



北京邮电大学  
Beijing University of Posts and Telecommunications

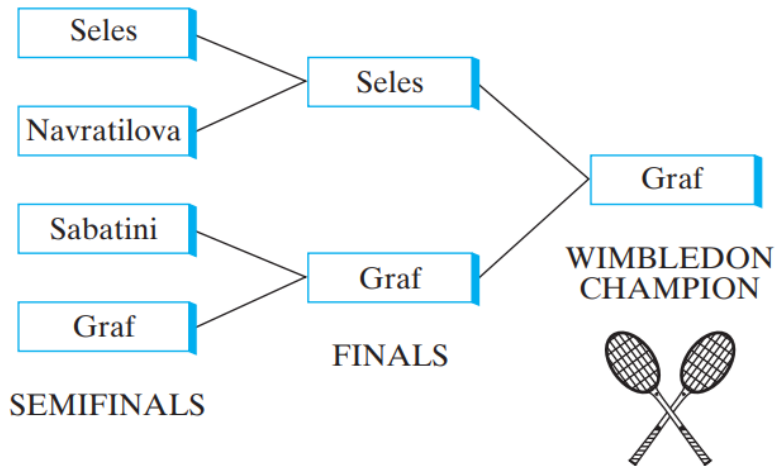
# Chapter 9 Trees 树

Lu Han

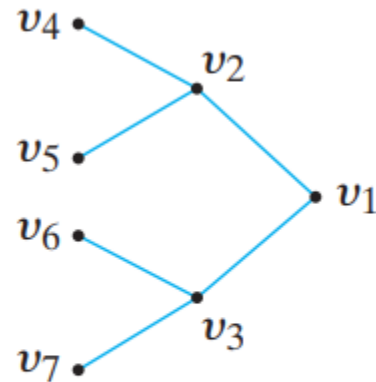
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## 9.1 Introduction 简介

Semifinals and finals of classic tennis competition Wimbledon.



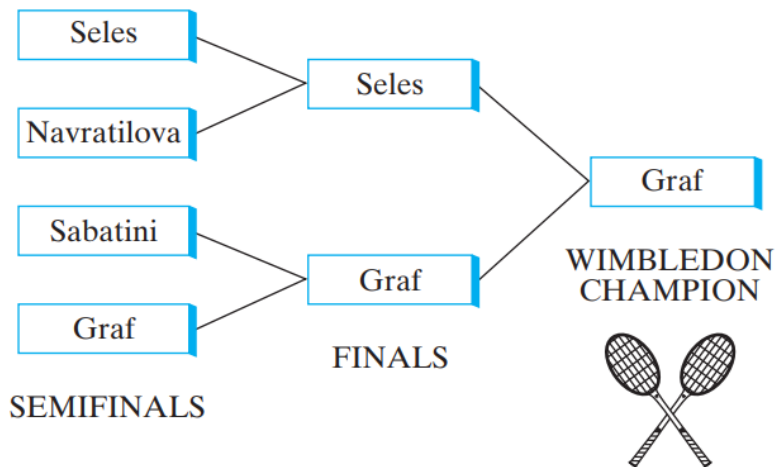
**Figure 9.1.1** Semifinals and finals at Wimbledon.



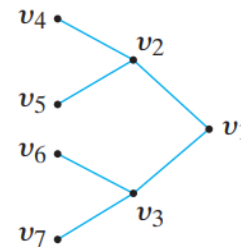
**Figure 9.1.2** The tournament of Figure 9.1.1 as a tree.

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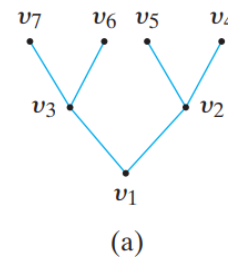
Semifinals and finals of classic tennis competition Wimbledon.



**Figure 9.1.1** Semifinals and finals at Wimbledon.



**Figure 9.1.2** The tournament of Figure 9.1.1 as a tree.



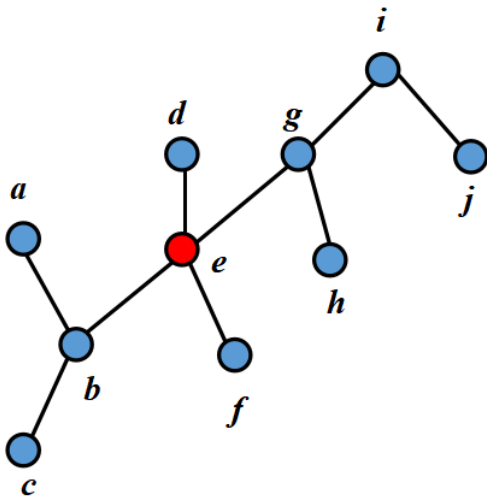
**Figure 9.1.3** The tree of Figure 9.1.2 rotated (a) compared with a natural tree (b).



