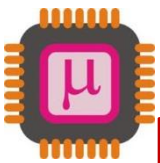


# Fall 2025/26– Lecture Notes # 3

- Introduction to Assembly Programming
- Program Segments



8086

# Introduction to Assembly Language Programming

- **ADD instruction**

**ADD** destination, source;    *dest = dest + source*

*mnemonic*

*operands*

**Example:**

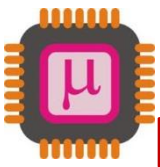
```
MOV    AL,24H    ;move 24H into  AL
MOV    DL,11H    ;move 11H into  DL
ADD    AL,DL      ;AL=AL+DL      (AL=35H)(DL =11H)
```

---

```
MOV    CH,24H    ;move 24H into  CH
MOV    BL,11H    ;move 11H into  BL
ADD    CH,BL      ;CH=CH+BL      (CH=35H)
```

---

```
MOV    CH,24H    ;load one operand into CH
ADD    CH,11H     ;add the second operand to CH (CH=35H)
```



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# Introduction to Assembly Language Programming

- **ADD instruction**

**ADD** destination, source;    *dest = dest + source*

*mnemonic*

*operands*

- ❖ If destination register is followed by an immediate data as the source, it is called the immediate operand.

```
MOV    CH,24H
ADD    CH,11H
```

- ❖ 8-bit registers can hold FFH (255) as the maximum value. Addition of larger numbers can be performed by the 16-bit nonsegment registers.

```
MOV    AX,34EH
MOV    DX,6A5H
ADD    DX,AX           ;DX=DX+AX  (DX=9F3H)
MOV    CX,34EH
ADD    CX,6A5H         ;CX=34EH+6A5=9F3H
```



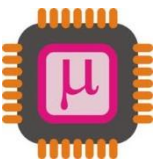
# Introduction to Program Segments

- **Segment**

- A **segment** is an area of memory that includes up to 64K bytes and begins an address evenly divisible by 16 (such an address ends in 0H).
- Assembly Language Program consists of three segments:
  - ❖ **code segment** : contains the program code (instructions)
  - ❖ **data segment** : used to store data to be processed by the program
  - ❖ **stack segment**: used to store information temporarily.

