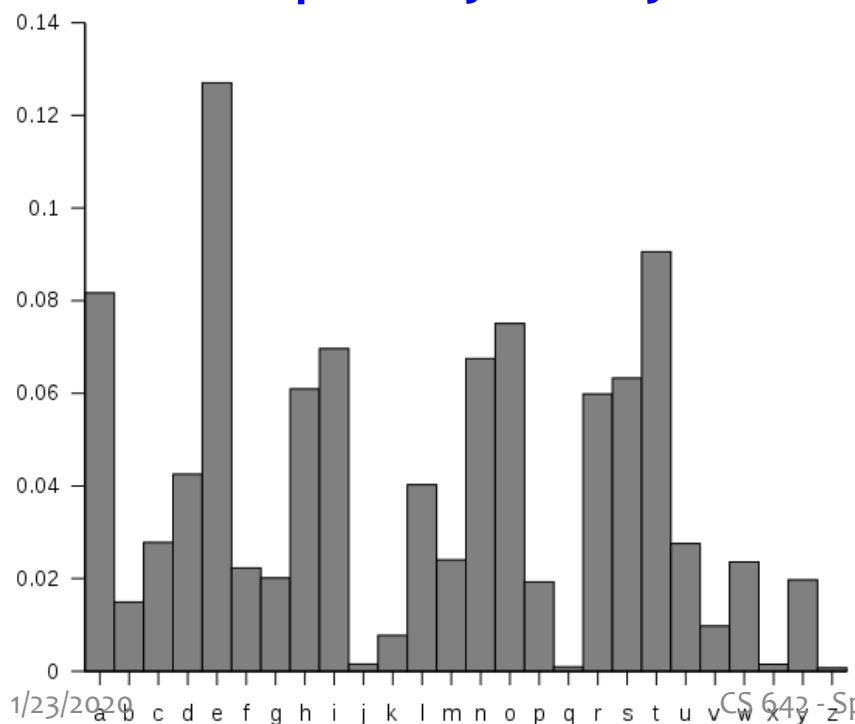
- Superset of shift ciphers: each letter is substituted for another one.
- Add a secret key
- Example:
 - Plaintext: ABCDEFGHIJKLMNOPQRSTUVWXYZ
 - Cipher: ZEBRAS CDEFGHIJKLMNOPQ TUVWXY
- “State of the art” for thousands of years

History: Substitution Cipher

- What is the key space? $26! \approx 2^{88}$

- How to attack?

— Frequency analysis.



Bigrams:

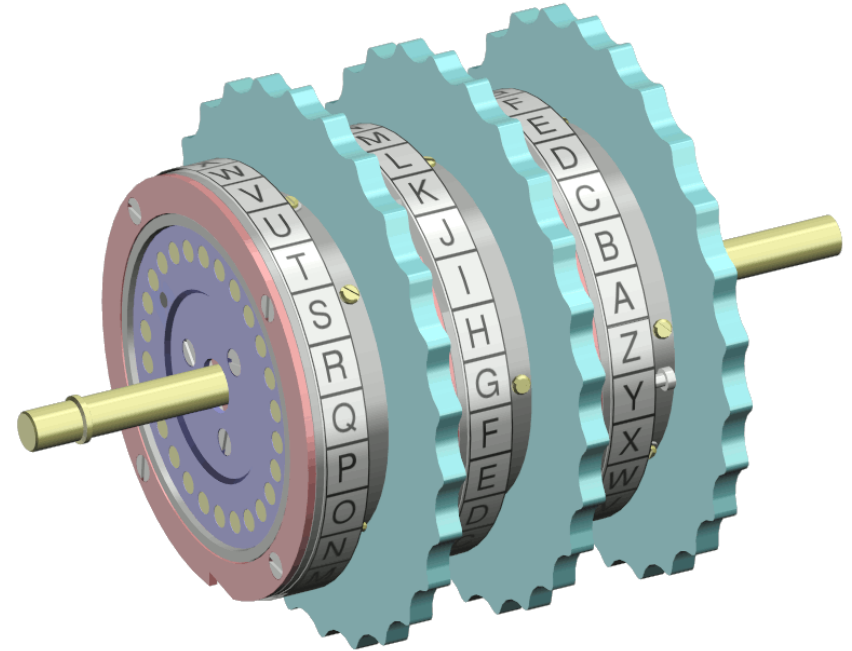
th 1.52%	en 0.55%	ng 0.18%
he 1.28%	ed 0.53%	of 0.16%
in 0.94%	to 0.52%	al 0.09%
er 0.94%	it 0.50%	de 0.09%
an 0.82%	ou 0.50%	se 0.08%
re 0.68%	ea 0.47%	le 0.08%
nd 0.63%	hi 0.46%	sa 0.06%
at 0.59%	is 0.46%	si 0.05%
on 0.57%	or 0.43%	ar 0.04%
nt 0.56%	ti 0.34%	ve 0.04%
ha 0.56%	as 0.33%	ra 0.04%
es 0.56%	te 0.27%	ld 0.02%
st 0.55%	et 0.19%	ur 0.02%