## 9.1 Introduction 简介

**Definition 9.1.1** A **(free) tree (自由树)**  $T$  is a simple graph satisfying the following: If  $v$  and  $w$  are vertices in  $T$ , there is a unique simple path from  $v$  to  $w$ .

A **rooted tree (有根树)** is a tree in which a particular vertex is designated the root.



Designate  $e$  as the root in the tree.

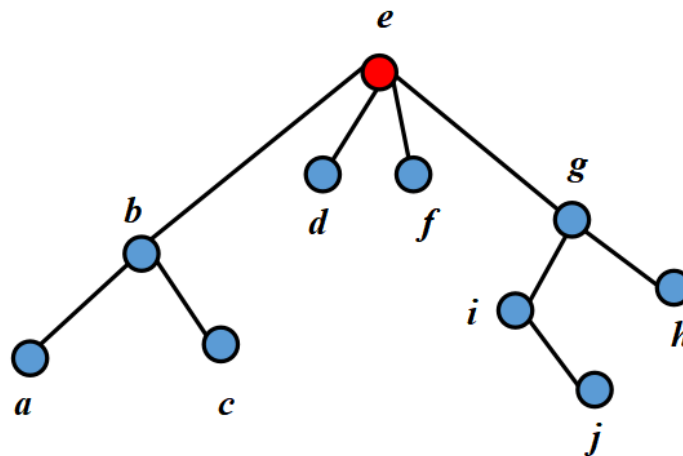
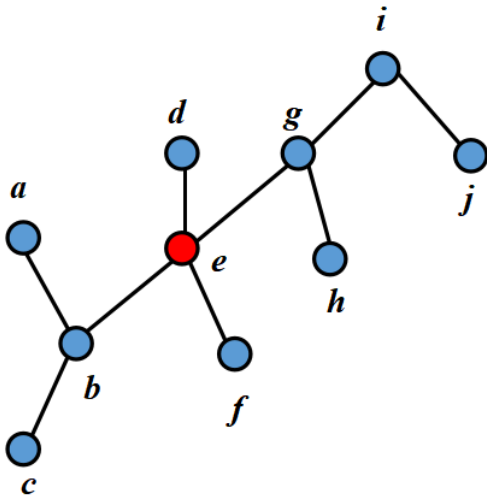


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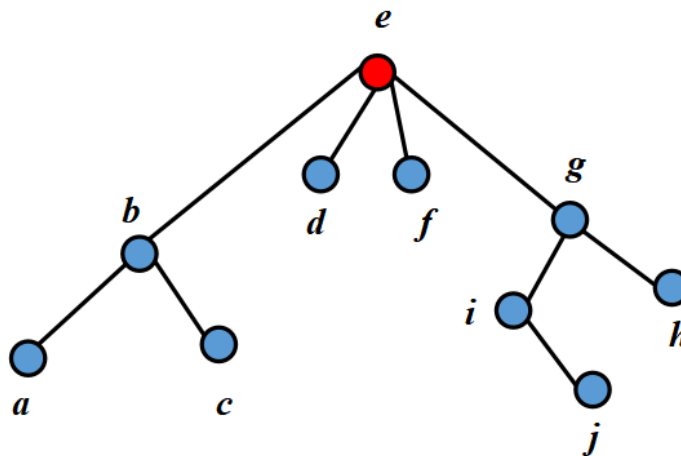
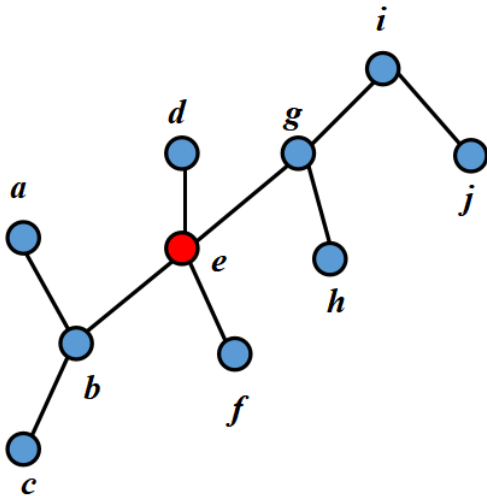


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- We call the level of the root level 0.
- The vertices under the root are said to be on level 1, and so on.