- **Logical and Physical Address**

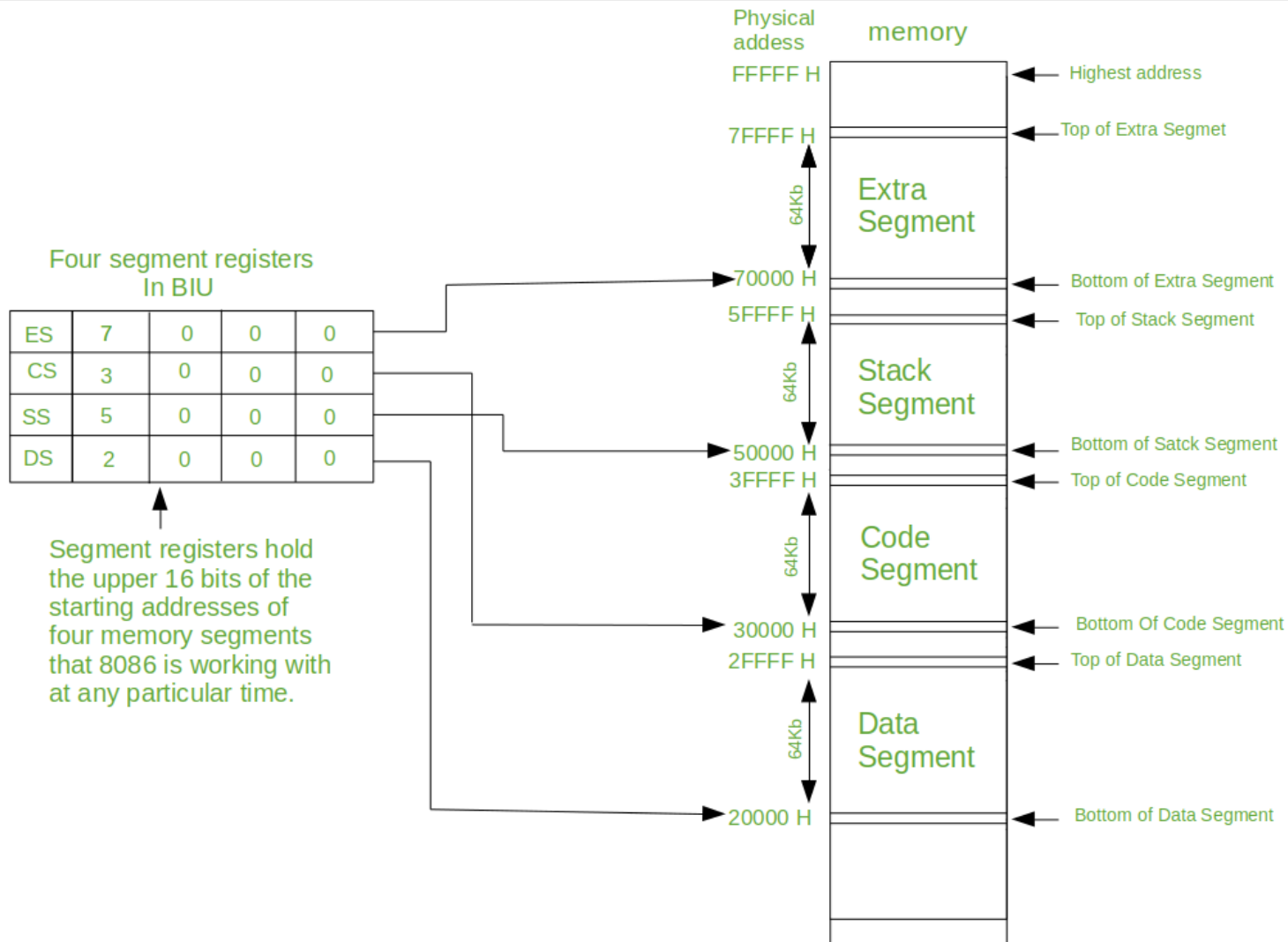
- **Physical Address** is the 20-bit address that actually put on the address bus. (in 8086)
  - ❖ Has a range of 00000H - FFFFFH
- **Segment Address** is a 16-bit address of the segment block.
  - Each segment is a **block** of 64 KB of memory space.
- **Offset Address** is a location **within** 64K byte segment range.
  - Has a range of 0000H - FFFFH
- **Logical Address** consists of **segment address** and **offset address**.

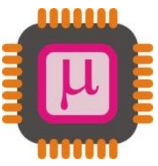


8086

# Introduction to Program Segments

- Memory Segmentation in 8086 Microprocessor**





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# Introduction to Program Segments

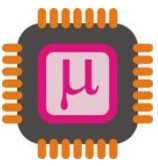
- Addressing in Code Segment
  - To execute a program, the 8086 fetches the instructions from the **code segment**.
  - The **logical address** of an instruction consists **CS** (Code Segment) and **IP**(instruction pointer).
  - **Logical Address in Code segment** is represented by using **segment address in CS** register and **Offset Address in IP** register as follows:

**CS:IP**

(16 bit CS and 16 bit IP making  
total of 32 bits)

**Example:** If CS register contains 2500H and IP register contains 95F3H. What is the **Logical Address** in the code segment?

CS:IP → **2500:95F3** (default in addressing is hex. You don't need H)



# Introduction to Program Segments

- Addressing in Code Segment

**Physical Address** is generated by shifting the CS one hex digit to the left and adding IP. Physical address is **20 bit address** which can be generated by using a logical address as follows.

1. Start with CS
2. Shift left CS (insert 0 as the Least significant digit)
3. Add IP

**Example:** If CS register contains 1980H and IP register contains 78FEH. What is the **Physical Address** in the code segment?