Trigrams:

1. the	6. ion	11. nce
2. and	7. tio	12. edt
3. tha	8. for	13. tis
4. ent	9. nde	14. oft
5. ing	10. has	15. sth

History: Enigma Machine

Uses rotors (substitution cipher) that change position after each key.



Key = initial setting of rotors

Key space?

26^n for n rotors

How Cryptosystems Work Today

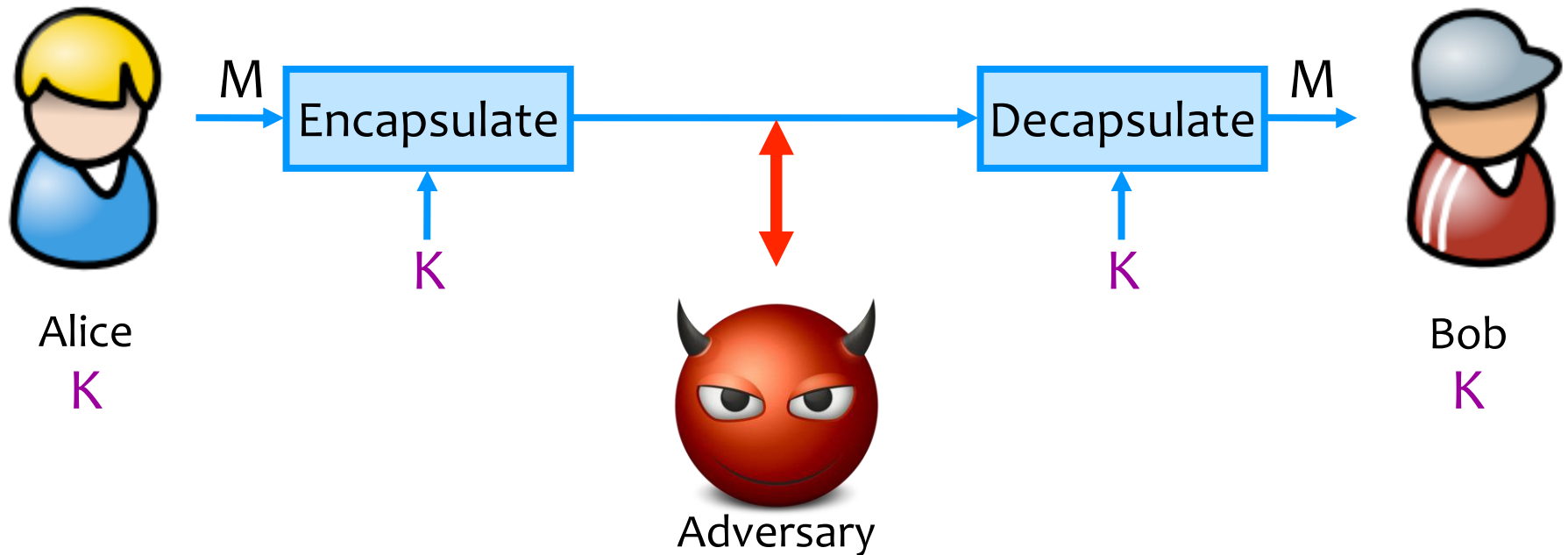
- Layered approach:
 - **Cryptographic primitives**, like block ciphers, stream ciphers, hash functions, and one-way trapdoor permutations
 - **Cryptographic protocols**, like CBC mode encryption, CTR mode encryption, HMAC message authentication
- Public algorithms (**Kerckhoff's Principle**)
- Security proofs based on assumptions (not this course)
- **Don't roll your own!**