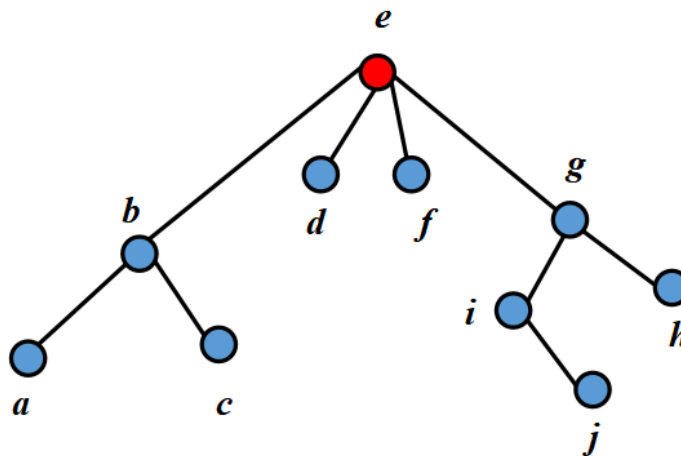
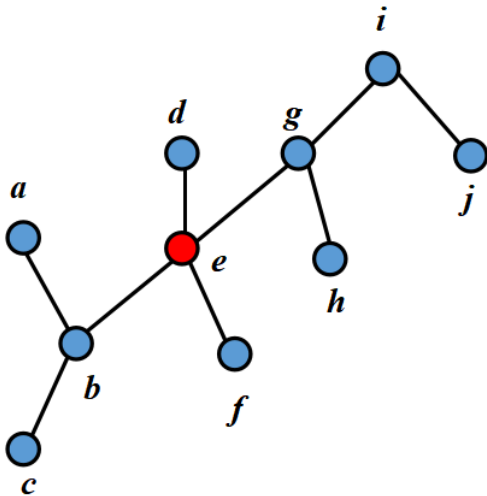
## 9.1 Introduction 简介

The **level of a vertex  $v$**  (顶点  $v$  所在的层次): the length of the simple path from the root to  $v$ .

The **height (高度)** of a rooted tree: the maximum level number that occurs.

A **rooted tree (有根树)** is a tree in which a particular vertex is designated the root.

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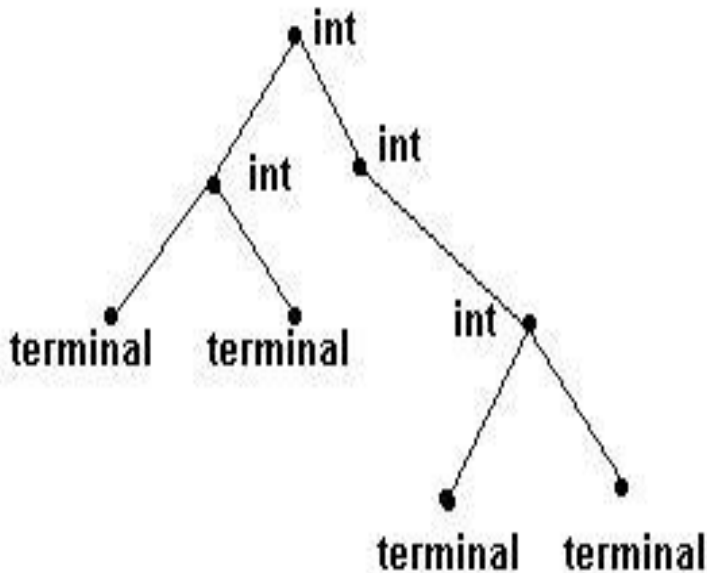
## 9.2 Terminology and Characterizations of Trees 树的术语和性质

**Definition 9.2.1** Let  $T$  be a tree with root  $v_0$ . Suppose that  $x, y$ , and  $z$  are vertices in  $T$  and that  $(v_0, v_1, \dots, v_n)$  is a simple path in  $T$ . Then

- (a)  $v_{n-1}$  is the **parent (父节点)** of  $v_n$ .
- (b)  $v_0, \dots, v_{n-1}$  are **ancestors (祖先节点)** of  $v_n$ .
- (c)  $v_n$  is a **child (子节点)** of  $v_{n-1}$ .
- (d) If  $x$  is an ancestor of  $y$ ,  $y$  is a **descendant (后代节点)** of  $x$ .
- (e) If  $x$  and  $y$  are children of  $z$ ,  $x$  and  $y$  are **siblings (兄弟节点)**.
- (f) If  $x$  has no children,  $x$  is a **terminal vertex (or a leaf) (终节点/叶节点)**.
- (g) If  $x$  is not a terminal vertex,  $x$  is an **internal (or branch) vertex (中间节点/枝节点)**.