*Logical address:* CS:IP → **1980:78FE**

- |    |               |               |                               |
|----|---------------|---------------|-------------------------------|
| 1. | Start with CS | 1980          |                               |
| 2. | Shift left CS | 1980 <b>0</b> |                               |
| 3. | Add IP        | 78FE          | (1980 <b>0</b> + 78FE =210FE) |

**Physical address:** The microprocessor will retrieve the instruction from the memory locations starting from **210FE (20 bit address)**.



# Introduction to Program Segments

- Addressing in Code Segment

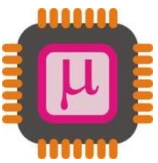
**Example:** If CS=24F6H and IP=634AH, determine:

- a) The logical address
- b) The offset address
- c) The physical address
- d) The lower range of the code segment
- e) The upper range of the code segment

**Solution:**

- a) The logical address is; **24F6:634A**
- b) The offset address is; **634A**
- c) The Physical address is; **24F60+634A= 2B2AA**
- d) The lower range of the code segment: **24F6:0000 → 24F60+0000 = 24F60**
- e) The upper range of the code segment: **24F6:FFFF → 24F60+FFFF = 34F5F**





# Introduction to Program Segments

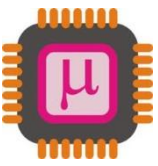
- Addressing in Data Segment
- The area of memory allocated strictly for data is called *data segment*.
- *Data segment* contains *variables* containing single values and arrays of values, where *code segment* only contain program *instructions*.
- *Logical Address in Data Segment* is represented by using *segment address in DS register* and *Offset Address in BX, SI or DI registers*.

DS:BX

DS:SI

DS:DI

- At any time *three locations* in the data segment are pointed with *DS:BX*, *DS:SI* and *DS:DI* respectively.



# Introduction to Program Segments

- Addressing in Data Segment

**Example:** If DS=7FA2H and the offset is 438EH, determine:

- a) The physical address
- b) The lower range of the data segment
- c) The upper range of the data segment
- d) Show the logical address

**Solution:**

- |                             |                        |
|-----------------------------|------------------------|
| a) The Physical address is; | $7FA20 + 438E = 83DAE$ |
| b) The lower range:         | $7FA20 + 0000 = 7FA20$ |
| c) The upper range:         | $7FA20 + FFFF = 8FA1F$ |
| d) The logical address is;  | $7FA2:438E$            |



8086

# Introduction to Program Segments

- Addressing in Data Segment

Why do we use data segment?

Assume that a program is needed to add 5 bytes of data (25H, 12H, 15H, 1FH and 2BH)

**One way:**

```
MOV AL,00H           ;initialize AL
ADD AL,25H
ADD AL,12H
ADD AL,15H
ADD AL,1FH
ADD AL,2BH           ; AL=25+12+15+1F+2B
```

code and data are mixed  
(bad programming practice)

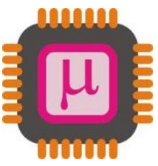
**Better way:** Assume that the Data segment contains the array of bytes starting from offset address 0200H.

```
MOV AL,0 ;clear AL
ADD AL,[0200];add the contents of DS:200 to AL
ADD AL,[0201];add the contents of DS:201 to AL
ADD AL,[0202];add the contents of DS:202 to AL
ADD AL,[0203];add the contents of DS:203 to AL
ADD AL,[0204];add the contents of DS:204 to AL
```

code and data are separated  
(good programming practice)

DS:01FF	?
DS:0200	25
DS:0201	12
DS:0202	15
DS:0203	1F
DS:0204	2B
DS:0205	?

Data Segment



8086

# Introduction to Program Segments

- Little endian convention**

Given 8-bit (1-byte) data, bytes are stored one after the other in the memory. However given 16-bit (2-bytes) of data how are data stored?

**Example:**            `MOV AX,35F3H`        ;load 35F3H into AX  
                      `MOV [1500],AX`       ; copy contents of AX to offset 1500H

In such a case the low byte goes to the low memory location and high byte goes to the high memory location.

DS:1500 = F3

DS:1501 = 35

This convention is called **little endian convention**: This convention is used by **Intel**.

**Big endian convention** is the opposite, where the high byte goes to the low address and low byte goes to the high address. **Motorola** microprocessor uses this convention.