Flavors of Cryptography

- Symmetric cryptography
 - Both communicating parties have access to a **shared random string K** , called the **key**.
- Asymmetric cryptography
 - Each party creates a public key **pk** and a secret key **sk** .

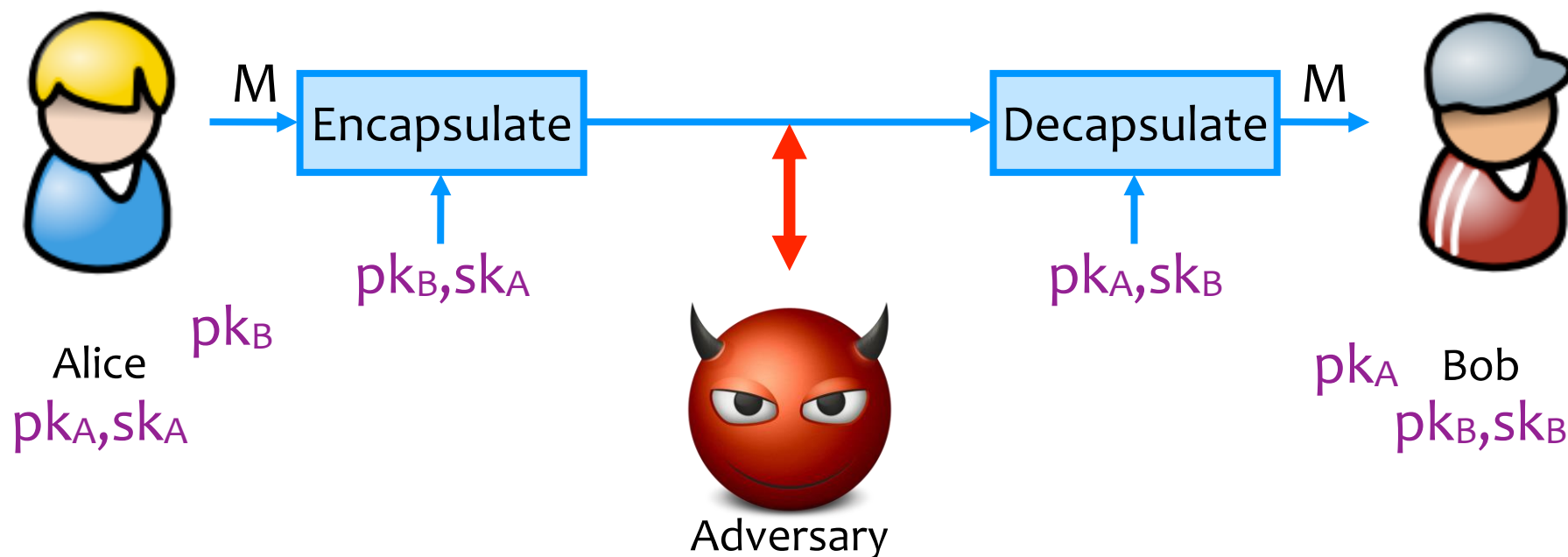
Symmetric Setting

Both communicating parties have access to a shared random string K , called the *key*.



Asymmetric Setting

Each party creates a public key pk and a secret key sk .



Flavors of Cryptography

- Symmetric cryptography
 - Both communicating parties have access to a **shared random string K** , called the **key**.
 - **Challenge: How do you privately share a key?**
- Asymmetric cryptography
 - Each party creates a public key **pk** and a secret key **sk** .
 - **Challenge: How do you validate a public key?**

Next Time

- Symmetric Encryption
 - One Time Pad
 - Block Ciphers
 - Modes of Operation