## 9.2 Terminology and Characterizations of Trees 树的术语和性质



If  $x$  is not a terminal vertex,  $x$  is an **internal (or branch) vertex** (中间节点/枝节点).

- An **internal vertex** (中间节点) is a vertex that has at least one child.
- A **terminal vertex** (终节点) is a vertex that has no children



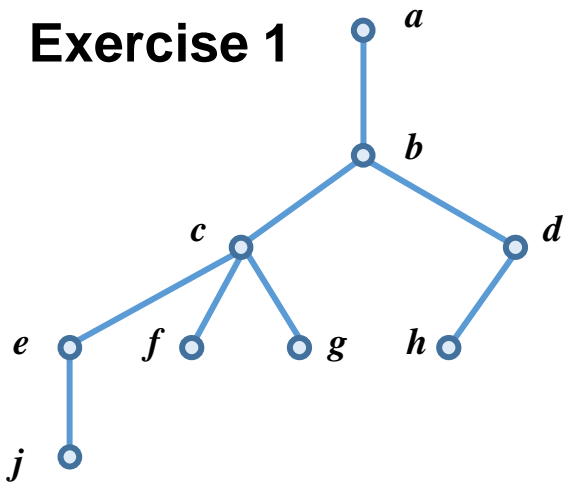
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(h) The **subtree (子树)** of  $T$  rooted at  $x$  is the graph with vertex set  $V$  and edge set  $E$ , where  $V$  is  $x$  together with the descendants of  $x$  and  $E = \{e \mid e \text{ is an edge on a simple path from } x \text{ to some vertex in } V\}$ .

### Exercise 1

Draw the subtree rooted at  $c$ .





## 9.2 Terminology and Characterizations of Trees 树的术语和性质

**Definition 9.1.1** A **(free) tree** (自由树)  $T$  is a simple graph satisfying the following: If  $v$  and  $w$  are vertices in  $T$ , there is a unique simple path from  $v$  to  $w$ .

A tree is connected.

A tree cannot contain a cycle.

A graph with no cycles is called an **acyclic graph** (非循环图).

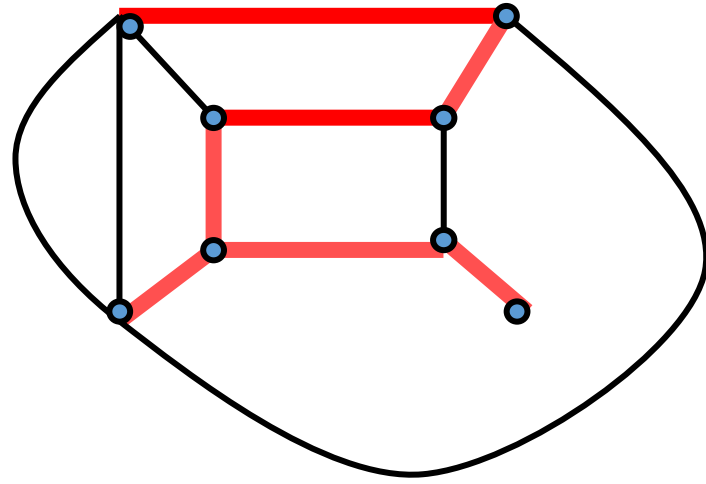
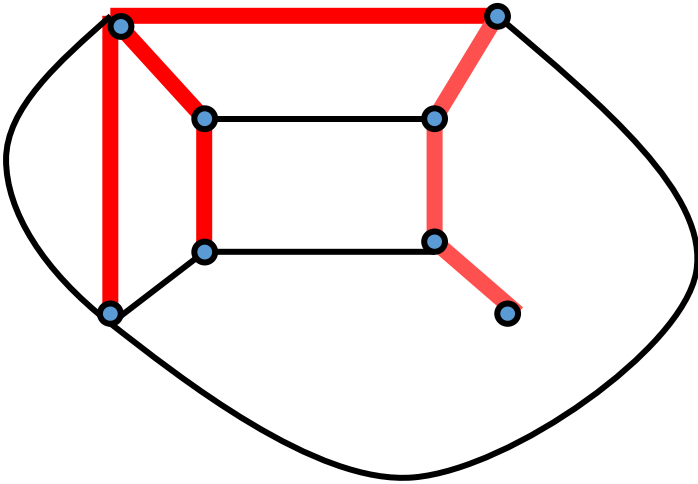
**Definition 9.2.3** Let  $T$  be a graph with  $n$  vertices. The following are equivalent.

- (a)  $T$  is a tree.
- (b)  $T$  is connected and acyclic.
- (c)  $T$  is connected and has  $n - 1$  edges.
- (d)  $T$  is acyclic and has  $n - 1$  edges.

## 9.3 Spanning Trees 生成树

**Definition 9.3.1** A tree  $T$  is a **spanning tree (生成树)** of a graph  $G$  if  $T$  is a subgraph of  $G$  that contains all of the vertices of  $G$ .

In general, a graph will have several spanning trees.





## 9.3 Spanning Trees 生成树

### Breadth-First Search 广度优先搜索

The idea of breadth-first search is to process all the vertices on a given level before moving to next-higher level.

### Depth-First Search 深度优先搜索 → Backtracking 回溯

The idea of depth-first search is to proceeds to successive levels in a tree at the earliest possible opportunity.



## 9.3 Spanning Trees 生成树

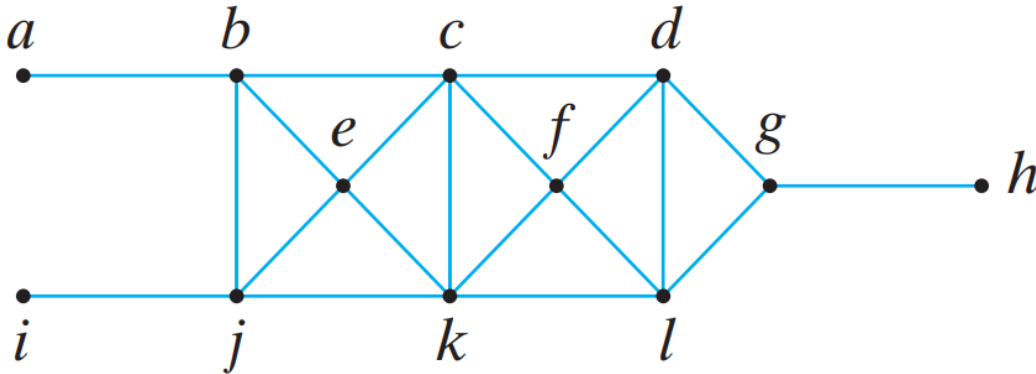
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**Exercise** Find a spanning tree for the graph using Breadth-First Search.





## 9.3 Spanning Trees 生成树

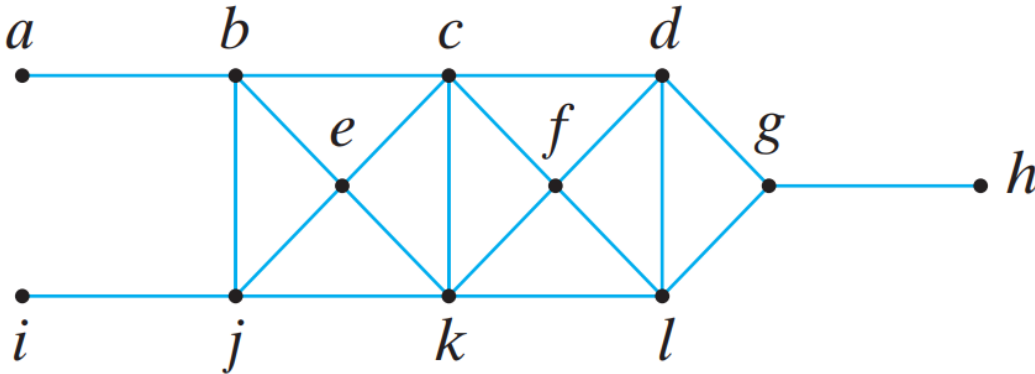
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### Depth-First Search 深度优先搜索 → Backtracking 回溯

The idea of depth-first search is to proceed to successive levels in a tree at the earliest possible opportunity.

**Exercise** Find a spanning tree for the graph using Depth-First Search.





## 9.4 Minimal Spanning Trees 最小生成树

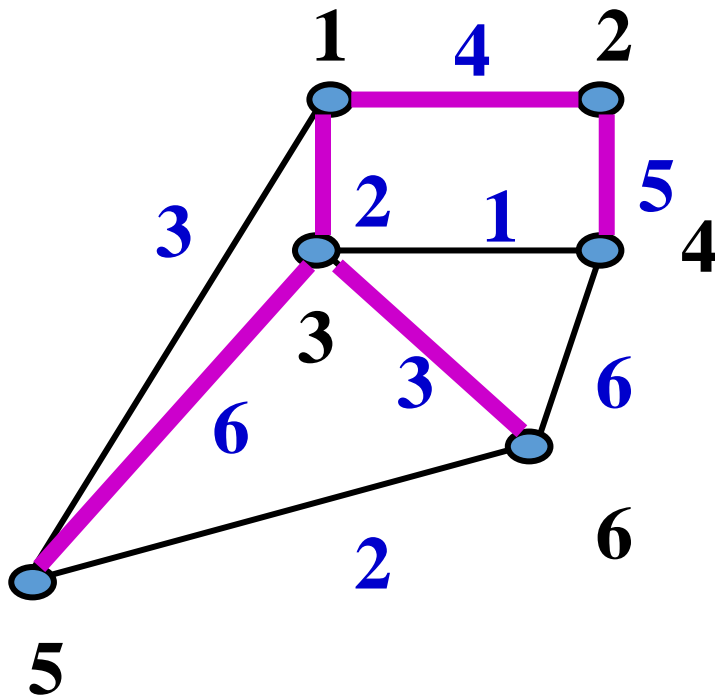
**Definition 9.4.1** Given a weighted graph  $G$ , a **minimal spanning tree** (最小生成树) of  $G$  is a spanning tree of  $G$  that has minimum weight.



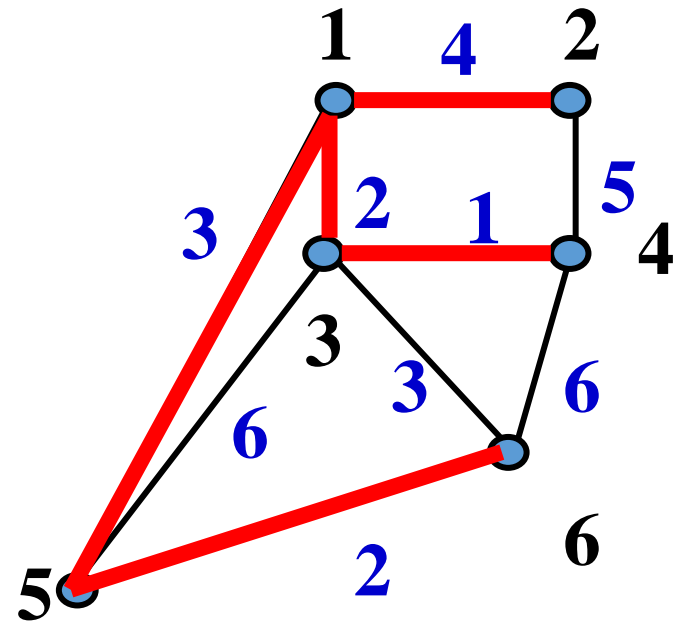


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Weight=20



Weight=12



## 9.4 Minimal Spanning Trees 最小生成树

**Definition 9.4.1** Given a weighted graph  $G$ , a **minimal spanning tree (最小生成树)** of  $G$  is a spanning tree of  $G$  that has minimum weight.

### **Prim's Algorithm** 普里姆算法

The algorithm begins with a single vertex. Then at each iteration, it adds to the current tree a minimum-weight edge that does not complete a cycle.

### **Kruskal's Algorithm** 克鲁斯卡尔算法

It is a greedy algorithm in graph theory as in each step it adds the next lowest-weight edge that will not form a cycle to the minimum spanning forest.



## 9.4 Minimal Spanning Trees 最小生成树

### Prim's Algorithm 普里姆算法

The algorithm begins with a single vertex. Then at each iteration, it adds to the current tree a minimum-weight edge that does not complete a cycle.

**Step 0:** Pick any vertex as a starting vertex (call it  $a$ ).  $T = \{a\}$ .

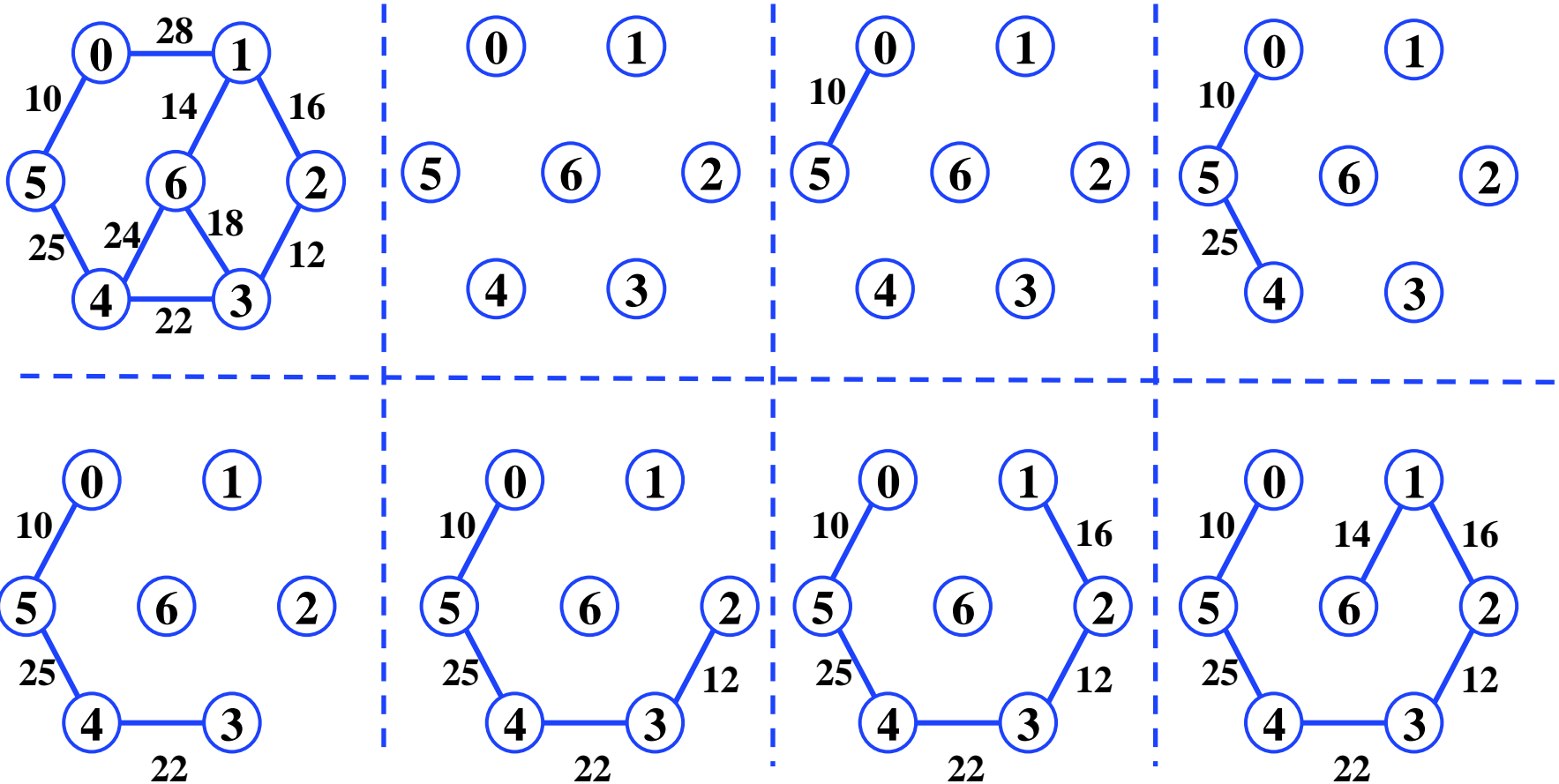
**Step 1:** Find the edge with smallest weight incident to  $a$ . Add it to  $T$ . Also include in  $T$  the next vertex and call it  $b$ .

**Step 2:** Find the edge of smallest weight incident to either  $a$  or  $b$ . Include in  $T$  that edge and the next incident vertex. Call that vertex  $c$ .

**Step 3:** Repeat Step 2, choosing the edge of smallest weight that does not form a cycle until all vertices are in  $T$ . The resulting subgraph  $T$  is a minimum spanning tree.



## Prim's Algorithm 普里姆算法





## 9.4 Minimal Spanning Trees 最小生成树

### Kruskal's Algorithm 克鲁斯卡尔算法

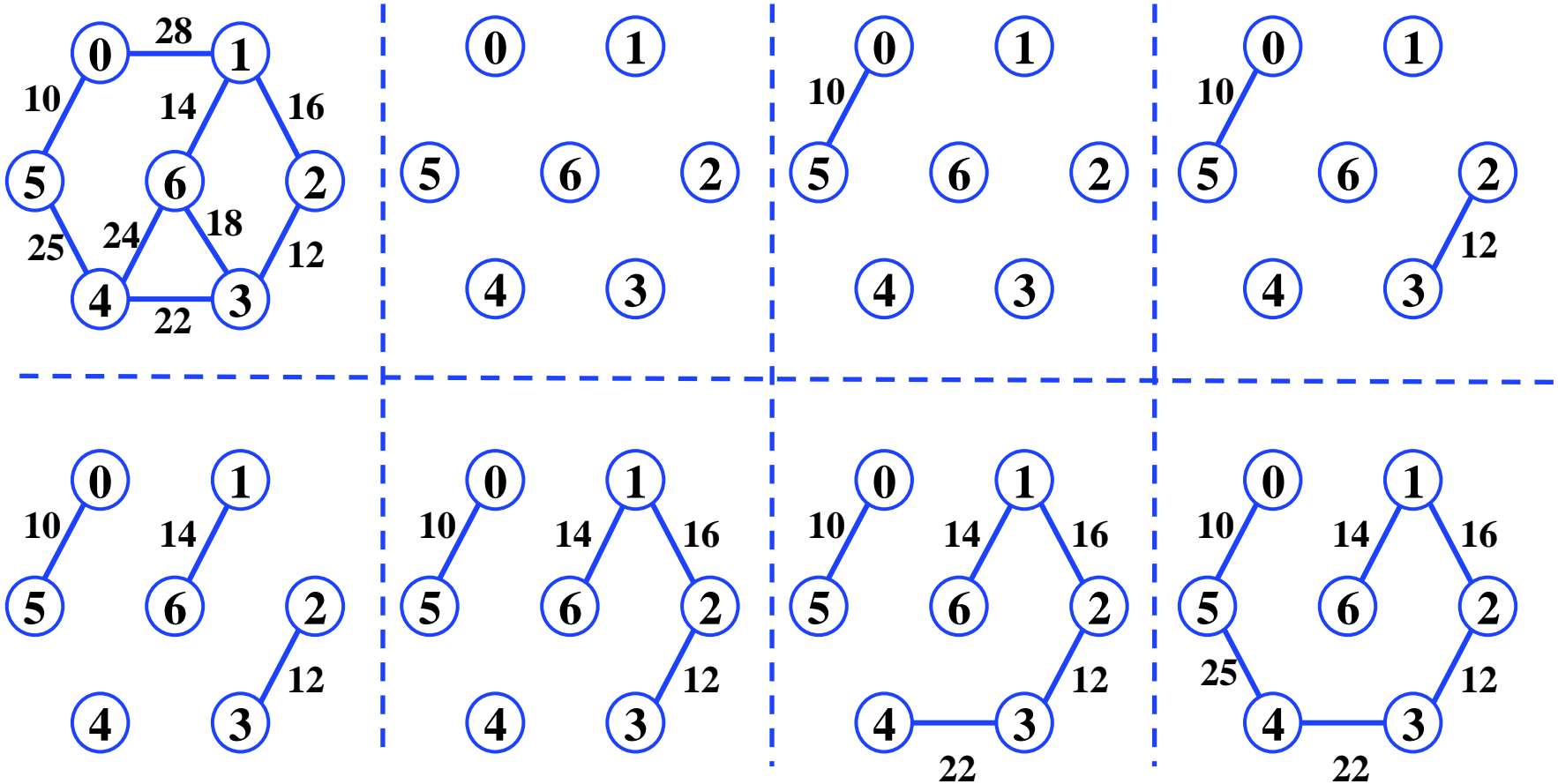
It is a greedy algorithm in graph theory as in each step it adds the next lowest-weight edge that will not form a cycle to the minimum spanning forest.

**Step 1:** Find the edge in the graph with smallest weight (if there is more than one, pick one at random).

**Step 2:** Find the next edge in the graph with smallest weight that doesn't close a cycle.

**Step 3:** Repeat Step 2 until you reach out to every vertex of the graph. The chosen edges form the desired minimum spanning tree.

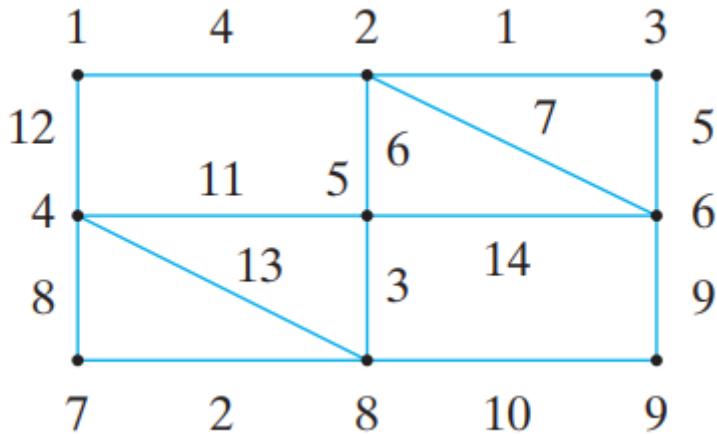
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